ORIGINAL RESEARCH

Analysis of Multi-delay and Piecewise Constant Delay Systems by Hybrid Functions Approximation

S. Mashayekhi · M. Razzaghi · M. Wattanataweekul

Published online: 26 April 2014 © Foundation for Scientific Research and Technological Innovation 2014

Abstract In this paper, a new numerical method for solving multi-delay and piecewise constant delay systems is presented. The method is based upon hybrid functions approximation. The properties of hybrid functions consisting of block-pulse functions and Bernoulli polynomials are presented. The operational matrices of integration, product and delay are given. These matrices are then utilized to reduce the solution of multi-delay systems and the piecewise constant delay systems to the solution of algebraic equations. Illustrative examples are included to demonstrate the validity and applicability of the technique.

Keywords Hybrid · Bernoulli polynomials · Block-pulse · Delay systems

Introduction

Delays occur frequently in biological, chemical, electronic and transportation systems. Timedelay systems are therefore a very important class of systems whose control and optimization have been of interest to many investigators $[1-4]$ $[1-4]$. It is well known that it is difficult to analytically solve a delay system. Several numerical methods have been used to obtain an approximate solution for delay differential equations [\[5\]](#page-18-2).

The available sets of orthogonal functions can be divided into three classes. The first class includes sets of piecewise constant basis functions (e.g., block-pulse, Haar, Walsh, etc.). The second class consists of sets of orthogonal polynomials (e.g., Chebyshev, Laguerre, Legendre, etc.). The third class is the set of sine–cosine functions in the Fourier series. Orthogonal functions have been used when dealing with various problems of the dynamical systems.

M. Wattanataweekul

Department of Mathematics, Statistics and Computer, Faculty of Science, Ubon Ratchathani University, Ubonratchathani 34190, Thailand

S. Mashayekhi · M. Razzaghi (\boxtimes)

Department of Mathematics and Statistics,

Mississippi State University, Mississippi State, MS 39762, USA

e-mail: razzaghi@math.msstate.edu

The main advantage of using orthogonal functions is that they reduce the dynamical system problems to those of solving a system of algebraic equations. The approach is based on converting the underlying differential equation into an integral equation through integration, approximating various signals involved in the equation by truncated orthogonal functions, and using the operational matrix of integration *P* to eliminate the integral operations. The matrix *P* can be uniquely determined based on the particular orthogonal functions. Typical examples are given in [\[6](#page-18-3)[–11\]](#page-18-4).

Among piecewise constant basis functions, block-pulse functions are found to be very attractive, in view of their properties of simplicity and disjointedness. Bernoulli polynomials and Taylor series are not based on orthogonal functions, nevertheless, they possess the operational matrix of integration. However, since the integration of the cross product of two Taylor series vectors is given in terms of a Hilbert matrix [\[12](#page-18-5)], which are known to be ill conditioned, the applications of Taylor series are limited. Furthermore, the operational matrix *P*, in Bernoulli polynomials $\beta_m(t)$, $m = 0, 1, 2, \ldots$, where $0 \le t \le 1$, has less errors than *P* for Taylor series in $\frac{1}{2} \le t \le 1$ and $1 < m < 8$. This is because for *P* in $\beta_m(t)$ we ignore the term $\frac{\beta_{m+1}(t)}{m+1}$ while for *P* in Taylor series we ignore the term *tm*+¹ $\frac{m+1}{m+1}$.

For approximating an arbitrary time function the advantages of Bernoulli polynomials $\beta_m(t)$, over shifted Chebyshev polynomials $T_m(t)$, and shifted Legendre polynomials *L_m*(*t*), *m* = 0, 1, 2, ..., where $0 \le t \le 1$, are:

- (a) the operational matrix *P*, in Bernoulli polynomials has less errors than *P* for shifted Chebyshev and shifted Legendre polynomials for $1 < m < 10$. This is because for *P* in $T_m(t)$ we ignore the term $\frac{T_{m+1}(t)}{4(m+1)}$; and for *P* in $L_m(t)$ we ignore the term $\frac{L_{m+1}(t)}{2(2m+1)}$;
- (b) the Bernoulli polynomials have less terms than shifted Chebyshev polynomials and shifted Legendre polynomials. For example $\beta_6(t)$ has 5 terms, while $T_6(t)$ and $L_6(t)$ have 7 terms, and this difference will increase by increasing *m*. Hence for approximating an arbitrary function we use less CPU time by applying Bernoulli polynomials as compared to shifted Chebyshev and shifted Legendre polynomials;
- (c) the coefficient of individual terms in Bernoulli polynomials are smaller than the coefficient of individual terms in the shifted Chebyshev and shifted Legendre polynomials. Since the computational errors in the product are related to the coefficients of individual terms, the computational errors are less by using Bernoulli polynomials.

In recent years different types of hybrid functions have been shown to be mathematical power tools for discretization of selected problems [\[13](#page-18-6)[–17](#page-18-7)]. In the present paper we introduce a new direct computational method to solve delay systems. The method consists of reducing the delay differential equations to a set of algebraic equations by first expanding the solution of delay differential equations as a hybrid function with unknown coefficients. These hybrid functions, which consist of block-pulse functions and Bernoulli polynomials, are introduced. The operational matrices of integration, product, and delay are given. These matrices are then used to evaluate the coefficients of the hybrid functions for the solution of the delay differential equations.

The outline of this paper is as follows: In "Properties of Hybrid Functions" section we introduce the basic properties of the hybrid functions of block-pulse and Bernoulli polynomials required for our subsequent development. Section "Problem Statement" is devoted to the problem statement. In "The Numerical Method" section we apply the proposed numerical method to approximate the multi-delay systems and piecewise constant delay systems. In "Discussion" section a discussion of the present method is given and in "Illustrative Examples" section we report our numerical findings and demonstrate the accuracy of the proposed numerical scheme by considering four numerical examples.

Properties of Hybrid Functions

Hybrid of Block-Pulse and Bernoulli Polynomials

Hybrid functions $b_{nm}(t)$, $n = 1, 2, ..., N$, $m = 0, 1, ..., M$ are defined on the interval $[0, t_f)$ as

$$
b_{nm}(t) = \begin{cases} \beta_m \left(\frac{N}{t_f} t - n + 1 \right), & t \in \left[\frac{n-1}{N} t_f, \frac{n}{N} t_f \right), \\ 0, & \text{otherwise,} \end{cases}
$$
(1)

where *n* and *m* are the order of block-pulse functions and Bernoulli polynomials, respectively. In Eq. [\(1\)](#page-2-0), $\beta_m(t)$, $m = 0, 1, 2, \ldots$, are the Bernoulli polynomials of order *m*, which can be defined by [\[18\]](#page-18-8)

$$
\beta_m(t) = \sum_{k=0}^m \binom{m}{k} \alpha_k t^{m-k},
$$

where α_k , $k = 0, 1, \ldots, m$, are Bernoulli numbers. These numbers are a sequence of signed rational numbers, which arise in the series expansions of trigonometric functions [\[19](#page-18-9)] and can be defined by the identity

$$
\frac{t}{e^t-1}=\sum_{n=0}^{\infty}\alpha_n\frac{t^n}{n!}.
$$

The first few Bernoulli numbers are

$$
\alpha_0 = 1, \quad \alpha_1 = \frac{-1}{2}, \quad \alpha_2 = \frac{1}{6}, \quad \alpha_4 = \frac{-1}{30},
$$

with $\alpha_{2k+1} = 0, k = 1, 2, 3, \ldots$

The first few Bernoulli polynomials are

$$
\beta_0(t) = 1, \quad \beta_1(t) = t - \frac{1}{2}, \quad \beta_2(t) = t^2 - t + \frac{1}{6}, \quad \beta_3(t) = t^3 - \frac{3}{2}t^2 + \frac{1}{2}t.
$$

According to [\[20\]](#page-18-10), Bernoulli polynomials form a complete basis over the interval [0,1].

Function Approximation

Suppose that $H = L^2[0, 1]$ and $\{b_{10}(t), b_{20}(t), \ldots, b_{NM}(t)\} \subset H$ be the set of hybrid of block-pulse and Bernoulli polynomials and

$$
Y = \text{span}\{b_{10}(t), b_{20}(t), \dots, b_{N0}(t), b_{11}(t), b_{21}(t), \dots, b_{N1}(t), \dots, b_{1M}(t), b_{2M}(t), \dots, b_{NM}(t)\},\
$$

and f be an arbitrary element in H . Since Y is a finite dimensional vector space, f has the unique best approximation out of *Y* such as $f_0 \in Y$, that is

$$
\forall y \in Y, \quad ||f - f_0|| \le ||f - y||.
$$

 $\hat{\mathfrak{D}}$ Springer

Since $f_0 \in Y$, there exist unique coefficients $c_{10}, c_{20}, \ldots, c_{NM}$ such that

$$
f \simeq f_0 = \sum_{m=0}^{M} \sum_{n=1}^{N} c_{nm} b_{nm}(t) = C^T B(t),
$$
 (2)

where

$$
BT(t) = [b10(t), b20(t), ..., bN0(t), b11(t), b21(t), ..., bN1(t), ..., b1M(t),b2M(t), ..., bNM(t)],
$$
\n(3)

and

$$
CT = [c10, c20, ..., cN0, c11, c21, ..., cN1, ..., c1M, c2M, ..., cNM].
$$
 (4)

Approximation Errors

In this section we obtain bounds for the error of best approximation in terms of Sobolev norms. This norm is defined in the interval (a, b) for $\mu \geq 0$ by

$$
||f||_{H^{\mu}(a,b)} = \left(\sum_{k=0}^{\mu} \int_{a}^{b} |f^{(k)}(x)|^{2} dx\right)^{\frac{1}{2}} = \left(\sum_{k=0}^{\mu} ||f^{(k)}||_{L^{2}(a,b)}^{2}\right)^{\frac{1}{2}},
$$
(5)

where $f^{(k)}$ denotes the *k*th derivative of f. The symbol $|f|_{H^{\mu;M}(0,1)}$ which is introduced in [\[21\]](#page-18-11) was defined by

$$
|f|_{H^{\mu;M}(0,1)} = \left(\sum_{k=\min(\mu,M+1)}^{\mu} ||f^{(k)}||^2_{L^2(0,1)}\right)^{\frac{1}{2}}.
$$

The seminorm [\[22](#page-18-12)] is

$$
|f|_{H^{r;\mu;M;N}(0,1)} = \left(\sum_{k=\min(\mu,M+1)}^{\mu} N^{2r-2k} ||u^{(k)}||_{L^{2}(0,1)}^{2}\right)^{\frac{1}{2}},
$$

where

$$
|f|_{H^{r,\mu;M;N}(0,1)} = N^{r-\mu}||f^{(\mu)}||_{L^{2}(0,1)},
$$
\n(6)

 $M \geq \mu - 1$, and if $N = 1$, $|\cdot|_{H^r; \mu; M; N}$ coincides with $|\cdot|_{H^{\mu; M}}$ which was introduced in [\[21\]](#page-18-11). To state our main results, the following theorem and lemma will be required.

Theorem 1 *Suppose that* $f \in H^{\mu}(0, 1)$ *with* $\mu \geq 0$ *. If* $P_M f = \sum_{m=0}^{M} c_m \beta_m$ *is the best approximation of f then*

$$
||f - P_M f||_{L^2(0,1)} \le cM^{-\mu} |f|_{H^{\mu;M}(0,1)},\tag{7}
$$

and for $1 \leq r \leq \mu$ *,*

$$
||f - P_M f||_{H^r(0,1)} \le cM^{2r - \frac{1}{2} - \mu} |f|_{H^{\mu;M}(0,1)},
$$
\n(8)

where c depends on μ*.*

² Springer

Proof Let $f \in H^{\mu}(0, 1)$ with $\mu \ge 0$ and $\sum_{m=0}^{M} c_m L_m$ be the best approximation of *f*, which is constructed by using shifted Legendre polynomials L_m , $m = 0, \ldots, M$ in the interval [0, 1]. Then [\[21\]](#page-18-11)

$$
\left| \left| f - \sum_{m=0}^{M} \zeta_m L_m \right| \right|_{L^2(0,1)} \le c M^{-\mu} |f|_{H^{\mu;M}(0,1)}, \tag{9}
$$

and for $1 \leq r \leq \mu$,

$$
\left| \left| f - \sum_{m=0}^{M} \zeta_m L_m \right| \right|_{H^r(0,1)} \le c M^{2r - \frac{1}{2} - \mu} |f|_{H^{\mu;M}(0,1)}.
$$
 (10)

Since the best approximation is unique [\[20](#page-18-10)], we have

$$
\left| \left| f - \sum_{m=0}^{M} \hat{c}_{m} L_{m} \right| \right|_{L^{2}(0,1)} = \left| \left| f - P_{M} f \right| \right|_{L^{2}(0,1)}, \quad \left| \left| f - \sum_{m=0}^{M} \hat{c}_{m} L_{m} \right| \right|_{H^{r}(0,1)} = \left| \left| f - P_{M} f \right| \right|_{H^{r}(0,1)},\tag{11}
$$

and by using Eqs. $(9-11)$ $(9-11)$ we can obtain Eqs. [\(7\)](#page-3-0) and [\(8\)](#page-3-1).

Lemma 1 *For n* = 1, 2, ..., *N* suppose $f_n : (\frac{n-1}{N}, \frac{n}{N}) \to R$ is a function in $H^{\mu}(\frac{n-1}{N}, \frac{n}{N})$. *Consider the function* $F_n f_n : (0, 1) \to R$ *such that* $(F_n f_n)(x) = f_n(\frac{1}{N}(x + n - 1))$ *for all* $x \in (0, 1)$, *then for* $0 \le l \le \mu$, *we have*

$$
||(F_nf_n)^{(l)}||_{L^2(0,1)} = N^{\frac{1}{2}-l}||f_n^{(l)}||_{L^2(\frac{n-1}{N},\frac{n}{N})}.
$$

Proof For $0 \le l \le \mu$, we have

$$
|| (F_n f_n)^{(l)} ||_{L^2(0,1)}^2 = \int_0^1 | (F_n f_n)^{(l)}(x)|^2 dx = \int_0^1 \left| f_n^{(l)} \left(\frac{1}{N} (x + n - 1) \right) \right|^2 dx
$$

=
$$
\int_0^{\frac{n}{N}} N^{-2l} | f_n^{(l)}(t) |^2 N dt = N^{1-2l} || f_n^{(l)} ||_{L^2 \left(\frac{n-1}{N}, \frac{n}{N} \right)}^2,
$$

where for the third equality, we used the change of variable rule by setting $t = \frac{1}{N}(x + n - 1)$. \Box

By using the following Theorem we can obtain error for the approximation function.

Theorem 2 *Suppose* $f \in H^{\mu}(0, 1)$ *with* $\mu \geq 0$ *, then*

$$
||f - P_M^N f||_{L^2(0,1)} \le c M^{-\mu} |f|_{H^{0;\mu;M;N}(0,1)},
$$
\n(12)

and for $1 \leq r \leq \mu$ *,*

$$
||f - P_M^N f||_{H^r(0,1)} \le c M^{2r - \frac{1}{2} - \mu} |f|_{H^{r;\mu;M;N}(0,1)}.
$$
\n(13)

Proof For $n = 1, 2, ..., N$ we consider the function $f_n : (\frac{n-1}{N}, \frac{n}{N}) \to R$ such that $f_n(x) =$ *f* (*x*) for all $x \in (\frac{n-1}{N}, \frac{n}{N})$. By using Lemma 1 and Eq. [\(7\)](#page-3-0) we have

$$
||f - P_{M}^{N} f||_{L^{2}(0,1)}^{2} = \sum_{n=1}^{N} \left| \left| f_{n} - \sum_{m=0}^{M} c_{nm} b_{nm} \right| \right|_{L^{2}(\frac{n-1}{N}, \frac{n}{N})}^{2} = cN^{-1} \sum_{n=1}^{N} \left| \left| F_{n} f_{n} - P_{M}(F_{n} f_{n}) \right| \right|_{L^{2}(0,1)}^{2}
$$

\n
$$
\le cN^{-1} M^{-2\mu} \sum_{n=1}^{N} \sum_{k=\min(\mu, M+1)}^{M} ||(F_{n} f_{n})^{(k)}||_{L^{2}(0,1)}^{2}
$$

\n
$$
= cN^{-1} M^{-2\mu} \sum_{k=\min(\mu, M+1)}^{M} \sum_{n=1}^{N} N^{-2k} ||f_{n}^{(k)}||_{L^{2}(\frac{n-1}{N}, \frac{n}{N})}^{2}
$$

\n
$$
= cM^{-2\mu} \sum_{k=\min(\mu, M+1)}^{M} N^{-2k} ||f^{(k)}||_{L^{2}(0,1)}^{2},
$$

Also, for $1 \le r \le \mu$, by using Lemma 1 and Eq. [\(8\)](#page-3-1) we have

$$
\begin{split} ||f - P_{M}^{N} f||_{H^{r}(0,1)}^{2} &= \sum_{n=1}^{N} \left| \left| f_{n} - \sum_{m=0}^{M} c_{nm} b_{nm} \right| \right|_{H^{r}\left(\frac{n-1}{N}, \frac{n}{N}\right)}^{2} \\ &= c \sum_{n=1}^{N} \sum_{p=0}^{r} N^{2p-1} ||(F_{n} f_{n})^{(p)} - (P_{M}(F_{n} f_{n}))^{(p)}||_{L^{2}(0,1)}^{2} \\ &\le c N^{2r-1} M^{4r-1-2\mu} \sum_{n=1}^{N} \sum_{k=\min(\mu, M+1)}^{n} ||(F_{n} f_{n})^{(k)}||_{L^{2}(0,1)}^{2} \\ &= c N^{2r-1} M^{4r-1-2\mu} \sum_{k=\min(\mu, M+1)}^{n} \sum_{n=1}^{N} N^{1-2k} ||f_{n}^{(k)}||_{L^{2}\left(\frac{n-1}{N}, \frac{n}{N}\right)}^{2} \\ &= c M^{4r-1-2\mu} \sum_{k=\min(\mu, M+1)}^{n} N^{2r-2k} ||f^{(k)}||_{L^{2}(0,1)}^{2} .\end{split}
$$

 \Box

Conclusion Suppose $f \in H^{\mu}(0, 1)$ with $\mu \ge 0$, and $M \ge \mu - 1$, then by using Eq. [\(6\)](#page-3-2) and Theorem 2 we get

$$
||f - P_M^N f||_{L^2(0,1)} \le c M^{-\mu} N^{-\mu} ||f^{(\mu)}||_{L^2(0,1)},
$$
\n(14)

and for $r \geq 1$,

$$
||f - P_M^N f||_{H^r(0,1)} \le cM^{2r - \frac{1}{2} - \mu} N^{r - \mu} ||f^{(\mu)}||_{L^2(0,1)}.
$$
 (15)

This result shows that in the case f is infinitely smooth, the rate of convergence of $P_M^N f$ to *f* is faster than $\frac{1}{N}$ to the power of $M + 1 - r$ and any power of $\frac{1}{M}$, which is superior to that for the classical spectral methods [\[21](#page-18-11)].

Operational Matrix of Integration and Product

The integration of the $B(t)$ defined in Eq. [\(3\)](#page-3-3) is given by

$$
\int_{0}^{t} B(t')dt' \simeq PB(t),\tag{16}
$$

,

where *P* is the $N(M + 1) \times N(M + 1)$ operational matrix of integration. The matrix *P* for $t_f = 1$ is given in [\[23](#page-19-0)] by

$$
P = \frac{1}{N} \begin{bmatrix} P_0 & I & O & \dots & O \\ \frac{-1}{2}\alpha_2 I & O & \frac{1}{2}I & \dots & O \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \frac{-1}{M}\alpha_M I & O & O & \dots & \frac{1}{M}I \\ \frac{-1}{M+1}\alpha_{M+1}I & O & O & \dots & O \end{bmatrix}
$$

where *I* and *O* are $N \times N$ identity and zero matrices respectively, and

$$
P_0 = \left[\begin{array}{ccccc} -\alpha_1 & 1 & \dots & 1 & 1 \\ 0 & -\alpha_1 & \dots & 1 & 1 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & -\alpha_1 & 1 \\ 0 & 0 & \dots & 0 & -\alpha_1 \end{array} \right].
$$

It is seen that *P* is more sparse than operational matrices of integration for the hybrid of blockpulse and Chebyshev polynomials [\[13\]](#page-18-6), the hybrid of block-pulse and Legendre polynomials [\[14\]](#page-18-13), and the hybrid of block-pulse and Taylor series [\[15\]](#page-18-14).

The product of two hybrid functions with the vector *C* is given by

$$
B(t)BT(t)C \simeq \tilde{C}B(t),
$$
\n(17)

where \tilde{C} is a $N(M + 1) \times N(M + 1)$ product operational matrix which is given in [\[24\]](#page-19-1).

The Operational Matrix of Multi-delay Systems

The delay function $B(t - \eta_i)$, $j = 1, 2, ..., r$, is the shift of the function $B(t)$ defined in Eq. [\(3\)](#page-3-3), along the time axis by η_j , where $\eta_1, \eta_2, \ldots, \eta_r$ are rational numbers in (0, 1). It is assumed that without loss of generality that $\eta_1 < \eta_2 < \cdots < \eta_r$. The general expression is given by

$$
B(t - \eta_j) = D_j B(t), \quad t > \eta_j, \quad 0 < t < 1,
$$
\n(18)

where D_i is the operational matrix of delay of hybrid functions corresponding to η_i . To find D_i , for $j = 1, 2, \ldots, r$, we first choose *N* in the following manner:

We define w as the smallest positive integer number for which $w\eta_i \in Z$ for $j =$ $1, 2, \ldots, r$, where *Z* is the set of all integer numbers. Next we choose λ as the greatest common divisor of the integers $w\eta_i$, $j = 1, 2, ..., r$. That is

$$
\lambda = \text{g.c.d}(w\eta_1, w\eta_2, \ldots, w\eta_r).
$$

 $\circled{2}$ Springer

Let

$$
N = \begin{cases} \frac{w}{\lambda}, & \frac{w}{\lambda} \in Z, \\ \left[\frac{w}{\lambda}\right] + 1, & \text{otherwise,} \end{cases}
$$
(19)

where $\lceil \cdot \rceil$ denotes the greatest integer value. Thus we have different subintervals given by

$$
\left[0, \frac{1}{N}\right), \left[\frac{1}{N}, \frac{2}{N}\right), \ldots, \left[\frac{N-1}{N}, 1\right).
$$

With the aid of Eq. [\(1\)](#page-2-0), it is noted that $b_{im}(t)$, $m = 0, 1, \ldots, M$, are non zero in the interval $\left[\frac{i-1}{N}, \frac{i}{N}\right)$. Thus $b_{im}(t-\eta_j)$ are non-zero in the $\left[\frac{i-1}{N} + \eta_j, \frac{i}{N} + \eta_j\right)$. If we expand $b_{im}(t-\eta_j)$ in terms of the element of $B(t)$ in Eq. [\(3\)](#page-3-3), we have $b_{im}(t-\eta_i) = b_{\beta_{ij}m}(t)$ where $\beta_{ij} = [N\eta_j + i]$. Thus, if we expand $B(t - \eta_i)$ in terms of $B(t)$ we find $N(M + 1) \times N(M + 1)$ delay matrix D_i as

$$
D_j = \begin{bmatrix} \psi_j & O & O & \cdots & O \\ O & \psi_j & O & \cdots & O \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ O & O & O & \psi_j & O \\ O & O & O & O & \psi_j \end{bmatrix},
$$

where *O* is $N \times N$ zero matrix, and ψ_j is $N \times N$ matrix which is given by

$$
\psi_j = \begin{bmatrix} 0 & 0 & \cdots & 1 & 0 & \cdots & 0 \\ 0 & 0 & \cdots & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \cdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 0 & \cdots & 1 \\ \vdots & \vdots & \vdots & \cdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 0 & 0 & 0 \end{bmatrix}.
$$

It is noted that the first one in the first row is located at the β_{1j} th column.

The Operational Matrix of Piecewise Constant Delay Systems

In this section we obtain the operational matrix of delay *Q* for piecewise constant delay systems. The general expression is given by

$$
B(t - a(t)) = QB(t), \quad 0 < t < 1,\tag{20}
$$

where $a(t)$ is considered as follows:

$$
a(t) = \begin{cases} \xi_1, & 0 \le t < T_1, \\ \xi_2, & T_1 \le t < T_2, \\ \vdots & \vdots \\ \xi_s, & T_{s-1} \le t < 1. \end{cases}
$$

It is assumed that $\xi_1, \xi_2, \ldots, \xi_r$ are rational numbers in [0, 1), and $\xi_1 < \xi_2 < \cdots < \xi_r$. We first choose *N* in the following manner: Let

$$
K=\{i:\,\xi_i\neq 0\},\,
$$

 \circledcirc Springer

and define γ as the smallest positive integer number for which

$$
\gamma \xi_i \in Z, \quad \gamma T_j \in Z, \quad i \in K, \quad j = 1, 2, \dots, s - 1.
$$

Next suppose $|K| = \alpha$, and choose δ as the greatest common divisor of the integers $\gamma \xi_i$ and γT_i , $i \in K$, $j = 1, 2, ..., s - 1$. That is

$$
\delta = g.c.d(\gamma\xi_1,\gamma\xi_2,\ldots,\gamma\xi_\alpha,\gamma T_1,\gamma T_2,\ldots,\gamma T_{s-1}).
$$

Let $h = \frac{\delta}{\gamma}$ and select *N* to minimize the number of subintervals as

$$
N = \begin{cases} \frac{1}{h}, & \frac{1}{h} \in Z, \\ \left[\frac{1}{h}\right] + 1, & \text{otherwise.} \end{cases}
$$
 (21)

Thus we have different subintervals given by

$$
[0, h), [h, 2h), \ldots, [(N-1)h, Nh).
$$

For these subintervals we can redefine $a(t)$ as

$$
a(t) = \xi_i \text{ for } h\left(\sum_{l=0}^{i-1} N_l\right) \le t < h\left(\sum_{l=0}^{i} N_l\right), \quad i = 1, 2, \dots, s,
$$
\nwhere

wh

$$
N_0 = 0, \quad N_j = \frac{T_j - T_{j-1}}{h}, \quad T_0 = 0, \quad T_s = 1,
$$

clearly $\sum_{j=1}^{s} N_j = N$. Therefore the problem has been reduced to find the delay operational matrix for the following delay function:

$$
B(t - a(t)) = \begin{cases} B(t - v_1 h), & 0 \le t < t_1, \\ B(t - v_2 h), & t_1 \le t < t_2, \\ \vdots & \vdots \\ B(t - v_N h), & t_{N-1} \le t < t_N. \end{cases}
$$

For evaluating v_i first we choose i as the smallest positive integer number such that

$$
1 + \sum_{l=0}^{i-1} N_l \le j \le \sum_{l=0}^{i} N_l.
$$

Then we choose v_i in such a way that

$$
t_i = ih
$$
, $v_j = \frac{\xi_i}{h}$, $i = 1, 2, ..., s$, $j = 1, 2, ..., N$.

In order to construct the matrix Q , we first find the matrix Q_i for $i = 1, 2, ..., N$ in such a way that the following relation is satisfied:

$$
B(t-v_i h) = Q_i B(t), \quad t_{i-1} \leq t < t_i.
$$

With the aid of Eq. [\(1\)](#page-2-0), it is noted that for the case $t_{i-1} \leq t < t_i$, the only terms with nonzero values are $b_{(i-v_i)m}(t-v_ih)$, for $m=0, 1, \ldots, M$. Since

$$
b_{(i-v_i)m}(t-v_ih) = b_{im}(t), \quad m = 0, 1, ..., M,
$$

we have

$$
Q_i = I_{M+1} \otimes S_i, \quad i = 1, 2, ..., N,
$$
 (22)

 $\hat{\mathfrak{D}}$ Springer

where I_{M+1} is the $(M + 1)$ dimensional identity matrix and S_i is an $N \times N$ matrix in which the only nonzero entry is equal to one and located at the $(i - v_i)$ th row and *i*th column and ⊗ denotes Kronecker product [\[25\]](#page-19-2). It is noted that if $i - v_i \le 0$, then S_i is a zero matrix of order *N* × *N*. Thus, if we expand $B(t - a(t))$ in terms of hybrid functions $B(t)$, we get

$$
Q=Q_1+Q_2+\cdots+Q_N.
$$

Problem Statement

Consider the following linear time-varying delay system:

$$
\dot{X}(t) = E(t)X(t) + \kappa_1 \sum_{j=1}^r F_j(t)X(t - \eta_j) + \kappa_2 H(t)X(t - a(t)) + G(t)U(t), \quad 0 \le t \le 1,
$$
 (23)

$$
X(0) = X_0,\tag{24}
$$

$$
X(t) = \phi(t), \quad t < 0,\tag{25}
$$

where $X(t) \in R^l$, $U(t) \in R^q$, $E(t)$, $F_j(t)$ $j = 1, 2, ..., r$, $H(t)$ and $G(t)$ are matrices of appropriate dimensions, X_0 is a constant specified vector, κ_1 and κ_2 are either 0 or 1 such that if $\kappa_1 = 1$ then $\kappa_2 = 0$ and vice versa, also if $\kappa_1 = 0$ then $\phi(t) = 0$. The problem is to find $X(t)$, $0 \le t \le 1$, in Eq. [\(23\)](#page-9-0) satisfying Eqs. [\(24\)](#page-9-0) and [\(25\)](#page-9-0). The method of this paper can also be extended to cases with delays in both state and control.

The Numerical Method

We approximate $X(t)$ in Eq. [\(23\)](#page-9-0) as follows:

$$
X(t) = [x_1(t), x_2(t), \dots, x_l(t)]^T, \quad U(t) = [u_1(t), u_2(t), \dots, u_q(t)]^T,
$$
 (26)

$$
\hat{B}(t) = I_l \otimes B(t), \quad \hat{B}_1(t) = I_q \otimes B(t), \tag{27}
$$

where I_l and I_q are the l - and q -dimensional identity matrices. Also, $B(t)$ and $B_1(t)$ are $l(M + 1)N \times l$ and $q(M + 1)N \times q$ matrices, respectively. By using Eq. [\(2\)](#page-3-4) each of $x_i(t)$ and each of $u_j(t)$, $i = 1, 2, ..., l$, $j = 1, 2, ..., q$, can be written in terms of hybrid functions as

$$
x_i(t) = B^T(t)X_i, \quad u_j(t) = B^T(t)U_j.
$$

From Eqs. (26) and (27) we get

$$
X(t) = \hat{B}^{T}(t)X, \quad U(t) = \hat{B}_{1}^{T}(t)U,
$$
\n(28)

where *X* and *U* are vectors of order $l(M + 1)N \times 1$ and $q(M + 1)N \times 1$, respectively, given by

$$
X = [X_1, X_2, ..., X_l]^T
$$
, $U = [U_1, U_2, ..., U_q]^T$.

Similarly we have

$$
X(0) = \hat{B}^{T}(t)d, \quad \phi_{j}(t - \eta_{j}) = \hat{B}^{T}(t)R_{j}, \tag{29}
$$

 \circledcirc Springer

where *d* and R_j , $j = 1, 2, ..., r$ are vectors of order $l(M + 1)N \times 1$. We expand $E(t)$, $F_i(t)$, $H(t)$ and $G(t)$ by hybrid functions as follows:

$$
E(t) = E^T \hat{B}(t), \quad F_j(t) = F_j^T \hat{B}(t), \quad H(t) = H^T \hat{B}(t), \quad G(t) = G^T \hat{B}_1(t),
$$

where E^T , F_{j}^T , H^T and G^T are of dimensions $l \times l(M+1)N$, $l \times l(M+1)N$ and $l \times q(M+1)N$, respectively. We can write $X(t - \eta_i)$ and $X(t - a(t))$ in terms of hybrid functions as

$$
X(t - \eta_j) = \begin{cases} \hat{B}^T(t)R_j, & 0 \le t \le \eta_j, \\ \hat{B}^T(t)\hat{D}_j^T X, & \eta_j \le t \le 1, \end{cases}
$$

and

$$
X(t - a(t)) = \hat{B}^T(t)\hat{Q}^T X,
$$

where

$$
\hat{D}_j = I_l \otimes D_j, \quad \hat{Q} = I_l \otimes Q,
$$

 D_i , and Q are the delay operational matrices given in Eqs. [\(18\)](#page-6-0) and [\(20\)](#page-7-0). Now we have

$$
E(t)X(t) = E^{T}\hat{B}(t)\hat{B}^{T}(t)X = \hat{B}^{T}(t)\tilde{E}^{T}X,
$$
\n(30)

$$
G(t)U(t) = G^T \hat{B}_1(t) \hat{B}_1^T(t)U = \hat{B}^T(t) \tilde{G}^T U,
$$
\n(31)

$$
H(t)X(t - a(t)) = H^T \hat{B}(t) \hat{B}^T(t) \hat{Q}^T X = \hat{B}^T(t) \tilde{H}^T \hat{Q}^T X,
$$
 (32)

where \tilde{E} , \tilde{G} and \tilde{H} can be calculated similarly to matrix \tilde{C} in Eq. [\(17\)](#page-6-1). Also we have

$$
\int_{0}^{t} \hat{B}^{T}(t')dt' = (I_{l} \otimes B^{T}(t))(I_{l} \otimes P^{T}) = \hat{B}^{T}(t)\hat{P}^{T},
$$
\n(33)

$$
\int_{0}^{t} F_{j}(t')X(t'-\eta_{j})dt' = \begin{cases} \hat{B}^{T}(t)\hat{P}^{T}\tilde{F}_{j}^{T}R_{j}, & 0 \leq t \leq \eta_{j}, \\ \hat{B}^{T}(t)Z_{j}\tilde{F}_{j}^{T}R_{j} + \hat{B}^{T}(t)\hat{P}^{T}\tilde{F}_{j}^{T}\hat{D}_{j}^{T}X, & \eta_{j} \leq t \leq 1, \end{cases}
$$
(34)

where P is the operational matrix of integration given in Eq. [\(16\)](#page-6-2) and

$$
\int\limits_{0}^{\tau_j} \hat{B}^T(t)dt = \hat{B}^T(t)Z_j,
$$

where Z_j is a constant matrix of order $l(M + 1)N \times l(M + 1)N$. By integrating Eq. [\(23\)](#page-9-0) from 0 to t and using Eqs. $(24-34)$ $(24-34)$ we have

$$
\hat{B}^{T}(t)X - \hat{B}^{T}(t)d = \hat{B}^{T}(t)\hat{P}^{T}\tilde{E}^{T}X + \kappa_{1}\sum_{j=1}^{r}(\hat{B}^{T}(t)\hat{P}^{T}\tilde{F}_{j}^{T}R_{j} \n+ \hat{B}^{T}(t)Z_{j}\tilde{F}_{j}^{T}R_{j} + \hat{B}^{T}(t)\hat{P}^{T}\tilde{F}_{j}^{T}\hat{D}_{j}^{T}X) \n+ \kappa_{2}\hat{B}^{T}(t)\hat{P}^{T}\tilde{H}^{T}\hat{Q}^{T}X + \hat{B}^{T}(t)\hat{P}^{T}\tilde{G}^{T}U.
$$
\n(35)

² Springer

From Eq. [\(35\)](#page-10-1) we get *X* as

$$
X = \left[I - \hat{P}^T \tilde{E}^T - \kappa_1 \sum_{j=1}^r \hat{P}^T \tilde{F}_j^T \hat{D}_j^T - \kappa_2 \hat{P}^T \tilde{H}^T \hat{Q}^T\right]^{-1}
$$

$$
\times \left[d + \kappa_1 \sum_{j=1}^r (\hat{P}^T \tilde{F}_j^T R_j + Z_j \tilde{F}_j^T R_j) + \hat{P}^T \tilde{G}^T U\right].
$$

Hence, $X(t)$ in Eq. [\(28\)](#page-9-2) can be obtained.

Discussion

Due to the nature of time-delay systems, the exact solutions of these systems are different functions on the distinct subintervals. In such situations, neither the continuous basis functions nor piecewise constant basis functions taken alone would form an efficient basis in the representation of such solutions. Datta and Mohan [\[26\]](#page-19-3) have correctly pointed out that, in general, the computed response of the delay systems via continuous or piecewise constant basis functions is not in good agreement with the exact response of the system. To meet these situations, we choose a suitable hybrid system of basis functions inherently possessing the required features of the solutions corresponding to delay systems. In the proposed method, with the aid of Eqs. [\(19\)](#page-7-1) and [\(21\)](#page-8-0) we determine the appropriate value for *N*, the order of block-pulse functions for multi-delay and piecewise constant delay systems respectively. To select *M*, we first choose an arbitrary number depending on the problem. If the exact solutions in each subinterval are polynomials we increase the value of *M* by 1 until two consecutive results are the same in each subinterval. When the exact solutions in each subinterval are not polynomials, we evaluate the results for two consecutive M for different t in [0, 1] until the results are similar up to a required number of decimal places for each subinterval.

Illustrative Examples

In this "The Numerical Method" section, examples are given to demonstrate the applicability and accuracy of our method. Example 1 is a two-dimensional time-varying multi-delay system which was considered in [\[26](#page-19-3)]. The exact solutions in Examples 1 are polynomials for each subinterval. Example 2 is a delay system with delay in both control and state. Example 2 was first considered in [\[26](#page-19-3)] and also solved in [\[13](#page-18-6)[–15\]](#page-18-14) by using the hybrid of block-pulse with Chebyshev polynomials, with Legendre polynomials, and with Taylor series respectively. In this example we compare the solution obtained using the proposed method with [\[13](#page-18-6)[–15\]](#page-18-14). The CPU time in each case are also given. Example 3 is a time-varying piecewise constant delay system whose solution is a polynomial for each subinterval and was considered in [\[5\]](#page-18-2). Example 4 is a piecewise constant delay system which was first considered in [\[26](#page-19-3)] and also solved in [\[5](#page-18-2)] by using the hybrid of block-pulse and Chebyshev polynomials. For example 4, we compare our findings with the numerical results in [\[5\]](#page-18-2) together with the CPU time.

Example 1

Consider the multi-delay systems described by [\[26](#page-19-3)]

$$
\begin{pmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \end{pmatrix} = \begin{pmatrix} t & 1 \\ t & 2t \end{pmatrix} \begin{pmatrix} x_1 \left(t - \frac{1}{3}\right) \\ x_2 \left(t - \frac{1}{3}\right) \end{pmatrix} + \begin{pmatrix} 2 & t \\ t^2 & 0 \end{pmatrix} \begin{pmatrix} x_1 \left(t - \frac{2}{3}\right) \\ x_2 \left(t - \frac{2}{3}\right) \end{pmatrix} + \begin{pmatrix} 0 \\ 1 \end{pmatrix} u(t), \quad (36)
$$

with

$$
x_1(t) = x_2(t) = u(t) = 0, \quad t \in \left[-\frac{2}{3}, 0 \right]. \tag{37}
$$

$$
u(t) = 2t + 1, \quad t > 0. \tag{38}
$$

The exact solutions are [\[26\]](#page-19-3)

$$
x_1(t) = \begin{cases} 0, & 0 \le t < \frac{1}{3}, \\ \frac{7}{162} - \frac{2}{9}t + \frac{1}{6}t^2 + \frac{1}{3}t^3, & \frac{1}{3} \le t < \frac{2}{3}, \\ \frac{11}{162} - \frac{58}{243}t + \frac{31}{162}t^2 + \frac{1}{9}t^3 + \frac{7}{72}t^4 + \frac{1}{6}t^5, & \frac{2}{3} \le t \le 1. \end{cases}
$$

$$
x_2(t) = \begin{cases} t + t^2, & 0 \le t < \frac{1}{3}, \\ \frac{5}{486} + t + \frac{7}{9}t^2 + \frac{2}{9}t^3 + \frac{1}{2}t^4, & \frac{1}{3} \le t < \frac{2}{3}, \\ \frac{1}{486} + t + \frac{200}{243}t^2 + \frac{20}{81}t^3 + \frac{29}{72}t^4 - \frac{1}{9}t^5 + \frac{1}{6}t^6, & \frac{2}{3} \le t \le 1. \end{cases}
$$

By using Eq. [\(19\)](#page-7-1) we have $\omega = 3$ and $\lambda = 1$, so we select $N = 3$. We also choose $M = 6$. Let

$$
x_1(t) = C_1^T B(t), \quad x_2(t) = C_2^T B(t).
$$
 (39)

By using Eq. [\(2\)](#page-3-4) we get

$$
u(t) = \left[\frac{4}{3}, \frac{6}{3}, \frac{8}{3}, \frac{2}{3}, \frac{2}{3}, \frac{2}{3}, \frac{0, \dots, 0}{15}\right]^T B(t) = U_1^T B(t)
$$
 (40)

and

$$
t = \left[\frac{1}{6}, \frac{1}{2}, \frac{5}{6}, \frac{1}{3}, \frac{1}{3}, \frac{1}{3}, \frac{0, \dots, 0}{15}\right]^T B(t) = Y_1^T B(t), \tag{41}
$$

$$
t^{2} = \left[\frac{1}{27}, \frac{7}{27}, \frac{19}{27}, \frac{1}{9}, \frac{1}{3}, \frac{1}{9}, \frac{5}{9}, \frac{1}{9}, \frac{1}{9}, \frac{1}{9}, \frac{0, \dots, 0}{12}\right]^{T} B(t) = Y_{2}^{T} B(t). \tag{42}
$$

Also we have

$$
x_1\left(t - \frac{1}{3}\right) = C_1^T \underline{D}_1 B(t), \quad x_2\left(t - \frac{1}{3}\right) = C_2^T \underline{D}_1 B(t), \tag{43}
$$

$$
x_1\left(t - \frac{2}{3}\right) = C_1^T \underline{D}_2 B(t), \quad x_2\left(t - \frac{2}{3}\right) = C_2^T \underline{D}_2 B(t), \tag{44}
$$

where \underline{D}_1 and \underline{D}_2 are the delay operational matrices given in Eq. [\(18\)](#page-6-0). By integrating Eq. (36) from 0 to *t* and using Eqs. $(37-44)$ $(37-44)$ we have

$$
C_1^T = C_1^T \underline{D}_1 \widetilde{Y}_1 P + C_2^T \underline{D}_1 P + 2C_1^T \underline{D}_2 P + C_2^T \underline{D}_2 \widetilde{Y}_1 P,\tag{45}
$$

$$
C_2^T = C_1^T \underline{D}_1 \widetilde{Y}_1 P + 2C_2^T \underline{D}_1 \widetilde{Y}_1 P + C_2^T \underline{D}_2 \widetilde{Y}_2 P + U_1^T P, \tag{46}
$$

where Y_1 and Y_2 can be obtained similarly to Eq. [\(17\)](#page-6-1). Solving Eqs. [\(45\)](#page-12-2) and [\(46\)](#page-12-3), the same values as the exact values of $x_1(t)$ and $x_2(t)$ would be obtained.

 \bigcirc Springer

Example 2

Consider the following delay system with delay in both control and state [\[26\]](#page-19-3)

$$
\dot{x}(t) = -x(t) - 2x\left(t - \frac{1}{4}\right) + 2u\left(t - \frac{1}{4}\right), \quad u(t) = 1, \quad t > 0,
$$
\n(47)

$$
x(t) = u(t) = 0, \quad -\frac{1}{4} \le t < 0. \tag{48}
$$

The exact solution is

$$
x(t) = \begin{cases} 0, & 0 \le t \le \frac{1}{4}, \\ 2 - 2 \exp\left[-\left(t - \frac{1}{4}\right)\right], & \frac{1}{4} \le t \le \frac{1}{2}, \\ -2 - 2 \exp\left[-\left(t - \frac{1}{4}\right)\right] + (2 + 4t) \exp\left[-\left(t - \frac{1}{2}\right)\right], & \frac{1}{2} \le t \le \frac{3}{4}, \\ 6 - 2 \exp\left[-\left(t - \frac{1}{4}\right)\right] + (2 + 4t) \exp\left[-\left(t - \frac{1}{2}\right)\right] \\ -\left(\frac{17}{4} + 2t + 4t^2\right) \exp\left[-\left(t - \frac{3}{4}\right)\right], & \frac{3}{4} \le t < 1. \end{cases}
$$

By using Eq. [\(19\)](#page-7-1) we have $\omega = 4$ and $\lambda = 1$, so we select $N = 4$. We also choose $M = 5$. Let

$$
x(t) = C_3^T B(t),\tag{49}
$$

by using Eq. [\(2\)](#page-3-4) we have

$$
u(t) = [1, 1, 1, 1, \underbrace{0, \dots, 0}_{20}]^{T} B(t) = U_{2}^{T} B(t).
$$
 (50)

Also we have

$$
x\left(t - \frac{1}{4}\right) = C_3^T \underline{D}_3 B(t),\tag{51}
$$

$$
u\left(t-\frac{1}{4}\right) = U_2^T \underline{D}_3 B(t),\tag{52}
$$

where \underline{D}_3 is the delay operational matrix given in Eq. [\(18\)](#page-6-0). By integrating Eq. [\(47\)](#page-13-0) from 0 to t and using Eqs. $(48-52)$ $(48-52)$ we have

$$
C_3^T = -C_3^T P - 2C_3^T \underline{D}_3 P + 2U_2^T \underline{D}_3 P.
$$
\n(53)

By solving Eq. [\(53\)](#page-13-2) we obtain *C*3. In Table [1,](#page-14-0) the solution obtained using the hybrid of block-pulse and Chebyshev polynomials in [\[13\]](#page-18-6) with $N = 4$ and $M_1 = 7$ and the hybrid of block-pulse and Legendre polynomials in $[14]$ $[14]$ with $N = 4$ and $M_2 = 7$ together with CPU time are given. In this table *M*¹ and *M*² denote the order of Chebyshev and Legendre polynomials respectively. In Table [2,](#page-14-1) the solution obtained using the hybrid of block-pulse and Taylor series in [\[15\]](#page-18-14) with $N = 4$ and $M_3 = 7$ and the proposed method with $N = 4$ and *M* = 5 together with CPU time and the exact solution of $x(t)$ for $\frac{1}{4} \le t \le 1$ are given. In this table M_3 , denotes the order of Taylor series. In Tables [1](#page-14-0) and [2,](#page-14-1) the approximate value of $x(t)$ on [0, $\frac{1}{4}$] is equal to zero, which is the same as the exact solution.

t	Block-pulse and Chebyshev $N = 4$ and $M_1 = 7$	CPU	Block-pulse and Legendre $N = 4$ and $M_2 = 7$	CPU	
0.25	$\mathbf{0}$	0.235	θ	0.171	
0.40	0.27858404	0.219	0.27858404	0.157	
0.55	0.51352714	0.234	0.51352714	0.166	
0.70	0.65465130	0.203	0.65465130	0.180	
0.85	0.70892964	0.204	0.70892964	0.183	
1.00	0.71174280	0.202	0.71174280	0.190	

Table 1 The solution obtained using the hybrid of block-pulse with Chebyshev and Legendre polynomials for Example 1

Table 2 The solution obtained using the hybrid of block-pulse with Taylor and Bernoulli polynomials for Example 1

\mathfrak{t}	Block-pulse and Taylor $N = 4$ and $M_3 = 7$	CPU	Present method $N = 4$ and $M = 5$	CPU	Exact value
0.25	Ω	0.188	Ω	0.126	θ
0.40	0.278584046	0.172	0.278584047	0.132	0.278584047
0.55	0.513527142	0.172	0.513527141	0.141	0.513527141
0.70	0.654651249	0.171	0.654651311	0.135	0.654651311
0.85	0.708929644	0.187	0.708929636	0.151	0.708929636
1.00	0.711741219	0.218	0.711742826	0.168	0.711742826

Example 3

Consider the following piecewise constant delay system [\[5](#page-18-2)]

$$
\dot{x}(t) = t^2 x(t - a(t)) + u(t), \quad 0 \le t \le 1,
$$
\n(54)

$$
x(0) = 1,\tag{55}
$$

$$
u(t) = 2t + 1, \quad t \ge 0,
$$
\n(56)

$$
a(t) = \begin{cases} 0.2, & 0 \le t < 0.3, \\ 0.7, & 0.3 \le t \le 1 \end{cases}
$$
 (57)

The exact solution is

$$
x(t) = \begin{cases} 1+t+t^2, & 0 \le t < 0.2, \\ \frac{62341}{62500} + t + t^2 + \frac{7}{25}t^3 + \frac{3}{20}t^4 + \frac{1}{5}t^5, & 0.2 \le t < 0.3, \\ \frac{1006717}{1000000} + t + t^2, & 0.3 \le t < 0.7, \\ \frac{2720369}{3000000} + t + t^2 + \frac{79}{300}t^3 - \frac{1}{10}t^4 + \frac{1}{5}t^5, & 0.7 \le t < 0.9, \\ \frac{76256550101}{84000000000} + t + t^2 + \frac{693817}{3000000}t^3 + \frac{459}{40000}t^4 \\ + \frac{167}{5000}t^5 + \frac{7}{50}t^6 - \frac{11}{140}t^7 + \frac{1}{40}t^8, & 0.9 \le t \le 1. \end{cases}
$$

² Springer

By using Eqs. [\(21\)](#page-8-0) and [\(57\)](#page-14-2) we have $\gamma = 10$ and $\delta = 1$, so we choose $N = 10$. We also choose $M = 8$. Let

$$
x(t) = C_4^T B(t),\tag{58}
$$

by using Eq. [\(2\)](#page-3-4) we have

$$
u(t) = \begin{bmatrix} \frac{11}{10}, \frac{13}{10}, \frac{3}{2}, \frac{17}{10}, \frac{19}{10}, \frac{21}{10}, \frac{23}{10}, \frac{5}{2}, \frac{27}{10}, \frac{29}{10}, \frac{1}{5}, \frac{1}{
$$

Also we have

$$
x(t - a(t)) = C_4^T \underline{D}_4 B(t),
$$
\n(62)

where \underline{D}_4 is the delay operational matrix given in Eq. [\(20\)](#page-7-0) and is obtained from

$$
\underline{D}_5 = \underline{Q}_1 + \underline{Q}_2 + \cdots + \underline{Q}_{10},
$$

where \underline{Q}_i , $1 \le i \le 10$ are calculated similarly to Eq. [\(22\)](#page-8-1). For example \underline{Q}_3 is

$$
\underline{Q}_3 = I_9 \otimes \underline{S}_3,
$$

where I_9 is the 9 dimensional identity matrix and S_3 is a 10 \times 10 matrix given by

$$
\underline{S}_3 = \begin{bmatrix} 0 & 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 0 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \dots & \vdots \\ 0 & 0 & 0 & 0 & \dots & 0 \end{bmatrix}.
$$

By integrating Eq. [\(54\)](#page-14-2) from 0 to *t* and using Eqs. [\(55](#page-14-2)[–62\)](#page-15-0) we have

$$
C_4^T - Y_4^T = C_4^T \underline{D}_4 \widetilde{Y}_3 P + U_3^T P, \qquad (63)
$$

where Y_3 can be obtained similarly to Eq. [\(17\)](#page-6-1). By solving Eq. [\(63\)](#page-15-1) the same value as the spectral value of ϵ (c) would be altriard. exact value of $x(t)$ would be obtained.

Example 4

Consider the following piecewise constant delay system [\[26](#page-19-3)]

$$
\begin{aligned}\n\dot{x}(\tau) &= -5x(\tau) - 5x(\tau - a(\tau)) + 2u(\tau), & 0 \le \tau \le 2, \\
x(0) &= 1, \\
u(\tau) &= 1, & \tau \ge 0, \\
a(\tau) &= \begin{cases}\n0, & 0 \le \tau < 0.8, \\
0.3, & 0.8 \le \tau < 1.4, \\
0.6, & 1.4 \le \tau < 1.7, \\
0.9, & 1.7 \le \tau \le 2.\n\end{cases}\n\end{aligned}
$$

The exact solution is

$$
x(\tau) = \begin{cases} 0.2 + 0.8e^{-10\tau}, & 0 \le \tau < 0.8, \\ 0.2 + 0.8e^{3-10\tau} + 0.8e^{-4}(1 - e^{3})e^{-5\tau}, & 0.8 \le \tau < 1.1, \\ 0.2 + 0.8e^{6-10\tau} - 4e^{-2.5}(1 - e^{3})\tau e^{-5\tau} & 1.1 \le \tau < 1.4, \\ + 0.8(1 - e^{3})(6.5e^{-2.5} + e^{-4})e^{-5\tau}, \\ 0.2 + 0.8e^{9-10\tau} - 4e^{-1}(1 - e^{3})\tau e^{-5\tau} \\ + 0.8(1 - e^{3})(8e^{-1} - 0.5e^{-2.5} + e^{-4})e^{-5\tau}, & 1.4 \le \tau < 1.7, \\ 0.2 + 0.8e^{12-10\tau} - 4e^{0.5}(1 - e^{3})\tau e^{-5\tau} \\ + 0.8(1 - e^{3})(9.5e^{0.5} - 0.5e^{-1} - 0.5e^{-2.5} + e^{-4})e^{-5\tau}, & 1.7 \le \tau \le 2. \end{cases}
$$

By using transformation $\tau = 2t$, $0 \le t \le 1$, we have

$$
\dot{x}(t) = 2(-5x(t) - 5x(t - a(t)) + 2u(t)), \quad 0 \le t \le 1,
$$
\n(64)

$$
x(0) = 1,\tag{65}
$$

$$
u(t) = 1, \quad t \ge 0,\tag{66}
$$

$$
a(t) = \begin{cases} 0, & 0 \le t < 0.4, \\ 0.3, & 0.4 \le t < 0.7, \\ 0.6, & 0.7 \le t < 0.85, \\ 0.9, & 0.85 \le t \le 1. \end{cases}
$$
(67)

From Eqs. [\(21\)](#page-8-0) and [\(67\)](#page-16-0) we have $\gamma = 100$ and $\delta = 5$, so we select $N = 20$. We also choose *M* an arbitrary positive integer number. Let $M = 1$ and

$$
x(t) = C_5^T B(t),\tag{68}
$$

Table 3 The solution obtained using the hybrid of block-pulse with Chebyshev and Bernoulli polynomials for Example 4

Block-pulse and Chebyshev with $N = 20$ and	$E_{\rm max}$	CPU	Present method with $N = 20$ and	$E_{\rm max}$	CPU
$M_1 = 1$	$2.7199e - 1$	0.150	$M=1$	$1.0156e - 1$	0.124
$M_1 = 2$	$3.9068e - 2$	0.275	$M=2$	$4.7036e - 3$	0.196
$M_1 = 4$	$1.7269e - 4$	0.567	$M = 3$	$1.9876e - 5$	0.492
$M_1 = 6$	$3.5338e - 7$	0.952	$M=5$	$1.5674e - 7$	0.780
$M_1 = 8$	$3.8955e - 10$	1.582	$M=7$	$2.8654e - 10$	1.092

by using Eq. [\(2\)](#page-3-4) we have

$$
1 = [\underbrace{1, 1, \dots, 1}_{20}]^T B(t) = Y_5^T B(t), \qquad u(t) = [\underbrace{1, 1, \dots, 1}_{20}]^T B(t) = U_4^T B(t). \tag{69}
$$

Also we have

$$
x(t - a(t)) = C_5^T \underline{D}_5 B(t),
$$
\n(70)

where D_5 is the delay operational matrix given in Eq. [\(20\)](#page-7-0). By integrating Eq. [\(64\)](#page-16-0) from 0 to t and using Eqs. $(65-70)$ $(65-70)$ we have,

$$
C_5^T - Y_5^T = 2\left(-5C_5^T P - 5C_5^T \underline{D}_5 P + 2U_4^T P\right).
$$
 (71)

By solving Eq. (71) we obtain C_5 . In Table [3,](#page-16-1) we compare maximum error obtained using the proposed method with the hybrid of block-pulse and Chebyshev polynomials in [\[5\]](#page-18-2). The approximate solution of $x(t)$, obtained by the proposed method for $M = 1$ and 2, are shown in Figs. [1](#page-17-2) and [2](#page-17-3) respectively.

Fig. 1 Exact and approximate solution of $x(t)$ for $M = 1$

Fig. 2 Exact and approximate solution of $x(t)$ for $M = 2$

Conclusion

In the present work, hybrid of block-pulse functions and Bernoulli polynomials are used to solve delay systems. The problem has been reduced to a problem of solving a system of algebraic equations. The matrix P given in Eq. [\(16\)](#page-6-2) has a large number of zero elements and is sparse. Hence, the present method is very attractive and reduces the CPU time and computer memory. Illustrative examples are given to demonstrate the validity and applicability of the proposed method. It is noted that the exact solutions obtained in the examples 1, 2 and 4 can not be obtained either with piecewise constant basis functions nor with continuous basis functions.

References

- 1. Malek-Zavarei, M., Jamshidi, M.: Time Delay Systems: Analysis Optimization and Applications North-Holland Systems and Control Series. Elsevier, New York (1987)
- 2. Cai, G., Huang, J.: Optimal control method with time delay in control. J. Sound Vib. **251**, 383–394 (2002)
- 3. Zhang, J., Knospe, C.R., Tsiotras, P.: New results for the analysis of linear systems with time-invariant delays. Int. J. Robust Nonlinear Control **13**, 1149–1175 (2003)
- 4. Abdulaziz, O., Hashim, I., Chowdhury, M.S.H.: Solving variational problems by homotopy perturbation method. Int. J. Numer. Methods Eng. **75**, 709–721 (2008)
- 5. Marzban, H.R., Shahsiah, M.: Solution of piecewise constant delay systems using hybrid of block-pulse and Chebyshev polynomials. Optim. Control Appl. Methods **32**, 647–659 (2010)
- 6. Hwang, C., Shih, Y.P.: Optimal control of delay systems via block pulse functions. J. Optim. Theory Appl. **45**, 101–112 (1985)
- 7. Palanisamy, K.R., Rao, G.P.: Optimal control of linear systems with delays in state and control via Wash functions. IEE Proc. D. Control Theory Appl. **130**, 300–312 (1983)
- 8. Hwang, C., Shih, Y.P.: Laguerre series direct method for variational problems. J. Optim. Theory Appl. **39**, 143–149 (1983)
- 9. Horng, I.R., Chou, J.H.: Shifted Chebyshev direct method for solving variational problems. Int. J. Syst. Sci. **16**, 855–861 (1985)
- 10. Chang, R.Y., Wang, M.L.: Shifted Legendre direct method for variational problems. J. Optim. Theory App. **39**, 299–307 (1983)
- 11. Razzaghi, M.: Optimal control of linear time-varying systems via Fourier series. J. Optim. Theory Appl. **65**, 375–384 (1990)
- 12. Razzaghi, M., Razzaghi, M.: Instabilities in the solution of a heat conduction problem using Taylor series and alternative approaches. J. Frankl. Inst. **326**, 683–690 (1989)
- 13. Razzaghi, M., Marzban, H.R.: A hybrid domain analysis for systems with delays in state and control. Math. Probl. Eng. **7**, 337–353 (2001)
- 14. Marzban, H.R., Razzaghi, M.: Solution of time-varying delay systems by hybrid functions. Math. Comput. Simul. **64**, 597–607 (2004)
- 15. Marzban, H.R., Razzaghi, M.: Analysis of time-delay systems via hybrid of block-pulse functions and Taylor series. J. Vib. Control **11**, 1455–1468 (2005)
- 16. Wang, X.T., Li, Y.M.: Numerical solutions of integro differential systems by hybrid of general block-pulse functions and the second Chebyshev polynomials. Appl. Math. Comput. **209**, 266–272 (2009)
- 17. Singh, V.K., Pandey, R.K., Singh, S.: A stable algorithm for Hankel transforms using hybrid of block-pulse and Legendre polynomials. Comput. Phy. Commun. **181**, 1–10 (2010)
- 18. Costabile, F., Dellaccio, F., Gualtieri, M.T.: A new approach to Bernoulli polynomials. Rend. Mat. Ser. VII **26**, 1–12 (2006)
- 19. Arfken, G.: Mathematical Methods for Physicists, 3rd edn. Academic press, San Diego (1985)
- 20. Kreyszig, E.: Introductory Functional Analysis with Applications. Wiley, New York (1978)
- 21. Canuto, C., Hussaini, M.Y., Quarteroni, A., Zang, T.A.: Spectral Methods, Fundamentals in Single Domains. Springer, Berlin (2006)
- 22. Marzban, H.R., Tabrizidooz, H.R., Razzaghi, M.: A composite collocation method for the nonlinear mixed Volterra–Fredholm–Hammerstein integral equations. Commun. Nonlinear Sci. Numer. Simul. **16**, 1186–1194 (2011)
- 23. Mashayekhi, S., Ordokhani, Y., Razzaghi, M.: Hybrid functions approach for nonlinear constrained optimal control problems. Commun. Nonlinear Sci. Numer. Simul. **17**, 1831–1843 (2012)
- 24. Mashayekhi, S., Ordokhani, Y., Razzaghi, M.: Hybrid functions approach for optimal control of systems described by integro-differential equations. Appl. Math. Model. **37**, 3355–3368 (2013)
- 25. Lancaster, P.: Theory of Matrices. Academic Press, New York (1969)
- 26. Datta, K.B., Mohan, B.M.: Orthogonal Functions in Systems and Control. World Scientific, Singapore (1995)