# SCIENTIFIC REPERTS

Received: 13 October 2016 Accepted: 11 November 2016 Published online: 23 December 2016

## **Lower and upper bounds for OPENentanglement of Rényi-***α* **entropy**

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**Entanglement Rényi-***α* **entropy is an entanglement measure. It reduces to the standard entanglement of formation when** *α* **tends to 1. We derive analytical lower and upper bounds for the entanglement Rényi-***α* **entropy of arbitrary dimensional bipartite quantum systems. We also demonstrate the application our bound for some concrete examples. Moreover, we establish the relation between entanglement Rényi-***α* **entropy and some other entanglement measures.**

Quantum entanglement is one the most remarkable features of quantum mechanics and is the key resource central to much of quantum information applications. For this reason, the characterization and quantifcation of entanglement has become an important problem in quantum-information science<sup>1</sup>. A number of entanglement measures have been proposed for bipartite states such as the entanglement of formation (EOF)<sup>2</sup>, concurrence<sup>[3](#page-8-2)</sup>, relative entropy<sup>[4](#page-8-3)</sup>, geometric entanglement<sup>5</sup>, negativity<sup>[6](#page-8-5)</sup> and squashed entanglement<sup>7[,8](#page-8-7)</sup>. Among them EOF is one of the most famous measures of entanglement. For a pure bipartite state  $|\psi\rangle_{AB}$  in the Hilbert space, the EOF is given by

$$
E_F(\left|\psi\right\rangle_{AB}) = S(\rho_A),\tag{1}
$$

where  $S(\rho_A) := -Tr \rho_A \log \rho_A$  is the von Neumann entropy of the reduced density operator of system *A*. Here "log" refers to the logarithm of base two. The situation for bipartite mixed states *ρ*<sub>*AB*</sub> is defined by the convex roof

$$
E_F(\rho_{AB}) = \min \sum_i p_i E_F(|\psi_i\rangle_{AB}),\tag{2}
$$

where the minimum is taken over all possible pure state decompositions of  $\rho_{AB} = \sum_i p_i |\psi_i\rangle_{AB} \langle \psi_i|$  with  $\sum_i p_i = 1$  and  $p_i$ >0. The EOF provides an upper bound on the rate at which maximally entangled states can be distilled from  $\rho$ and a lower bound on the rate at which maximally entangled states needed to prepare copies of *ρ*<sup>[9](#page-8-8)</sup>. For two-qubit systems, an elegant formula for EOF was derived by Wootters in ref. [3.](#page-8-2) However, for the general highly dimensional case, the evaluation of EOF remains a nontrivial task due to the the difficulties in minimization procedures<sup>10</sup>. At present, there are only a few analytic formulas for EOF including the isotropic states<sup>11</sup>, Werner states<sup>12</sup> and Gaussian states with certain symmetries<sup>13</sup>. In order to evaluate the entanglement measures, many efforts have also been devoted to the study of lower and upper bounds of different entanglement measures $14-32$  $14-32$ . Especially, Chen *et al*. [18](#page-8-15) derived an analytic lower bound of EOF for an arbitrary bipartite mixed state, which established a bridge between EOF and two strong separability criteria. Based on this idea, there are several improved lower and upper bounds for EOF presented in refs [33–](#page-8-16)[36](#page-8-17). While the entanglement of formation is the most common measure of entanglement, it is not the unique measure. Tere are other measures such as entanglement Rényi-*α* entropy (ΕRαΕ) which is the generalization of the entanglement of formation. The ΕRαΕ has a continuous spectrum parametrized by the non-negative real parameter *α*. For a bipartite pure state  $|\psi\rangle_{AB}$ , the ER*α*E is defined as<sup>37</sup>

$$
E_{\alpha}(|\psi\rangle_{AB}) := S_{\alpha}(\rho_A) := \frac{1}{1 - \alpha} \log(\text{tr}\rho_A^{\alpha}), \tag{3}
$$

where  $S_{\alpha}(\rho_A)$  is the Rényi- $\alpha$  entropy. Let  $\mu_1, \dots, \mu_m$  be the eigenvalues of the reduced density matrix  $\rho_A$  of  $|\psi\rangle_{AB}$ . We have

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$$
S_{\alpha}(\rho_A) = \frac{1}{1 - \alpha} \log \biggl( \sum_{i} \mu_i^{\alpha} \biggr) := H_{\alpha}(\vec{\mu}), \tag{4}
$$

where  $\vec{\mu}$  is called the Schmidt vector ( $\mu_1$ ,  $\mu_2$ ,  $\cdots$ ,  $\mu_m$ ). The Rényi- $\alpha$  entropy is additive on independent states and has found important applications in characterizing quantum phases with differing computational power<sup>38</sup>, ground state properties in many-body systems<sup>[39](#page-8-20)</sup>, and topologically ordered states<sup>40,41</sup>. Similar to the convex roof in (2), the ER $\alpha$ E of a bipartite mixed state  $\rho_{AB}$  is defined as

$$
E_{\alpha}(\rho_{AB}) = \min \sum_{i} p_i E_{\alpha}(|\psi_i\rangle_{AB}). \tag{5}
$$

It is known that the Rényi-*α* entropy converges to the von Neumann entropy when *α* tends to 1. So the ER*α*E reduces to the EOF when *α* tends to 1. Further ER*α*E is not increased under local operations and classical communications (LOCC)<sup>37</sup>. So the ER $\alpha$ E is an entanglement monontone, and becomes zero if and only if  $\rho_{AB}$  is a separable state.

An explicit expression of ER $\alpha$ E has been derived for two-qubit mixed state with  $\alpha \ge (\sqrt{7} - 1)/2 \simeq 0.823^{37/42}$ . Recently, Wang *et al.*[42](#page-8-23) further derived the analytical formula of ER*α*E for Werner states and isotropic states. However, the general analytical results of ER*α*E even for the two-qubit mixed state with arbitrary parameter *α* is still a challenging problem.

The aim of this paper is to provide computable lower and upper bounds for ERαE of arbitrary dimensional bipartite quantum systems, and these results might be utilized to investigate the monogamy relation<sup>[43–](#page-8-24)46</sup> in high-dimensional states. The key step of our work is to relate the lower or upper bounds with the concurrence which is relatively easier to dealt with. We also demonstrate the application of these bounds for some examples. Furthermore, we derive the relation of ER*α*E with some other entanglement measures.

**Lower and upper bounds for entanglement of Rényi-***α* **entropy.** For a bipartite pure state with Schmidt decomposition  $|\psi\rangle = \sum_{i=1}^{m} \sqrt{\mu_i} |ii\rangle$ , the concurrence of  $|\psi\rangle$  is given by  $c(\ket{\psi}) := \sqrt{2(1 - \text{Tr}\rho_A^2)} = \sqrt{2(1 - \sum_{i=1}^m \mu_i^2)}$ . The expression  $1 - \text{Tr}\rho_A^2$  is also known as the mixedness and linear entropy<sup>47,48</sup>. The concurrence of a bipartite mixed state  $\rho$  is defined by the convex roof  $c(\rho) = \min_{\rho} \sum_i \rho_i c(|\psi_i\rangle)$ for all possible pure state decompositions of  $\rho = \sum_i p_i |\psi_i\rangle \langle \psi_i|$ . A series of lower and upper bounds for concurrence have been obtained in refs [19](#page-8-28)[,24,](#page-8-29)[25.](#page-8-30) For example, Chen *et al*.<sup>[19](#page-8-28)</sup> provides a lower bound for the co making a connection with the known strong separability criteria<sup>[49](#page-8-31),[50](#page-8-32)</sup>, i.e.,

$$
c(\rho) \ge \sqrt{\frac{2}{m(m-2)}} (\max(||\rho^{T_A}||, ||\mathcal{R}(\rho)||) - 1),
$$
\n(6)

for any  $m \otimes n$ ( $m \le n$ ) mixed quantum system. The  $\|\cdot\|$  denotes the trace norm and  $T_A$  denotes the partial transpose. Another important bound of squared concurrence used in our work is given by refs [24,](#page-8-29)[25.](#page-8-30)

$$
\operatorname{Tr}(\rho \otimes \rho V_i) \leq [C(\rho)]^2 \leq \operatorname{Tr}(\rho \otimes \rho K_i),\tag{7}
$$

 $\text{with}\, V_1=4(P_-^{(1)}-P_+^{(1)})\otimes P_-^{(2)}\!, V_2=4P_-^{(1)}\otimes (P_-^{(2)}-P_+^{(2)})\!, K_1=4(P_-^{(1)}\otimes I^{(2)})\!, K_2=4(I^{(1)}\otimes P_-^{(2)})\text{and}\, P_-^{(i)}\!(P_+^{(i)})$ is the projector on the antisymmetric (symmetric) subspace of the two copies of the *i*th system. These bounds can be directly measured and can also be written as

$$
\operatorname{Tr}(\rho \otimes \rho V_1) = 2(\operatorname{Tr}\rho^2 - \operatorname{Tr}\rho_A^2),\tag{8}
$$

$$
\operatorname{Tr}(\rho \otimes \rho V_2) = 2(\operatorname{Tr}\rho^2 - \operatorname{Tr}\rho_B^2),\tag{9}
$$

$$
\operatorname{Tr}(\rho \otimes \rho K_1) = 2(1 - \operatorname{Tr}\rho_A^2),\tag{10}
$$

$$
\operatorname{Tr}(\rho \otimes \rho K_2) = 2(1 - \operatorname{Tr}\rho_B^2). \tag{11}
$$

Below we shall derive the lower and upper bounds of ER*α*E based on these existing bounds of concurrence. Different states may have the same concurrence. Thus the value of  $H_{\alpha}(\vec{\mu})$  varies with different Schmidt coefficients  $\mu_i$  for fixed concurrence. We define two functions

$$
R_U(c) = \max\bigg\{H_{\alpha}(\vec{\mu})\bigg|\sqrt{2\big(1-\sum_{i=1}^m \mu_i^2\big)} \equiv c\bigg\},\tag{12}
$$

$$
R_L(c) = \min\bigg\{H_\alpha(\vec{\mu})\bigg|\sqrt{2\big(1-\sum_{i=1}^m\mu_i^2\big)}\bigg\} = c\bigg\}.\tag{13}
$$

The derivation of them is equivalent to finding the maximal and minimal of  $H_\alpha(\vec{\mu})$ . Notice that the definition of  $H_{\alpha}(\vec{\mu})$ , it is equivalent to find the maximal and minimal of  $\sum_{i=1}^{m}\mu_{i}^{\alpha}$  under the constraint  $\sqrt{2(1-\sum_{i=1}^{m}\mu_{i}^{2})} \equiv c$ 

since the logarithmic function is a monotonic function . With the method of Lagrange multipliers we obtain the necessary condition for the maximum and minimum of  $\sum_{i=1}^{m}\mu_{i}^{\alpha}$  as follows

$$
\alpha \mu_i^{\alpha - 1} = 2\lambda_i \mu_i - \lambda_2,\tag{14}
$$

where  $\lambda_1$ ,  $\lambda_2$  denote the Lagrange multipliers. This equation has maximally two nonzero solutions  $\gamma$  and *δ* for each *μ*<sub>*i*</sub> Let *n*<sub>1</sub> be the number of entries where *μ*<sub>*i*</sub> = γ and *n*<sub>2</sub> be the number of entries where *μ*<sub>*i*</sub> = *δ*. Thus the derivation is reduced to maximizes or minimizes the function

$$
R_{n_1 n_2}(c) = \frac{1}{1 - \alpha} \log(n_1 \gamma^{\alpha} + n_2 \delta^{\alpha}), \qquad (15)
$$

<span id="page-2-1"></span>under the constrains

$$
n_1\gamma + n_2\delta = 1, 2(1 - n_1\gamma^2 - n_2\delta^2) = c^2,
$$
\n(16)

<span id="page-2-0"></span>where  $n_1 + n_2 = d \le m$ . From Eq. [\(16\)](#page-2-0) we obtain two solutions of  $\gamma$ 

$$
\gamma_{n_1 n_2}^{\pm} = \frac{n_1 \pm \sqrt{n_1^2 - n_1(n_1 + n_2)[1 - n_2(1 - c^2/2)]}}{n_1(n_1 + n_2)},\tag{17}
$$

$$
\delta_{n_1 n_2}^{\pm} = \frac{1 - n_1 \gamma_{n_1 n_2}^{\pm}}{n_2},\tag{18}
$$

with max  $\{\sqrt{2(n_1-1)/n_1}, \sqrt{2(n_2-1)/n_2}\} \le c \le \sqrt{2(d-1)/d}$ . Because  $\gamma_{n_2n_1}^- = \delta_{n_1n_2}^+$ ,  $\delta_{n_2n_1}^- = \gamma_{n_1n_2}^+$ , we should only consider the case for  $\gamma_{n_1n_2}^+$ . When  $n_2 = 0$ ,  $\gamma$  can be uniquely determined case.

When  $m=3$ , the solution of Eq. ([15\)](#page-2-1) is  $R_{12}(c)$  and  $R_{21}(c)$  for  $1 < c \leq 2/\sqrt{3}$ . After a direct calculation we find  $R_{12}(c)$  and  $R_{21}(c)$  are both monotonically function of the concurrence *c*, and  $R_{12}(2/\sqrt{3}) = R_{21}(2/\sqrt{3})$ . In order to compare the value of  $R_{12}(c)$  and  $R_{21}(c)$  we only need to compare the value of them at the endpoint  $c=1$ . For convenience we divide the problem into three cases. If  $0 < \alpha < 2$ , then  $R_{12}(1) > R_{21}(1)$ ; If  $\alpha = 2$ , then  $R_{12}(1) = R_{21}(1)$ ; If  $\alpha > 2$ , then  $R_{12}(1) < R_{21}(1)$ . Thus we conclude that the maximal and minimal function of  $H_{\alpha}(\vec{\mu})$  is given by *R*<sub>21</sub>(*c*) and *R*<sub>12</sub>(*c*) respectively for  $\alpha > 2$ . When  $\alpha < 2$ , the maximal and minimal function of  $H_{\alpha}(\vec{\mu})$  is *R*<sub>12</sub>(*c*) and *R*<sub>12</sub>(*c*) and  $R_{21}(c)$  respectively. When  $\alpha = 2$ , we can check that the two functions  $R_{21}(c)$  and  $R_{12}(c)$  always have the same value for  $1 < c \leq 2/\sqrt{3}$ . In the general case for  $m = d$ , numerical calculation shows the following results

(i) When  $\alpha > 2$ ,

$$
R_L(c) = \frac{\log\left[ (\gamma_{1,d-1}^+)^{\alpha} + (d-1)^{1-\alpha} (1-\gamma_{1,d-1}^+)^{\alpha} \right]}{1-\alpha},\tag{19}
$$

$$
R_U(c) = \frac{\log\left[ (\gamma_{1,d-1}^-)^{\alpha} + (d-1)^{1-\alpha} (1-\gamma_{1,d-1}^-)^{\alpha} \right]}{1-\alpha},\tag{20}
$$

with 
$$
\sqrt{2(d-2)/(d-1)} < c \le \sqrt{2(d-1)/d}
$$
,  $1 \le d \le m-1$  and  
 $\gamma_{1,d-1}^{\pm} = (2 \pm \sqrt{2(d-1)[d(2-c^2)-2]})/2d$ .

(ii) When  $\alpha < 2$ ,

$$
R_L(c) = \frac{\log\left[ (\gamma_{1,d-1}^{\dagger})^{\alpha} + (d-1)^{1-\alpha} (1-\gamma_{1,d-1}^{\dagger})^{\alpha} \right]}{1-\alpha},\tag{21}
$$

$$
R_U(c) = \frac{\log\left[ (\gamma_{1,d-1}^+)^{\alpha} + (d-1)^{1-\alpha} (1-\gamma_{1,d-1}^+)^{\alpha} \right]}{1-\alpha}.
$$
\n(22)

(iii) When  $\alpha = 2$ , these lower and upper bounds give the same value.

We use the denotation co(*g*) to be the convex hull of the function *g*, which is the largest convex function that is bounded above by *g*, and ca(*g*) to be the smallest concave function that is bounded below by *g*. Using the results presented in Methods, we can prove the main result of this paper.

*Theorem.* For any  $m \otimes n$ ( $m \leq n$ ) mixed quantum state  $\rho$ , its ER $\alpha$ E satisfies

$$
co[R_{L}(C)] \leq E_{\alpha}(\rho) \leq ca[R_{U}(\overline{C})], \qquad (23)
$$

where



<span id="page-3-0"></span>**Figure 1.** The plot of lower bound (dashed line) and upper bound (dotted line) for  $\alpha = 3$ , m = 3. The upper bound consists of two segments and the lower bound consists of three segments. The solid line corresponds to  $R_{11}$ ,  $R_{12}$  and  $R_{21}$ .



<span id="page-3-1"></span>**Figure 2.** The plot of the second derivative of  $R_{12}$  for  $1 \lt c \leq 2/\sqrt{3}$ . 

$$
\overline{C} = \min \left\{ \sqrt{2(1 - \text{Tr}\rho_A^2)}, \sqrt{2(1 - \text{Tr}\rho_B^2)} \right\},\tag{24}
$$

and

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$$
\underline{C}^2 = \max\{0, 2/m(m-1)(\|\rho^{T_A}\| - 1)^2, 2/m(m-1)(\|\mathcal{R}(\rho)\| - 1)^2, 2(\text{Tr}\rho^2 - \text{Tr}\rho_A^2), 2(\text{Tr}\rho^2 - \text{Tr}\rho_B^2)\}.
$$
\n(25)

Next we consider how to calculate the expressions of  $co(R_L(c))$  and  $ca(R_U(c))$ . As an example, we only consider the case  $m=3$ . In order to obtain  $co(R<sub>L</sub>(c))$ , we need to find the largest convex function which bounded above by *R<sub>L</sub>*(*c*). We first set the parameter  $\alpha = 3$ , then we can derive

$$
R_{L}(c) = \begin{cases} R_{11}, & 0 < c \le 1 \\ R_{12}, & 1 < c \le 2/\sqrt{3} \\ R_{11}, & 0 < c \le 1 \end{cases}
$$
\n
$$
R_{U}(c) = \begin{cases} R_{11}, & 0 < c \le 1 \\ R_{21}, & 1 < c \le 2/\sqrt{3} \end{cases} \tag{26}
$$

We plot the function  $R_{11}$ ,  $R_{12}$  and  $R_{21}$  in Fig. [1](#page-3-0) which illustrates our result. It is direct to check that  $R''_{11} \ge 0$ , therefore  $co(R_{11}) = R_{11}$  for  $0 < c \le 1$ . The second derivative of  $R_{12}$  is not convex near  $c = 2/\sqrt{3}$  as shown in Fig. [2.](#page-3-1) In order to calculate  $co(R_{12})$ , we suppose  $l_1(c) = k_1(c - 2/\sqrt{3}) + \log 3$  to be the line crossing through the point  $[2/\sqrt{3}, R_{12}(2/\sqrt{3})]$ . Then we solve the equations  $l_1(c) = R_{12}(c)$  and  $dl_1(c)/dc = dR_{12}(c)/dc = k_1$  and the solution is  $k_1 = 5.2401$ ,  $c = 1.1533$ . Combining the above results, we get

$$
co(R_L(c)) = \begin{cases} R_{11}(0 < c \le 1) \\ R_{12}(1 < c \le 1.1533) \\ 5.2401(c - 2/\sqrt{3}) + \log 3 \\ (1.1533 < c \le 2/\sqrt{3}). \end{cases}
$$
(27)



<span id="page-4-0"></span>**Figure 3.** The plot of lower bound (dashed line) and upper bound (dotted line) for  $\alpha = 0.6$ , m = 3. The upper bound consists of two segments and the lower bound also consists of two segments. The solid line corresponds to  $R_{11}$ ,  $R_{12}$  and  $R_{21}$ .

Similarly, we can calculate that  $R''_{11} \geq 0$  and  $R''_{21} \geq 0$ , thus  $co(R_U(c))$  is the broken line connecting the following points:  $[0, 0]$ ,  $[1, \log 2]$ ,  $[2/\sqrt{3}, \log 3]$  $[2/\sqrt{3}, \log 3]$  $[2/\sqrt{3}, \log 3]$ . In Fig. 3 we have plotted the lower and upper bounds with dashed and dotted line respectively.

Then we choose the parameter  $\alpha$  = 0.6, and we get

. . . . . . . . . . . . . . . . . . . .

$$
R_L(c) = \begin{cases} R_{11}(0 < c \le 1) \\ R_{21}(1 < c \le 2/\sqrt{3}), \end{cases} \tag{28}
$$

$$
R_U(c) = \begin{cases} R_{11}(0 < c \le 1) \\ R_{12}(1 < c \le 2/\sqrt{3}). \end{cases} \tag{29}
$$

Since  $R_{11}'' \leq 0$ ,  $R_{21}'' \leq 0$ , we have that  $co(R_L(c))$  is the broken line connecting the points: [0, 0], [1,  $\log 2$ ],  $[2/\sqrt{3}$ ,  $\log 3$ ]. In order to obtain *ca*( $R_U(c)$ ), we need to find the smallest concave function which bounded below by  $R_U(c)$ . We find  $R_{11}'' \le 0$ ,  $R_{12}'' \ge 0$ , therefore  $ca(R_U(c))$  is the curve consisting  $R_{11}$  for  $0 < c \le 1$ and the line connecting points  $[1, R_{12}(1)]$  and  $[2/\sqrt{3}, R_{12}(2/\sqrt{3})]$  for  $1 < c \leq 2/\sqrt{3}$ . As shown in Fig. [3,](#page-4-0) the lower and upper bound both consists of two segments in this case.

Generally, we can get the expression of  $co(R<sub>L</sub>(c))$  and  $ca(R<sub>U</sub>(c))$  for other parameters  $\alpha$  and  $m$  using similar method.

**Examples.** In the following, we give two examples as applications of the above results.

*Example 1*. We consider the *d*⊗*d* Werner states

$$
\rho_f = \frac{1}{d^3 - d} [(d - f)I + (df - 1)\mathcal{F}],
$$
\n(30)

where  $-1 \le f \le 1$  and  $\mathcal F$  is the flip operator defined by  $\mathcal F(\phi \otimes \psi) = \psi \otimes \phi$ . It is shown in ref. [51](#page-8-33) that the concurrence  $C(\rho_f) = -f$  for  $f \le 0$  and  $C(\rho_f) = 0$  for  $f \ge 0$ . According to the theorem we obtain that  $1/(1 - \alpha)\log[((1 + \sqrt{1 - f^2})/2)^{\alpha} + ((1 - \sqrt{1 - f^2})/2)^{\alpha}] \le E_a(r_f) \le -f$  for  $-1 \le f \le 0$  when  $m = 3$ .

*Example 2.* The second example is the  $3 \otimes 3$  isotropic state  $\rho = (x/9)I + (1 - x)|\psi\rangle\langle\psi|$ , where  $\ket{\psi} = (a, 0, 0, 0, 1/\sqrt{3}, 0, 0, 0, 1/\sqrt{3})^t/\sqrt{a^2 + 2/3}$  with  $0 \le a \le 1$ . We choose  $x = 0.1$ , it is direct to calculate that

$$
C_1 = \sqrt{2(Tr\rho^2 - Tr\rho_A^2)} = \sqrt{2(Tr\rho^2 - Tr\rho_B^2)} = \frac{2\sqrt{6.53 + 41.46a^2 - 1.71a^4}}{3(2 + 3a^2)},
$$
\n(31)

$$
C_2 = \frac{1}{\sqrt{3}} (\|\rho^{T_A}\| - 1) = \frac{2(5 + 6.9a^2 - 0.9a^4 + 9.353a(2 + 3a^2))}{3(2 + 3a^2)^2},
$$
\n(32)

$$
C_3 = \frac{1}{\sqrt{3}} (\|R(\rho)\| - 1) = \frac{0.346 + 1.2a}{0.667 + a^2},
$$
\n(33)



<span id="page-5-0"></span>**Figure 4.** Lower and upper bounds of E $\alpha(\rho)$  for  $\alpha = 0.6$  where we have set x = 0.1. Red solid line is obtained by *C*<sub>1</sub>, the dash-dotted and dashed line is obtained by *C*<sub>2</sub> and *C*<sub>3</sub>, respectively. The blue solid line is the upper bound of  $E_{\alpha}(\rho)$ .

$$
\overline{C} = \sqrt{2(1 - Tr\rho_A^2)} = \sqrt{2(1 - Tr\rho_B^2)} = \frac{\sqrt{6(6.38 + 33.72a^2 + 3.42a^4)}}{3(2 + 3a^2)}.
$$
\n(34)

When  $\alpha$  = 0.6, we can calculate the lower and upper bounds and the results is shown in Fig. [4.](#page-5-0) The solid red line corresponds to the lower bound of *Eα* by choosing the lower bound of concurrence is *C*1, and the dash-dotted and dashed line correspond to the cases when we choose the lower bound of concurrence is  $C_2$  and  $C_3$ , respectively. We can choose the maximum value of the three curves as the lower bound of *Eα*. Te blue solid line is the upper bound of *Eα*.

**Relation with other entanglement measures.** In this section we establish the relation between ER*α*E and other well-known entanglement measures, such as the entanglement of formation, the geometric measure of entanglement $52$ , the logarithmic negativity and the G-concurrence.

**Entanglement of formation.** Let  $\rho$  be a bipartite pure state with Schmidt coefficients ( $\mu_1, \mu_2, \ldots$ ). We investigate the derivative of ER*α*E w.r.t. *α* as follows.

$$
\frac{dE_{\alpha}(\rho)}{d\alpha} = \frac{1}{(1-\alpha)^{2}} \left[ \sum_{j} \frac{\mu_{j}^{\alpha}}{\sum_{k} \mu_{k}^{\alpha}} \log \mu_{j}^{1-\alpha} + \log \sum_{k} \mu_{k}^{\alpha} \right]
$$
\n
$$
\leq \frac{1}{(1-\alpha)^{2}} \left[ \log \frac{\sum_{j} \mu_{j}}{\sum_{k} \mu_{k}^{\alpha}} + \log \sum_{k} \mu_{k}^{\alpha} \right]
$$
\n
$$
= 0. \tag{35}
$$

The inequality follows from the concavity of logarithm function. The last equality follows from the fact  $\sum_j \mu_j = 1$ . Hence the ER $\alpha$ E is monotonically non-increasing with  $\alpha \geq 0$ . Since it becomes the von Neumann entropy when *α* tends to one, we have

$$
E_{\alpha}(\rho) \ge E_{F}(\rho) \ge E_{\beta}(\rho) \tag{36}
$$

where 0≤*α*≤1 and *β*≥1. Using the convex roof, one can show that (36) also holds for mixed bipartite states *ρ*.

**Geometric measure of entanglement.** The geometric measure (GM) of entanglement measures the closest distance between a quantum state and the set of separable states<sup>52</sup>. The GM has many operational interpretations, such as the usability of initial states for Grovers algorithm, the discrimination of quantum states under LOCC and the additivity and output purity of quantum channels, see the introduction of ref. [48](#page-8-27) for a recent review on GM. For pure state  $|\psi\rangle$  we define  $G_l(\psi)=-\log\max|\langle\varphi|\psi\rangle|^2$ , where the maximum runs over all product states  $|\varphi\rangle$ . it is easy to see that max  $|\langle\varphi|\psi\rangle|^2$  is equal to the square of the maximum of Schmidt coefficients of |*ψ*〉. For mixed states *ρ* we defne

$$
G_{I}^{c}(\rho) := \min \sum_{i} p_{i} G_{I}(|\psi_{i}\rangle), \qquad (37)
$$

where the minimum runs over all decompositions of  $\rho = \sum_i p_i |\psi_i\rangle\langle\psi_i|^{48}$ . We construct the linear relation between the GM and ER*α*E as follows.

*Lemma*. If  $\alpha > 1$  then

$$
\frac{\alpha}{2(\alpha-1)}G_{I}^{c}(\rho)\geq E_{\alpha}(\rho).
$$
\n(38)

If  $\alpha = 1$  and  $\rho$  is a pure state then

$$
G_l^c(\rho) \le E_\alpha(\rho). \tag{39}
$$

If  $\alpha$  < 1 then

$$
E_{\alpha}(\rho) + \frac{\alpha}{2(1-\alpha)} G_{\mathbf{l}}^{c}(\rho) \le \frac{1}{1-\alpha} \log d, \tag{40}
$$

where *d* is the minimum dimension of  $H_A$  and  $H_B$ . The details for proving the lemma can be seen from Methods.

**logarithmic negativity.** In this subsection we consider the logarithmic negativity<sup>53</sup>. It is the lower bound of the PPT entanglement cost<sup>53</sup>, and an entanglement monotone both under general LOCC and PPT operations<sup>54</sup>. The logarithmic negativity is defined as

$$
LN(\rho) = \log \|\rho^{T_A}\|.\tag{41}
$$

Suppose  $\rho = \sum_i p_i |\psi_i\rangle \langle \psi_i|$  is the optimal decomposition of ER $\alpha E E_\alpha(\rho)$ , and the pure state  $|\psi_i\rangle$  has the standard  $\text{Schmidt form} |\psi_i\rangle = \sum_j \sqrt{\mu_{i,j}} |a_{i,j},\,b_{i,j}\rangle.$  For  $1/2$   $\leq$   $\alpha$   $\leq$   $(2n-1)/2n$  and  $n$   $>$  1, we have

$$
n \times LN(\rho) = n \log ||\rho^{T_A}|| \ge n \log \sum_{i} p_i ||(\psi_i) \langle \psi_i|)^{T_A}|| \ge n \sum_{i} p_i \log ||(\psi_i) \langle \psi_i|)^{T_A}||
$$
  
\n
$$
= 2n \sum_{i} p_i \log \left( \sum_{j} \sqrt{\mu_{i,j}} \right) \ge 2n \sum_{i} p_i \log \sum_{j} \mu_{i,j}^{\alpha} \ge \frac{1}{1 - \alpha} \sum_{i} p_i \log \sum_{j} \mu_{i,j}^{\alpha}
$$
  
\n
$$
= \sum_{i} p_i E_{\alpha}(|\psi_i\rangle) = E_{\alpha}(\rho) \tag{42}
$$

where the first inequality is due to the property proved in ref. [54,](#page-8-36) the second inequality is due to the concavity of logarithm function, and in the last inequality we have used the inequality  $2n \ge 1/(1-\alpha)$  for  $1/2 \le \alpha \le (2n-1)/2n$ ,  $n \ge 1$ .

**G-concurrence.** The G-concurrence is one of the generalizations of concurrence to higher dimensional case. It can be interpreted operationally as a kind of entanglement capacity<sup>[55](#page-8-37),[56](#page-8-38)</sup>. It has been shown that the G-concurrence plays a crucial role in calculating the average entanglement of random bipartite pure states<sup>57</sup> and demonstration of an asymmetry of quantum correlations<sup>58</sup>. Let  $|\psi\rangle$  be a pure bipartite state with the Schmidt decomposition  $\ket{\psi}=\sum_{i=1}^d\sqrt{\mu_i}|\vec{u}\rangle$ . The G-concurrence is defined as the geometric mean of the Schmidt coefficients<sup>55,[56](#page-8-38)</sup>

$$
G(\left|\psi\right) := d(\mu_1 \mu_2 \cdots \mu_d)^{1/d}.
$$
\n(43)

For  $\alpha > 1$ , we have

$$
E_{\alpha}(|\psi\rangle) = \frac{1}{1-\alpha} \log \sum_{i} \mu_{i}^{\alpha}
$$
  
\n
$$
\leq \frac{1}{1-\alpha} \log \left[ d(\mu_{1}^{\alpha} \cdots \mu_{d}^{\alpha}) \frac{1}{d} \right]
$$
  
\n
$$
= \frac{1}{(1-\alpha)} [\alpha \log d + \log (\mu_{1} \cdots \mu_{d}) \frac{\alpha}{d}]
$$
  
\n
$$
- (\alpha - 1) \log d]
$$
  
\n
$$
= \frac{\alpha}{1-\alpha} \log G(|\psi\rangle) + \log d.
$$
 (44)

For  $0 < \alpha < 1$ , we have

$$
E_{\alpha}(|\psi\rangle) \ge \frac{\alpha}{1-\alpha} \log G(|\psi\rangle) + \log d. \tag{45}
$$

#### **Discussion and Conclusion**

Entanglement Rényi-*α* entropy is an important generalization of the entanglement of formation, and it reduces to the standard entanglement of formation when  $\alpha$  approaches to 1. Recently, it has been proved<sup>59</sup> that the squared ER*α*E obeys a general monogamy inequality in an arbitrary *N*-qubit mixed state. Correspondingly, we can construct the multipartite entanglement indicators in terms of ER*α*E which still work well even when the indicators based on the concurrence and EOF lose their efficacy. However, the difficulties in minimization procedures restrict the application of ER*α*E. In this work, we present the frst lower and upper bounds for the ER*α*E of arbitrary dimensional bipartite quantum systems based on concurrence, and these results might provide an alternative method to investigate the monogamy relation in high-dimensional states. We also demonstrate the application our bound for some examples. Furthermore, we establish the relation between ER*α*E and some other entanglement measures. These lower and upper bounds can be further improved for other known bounds of concurrence<sup>[60,](#page-9-1)61</sup>. Afer completing this manuscript, we became aware of a recently related paper by Leditzky *et al*. in which they also obtained another lower bound of ΕRαΕ in terms of Rényi conditional entropy<sup>62</sup>.

### **Methods**

**Proof of the theorem.** Suppose  $\rho = \sum_j p_j |\psi_j\rangle \langle \psi_j|$  is the optimal decomposition of ER $\alpha E E_a(\rho)$ , and the concurrence of  $|\psi_j\rangle$  is denoted as  $c_j$ . Thus we have

$$
E_{\alpha}(\rho) = \sum_{j} p_{j} E_{\alpha}(\left|\psi_{j}\right\rangle) = \sum_{j} p_{j} H_{\alpha}(\overrightarrow{\mu})
$$
  
\n
$$
\geq \sum_{j} p_{j} co(R_{L}(c_{j})) \geq co\left[R_{L}\left(\sum_{j} p_{j} c_{j}\right)\right]
$$
  
\n
$$
\geq co[R_{L}(C)], \qquad (46)
$$

where the first inequality is due to the definition of  $co(g)$ ; in the second inequality we have used the monotonically increasing and convex properties of  $co(R_L(c_j))$  as a function of concurrence  $c_j$ ; and in the last inequality we have used the lower bound of concurrence. On the other hand, we have

$$
E_{\alpha}(\rho) = \sum_{j} p_{j} E_{\alpha}(|\psi_{j}\rangle) = \sum_{j} p_{j} H_{\alpha}(\vec{\mu})
$$
  
\n
$$
\leq \sum_{j} p_{j} ca(R_{U}(c_{j})) \leq ca \bigg[R_{U} \bigg(\sum_{j} p_{j} c_{j}\bigg)\bigg]
$$
  
\n
$$
\leq ca[R_{U}(\overline{C})], \qquad (47)
$$

where the first inequality is due to the definition of  $ca(g)$ ; the second inequality is due to the monotonically increasing and concave properties of  $ca(R_U(c_j))$  as a function of concurrence  $c_j$ ; and in the last inequality we have used the upper bound of concurrence. Tus we have completed the proof of the theorem.

**Proof of the lemma.** Suppose the minimum in (37) is reached at  $\rho = \sum_i p_i |\psi_i\rangle \langle \psi_i|$ . Let the Schmidt decom- $\text{position of }|\psi_i\rangle\text{ be }|\psi_i\rangle=\sum_j\!\!\sqrt{\mu_{i,j}}|a_{i,j},\,b_{i,j}\rangle\,\text{where}\,\mu_{i,1}\,\text{is the maximum Schmidt coefficient.}\,\text{For }\alpha>1\text{, we have}\,\text{for all }i\geq 1\text{, we have}\,\left\{\begin{aligned} &\int_{0}^{T}|\psi_i\rangle\,\text{d}\psi_i\rangle=\sum_j\sqrt{\mu$ 

$$
\frac{\alpha}{2(\alpha - 1)} G_1^c(\rho) = \frac{-\alpha}{2(\alpha - 1)} \sum_i p_i \log \mu_{i,1}^2
$$
  
\n
$$
= \frac{-1}{\alpha - 1} \sum_i p_i \log \mu_{i,1}^{\alpha}
$$
  
\n
$$
\geq \frac{-1}{\alpha - 1} \sum_i p_i \log \left( \sum_j \mu_{i,j}^{\alpha} \right)
$$
  
\n
$$
= \sum_i p_i E_{\alpha}(|\psi_i|)
$$
  
\n
$$
\geq E_{\alpha}(\rho).
$$
 (48)

We have proved (38). For  $\alpha = 1$ , let  $\mu_i$  be the Schmidt coefficients of  $\rho$ , we have

$$
E_{\alpha}(\rho) = S(\rho) = -\sum_{i} \mu_{i} \log \mu_{i}
$$
  
\n
$$
\geq -\sum_{i} \mu_{i} \log \max_{j} {\mu_{j}} = -\log \max_{j} {\mu_{j}}
$$
  
\n
$$
= G_{1}^{c}(\rho).
$$
 (49)

We have proved (39). For  $\alpha$  < 1, we have

$$
E_{\alpha}(\rho) + \frac{\alpha}{2(1-\alpha)} G_{1}^{c}(\rho) = E_{\alpha}(\rho) - \frac{\alpha}{2(1-\alpha)} \sum_{i} p_{i} \log \mu_{i,1}^{2}
$$
  
\n
$$
= E_{\alpha}(\rho) - \frac{1}{1-\alpha} \sum_{i} p_{i} \log \mu_{i,1}^{\alpha}
$$
  
\n
$$
= E_{\alpha}(\rho) - \frac{1}{1-\alpha} (\sum_{i} p_{i} \log(d\mu_{i,1}^{\alpha}) - \log d)
$$
  
\n
$$
\leq E_{\alpha}(\rho) - \frac{1}{1-\alpha} \sum_{i} p_{i} \log \left(\sum_{j} \mu_{i,j}^{\alpha}\right) + 11 - \alpha \log d
$$
  
\n
$$
\leq \frac{1}{1-\alpha} \log d.
$$
 (50)

The inequality holds because the pure state  $|\psi_i\rangle$  is in the  $d\times d$  space. So we have proved (40).

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#### **Acknowledgements**

WS was supported by NSF-China under Grant Nos 11374085, 11274010, the discipline top-notch talents Foundation of Anhui Provincial Universities, the Excellent Young Talents Support Plan of Anhui Provincial Universities, the Anhui Provincial Natural Science Foundation and the 136 Foundation of Hefei Normal University under Grant No. 2014136KJB04. LC was supported by the NSF-China (Grant No. 11501024), and the Fundamental Research Funds for the Central Universities (Grant Nos 30426401, 30458601 and 29816133).

#### **Author Contributions**

W. Song and L. Chen carried out the calculations. W. Song and L. Chen conceived the idea. All authors contributed to the interpretation of the results and the writing of the manuscript. All authors reviewed the manuscript.

#### **Additional Information**

Competing financial interests: The authors declare no competing financial interests.

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