

# A Farkas-type theorem for interval linear inequalities

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**Abstract** We describe a Farkas-type condition for strong solvability of interval linear inequalities. The result is used to derive several descriptions of the set of strong solutions and to show that this set forms a convex polytope.

**Keywords** Linear inequalities · Interval data · Strong solvability · Farkas-type theorem

## 1 Introduction

The famous theorem (often called a lemma) proved by Julius Farkas 111 years ago [1] asserts that a system of linear equations

$$Ax = b \tag{1}$$

(with  $A \in \mathbb{R}^{m \times n}$  and  $b \in \mathbb{R}^m$ ) has a nonnegative solution if and only if for each  $p \in \mathbb{R}^m$ ,  $A^T p \geq 0$  implies  $b^T p \geq 0$ . Several consequences for systems of other types than (1) can be drawn from this result. For instance, since a system of linear inequalities

$$Ax \leq b \tag{2}$$

has a solution if and only if the system of linear equations

$$Ax_1 - Ax_2 + x_3 = b$$

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has a nonnegative solution, from the Farkas theorem we obtain that (2) is solvable if and only if for each  $p \geq 0$ ,  $A^T p = 0$  implies  $b^T p \geq 0$ .

Rohn [4] and recently independently Karademir and Prokopyev [3] (see also [6]) formulated a Farkas-type theorem for interval linear equations. Given an  $m \times n$  interval matrix  $A = [\underline{A}, \overline{A}] = [A_c - \Delta, A_c + \Delta] = \{ A \mid A_c - \Delta \leq A \leq A_c + \Delta \}$  (where  $A_c = (\underline{A} + \overline{A})/2$  and  $\Delta = (\overline{A} - \underline{A})/2$  are the midpoint matrix and the radius matrix of  $A$ , respectively) and an interval  $m$ -vector  $b = [\underline{b}, \overline{b}] = [b_c - \delta, b_c + \delta] = \{ b \mid b_c - \delta \leq b \leq b_c + \delta \}$  (where again  $b_c = (\underline{b} + \overline{b})/2$  and  $\delta = (\overline{b} - \underline{b})/2$ ), they proved that the system

$$Ax = b$$

has a nonnegative solution for each  $A \in A, b \in b$  if and only if for each  $p, A_c^T p + \Delta^T |p| \geq 0$  implies  $b_c^T p - \delta^T |p| \geq 0$ . The presence of absolute values renders, however, this result practically useful for small values of  $m$  only: checking this Farkas-type condition is NP-hard [4].

In this paper we formulate a Farkas-type condition for interval linear inequalities. Given an interval matrix  $A$  and an interval vector  $b$  as above, we are interested in solvability of each system

$$Ax \leq b \tag{3}$$

with data satisfying

$$A \in A, b \in b. \tag{4}$$

This type of solvability is called *strong solvability* of a formally written system of interval linear inequalities  $Ax \leq b$ , see Chapter 2 in [2] for a survey of results. In Theorem 1 we prove a Farkas-type condition for strong solvability which we then use to obtain another proof of the result by Rohn and Kreslová [7] saying that if  $Ax \leq b$  is strongly solvable, then all the systems (3), (4) have a common solution which is called a *strong* (it could also be termed “universal”) *solution* of  $Ax \leq b$ . As the main result of this paper we give in Theorem 4 four alternative descriptions of the set of strong solutions, and in the concluding Sect. 5 we show an interconnection between strong solvability of interval linear equations and that of interval linear inequalities.

## 2 The Farkas-type theorem

We have the following Farkas-type characterization.

**Theorem 1** *A system  $Ax \leq b$  is strongly solvable if and only if for each  $p \geq 0, \underline{A}^T p \leq 0 \leq \overline{A}^T p$  implies  $\underline{b}^T p \geq 0$ .*

*Proof* “If”: Let for each  $p \geq 0, \underline{A}^T p \leq 0 \leq \overline{A}^T p$  imply  $\underline{b}^T p \geq 0$ . Assume to the contrary that some system (3) with data satisfying (4) is not solvable. Then according

to the above-quoted Farkas condition for (3) there exists a vector  $p_0 \geq 0$  such that  $A^T p_0 = 0$  and  $b^T p_0 < 0$ . But from  $\underline{A} \leq A \leq \bar{A}$  and  $p_0 \geq 0$  we obtain  $\underline{A}^T \leq A^T \leq \bar{A}^T$  and  $\underline{A}^T p_0 \leq A^T p_0 \leq \bar{A}^T p_0$ , hence  $\underline{A}^T p_0 \leq 0 \leq \bar{A}^T p_0$ . Thus our assumption implies  $\underline{b}^T p_0 \geq 0$ , but we also have  $\underline{b}^T p_0 \leq b^T p_0 < 0$ , which is a contradiction.

“Only if”: Let each system (3) with data satisfying (4) be solvable. Assume to the contrary that the Farkas-type condition does not hold, i.e., that there exists a vector  $p_1 \geq 0$  such that

$$\underline{A}^T p_1 \leq 0 \leq \bar{A}^T p_1$$

and

$$\underline{b}^T p_1 < 0. \tag{5}$$

For each  $i = 1, \dots, n$  define a real function of one real variable by

$$f_i(t) = (t\underline{A}_{\bullet i} + (1 - t)\bar{A}_{\bullet i})^T p_1.$$

Then  $f_i(0) = (\bar{A}_{\bullet i})^T p_1 = (\bar{A}^T p_1)_i \geq 0$  and  $f_i(1) = (\underline{A}_{\bullet i})^T p_1 = (\underline{A}^T p_1)_i \leq 0$ , hence by continuity there exists a  $t_i \in [0, 1]$  such that  $f_i(t_i) = 0$ . Now define  $A$  columnwise by

$$A_{\bullet i} = t_i \underline{A}_{\bullet i} + (1 - t_i) \bar{A}_{\bullet i} \quad (i = 1, \dots, n).$$

Then  $A_{\bullet i}$ , as a convex combination of  $\underline{A}_{\bullet i}$  and  $\bar{A}_{\bullet i}$ , belongs to  $[\underline{A}_{\bullet i}, \bar{A}_{\bullet i}]$  for each  $i$ , hence  $A \in \mathcal{A}$ . Moreover, from the definition of  $t_i$  we have  $(A^T p_1)_i = (A_{\bullet i})^T p_1 = f_i(t_i) = 0$  for each  $i$ , hence  $A^T p_1 = 0$  which together with (5) implies that the system  $Ax \leq \underline{b}$  has no solution in contradiction to our assumption.  $\square$

### 3 Strong solvability of interval linear inequalities

Let us now have a closer look at the Farkas-type condition of Theorem 1. If we write it in a slightly different form: for each  $p$ ,

$$p \geq 0, \bar{A}^T p \geq 0, -\underline{A}^T p \geq 0 \text{ implies } \underline{b}^T p \geq 0,$$

we can see that it is just the Farkas condition for nonnegative solvability of the system

$$\bar{A}x_1 - \underline{A}x_2 \leq \underline{b}. \tag{6}$$

In this way we have found another proof of a theorem by Rohn and Kreslová [7] formulated, unlike the Farkas condition, in primal terms:

**Theorem 2** *A system of interval linear inequalities  $Ax \leq b$  is strongly solvable if and only if the system (6) has a nonnegative solution.*

This result has a nontrivial consequence which is also contained in [7]: namely, the existence of strong solutions.

**Theorem 3** *If a system of interval linear inequalities  $Ax \leq b$  is strongly solvable, then there exists an  $x_0$  such that*

$$Ax_0 \leq b \tag{7}$$

holds for each  $A \in \mathbf{A}$  and  $b \in \mathbf{b}$ .

*Proof* According to Theorem 2, strong solvability of  $Ax \leq b$  implies nonnegative solvability of (6). If  $x_1 \geq 0$  and  $x_2 \geq 0$  solve (6), then for each  $A \in \mathbf{A}$  and  $b \in \mathbf{b}$  we have  $A(x_1 - x_2) \leq \overline{A}x_1 - \underline{A}x_2 \leq \underline{b} \leq b$ , so that  $x_0 = x_1 - x_2$  possesses the required property.  $\square$

In order to be able to convert solutions of (6) into complementary ones, we need the following auxiliary result. Minimum of two vectors is understood entrywise.

**Proposition 1** *If  $x_1, x_2$  is a nonnegative solution of (6), then for each  $d$  with*

$$0 \leq d \leq \min\{x_1, x_2\},$$

$x'_1 = x_1 - d, x'_2 = x_2 - d$  is also a nonnegative solution of (6).

*Proof* Since  $d \leq \min\{x_1, x_2\} \leq x_1$ , we have  $x'_1 \geq 0$ , and similarly  $x'_2 \geq 0$ . Next,  $\overline{A}x'_1 - \underline{A}x'_2 = \overline{A}x_1 - \underline{A}x_2 + (\overline{A} - \underline{A})d \leq \overline{A}x_1 - \underline{A}x_2 \leq \underline{b}$ , so that  $x'_1, x'_2$  solve (6).  $\square$

A vector  $x_0$  satisfying (7) for each  $A \in \mathbf{A}, b \in \mathbf{b}$  is called a *strong solution* of  $Ax \leq b$ . Denote by  $X_S(\mathbf{A}, \mathbf{b})$  the set of strong solutions. For  $x = (x_i)_{i=1}^n$  let

$$\begin{aligned} x^+ &= (\max\{x_i, 0\})_{i=1}^n, \\ x^- &= (\max\{-x_i, 0\})_{i=1}^n, \end{aligned}$$

and

$$|x| = (|x_i|)_{i=1}^n,$$

then  $x = x^+ - x^-, |x| = x^+ + x^-, (x^+)^T x^- = 0, x^+ \geq 0$  and  $x^- \geq 0$ . The following theorem brings several alternative descriptions of the set of strong solutions.

**Theorem 4** *We have*

$$X_S(\mathbf{A}, \mathbf{b}) = \{x_1 - x_2 \mid \overline{A}x_1 - \underline{A}x_2 \leq \underline{b}, x_1 \geq 0, x_2 \geq 0\} \tag{8}$$

$$= \{x \mid \overline{A}x^+ - \underline{A}x^- \leq \underline{b}\} \tag{9}$$

$$= \{x \mid A_c x + \Delta|x| \leq \underline{b}\} \tag{10}$$

$$= \{x \mid A_c x + \Delta t \leq \underline{b}, -t \leq x \leq t\}. \tag{11}$$

*Proof* The equality (8) was proved in [7, Prop. 1]. Denote by  $X_1, X_2, X_3, X_4$  the right-hand side sets in (8)–(11), respectively. We shall prove that  $X_1 \subseteq X_2 \subseteq X_3 \subseteq X_4 \subseteq X_1$ .

“ $X_1 \subseteq X_2$ ”: Let  $x_1 \geq 0, x_2 \geq 0$  satisfy  $\overline{A}x_1 - \underline{A}x_2 \leq \underline{b}$ . Put  $d = \min\{x_1, x_2\}$  (entrywise) and  $x = x_1 - x_2$ . Then  $d \geq 0, x^+ = x_1 - d, x^- = x_2 - d$ , and  $\overline{A}x^+ - \underline{A}x^- \leq \underline{b}$  by Proposition 1 which proves that  $X_1 \subseteq X_2$ .

“ $X_2 \subseteq X_3$ ”: From  $x = x^+ - x^-, |x| = x^+ + x^-$  we have  $x^+ = (|x| + x)/2, x^- = (|x| - x)/2$ . Substituting these quantities into  $\overline{A}x^+ - \underline{A}x^- \leq \underline{b}$  leads to  $A_c x + \Delta|x| \leq \underline{b}$ .

“ $X_3 \subseteq X_4$ ”: If  $A_c x + \Delta|x| \leq \underline{b}$ , then for  $t = |x|$  we have  $A_c x + \Delta t \leq \underline{b}$  and  $-t \leq x \leq t$ .

“ $X_4 \subseteq X_1$ ”: If  $A_c x + \Delta t \leq \underline{b}$  and  $-t \leq x \leq t$ , then  $|x| \leq t$  and nonnegativity of  $\Delta$  implies  $A_c x + \Delta|x| \leq A_c x + \Delta t \leq \underline{b}$ . Substituting  $x = x^+ - x^-$  and  $|x| = x^+ + x^-$ , we obtain  $\overline{A}x^+ - \underline{A}x^- \leq \underline{b}$ , hence  $x_1 = x^+, x_2 = x^-$  satisfy  $x_1 - x_2 = x, x_1 \geq 0, x_2 \geq 0$  and  $\overline{A}x_1 - \underline{A}x_2 \leq \underline{b}$ . □

In the proof we have relied on the equality (8) proved in [7]. But we can also pursue another path starting from the following general-purpose theorem.

**Theorem 5** *Let  $A = [A_c - \Delta, A_c + \Delta]$  be an  $m \times n$  interval matrix,  $b = [b_c - \delta, b_c + \delta]$  an interval  $m$ -vector, and let  $x \in \mathbb{R}^n$ . Then we have*

$$\{Ax - b \mid A \in \mathbf{A}, b \in \mathbf{b}\} = [A_c x - b_c - (\Delta|x| + \delta), A_c x - b_c + (\Delta|x| + \delta)]. \tag{12}$$

*Proof* According to [2, Prop. 2.27], there holds

$$\{Ax \mid A \in \mathbf{A}\} = [A_c x - \Delta|x|, A_c x + \Delta|x|].$$

Then it is sufficient to apply this result to the augmented interval matrix  $A' = (A \mid -b)$  and to the augmented vector  $x' = (x^T \mid 1)^T$  to obtain (12). □

Now, if  $x$  is a strong solution of  $Ax \leq b$ , then  $Ax - b \leq 0$  for each  $A \in \mathbf{A}, b \in \mathbf{b}$  which is equivalent to nonpositivity of the upper bound of the right-hand side interval in (12) which is the case if and only if  $A_c x + \Delta|x| \leq \underline{b}$ . This proves (10), and the other three descriptions can be proved in the same way as in the main proof.

As a direct consequence of (11) we obtain the next result. As the terminology varies, we mention explicitly that a polytope can be bounded as well as unbounded.

**Corollary 1** *The set  $X_S(\mathbf{A}, \mathbf{b})$  is a convex polytope.*

*Proof* The above set  $X_4$ , described by a set of linear inequalities (11), is a convex polytope in  $\mathbb{R}^{2n}$ , hence  $X_S(\mathbf{A}, \mathbf{b})$ , as an  $x$ -projection of it, is a convex polytope as well. □

#### 4 Computation of a strong solution

The inequalities in (9), (10) provide for direct checks of whether a given  $x$  is a strong solution, or not, whereas the description in (8) makes it possible to find a strong solution by solving the linear program

$$\min\{e^T x_1 + e^T x_2 \mid \bar{A}x_1 - \underline{A}x_2 \leq \underline{b}, x_1 \geq 0, x_2 \geq 0\}, \quad (13)$$

where  $e = (1, 1, \dots, 1)^T \in \mathbb{R}^n$ . Indeed, solve (13); if an optimal solution  $x_1^*, x_2^*$  is found, then  $Ax \leq b$  is strongly solvable and  $x = x_1^* - x_2^*$  is a strong solution of it; if (13) is infeasible, then  $Ax \leq b$  is not strongly solvable (the problem (13) cannot be unbounded since its objective is nonnegative).

Finally we show that the very form of (6) enforces complementarity of any optimal solution.

**Theorem 6** *Each optimal solution  $x_1^*, x_2^*$  of (13) satisfies  $(x_1^*)^T x_2^* = 0$ .*

*Proof* Assume that it is not so, so that some optimal solution  $x_1^*, x_2^*$  of (13) satisfies  $(x_1^*)^T x_2^* > 0$ . Then  $(x_1^*)_i (x_2^*)_i > 0$  for some  $i$ . Set  $d = \min\{x_1^*, x_2^*\}$ , then  $d \geq 0$  and  $d_i > 0$ , and by Proposition 1,  $x_1^* - d, x_2^* - d$  is a feasible solution of (13) whose objective value  $e^T (x_1^* - d) + e^T (x_2^* - d) = e^T x_1^* + e^T x_2^* - 2e^T d < e^T x_1^* + e^T x_2^*$  is less than the optimal value, a contradiction.  $\square$

#### 5 Strong solvability of interval linear equations and inequalities

For each  $y \in \{-1, 1\}^m$  (a  $\pm 1$ -vector) denote by  $T_y = \text{diag}(y)$  the  $m \times m$  diagonal matrix with diagonal vector  $y$ . Given an  $m \times n$  interval matrix  $A$ , for each  $y \in \{-1, 1\}^m$  let

$$A_y = \{T_y A \mid A \in A\}. \quad (14)$$

It can be easily seen that if  $A = [A_c - \Delta, A_c + \Delta]$ , then

$$A_y = [T_y A_c - \Delta, T_y A_c + \Delta],$$

i.e.,  $A_y$  is again an interval matrix. The same notation also applies to interval vectors that are special cases of interval matrices.

Given an  $m \times n$  interval matrix  $A$  and an interval  $m$ -vector  $b$ , the system of interval linear equations  $Ax = b$  is called strongly solvable [2] if each system

$$Ax = b$$

with data satisfying

$$A \in A, b \in b$$

is solvable.

The following theorem establishes an interconnection between strong solvability of interval linear equations and inequalities. It shows that with each strongly solvable system of interval linear equations  $Ax = b$  with an  $m \times n$  interval matrix  $A$  we may associate  $2^m$  strongly solvable systems of interval linear inequalities.

**Theorem 7** *A system  $Ax = b$  is strongly solvable if and only if  $A_y x \leq b_y$  is strongly solvable for each  $y \in \{-1, 1\}^m$ .*

*Proof* “Only if”: Let  $y \in \{-1, 1\}^m$ ,  $A' \in A_y$  and  $b' \in b_y$ . In view of (14),  $A' = T_y A$  and  $b' = T_y b$  for some  $A \in A$  and  $b \in b$ . Since  $Ax = b$  is strongly solvable,  $Ax = b$  and thus also  $T_y Ax = T_y b$  and  $T_y Ax \leq T_y b$  are solvable, so that  $A_y x \leq b_y$  is strongly solvable.

“If”: Conversely, let each  $A_y x \leq b_y$ ,  $y \in \{-1, 1\}^m$ , have a strong solution  $x_y$ , so that

$$T_y A x_y \leq T_y b,$$

which we can write as

$$T_y(b - Ax_y) \geq 0, \quad (15)$$

holds for each  $A \in A$ ,  $b \in b$ . Take some fixed  $A \in A$  and  $b \in b$ . Then (15) shows that the residual set

$$\{b - Ax \mid x \in \mathbb{R}^n\}$$

intersects all orthants, hence Theorem 3 in [5] implies that  $Ax = b$  has a solution. Since  $A \in A$  and  $b \in b$  were arbitrary, the system  $Ax = b$  is strongly solvable, which was to be proved.  $\square$

This theorem offers some answer to the question why checking strong solvability on interval linear inequalities is essentially easier than checking that of interval linear equations: in the former case only one linear program (13) needs to be solved whereas  $2^m$  of them are required in the latter case.

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