



Trace representation of Legendre sequences over non-binary fields

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Abstract

For distinct odd primes N and p , we view the N -periodic binary Legendre sequence as a p -ary sequence and present its trace representation via trace functions over \mathbb{F}_p . We use a skill to calculate the Mattson–Solomon polynomials of Legendre sequences and then describe the Mattson–Solomon polynomials by means of trace functions over \mathbb{F}_p .

Keywords Legendre sequence · Trace representation · Mattson–Solomon polynomial

Mathematics Subject Classification 94A55 · 94A60 · 65C10

1 Introduction

For an odd prime number N , the N -periodic Legendre sequence is defined as

$$s_u = \begin{cases} \frac{1+(\frac{u}{N})}{2}, & \text{if } \gcd(u, N) = 1, \\ 0, & \text{otherwise,} \end{cases} \quad u \geq 0, \quad (1)$$

where $(\frac{\cdot}{N})$ is the Legendre symbol. Let g be a (fixed) primitive root modulo N , one can define the *cyclotomic classes*

$$D_0 = \{g^{2k} \pmod{N} : 0 \leq k < (N-1)/2\}$$

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and

$$D_1 = gD_0 = \{g^{2k+1} \pmod N : 0 \leq k < (N - 1)/2\}.$$

Then we get an equivalent definition of the Legendre sequence

$$s_u = \begin{cases} 0, & \text{if } u \pmod N \in D_1 \cup \{0\}, \\ 1, & \text{if } u \pmod N \in D_0, \end{cases} \quad u \geq 0. \tag{2}$$

The Legendre sequences (s_u) have been extensively studied in the literature. They have strong pseudorandomness properties: equidistribution, optimal correlation, high linear complexity, etc., see [3,4,6,9,10,14,18,23]. Aly, Winterhof [1] studied the k -error linear complexity (over \mathbb{F}_N) by viewing the N -periodic (s_u) as a sequence over \mathbb{F}_N .

In particular, for certain applications to coding theory, some binary sequences are discussed over different finite fields (not in \mathbb{F}_2) [7,8]. Partially motivated by the study, Wang et al considered the N -periodic Legendre sequence (s_u) in \mathbb{F}_p , where p is an odd prime (or a prime-power) with $\gcd(p, N) = 1$, and investigated the linear complexity and minimal polynomials over \mathbb{F}_p in [11,21,22]. Certain work had actually been done by He in [13]. In this work, we will continue this project to investigate the trace representation of N -periodic Legendre sequence (s_u) in \mathbb{F}_p (not in \mathbb{F}_2). We should remark that, the trace representation of (s_u) of Mersenne prime period and of any prime period have been described via trace functions from \mathbb{F}_{2^n} to \mathbb{F}_2 , where n is the order of 2 modulo N , by No et al in [19] and by Kim et al in [15], sequentially. Some special cases have been studied in [20] recently.

We will compute the *Mattson–Solomon polynomial* (see definition below) of (s_u) and present the trace representation by using trace functions over \mathbb{F}_p . For any N -periodic p -ary sequence (t_u) , there always exists a polynomial $G(X)$ defined over finite fields of characteristic p such that

$$t_u = G(\beta^u), \quad u \geq 0,$$

where β is an N th root of unity in an extension field of \mathbb{F}_p . $G(X)$ is unique if its degree is smaller than N , see [16]. Such $G(X)$ is called the *Mattson–Solomon polynomial* of (t_u) in coding theory [17]. Dai et al called $G(X)$ as a *defining polynomial* and $(G(X), \beta)$ as the *defining pair* of (t_u) in [5], where they discussed trace representation and linear complexity of certain binary sequences.

Throughout the work, we always let p be an odd prime and co-prime to N , the period of Legendre sequences.

2 Mattson–Solomon polynomials

Define polynomials

$$d_l(X) = \sum_{u \in D_l} X^u \in \mathbb{F}_p[X], \quad l = 0, 1.$$

We need the following technical lemma.

Lemma 1 *Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p . For any fixed pair of integers i, j with $0 \leq i, j < 2$, we have*

$$d_i(\beta)d_j(\beta) + d_{i+1}(\beta)d_{j+1}(\beta) + \frac{N-1}{2} = \begin{cases} N, & \text{if } \frac{N-1}{2} + i - j \equiv 0 \pmod{2}, \\ 0, & \text{otherwise.} \end{cases}$$

Here and hereafter, the subscript of d is performed modular 2.

Proof We calculate

$$\begin{aligned} d_i(\beta)d_j(\beta) + d_{i+1}(\beta)d_{j+1}(\beta) &= \sum_{k=0}^1 \sum_{u \in D_0} \beta^{ug^{i+k}} \sum_{v \in D_0} \beta^{vg^{j+k}} \\ &= \sum_{k=0}^1 \sum_{u \in D_0} \beta^{ug^{i+k}} \sum_{w \in D_0} \beta^{uwg^{j+k}} \\ &\quad (\text{we use } v = uw) \\ &= \sum_{k=0}^1 \sum_{u \in D_0} \sum_{w \in D_0} \beta^{ug^{j+k}(g^{i-j}+w)} \\ &= \sum_{w \in D_0} \sum_{k=0}^1 \sum_{z \in D_{j+k}} \gamma_w^z \\ &\quad (\text{we use } z = ug^{j+k}, \gamma_w = \beta^{g^{i-j}+w}) \\ &= \sum_{w \in D_0} \sum_{z=1}^{N-1} \gamma_w^z. \end{aligned}$$

Let $\text{ord}(\gamma_w)$ denote the order of γ_w . We note that $\text{ord}(\gamma_w) | N$ since β is a primitive N th root of unity. If $\text{ord}(\gamma_w) = N$, then we have

$$\sum_{z=1}^{N-1} \gamma_w^z = \sum_{z=0}^{N-1} \gamma_w^z - 1 = \frac{1 - \gamma_w^N}{1 - \gamma_w} - 1 = -1 \in \mathbb{F}_p.$$

If $\text{ord}(\gamma_w) = 1$, then we have

$$\sum_{z=1}^{N-1} \gamma_w^z = N - 1 \in \mathbb{F}_p.$$

Now we need to determine the number of $w \in D_0$ with $\text{ord}(\gamma_w) = 1$ and the number of $w \in D_0$ with $\text{ord}(\gamma_w) = N$.

We have $\text{ord}(\gamma_w) = 1$ if and only if $g^{i-j} + w \equiv 0 \pmod{N}$, which is equivalent to $w \equiv g^{(N-1)/2+i-j} \pmod{N}$. This implies that $2 \mid ((N-1)/2 + i - j)$ since $w \in D_0$. That is to say, there exists an $w \in D_0$ such that $g^{i-j} + w \equiv 0 \pmod{N}$, which holds if and only if $2 \mid ((N-1)/2 + i - j)$. In this case w is unique. We conclude that if $2 \mid ((N-1)/2 + i - j)$, then there are $(N-1)/2 - 1$ elements $w \in D_0$ such that $\text{ord}(\gamma_w) = N$ and one $w \in D_0$ such that $\text{ord}(\gamma_w) = 1$, while if $2 \nmid ((N-1)/2 + i - j)$, all $w \in D_0$ satisfy $\text{ord}(\gamma_w) = N$.

Putting everything together, we derive

$$d_i(\beta)d_j(\beta) + d_{i+1}(\beta)d_{j+1}(\beta) = \begin{cases} \frac{N+1}{2}, & \text{if } 2 \mid \left(\frac{N-1}{2} + i - j\right), \\ -\frac{N-1}{2}, & \text{otherwise.} \end{cases}$$

This completes the proof. □

Theorem 1 *Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p . Then the Mattson–Solomon polynomial of (s_u) defined in Eq. (1) or Eq. (2) is*

$$G(X) = N^{-1} \left(d_0(\beta)d_0(X) + d_1(\beta)d_1(X) + \frac{N-1}{2} \right)$$

if $N \equiv 1 \pmod{4}$, and otherwise

$$G(X) = N^{-1} \left(d_0(\beta)d_1(X) + d_1(\beta)d_0(X) + \frac{N-1}{2} \right).$$

Proof We get from Lemma 1 that

$$(d_0(\beta))^2 + (d_1(\beta))^2 + \frac{N-1}{2} = \begin{cases} N, & \text{if } N \equiv 1 \pmod{4}, \\ 0, & \text{if } N \equiv -1 \pmod{4}, \end{cases}$$

and

$$2d_0(\beta)d_1(\beta) + \frac{N-1}{2} = \begin{cases} 0, & \text{if } N \equiv 1 \pmod{4}, \\ N, & \text{if } N \equiv -1 \pmod{4}. \end{cases}$$

Note that $d_i(\beta^u) = d_{i+j}(\beta)$ if $u \in D_j$, where $i, j \in \{0, 1\}$ and the subscript of d is performed modulo 2. Now, we discuss the Mattson–Solomon polynomial of (s_u) .

Case 1 $N \equiv 1 \pmod{4}$.

For $u \in D_0$, we have

$$\begin{aligned} G(\beta^u) &= N^{-1} \left(d_0(\beta)d_0(\beta^u) + d_1(\beta)d_1(\beta^u) + \frac{N-1}{2} \right) \\ &= N^{-1} \left((d_0(\beta))^2 + (d_1(\beta))^2 + \frac{N-1}{2} \right) \\ &= N^{-1} \cdot N = 1 = s_u. \end{aligned}$$

For $u \in D_1$, we have

$$\begin{aligned} G(\beta^u) &= N^{-1} \left(d_0(\beta)d_0(\beta^u) + d_1(\beta)d_1(\beta^u) + \frac{N-1}{2} \right) \\ &= N^{-1} \left(d_0(\beta)d_1(\beta) + d_1(\beta)d_0(\beta) + \frac{N-1}{2} \right) \\ &= N^{-1} \left(2d_0(\beta)d_1(\beta) + \frac{N-1}{2} \right) \\ &= N^{-1} \cdot 0 = 0 = s_u. \end{aligned}$$

For $u = 0$, we note that

$$d_0(1) = d_1(1) = \frac{N-1}{2},$$

and

$$d_0(\beta) + d_1(\beta) = \sum_{u=1}^{N-1} \beta^u = \sum_{u=0}^{N-1} \beta^u - 1 = \frac{1-\beta^N}{1-\beta} - 1 = -1.$$

Then, we get

$$\begin{aligned} G(\beta^0) &= N^{-1} \left(d_0(\beta)d_0(1) + d_1(\beta)d_1(1) + \frac{N-1}{2} \right) \\ &= N^{-1} \left(-\frac{N-1}{2} + \frac{N-1}{2} \right) = 0 = s_u. \end{aligned}$$

Putting everything together, we derive that

$$G(X) = N^{-1} \left(d_0(\beta)d_0(X) + d_1(\beta)d_1(X) + \frac{N-1}{2} \right)$$

is the Mattson–Solomon polynomial of (s_u) when $N \equiv 1 \pmod{4}$.

Case 2 $N \equiv -1 \pmod{4}$.

It can be verified in a similar way. □

Now we further consider the values of $d_0(\beta)$ and $d_1(\beta)$ in Theorem 1.

Lemma 2 *Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p and p a quadratic residue class modulo N (i.e., $p \in D_0$). If N satisfies one of the following two conditions*

- (1) $N \equiv 1 \pmod{4}$ and $N \equiv 1 \pmod{p}$,
- (2) $N \equiv -1 \pmod{4}$ and $N \equiv -1 \pmod{p}$,

then we have

$$\begin{cases} d_0(\beta) = 0, \\ d_1(\beta) = -1, \end{cases} \quad \text{or} \quad \begin{cases} d_0(\beta) = -1, \\ d_1(\beta) = 0, \end{cases}$$

and otherwise we have

$$d_0(\beta), d_1(\beta) \in \mathbb{F}_p \setminus \{0\},$$

which means that both $d_0(\beta)$ and $d_1(\beta)$ are non-zero.

Proof Firstly, we have $d_0(\beta) = d_0(\beta^p) = (d_0(\beta))^p$ since $p \in D_0$. That is to say $d_0(\beta) \in \mathbb{F}_p$. Similarly, we have $d_1(\beta) \in \mathbb{F}_p$.

For $N \equiv 1 \pmod{4}$, we see that in the proof of Theorem 1

$$(d_0(\beta))^2 + (d_1(\beta))^2 = \frac{N+1}{2} = 1$$

if and only if $N \equiv 1 \pmod{p}$. So together with $d_0(\beta) + d_1(\beta) = -1$, we get for $N \equiv 1 \pmod{p}$

$$2d_0(\beta)d_1(\beta) = 0,$$

which derives that either $d_0(\beta)$ or $d_1(\beta)$ is zero. Then, it is easy to get that

$$\begin{cases} d_0(\beta) = 0, \\ d_1(\beta) = -1, \end{cases} \quad \text{or} \quad \begin{cases} d_0(\beta) = -1, \\ d_1(\beta) = 0. \end{cases}$$

For $N \equiv -1 \pmod{4}$, we get similarly $2d_0(\beta)d_1(\beta) = \frac{N+1}{2} = 0$ if and only if $N \equiv -1 \pmod{p}$ and then the result is derived.

The proof above also tells us that

$$2d_0(\beta)d_1(\beta) \neq 0$$

for other N . □

Lemma 3 Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p and p a quadratic non-residue class modulo N (i.e., $p \in D_1$). Then both $d_0(\beta)$ and $d_1(\beta)$ are non-zero.

Proof Since $p \in D_1$, we have for $i = 0, 1$

$$(d_i(\beta))^p = d_i(\beta^p) = d_{i+1}(\beta) = -1 - d_i(\beta),$$

which indicates both $d_0(\beta)$ and $d_1(\beta)$ are non-zero. □

From Theorem 1 and Lemmas 2 and 3, we immediately get the following results.

Theorem 2 Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p , p a quadratic residue class modulo N (i.e., $p \in D_0$) and (s_u) defined in Eq. (1) or Eq. (2).

- (1) For N satisfying $N \equiv 1 \pmod{4}$ and $N \equiv 1 \pmod{p}$, if we suppose $d_0(\beta) = 0$ (of course we can also suppose $d_1(\beta) = 0$), then the Mattson–Solomon polynomial of (s_u) is

$$G(X) = -N^{-1}d_1(X).$$

- (2) For N satisfying $N \equiv -1 \pmod{4}$ and $N \equiv -1 \pmod{p}$, if we suppose $d_0(\beta) = 0$ (of course we can also suppose $d_1(\beta) = 0$), then the Mattson–Solomon polynomial of (s_u) is

$$G(X) = -N^{-1}d_1(X) + N^{-1}.$$

- (3) For other N , the Mattson–Solomon polynomial of (s_u) is

$$G(X) = N^{-1}\left(\rho d_1(X) - (1 + \rho)d_0(X) + \frac{N-1}{2}\right).$$

where $\rho = d_0(\beta)$ and $\rho(1 + \rho) \neq 0$.

Theorem 3 Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p and p a quadratic non-residue class modulo N (i.e., $p \in D_1$). Then the Mattson–Solomon polynomial of (s_u) defined in Eq. (1) or Eq. (2) is

$$G(X) = N^{-1}\left(\rho d_1(X) - (1 + \rho)d_0(X) + \frac{N-1}{2}\right),$$

where $\rho = d_0(\beta)$ and $\rho(1 + \rho) \neq 0$.

3 Trace representation

In this section, we describe the trace representation of (s_u) . For $n|m$, the trace function from finite field \mathbb{F}_{p^m} to \mathbb{F}_{p^n} is defined as

$$\text{Tr}_n^m(X) = X + X^{p^n} + X^{p^{2n}} + \cdots + X^{p^{(m/n-1)n}}.$$

The trace functions play an important role in sequences design [12].

Theorem 4 Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p , p a quadratic residue class modulo N (i.e., $p \in D_0$) and (s_u) defined in Eq. (1) or Eq. (2). Let ℓ be the order of p modulo N .

(1) For N satisfying $N \equiv 1 \pmod{4}$ and $N \equiv 1 \pmod{p}$, if we suppose $d_0(\beta) = 0$, then the trace representation of (s_u) is

$$s_u = -N^{-1} \sum_{j=0}^{\frac{N-1}{2\ell}-1} \text{Tr}_1^\ell \left(\beta^{g^{2j+1}} \right).$$

(2) For N satisfying $N \equiv -1 \pmod{4}$ and $N \equiv -1 \pmod{p}$, if we suppose $d_0(\beta) = 0$, then the trace representation of (s_u) is

$$s_u = -N^{-1} \sum_{j=0}^{\frac{N-1}{2\ell}-1} \text{Tr}_1^\ell \left(\beta^{g^{2j+1}} \right) + N^{-1}.$$

(3) For other N , the trace representation of (s_u) is

$$s_u = N^{-1} \left(\rho \sum_{j=0}^{\frac{N-1}{2\ell}-1} \text{Tr}_1^\ell \left(\beta^{g^{2j+1}} \right) - (1 + \rho) \sum_{j=0}^{\frac{N-1}{2\ell}-1} \text{Tr}_1^\ell \left(\beta^{ug^{2j}} \right) + \frac{N-1}{2} \right).$$

where $\rho = d_0(\beta)$ and $\rho(1 + \rho) \neq 0$.

Proof To get the trace presentation of $s(u)$, we only need to describe $d_0(X)$ and $d_1(X)$ in Theorem 2 using trace functions.

Let U be set generated by p modulo N , i.e.,

$$U = \langle p \rangle = \{p^k \pmod{N} : 0 \leq k < \ell\}.$$

Since $p \in D_0$, we see that U is a subgroup of D_0 (under the multiplication). Then D_0, D_1 can be written as the union

$$D_0 = \bigcup_{k=0}^{\frac{N-1}{2\ell}-1} g^{2k}U, \quad D_1 = \bigcup_{k=0}^{\frac{N-1}{2\ell}-1} g^{2k+1}U.$$

Write polynomial

$$u(X) = \sum_{u \in U} X^u.$$

Using the fact that

$$\text{Tr}_1^\ell(X) = X + X^p + X^{p^2} + \dots + X^{p^{\ell-1}} \equiv u(X) \pmod{X^N - 1},$$

we derive

$$d_0(X) = \sum_{j=0}^{\frac{N-1}{2\ell}-1} u(X^{g^{2j}}) \equiv \sum_{j=0}^{\frac{N-1}{2\ell}-1} \text{Tr}_1^\ell(X^{g^{2j}}) \pmod{X^N - 1}$$

and

$$d_1(X) = \sum_{j=0}^{\frac{N-1}{2\ell}-1} u(X^{g^{2j+1}}) \equiv \sum_{j=0}^{\frac{N-1}{2\ell}-1} \text{Tr}_1^\ell(X^{g^{2j+1}}) \pmod{X^N - 1}.$$

Then, replacing $d_0(X)$ and $d_1(X)$ in Theorem 2 and noting that $s_u = G(\beta^u)$, we finish the proof. \square

Theorem 5 *Let β be a primitive N th root of unity in an extension field of \mathbb{F}_p and p a quadratic non-residue class modulo N (i.e., $p \in D_1$). Let ℓ be the order of p modulo N . Then, the trace representation of (s_u) defined in Eq. (1) or Eq. (2) is*

$$s_u = N^{-1} \left(\rho \sum_{j=0}^{\frac{N-1}{\ell}-1} \text{Tr}_2^\ell(\beta^{ug^{2j+1}}) - (1 + \rho) \sum_{j=0}^{\frac{N-1}{\ell}-1} \text{Tr}_2^\ell(\beta^{ug^{2j}}) + \frac{N-1}{2} \right).$$

where $\rho = d_0(\beta)$ and $\rho(1 + \rho) \neq 0$.

Proof The proof is similar to that of Theorem 4. From the condition $p \in D_1$, we see that $p^2 \in D_0$ and the order of p^2 modulo N is $\frac{\ell}{2}$. We remark here that ℓ is even. Indeed, if $p \equiv g^{2k+1} \pmod{N}$ for some k , we get $p^\ell \equiv g^{(2k+1)\ell} \equiv 1 \pmod{N}$, which indicates that $(N-1) | \ell(2k+1)$. Then ℓ is even since $N-1$ is even.

Now write

$$V = \langle p^2 \rangle = \left\{ p^{2k} \pmod{N} : 0 \leq k < \frac{\ell}{2} \right\}.$$

Then V is a subgroup of D_0 and D_0, D_1 can be represented as

$$D_0 = \bigcup_{k=0}^{\frac{N-1}{\ell}-1} g^{2k} V, \quad D_1 = \bigcup_{k=0}^{\frac{N-1}{\ell}-1} g^{2k+1} V.$$

Similar to the proof of Theorem 4, we have

$$\text{Tr}_2^\ell(X) = X + X^{p^2} + X^{p^{2 \times 2}} + \dots + X^{p^{2 \times (\frac{\ell}{2}-1)}} \equiv v(X) \pmod{X^N - 1},$$

where $v(X) = \sum_{u \in V} X^u$. Then we describe $d_0(X)$ and $d_1(X)$ as follows

$$d_0(X) = \sum_{j=0}^{\frac{N-1}{\ell}-1} v\left(X^{g^{2j}}\right) \equiv \sum_{j=0}^{\frac{N-1}{\ell}-1} \text{Tr}_2^\ell\left(X^{g^{2j}}\right) \pmod{X^N - 1}$$

and

$$d_1(X) = \sum_{j=0}^{\frac{N-1}{\ell}-1} v\left(X^{g^{2j+1}}\right) \equiv \sum_{j=0}^{\frac{N-1}{\ell}-1} \text{Tr}_2^\ell\left(X^{g^{2j+1}}\right) \pmod{X^N - 1}.$$

Then, replacing $d_0(X)$ and $d_1(X)$ in Theorem 3 and noting that $s_u = G(\beta^u)$, we finish the proof. □

4 Remarks and conclusions

In this work, we view N -periodic Legendre sequences in \mathbb{F}_2 as in \mathbb{F}_p and considered their trace representation by calculating Mattson–Solomon polynomials. The results extended the early work of No et al and Kim et al on trace representation over \mathbb{F}_2 .

The way in this work also can be used to consider the trace representation if we put N -periodic Legendre sequences in rings, for example in \mathbb{Z}_4 , the residue class ring modulo 4.

We finally remark that, there is a relationship between Mattson–Solomon polynomials of prime periodic sequences and their linear complexity[12, Theorem 6.3]. The *linear complexity* $LC(t_u)$ of an N -period sequence (t_u) over \mathbb{F}_p is the least order L of a linear recurrence relation over \mathbb{F}_p

$$t_{u+L} + c_1 t_{u+L-1} + \dots + c_{L-1} t_{u+1} + c_L t_u = 0 \quad \text{for } u \geq 0,$$

where $c_1, c_2, \dots, c_L \in \mathbb{F}_p$. By [2], $LC(t_u)$ equals the number of nonzero coefficients of the Mattson–Solomon polynomial $G(x)$ of degree $< N$. So from Theorems 2 and 3, we immediately derive the linear complexity of N -periodic Legendre sequences studied in [13,22].

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