Multi-dimensional color image storage and retrieval for a normal arbitrary quantum superposition state

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Abstract Multi-dimensional color image processing has two difficulties: One is that a large number of bits are needed to store multi-dimensional color images, such as, a three-dimensional color image of $1024 \times 1024 \times 1024$ needs $1024 \times 1024 \times 1024 \times$ 24 bits. The other one is that the efficiency or accuracy of image segmentation is not high enough for some images to be used in content-based image search. In order to solve the above problems, this paper proposes a new representation for multidimensional color image, called a (n + 1)-qubit normal arbitrary quantum superposition state (NAQSS), where *n* qubits represent colors and coordinates of 2^n pixels (e.g., represent a three-dimensional color image of $1024 \times 1024 \times 1024$ only using 30 qubits), and the remaining 1 qubit represents an image segmentation information to improve the accuracy of image segmentation. And then we design a general quantum circuit to create the NAOSS state in order to store a multi-dimensional color image in a quantum system and propose a quantum circuit simplification algorithm to reduce the number of the quantum gates of the general quantum circuit. Finally, different strategies to retrieve a whole image or the target sub-image of an image from a quantum system are studied, including Monte Carlo sampling and improved Grover's algorithm which can search out a coordinate of a target sub-image only running in $O(\sqrt{N/r})$ where N and r are the numbers of pixels of an image and a target sub-image, respectively.

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H.-S. Li · R.-G. Zhou · L. Song College of Information Engineering, East China JiaoTong University, Nanchang 330013, Jiangxi, People's Republic of China e-mail: lhs@ecjtu.jx.cn **Keywords** Quantum image processing · Image storage · Image retrieval · Grover search algorithms · Quantum computing · Quantum circuit

1 Introduction

Quantum computing [1] which has the unique computing performance of quantum coherence, entanglement, superposition of quantum states and other inherent characteristics quickly becomes an international research focus. In fact, utilizing the unique properties, Shor's discrete logarithms and integer factoring algorithms in polynomial time [2], Deutsh's parallel computing algorithm with quantum parallelism and coherence [3] and Grover's quadratic speedup for unordered database search algorithm [4] are insurmountable so far by any known classical algorithms.

Since the amplitude or phase of a quantum state can be used to store information [16], a quantum system is superior to a classical computer for information storage (e.g., we may consider a system of n qubits which can store 2^n complex numbers. For n = 500, 2^{500} , this number is lager than the estimated number of atoms in the Universe! Trying to store all these complex numbers would not possible on any conceivable computer [12]). In a quantum system, frequency of the physical nature of color could represent a color instead of the RGB model or the HIS model, so a color was represented by only a 1-qubit quantum state [5] and an image was stored in a quantum array [5,6]. A quantum composite state (FRQI state) stored the colors and the coordinates of a 2D gray image of 2^n pixels with (n + 1) qubits [7]. A set of quantum states (QSMC) was proposed to represent *M* colors in an image and discussed how to retrieve images stored in a quantum system [8]. The phase-space distribution functions (Wigner and Husimi functions) were used to store an image in a quantum system [10].

Quantum computing can be realized by quantum gate operations. A finite set of basic gate operations was used to construct any quantum gate operation [11]. Universal quantum gates were expressed as the combination of one-bit gate and two-bit gates [12,13], and two-bit gates were universal for quantum computing [14]. An efficient scheme was proposed for initializing a quantum register with an arbitrary superposed state, and application of the scheme in three special cases was discussed [15].

Many scholars have studied Grover search algorithms. Grover's algorithm was generalized to deal with an arbitrary initial amplitude distribution, and a bound of the success probability of searching a marked state was derived [17]. For the known number of solutions, the rotation phase π in Grover's algorithm was replaced with an arbitrary phase of the rotation and modified algorithm's properties were analyzed by recursion equations [18]. In order to search out marked states with the maximum success probability, rotational phases of Grover algorithm must satisfy certain matching conditions [9, 19, 20].

Content-based image search is an emerging technology [21], and the image segmentation is a key technology. Significant progress had been made in field of image segmentation, for example, snake- and region growing were combined with a principled framework [21]. But for content-based image search, the accuracy or efficiency of image segmentation is not high enough for some images by employing classical segmentation algorithms. Such as, snake [22], region growing [23], clustering [24], these algorithms search out a coordinate of a target object in o(N) where N is the number of pixels. For some special images (e.g., overlapping target objects [6], the target object and the background whose colors are similar [25]), classical algorithms are difficult to segment target objects. In order to improve the efficiency and accuracy of content-based image search, image segmentation information is extracted from an original image by using classic segmentation algorithms in advance. For some special images (e.g., overlapping target objects), we can get segmentation information of target objects from the image with human-computer interaction (e.g., using Photoshop software). And then, we store the colors, coordinates and segmentation information of an image in a quantum system. It is a motivation to propose that a (n + 1)-qubit normal arbitrary quantum superposition state (NAQSS) represents a multi-dimensional color image (including segmentation information) in this paper. On the basis of the color treatment strategy in the paper [8], we design a generic quantum circuit to transform a simple initial state $|0\rangle^{\otimes (n+1)}$ into a NAQSS state to realize an image storage. And the generic quantum circuit is optimized by a quantum circuit simplification (QCS) algorithm. For the purpose of retrieving the entire image from a quantum system, we measure directly the NAQSS state with a set of projection measurement operators and calculate the maximum number of measurements (namely, Monte Carlo sampling). We adopt an improved Grover search algorithm with arbitrary rotation phases and an arbitrary initial state to efficiently retrieve a target sub-image of an image (e.g., search out a coordinate of a target sub-image only running in $O(\sqrt{N/r})$).

The paper is organized as follows: An image representation are presented in Sect. 2. The realizations of image storage and quantum circuit optimization are discussed in Sect. 3. Multi-dimensional color image retrievals are described in detail in Sect. 4. Conclusion is shown in Sect. 5.

2 Basic quantum gates and representation of a quantum image

2.1 Representation of a quantum image

We describe briefly an angle how to represent a color [8] as follows.

A bijective function F_1 which sets up a one-to-one relationship between color and angle is created

$$F_1: Color \leftrightarrow \phi \tag{1}$$

where $Color = \{color_1, color_2, \dots, color_M\}$, $color_i$ corresponds to the *ith* color in ordered *M* colors, $\phi = \{\phi_1, \phi_2, \dots, \phi_M\}$, $\phi_i = \frac{\pi(i-1)}{2(M-1)}$, $i \in \{1, 2, \dots, M\}$. For grayscale images, M = 256, colors are sorted in ascending order by the grayscale values. For example, $color_1$ and $color_{256}$ correspond to grayscale values 0 and 255, respectively. For color images, $M = 2^{24}$, we make x, y, z be the values of *R*, *G*, *B* in 24-bit RGB true color, and let $i = x \times 256 \times 256 + y \times 256 + z + 1$, thus, $color_i$ corresponds to the value (x, y, z) of *RGB*. For instance, $color_1$ and $color_{16777216}$ correspond to RGB values (0,0,0) and (255,255,255).

If V is a k-dimensional Euclidean space spanned by the orthogonal basis vectors $b_1, b_2, \ldots b_k$, a k-dimensional digital image is represented as the function $f: V \to R$

where *V* notates the position information of an image and f(V) is a color set of pixels corresponding to the position *V*. Let $f(V) \subset color$ in this paper, we obtain that the angle set ϕ represents colors of a k-dimensional digital image by (1).

A quantum superposition state in 2^n dimensional Hilbert space may be expressed as $|\psi_a\rangle = \sum_{i=0}^{2^n-1} a_i |i\rangle$ where $|i\rangle$, $i = 0, 1, \dots 2^n - 1$, is a set of orthogonal basis and a_i is an arbitrary real.

In order to represent a k-dimensional color digital image, $|\psi_a\rangle = \sum_{i=0}^{2^n-1} a_i |i\rangle$ is changed

$$\left|\psi_{\phi}\right\rangle = \sum_{i=0}^{2^{n}-1} a_{i} \left|v_{1}\right\rangle \left|v_{2}\right\rangle \cdots \left|v_{k}\right\rangle$$

$$(2)$$

where $i = i_1 \cdots i_j i_{j+1} \cdots i_l \cdots i_m \cdots i_n$, $v_1 = i_1 \cdots i_j$, $v_2 = i_{j+1} \cdots i_l$ and $v_k = i_m \cdots i_n$ are the binary expansions for i, v_1 , v_2 and v_k , respectively; $|i\rangle = |v_1\rangle |v_2\rangle \cdots |v_k\rangle$ is a coordinate (v_1, v_2, \dots, v_k) in the k-dimensional space V and $a_i \in \phi$ (see (1)) notates the color of pixel corresponding to the coordinate.

To normalize the state $|\psi_{\phi}\rangle$ in (2), we set

$$\theta_i = \frac{a_i}{\sqrt{\sum_{y=0}^{2^n - 1} a_y^2}}$$
(3)

Substituting (3) for a_i in (2), we obtain a state $|\psi_N\rangle$

$$|\psi_N\rangle = \sum_{i=0}^{2^n - 1} \theta_i |v_1\rangle |v_2\rangle \cdots |v_k\rangle$$
(4)

where $(\sum_{i=0}^{2^{n-1}} \theta_{i}^{2}) = 1$ and $i = v_{1}v_{2}\cdots v_{k}$.

Thus, the state $|\psi_N\rangle$ can represent a k-dimensional color image. For the sake of storing image segmentation information in a quantum state, we create a bijective function F_2 which sets up a one-to-one relationship between an angle and an integer

$$F_2: Number \leftrightarrow Angle \tag{5}$$

where $Angle = \{\beta_0, \beta_2, ..., \beta_{m-1}\}$ is an angle set, and $Number = \{1, 2, ..., m\}$ is a set of integers, $\beta_i = \frac{i\pi}{2(m-1)}, i \in \{0, 1, ..., m-1\}$; if m = 1, let $\beta_0 = 0$.

Assuming that an image is divided into *m* sub-images numbered as 1, 2, ..., m, we define a quantum state $|\psi\rangle$ by changing (4) to represent the whole image (including segmentation information)

$$|\psi\rangle = \sum_{i=0}^{2^{n}-1} \theta_{i} |v_{1}\rangle |v_{2}\rangle \cdots |v_{k}\rangle |\chi_{i}\rangle$$
(6)

where $|\chi_i\rangle = \cos \gamma_i |0\rangle + \sin \gamma_i |1\rangle$, $\gamma_i \in Angle$ corresponding to an integer from (5), represents the serial number of the sub-image which contains the pixel corresponding to the coordinate $|v_1\rangle |v_2\rangle \cdots |v_k\rangle$.

Since $\||\psi\rangle\| = \sqrt{\sum_{i=0}^{2^n-1} \theta_i^2(\cos^2 \gamma_i + \sin^2 \gamma_i)} = 1$, $|\psi\rangle$ is called a normal arbitrary quantum superposition state (NAQSS).

3 Realizations of image storage and circuit optimization

Two unitary matrixes are defined as

$$R_{x}(\theta) = \begin{bmatrix} \cos\theta \sin\theta\\ \sin\theta - \cos\theta \end{bmatrix}, \theta \in [0, 2\pi]$$
(7)

$$R_{y}(\theta) = \begin{bmatrix} \cos \theta - \sin \theta\\ \sin \theta & \cos \theta \end{bmatrix}, \theta \in [0, 2\pi]$$
(8)

 $R_x(\theta)$ and $R_y(\theta)$ have the following properties

$$R_{y}(\theta) \cdot R_{y}(-\theta) = R_{y}(-\theta) \cdot R_{y}(\theta) = I$$

$$R_{y}(\theta) |0\rangle = \cos \theta |0\rangle + \sin \theta |1\rangle$$

$$R_{y}(0) = I, \ R_{y}\left(\frac{\pi}{2}\right) |0\rangle = |1\rangle$$

$$R_{x}(\theta) \cdot R_{x}(\theta) = I, \ R_{x}\left(\frac{\pi}{4}\right) = H, \ and \ R_{x}\left(\frac{\pi}{2}\right) = X$$
(9)

where *I*, *H*, *X* are an identity matrix, a Hadamard matrix and a pauli-x matric, respectively.

A sequence of angles is given as follows

$$\alpha_1 = \arctan \sqrt{\frac{\sum_{i_2 \cdots i_n} |\theta_{1i_2 \cdots i_n}|^2}{\sum_{i_2 \cdots i_n} |\theta_{0i_2 \cdots i_n}|^2}}$$
(10)

$$\alpha_{j,i_1\cdots i_{j-1}} = \arctan \sqrt{\frac{\sum_{i_{j+1}\cdots i_n} \left|\theta_{i_1\cdots i_{j-1}1i_{j+1}\cdots i_n}\right|^2}{\sum_{i_{j+1}\cdots i_n} \left|\theta_{i_1\cdots i_{j-1}0i_{j+1}\cdots i_n}\right|^2}}$$
(11)

where j = 2, 3, ..., n and $i_1 i_2 \cdots i_{j-1}$ is the binary expansion of an integer. For instance, when j = 2, (11) is rewritten

$$\alpha_{2,0} = \arctan \sqrt{\frac{\sum_{i_3 \cdots i_n} |\theta_{01i_3 \cdots i_n}|^2}{\sum_{i_3 \cdots i_n} |\theta_{00i_3 \cdots i_n}|^2}}$$
(12)

$$\alpha_{2,1} = \arctan \sqrt{\frac{\sum_{i_3 \cdots i_n} |\theta_{11i_3 \cdots i_n}|^2}{\sum_{i_3 \cdots i_n} |\theta_{10i_3 \cdots i_n}|^2}}$$
(13)

A controlled- R_{xyi} operation is defined

$$R_{xji} = \left(\sum_{k=0, k\neq i}^{2^{j-1}-1} |k\rangle \langle k|\right) \otimes I + |i\rangle \langle i| \otimes R_x(\alpha_{j,i}), j \ge 2$$
(14)



Fig. 1 The realization of a normal arbitrary quantum superposition state. Dashed box i(i = 1, 2, n) corresponds to $R_{xj}(j = 1, 2, n)$, Dashed box n + 1 corresponds to $\prod_{i=0}^{2^n-1} R_{yi}$. $\gamma_0, \gamma_1, \ldots, \gamma_{2^n-1} \in \{\beta_1, \beta_2, \ldots, \beta_m\}$ (see (5))

The controlled- R_{xyi} operation is a unitary matrix, since $R_{xyi}R_{xyi}^+ = I^{\otimes j}$. A unitary matrix R_{xj} is given for j = 1 and $j \ge 2$

$$R_{xj} = R_x(\alpha_1) \otimes I^{\otimes (n-1)}, \, j = 1$$
(15)

and

$$R_{xj} = \prod_{i=0}^{2^{j-1}} -1 \left(R_{xji} \otimes I^{\otimes (n-j)} \right), j \ge 2$$
(16)

Applying successively the unitary matrix R_{xj} on the initial state $|0\rangle^{\otimes n}$, we acquire that $|\psi_N\rangle = (\prod_{j=1}^n R_{xj})|0\rangle^{\otimes n} = \sum_{i=0}^{2^n-1} \theta_i |v_1\rangle |v_2\rangle \cdots |v_k\rangle$. A controlled- R_{yi} operation is defined

$$R_{yi} = \left(\sum_{j=0, j \neq i}^{2^n - 1} |j\rangle \langle j|\right) \otimes I + |i\rangle \langle i| \otimes R_y(\gamma_i)$$
(17)

Applying successively the unitary matrix R_{yi} on the initial state $|\psi_N\rangle \otimes |0\rangle$, we achieve a NAQSS state $|\psi\rangle$

$$|\psi\rangle = \left(\prod_{i=0}^{2^n-1} R_{yi}\right) (|\psi_N\rangle \otimes |0\rangle) = \sum_{i=0}^{N-1} \theta_i |v_1\rangle |v_2\rangle \cdots |v_k\rangle (\cos \gamma_i |0\rangle + \sin \gamma_i |1\rangle)$$
(18)

 $|\psi\rangle$ is realized by the quantum circuit shown in Fig. 1. The circuit is built with $(2^{n+1} - 1)$ quantum gates which can be constructed by one-bit and two-bit gates whose total number is $o(Nlog^2N)$ where $N = 2^n$ is the pixel's number of an image.

Usually, an image may have many same colors or be locally symmetric, and a large number of pixels belong to a sub-image, so we design QCS algorithm (see Algorithm 1) to reduce the quantum gates in Fig. 1. The computational complexity of this algorithm is $o(N^2 \log N)$.

Algorithm 1 Quantum circuit simplification (QCS) algorithm

- 1. Analyzing 2^{j-1} angles $\alpha_{j,i_1\cdots i_{j-1}}$ (see Eq. (11)), divide angles into *t* groups according to the values of angles and notate $G_1, G_2, \cdots G_t$, respectively.
- 2. Regard the second subscript $i_1 \cdots i_{j-1}$ of $\alpha_{j,i_1 \cdots i_{j-1}}$ in G_i as a set of code and merge angles of G_i by the following strategy:

while (there are at least two codes whose Hamming distance equals 1 in G_i) do

Find two codes whose hamming distance is equal to 1 and merge them, e.g., $i_1 \cdots 0 \cdots i_{j-1}$ and $i_1 \cdots 1 \cdots i_{j-1}$ can be merged into $i_1 \cdots x \cdots i_{j-1}$ where x means there is no gate operator on the corresponding qubit (see Fig. (3).

end while

- 3. Successively take $i = 1, 2, \dots, t$, repeating Step 2.
- 4. Successively take $j = 1, 2, \dots, n$, repeating Step 1 to Step 3.
- 5. Analyzing 2^n angles $\gamma_{i_1i_2\cdots i_n}$, divide angles into *m* groups according to the values of angles and notate $G_1, G_2, \cdots G_m$, respectively.
- 6. Regard the subscript $i_1 i_2 \cdots i_n$ of $\gamma_{i_1 i_2 \cdots i_n}$ in G_i as a set of code, and merge angles of G_i by the following strategy

while (there are at least two codes whose Hamming distance equals 1 in G_i) do

Find two codes whose hamming distance is equal to 1 and merge them, e.g., $i_1 \cdots 0 \cdots i_n$ and $i_1 \cdots 1 \cdots i_n$ can be merged into $i_1 \cdots x \cdots i_n$.

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end while
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7. Successively take $i = 1, 2, \dots, m$, repeating Step 6.

In order to make the above algorithm explicit, let us consider a $4 \times 4 \times 2$ image shown in Fig. 2 as an example. From formulas (7)–(18) and Fig. 2, values of $\alpha_{j,i_1\cdots i_{j-1}}$ are given and QCS algorithm implementation is shown in Table 1. The optimized quantum circuit for the image is shown in Fig. 3.

4 Image retrieval

In order to retrieve images from a quantum system, we adopt two strategies: direct measurement (namely retrieve a whole image by Monte Carlo sampling), retrieve target



Fig. 2 Left a $4 \times 4 \times 2$ image which is divided into two sub-images (numbered 0 and 1) and on the *right* for lower sub-image. The image is represented as $|\psi\rangle = \sum_{i=0}^{31} \theta_i |v_1\rangle |v_2\rangle |v_3\rangle (\cos \beta_i |0\rangle + \sin \beta_i |1\rangle)$ where $|v_1\rangle |v_2\rangle |v_3\rangle = |x\rangle |y\rangle |z\rangle = |x_1x_2\rangle |y_1y_2\rangle |z_1\rangle$, x_i , y_i , $z_i \in \{0, 1\}$



Fig. 3 The optimized quantum circuit for the $4 \times 4 \times 2$ image. *Dashed box i* ($i = 1, 2, \dots, 6$) corresponds, respectively, to $\alpha_1 = \frac{\pi}{4}(R_x(\frac{\pi}{4}) = H), \alpha_{2,x} = \frac{\pi}{4}, \alpha_{3,xx}, \alpha_{4,xxx} = \frac{\pi}{4}, \alpha_{5,xxxx} = \frac{\pi}{4}, \gamma_{xx0xx} = 0(R_y(0) = I), \gamma_{xx1xx} = \frac{\pi}{2}(R_y(\frac{\pi}{2}))$

sub-images whose amplitudes are amplified by applying improved improved Grover's algorithm. To be simply, we substitute $|i\rangle$ for $|v_1\rangle \cdots |v_k\rangle$ in (6) in the section, i.e., the NAQSS state is expressed as

$$|\psi\rangle = \sum_{i=0}^{2^{n}-1} \theta_{i} |i\rangle |\chi_{i}\rangle$$
(19)

We define dim $(|v_i\rangle)$ as the dimensions of the subspace $|v_i\rangle$, i.e., $|v_i\rangle$ is expressed with dim $(|v_i\rangle)$ 1-qubit states. To successfully retrieve images, suppose we have a documentation, including the following known information: the values of k and dim $(|v_i\rangle)$ of $|i\rangle = |v_1\rangle \cdots |v_k\rangle$, the value of $\sqrt{\sum_{i=0}^{2^n-1} a_i^2}$, the number of sub-images, the number of pixels of each sub-image, average and variance of color values of each sub-image, projection measurement operator.

4.1 Retrieving a whole image from a quantum system

We define the observable operator $M = \sum_{j=0}^{2^{n+1}-1} m_j P_j$ where $P_j = |j\rangle \langle j|$, apply M to measure the quantum state $|\psi\rangle$ in Eq. (19), and obtain result m_j with probability $p(m_j) = \langle \psi | P_j | \psi \rangle = \theta_j^2$, i.e.,

$$\theta_j = \sqrt{p(m_j)} \tag{20}$$

Define n_s as the number of samples and suppose that the results of n_s measurements are n_j times of m_j , set $\hat{p}(m_j) = n_j / n_s$, so we know that $\hat{p}(m_j)$, $\hat{\theta}_j$ are the estimates of $p(m_j)$ and θ_j , i.e.,

$$\widehat{\theta}_j = \sqrt{\widehat{p}(m_j)} \tag{21}$$

Assuming that a quantum state is detected with the probability of $1 - \alpha$ in n_{max} measurements, now, we describe how to solve the value of n_{max} (i.e., how many measurements are necessary to acquire the correct value of θ_j) by appropriately changing the method in Sect. 2.2 of paper [8].

Set

$$Z = \begin{cases} 0, \text{ result of measurment is not } m_j \\ 1, \text{ result of measurment is } m_j \end{cases}$$
(22)

Since Z is either 1 or 0, Z is a Bernoulli random variable. The probability mass function of random variable Z is given by

$$\begin{cases} p(\widetilde{m_j}) = P\{Z = 0\} = 1 - p \\ p(m_j) = P\{Z = 1\} = p \end{cases}$$
(23)

where $\widetilde{m_i}$ notates result of measurement is not m_i .

The expectation μ and variance σ of Z are $\mu = p$ and $\sigma^2 = p(1-p)$, respectively. Suppose that $Z_1, Z_2, \dots Z_{n_s}$ are n_s samples of Z and n_s is sufficiently large, then $(\sum_{i=1}^{n_s} Z_i - n_s p) / \sqrt{n_s p(1-p)} = (n_s \overline{Z} - n_s p) / \sqrt{n_s p(1-p)}$ has approximately a standard normal distribution by the Central Limit Theorem. Thus, $P\{|(n_s \overline{Z} - n_s p) / \sqrt{n_s p(1-p)}| < Z_{\alpha/2}^2\} \approx 1 - \alpha$ where $1 - \alpha$ is a confidence level (the value of $Z_{\alpha/2}$ can be found in standard normal distribution lookup tables: e.g., when $\alpha = 0.05$, we can obtain $Z_{\alpha/2} = 1.96$). Solving the inequality $|(n_s \overline{Z} - n_s p) / \sqrt{n_s p(1-p)}| < Z_{\alpha/2}^2$, we obtain that the confidence interval of p is $[p_{\min}, p_{\max}]$ with approximate confidence level $1 - \alpha$. p_{\min} and p_{\max} are expressed as follows

$$p_{\min} = \frac{2n_s \overline{Z} + \lambda - \sqrt{(2n_s \overline{Z} + \lambda)^2 - 4(n_s + \lambda)n_s \overline{Z}^2}}{2n_s + 2\lambda}$$
(24)

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and

$$p_{\max} = \frac{2n_s\overline{Z} + \lambda + \sqrt{(2n_s\overline{Z} + \lambda)^2 - 4(n_s + \lambda)n_s\overline{Z}^2}}{2n_s + 2\lambda}$$
(25)

where $\lambda = Z_{\alpha/2}^2$. The size of the confidence interval $[p_{\min}, p_{\max}]$ is

$$\Delta p = p_{\max} - p_{\min} = \frac{\sqrt{4n_s\lambda\overline{Z} - 4n_s\lambda\overline{Z}^2 + \lambda^2}}{n_s + \lambda}$$
(26)

Define

$$\theta_{\max} = \sqrt{p_{\max}}, \, \theta_{\min} = \sqrt{p_{\min}}$$
(27)

From (26) and (27), we infer

$$\Delta \theta = \sqrt{p_{\text{max}}} - \sqrt{p_{\text{min}}} \le \sqrt{p_{\text{max}} - p_{\text{min}}} = \sqrt{\Delta p}$$
(28)

where $\Delta \theta = \theta_{\text{max}} - \theta_{\text{min}}$.

Since $\overline{Z} = \sum_{i=1}^{n_s} Z_i / n_s$ and $\hat{p}(m_j) = \overline{Z}$, we know that $\hat{\theta}_j = \sqrt{\overline{Z}}$ and $\overline{Z} \approx p$ by Law of Large Numbers. The confidence interval of p is $[p_{\min}, p_{\max}]$; therefore, $\overline{Z} \in [p_{\min}, p_{\max}]$ and $p \in [p_{\min}, p_{\max}]$ are with the approximate probability of $1 - \alpha$. Thus

$$\theta_j \in [\theta_{\min}, \theta_{\max}], \, \widehat{\theta}_j \in [\theta_{\min}, \theta_{\max}]$$
(29)

By (29), we find

$$\left|\widehat{\theta}_{j} - \theta_{j}\right| \le \Delta\theta \tag{30}$$

Since $a_j = \theta_j \sqrt{\sum_{i=0}^{2^n - 1} a_i^2}$ (see (3)), from (29) and (30), we yield

$$\frac{\left|\widehat{a}_{j}-a_{j}\right|}{G_{\phi}} \leq \Delta\theta \tag{31}$$

and

$$a_j, \hat{a}_j \in \left[G_\phi \theta_{\min}, G_\phi \theta_{\max}\right]$$
 (32)

where $G_{\phi} = \sqrt{\sum_{i=0}^{2^n - 1} a_i^2}$ and $\widehat{a}_j = \widehat{\theta}_j \sqrt{\sum_{i=0}^{2^n - 1} a_i^2}$. According to (1), we calculate

$$\Delta \phi = |\phi_{i+1} - \phi_i| = \frac{\pi}{2(M-1)}$$
(33)

where $i \in \{1, 2, \dots, (M-1)\}, M = 2^{24}$ or M = 256.



Fig.4 The relation of $\Delta \hat{\phi}$, $\Delta \phi$, a_j , \hat{a}_j , $G_{\phi} \theta_{\min}$, $G_{\phi} \theta_{\max}$ and $G_{\phi} \Delta \theta$. When $\hat{\theta}_j \in \Delta \theta$ (i.e., $\hat{a}_j \in G_{\phi} \Delta \theta$), we derive that $a_i = \phi_i$ from the figure

Suppose

$$\sqrt{\Delta p} < \frac{\Delta \phi}{G_{\phi}} \tag{34}$$

Let $\Delta \hat{\phi} = |a_i - \hat{a}_i|$, from (28), (31), (32) and (34), we see

$$\Delta \phi \le G_{\phi} \Delta \theta < \Delta \phi \tag{35}$$

 $\Delta \hat{\phi}, \ \Delta \phi, \ a_j, \ \hat{a}_j, \ G_{\phi} \theta_{\min}, \ G_{\phi} \theta_{\max} \text{ and } G_{\phi} \Delta \theta \text{ are shown in Fig. 4.}$ Formula (34) is equivalent to

$$y^4 n_s^2 + \left[2\lambda y^4 - 4\lambda \left(\overline{Z} - \overline{Z}^2\right)\right] n_s + y^4 \lambda^2 - \lambda^2 > 0$$
(36)

where $y = \frac{\Delta \phi}{G_{\phi}}$. Solving (36), two solutions are

$$n_{-} = \lambda \left[\frac{2(\overline{Z} - \overline{Z}^{2})}{y^{4}} - 1 - \sqrt{\left(\frac{2(\overline{Z} - \overline{Z}^{2})}{y^{4}} - 1\right)^{2} + (1 - y^{8})} \right]$$
(37)

$$n_{+} = \lambda \left[\frac{2\left(\overline{Z} - \overline{Z}^{2}\right)}{y^{4}} - 1 + \sqrt{\left(\frac{2(\overline{Z} - \overline{Z}^{2})}{y^{4}} - 1\right)^{2} + (1 - y^{8})} \right] < \frac{4\lambda(\overline{Z} - \overline{Z}^{2})}{y^{4}}$$
(38)

where $\lambda = Z_{\alpha/2}^2$, $y = \frac{\Delta \phi}{G_{\phi}}$, $\Delta \phi = \frac{\pi}{2(M-1)}$, $G_{\phi} = \sqrt{\sum_{i=0}^{2^n-1} a_i^2}$ and $\overline{Z} = \sum_{i=1}^{n_s} Z_i/n_s$. Let

$$n_{\max} = \left\lceil \frac{4\lambda(\overline{Z} - \overline{Z}^2)}{y^4} \right\rceil \tag{39}$$

where $\lceil \cdot \rceil$ means rounding up.

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Since $n_{-} \le 0$ in (37), n_{-} is not selected. Thus, we can acquire the correct quantum state after at most n_{max} (see Eq. (39)) measurements with the approximate probability of $1 - \alpha$.

Defining $\overline{a} = \frac{\sum_{i=0}^{2^n - 1} a_i^2}{2^n}$, we have

$$G_{\phi} = \sqrt{2^n} \sqrt{\overline{a}} \tag{40}$$

Since $a_i \in [0, \frac{\pi}{2}]$ (see (2)), we yield that $\sqrt{\overline{a}} \in [0, \frac{\pi}{2}]$ and $\sqrt{\overline{a}}$ can be seen as a number which is irrelevant to the size of the number *n*. From (22), (23) and (40), we infer

$$\bar{z} \approx p = \theta_j^2 = \left(\frac{a_j}{G_{\phi}}\right)^2 = \frac{a_j^2}{2^n \bar{a}}$$
(41)

Substitute $\Delta \phi = \frac{\pi}{2(M-1)}$, $G_{\phi} = \sqrt{2^n} \sqrt{\overline{a}}$ and $\overline{z} \approx \frac{a_j^2}{2^n \overline{a}}$ into (39), so that

$$n_{\max} = \left\lceil N^2 (M-1)^4 \frac{2^6 \lambda (\overline{Z} - \overline{Z}^2) \overline{a}^2}{\pi^4} \right\rceil$$
$$\approx N(M-1)^4 \frac{2^6 \lambda a_j^2 \left(\overline{a} - \frac{a_j^2}{N}\right)}{\pi^4}$$
$$\approx N(M-1)^4 \frac{2^6 \lambda a_j^2 \overline{a}}{\pi^4}$$
(42)

where $N = 2^n$ (suppose $N \gg 1$) is the number of pixels and M is the number of colors ($M = 2^{24}$ for color images and M = 256 for grayscale images).

We can conclude that n_{max} is proportional to $N(M-1)^4$ by (42).

4.2 Retrieving a target sub-image from a quantum system

Suppose an image is only divided into two sub-images (i.e., target and background), in the case, $Angle = \{\beta_0 = 0, \beta_1 = \pi/2\}$ in (5), let $\gamma_i = \beta_0$ for $i \in B$ and $\gamma_i = \beta_1$ for $i \in A$, the NAQSS state $|\psi\rangle$ in (19) is rewritten

$$|\psi(r)\rangle = \sum_{i=0,i\in B}^{2^n-1} \theta_i |i\rangle \otimes |0\rangle + \sum_{i=0,i\in A}^{2^n-1} \theta_i |i\rangle \otimes |1\rangle$$
(43)

where *A* and *B* are defined as sets contain coordinates of the target and background in an image, respectively. That is to say, *A* is a set of marked states and *B* is a set of unmarked states. We employ Grover search algorithm to efficiently search out the target sub-image corresponding to marked states. Synthesizing papers [9] and [17– 20], we propose an improved Grover's algorithm which is suitable for the NAQSS state $|\psi(r)\rangle$ in (43) and describe in detail the Grover search algorithm how to search out the marked states with the minimum number of iterations and the largest success probability.

We assume : $N = 2^n$ is a total of states, and *r* is a number of marked states; *A* and *B* are sets of marked and unmarked states, respectively; I_r is the rotation operator of marked states, and I_0 is the rotation operator of the initial state $|\psi(r)\rangle$ in (43). So Grover search operator is expressed as $G_{\psi} = G \otimes I = -H^{\otimes n}I_0(H^{\otimes n})^{-1}I_r \otimes I$ where *H* and *I* are Hadamard and identity matrixes, respectively. *A*, *B*, I_0 and I_r are listed as follows

$$A = \{i | f(i_1 i_2 \cdots i_n j) = 1\}, B = \{i | f(i_1 i_2 \cdots i_n j) = 0\}$$
(44)

where $f(\cdot)$ is a quantum oracle (a blank box) of Grover's algorithm, $f(i_1i_2\cdots i_n1) = 1$, $f(i_1i_2\cdots i_n0) = 0$, $i = i_1i_2\cdots i_n$ and $i_1, i_2, \ldots i_n, j \in \{0, 1\}$.

$$I_r = I^{\otimes n} - (1 - e^{i\alpha}) \sum_{\tau \in A} |\tau\rangle \langle \tau|$$
(45)

$$I_0 = I^{\otimes n} - (1 - e^{i\alpha}) |0\rangle \langle 0|$$
(46)

where α is a rotation angle of marked and initial states in (45) and (46).

Applying Grover search operator G_{ψ} on $|\psi(r)\rangle$

$$G_{\psi} |\psi(r)\rangle = (G \otimes I) \left(\sum_{i=0,i\in B}^{2^{n}-1} \theta_{i} |i\rangle \otimes |0\rangle + \sum_{i=0,i\in A}^{2^{n}-1} \theta_{i} |i\rangle \otimes |1\rangle \right)$$
$$= \sum_{i=0,i\in B}^{N} G(\theta_{i} |i\rangle) \otimes |0\rangle + \sum_{i=0,i\in A}^{N} (G\theta_{i} |i\rangle) \otimes |1\rangle)$$
(47)

The state $|\psi(r)\rangle$ after t iterations of G_{ψ} is changed into

$$|\psi(t)\rangle = \sum_{i \in B} l_i(t) |i\rangle |0\rangle + \sum_{i \in A} k_i(t) |i\rangle |1\rangle$$
(48)

We denote the averages of the amplitudes as follows

$$\bar{l}(t) = \frac{1}{N-r} \sum_{i \in B} l_i(t)$$
(49)

$$\overline{k}(t) = \frac{1}{r} \sum_{i \in A} k_i(t) \tag{50}$$

$$C(t) = \frac{r(1 - e^{i\alpha})e^{i\alpha}}{N}\overline{k}(t) + \frac{(1 - e^{i\alpha})(N - r)}{N}\overline{l}(t)$$
(51)

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The recursion equations of the amplitudes take

$$k_i(t+1) = C(t) - e^{i\alpha}k_i(t)$$
(52)

and

$$l_i(t+1) = C(t) - l_i(t)$$
(53)

Accumulating (52) and (53) from i = 0 to i = N - 1, we gain another recursion equations

$$\overline{k}(t+1) = C(t) - e^{i\alpha}\overline{k}(t)$$
(54)

and

$$\bar{l}(t+1) = C(t) - \bar{l}(t)$$
 (55)

Solving the recurrence Eqs. (54) and (55) to give

$$\overline{k}(t) = (-1)^{t} e^{i\alpha t} \left[\frac{d \sin(t\beta) - \sin(t\beta - \beta)}{\sin \beta} \overline{k}(0) + \frac{(1 - e^{-i\alpha}) \cos^{2} \theta \sin(t\beta)}{\sin \beta} \overline{l}(0) \right]$$
(56)

and

$$\bar{l}(t) = (-1)^{t} e^{i\alpha t} \left[\frac{\sin(t\beta + \beta) - d\sin(t\beta)}{\sin\beta} \bar{l}(0) - \frac{(d^{2} + 1)\sin(t\beta) - 2d\cos\beta\sin(t\beta)}{(1 - e^{-i\alpha})\cos^{2}\theta\sin\beta} \bar{k}(0) \right]$$
(57)

where $d = e^{i\alpha}\sin^2\theta + \cos^2\theta$, $\sin\theta = \sqrt{\frac{r}{N}}$, $\cos\theta = \sqrt{\frac{N-r}{N}}$ ($0 \le \theta \le \pi/2$), and $\cos\beta = 1 - 2\sin^2\theta \sin^2(\alpha/2) \ (0 \le \beta \le 2\pi, 0 < \alpha < 2\pi).$

Subtracting (52) from (54) and subtracting (53) from (55), we find

$$k_i(t) - \bar{k}(t) = k_i(t-1) - \bar{k}(t-1) = \dots = k_i(0) - \bar{k}(0)$$
(58)

and

$$l_i(t) - \bar{l}(t) = -\left[l_i(t-1) - \bar{l}(t-1)\right] = \dots = (-1)^t \left[l_i(0) - \bar{l}(0)\right]$$
(59)

So, we acquire the solutions of (52) and (53) from (58) and (59)

$$k_i(t) = \bar{k}(t) + k_i(0) - \bar{k}(0)$$
(60)

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and

$$l_i(t) = \bar{l}(t) + (-1)^t \left[l_i(0) - \bar{l}(0) \right]$$
(61)

From (58) and (59), we can denote the variances of the amplitudes as follows

$$\sigma_k^2(t) = \frac{1}{r} \sum_{i \in A} \left| k_i(t) - \bar{k}(t) \right|^2 = \sigma_k^2(0)$$
(62)

and

$$\sigma_l^2(t) = \frac{1}{N - r} \sum_{i \in B} \left| l_i(t) - \bar{l}(t) \right|^2 = \sigma_l^2(0)$$
(63)

Set P_t is the probability of success after t iterations to search out marked states

$$P_t = \sum_{i \in A} |k_i(t)|^2 = 1 - \sum_{i \in B} |l_i(t)|^2 = 1 - (N - r)\sigma_l^2(t) - (N - r)\left|\bar{l}(t)\right|^2$$
(64)

Since $\sigma_l(t) = \sigma_l(0)$, (64) can be rewritten as

$$P_t = 1 - (N - r)\sigma_l^2(0) - (N - r)\left|\bar{l}(t)\right|^2$$
(65)

When $\bar{l}(t) = 0$, we define

$$P_{\max} = P_t = 1 - (N - r)\sigma_l^2(0) \tag{66}$$

Simplifying the equation $\overline{l}(t) = 0$, results are

$$\sin(t\beta) \left[(1+\cos\alpha)(1-\cos\beta)(\cos^2\theta)\bar{l}(0) - 2\cos^2\theta(1-\cos\beta)\bar{k}(0) \right] + \cos(t\beta)\sin\beta(1-\cos\alpha)(\cos^2\theta)\bar{l}(0) = 0$$
(67)

and

$$\sin(t\beta)(\cos\beta - 1)\sin\alpha(\cos^2\theta)\bar{l}(0) + \cos(t\beta)\sin\beta\sin\alpha(\cos^2\theta)\bar{l}(0)$$
(68)

Since $0 < \alpha < 2\pi$ and $\cos \theta \neq 0$ (otherwise, the search space is composed by the marked states), we rewrite (67) as

$$\frac{\sin(t\beta)(1+\cos\alpha)\sin^2\theta+\cos(t\beta)\sin\beta}{2\sin^2\theta\sin(t\beta)} = \frac{\overline{k}(0)}{\overline{l}(0)}$$
(69)

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Solving (69), the optimal iterations of Grover is

$$t = \left\lfloor \frac{\frac{\pi}{2} - \arctan\left[\left(\frac{\overline{k}(0)}{\overline{l}(0)} - \frac{1 + \cos\alpha}{2}\right)\frac{2\sin^2\theta}{\sin\beta}\right]}{\beta}\right\rfloor$$
(70)

Where $\lfloor \rfloor$ means rounding down.

Substitute t = 1 in (69), we derive that the rotation angle α satisfies

$$\cos \alpha = \frac{\overline{k}(0)}{\overline{l}(0)} - \frac{1}{2\sin^2\theta}$$
(71)

Thus, $P_t = P_{\text{max}}$ after applying just one iteration when $-1 \le \cos \alpha < 1$, i.e.,

$$-1 \le \frac{\overline{k}(0)}{\overline{l}(0)} - \frac{1}{2\sin^2\theta} < 1$$
(72)

Substitute t = 2 in (69), we derive that the rotation angle α satisfies

$$\cos \alpha = \frac{1}{2} \left(1 + \frac{\bar{k}(0)}{\bar{l}(0)} \right) - \frac{3 \pm \sqrt{\left[2\sin^2 \theta (1 - \frac{\bar{k}(0)}{\bar{l}(0)}) - 1 \right]^2 + 4}}{4\sin^2 \theta}$$
(73)

Thus, $P_t = P_{\text{max}}$ after applying just two iterations when

$$-1 \le \frac{1}{2} \left(1 + \frac{\bar{k}(0)}{\bar{l}(0)} \right) - \frac{3 \pm \sqrt{\left[2\sin^2\theta \left(1 - \frac{\bar{k}(0)}{\bar{l}(0)} \right) - 1 \right]^2 + 4}}{4\sin^2\theta} < 1$$
(74)

If (72) or (74) are not satisfied, we take $\alpha = \pi$ in (70) and $P_t \approx P_{\text{max}}$ after

$$t = \left\lfloor \frac{\frac{\pi}{2} - \arctan\left[\frac{\overline{k}(0)}{\overline{l}(0)}\sqrt{\frac{r}{N-r}}\right]}{\arccos\left(1 - 2\frac{r}{N}\right)} \right\rfloor$$
(75)

iterations.

When $r \ll N$, t is in maximum. Setting $r \ll N$, from (75), we infer

$$t < \frac{\pi}{2} \frac{1}{\arccos(1 - 2\frac{r}{N})} < \frac{\frac{\pi}{4}\sqrt{N/r}}{\sqrt{1 - \frac{r}{N}}} \approx \frac{\pi}{4}\sqrt{N/r}$$
(76)

Therefore, the minimum number of iterations t is in $O(\sqrt{N/r})$ from (71), (73) and (76).

Suppose that

$$|k_i(t)|^2 = P_t^i, \ |k_i(0)|^2 = P_0^i$$
(77)

So $\sum_{i \in A} P_t^i = P_t$, $\sum_{i \in A} P_0^i = P_0$ from (64). We rewrite (60) as

$$k_i(0) = k_i(t) - \bar{k}(t) + \bar{k}(0)$$
(78)

By (72), (74) or (75), taking appropriate t and α in (56), we calculate $\overline{k}(t)$. Since $k_i(0)$ and $\overline{k}(0)$ are real, by (78), we can suppose

$$k(t) = a + bi, \ k_i(t) = c_i + bi$$
 (79)

For $k_i(0) > 0(i = 0, 1, ..., (N - 1))$, Grover search algorithm increases the total probability of the marked states and keeps the relative amplitude among the marked states, so $a \times c_i > 0$ (i.e., the signs of *a* and c_i are the same). Hence, from (77) and (79), we yield

$$c_i = \sqrt{P_t^i - b^2}, when \quad a > 0 \tag{80}$$

$$c_i = -\sqrt{P_t^i - b^2}, when \quad a < 0 \tag{81}$$

From (78) and (79), we obtain

$$k_i(0) = c_i - a + \bar{k}(0) \tag{82}$$

Therefore, select α , after t iterations, $P_t = P_{\max}$ or $P_t \approx P_{\max}$. We detect the state $(G_{\phi})^t(|\psi\rangle)$ to acquire values of $|k_i(t)|^2(i = 0, 1, ..., (N - 1))$ by the method shown in Sect. 4.1 in this paper, and further get values of $k_i(0)$ (i.e., amplitudes of the initial state $|\psi(r)\rangle$) from (80), (81) and (82). When $\sigma_l^2(0) \approx 0$ (colors of the background sub-image are similar), $P_t \approx 1$, in this case, measuring $(G_{\phi})^t(|\psi\rangle) \approx \sum_{i \in A} k_i(t) |i\rangle |1\rangle$, we retrieve a coordinate of the marked sub-image with approximate probability 100 %

every time. In order to make the above description explicit, let us consider a 512×512 image shown on the left of Fig. 5 as an example. We consider two cases: Only retrieve the coordinates of the edge of the airplane (see the middle of Fig. 5) and retrieve the whole

airplane including colors and coordinates (see the right of Fig. 5). The process of the example consists of the following steps.

- Store the original image on the left of Fig. 5 in quantum systems as a quantum state |ψ(1)⟩ or |ψ(2)⟩.
- (2) Apply improved Grover's algorithm on a quantum state $|\psi(1)\rangle$ or $|\psi(2)\rangle$.
- (3) Measure the quantum state after *t* iterations of Grover and retrieve the target sub-image.



Fig. 5 (*left*) The original image is segmented into two sub-images, one of which is marked '0' as a background sub-image and the other one is marked '1' as a target sub-image. (*middle*) The target sub-image is the edge of the airplane, and the background sub-image is the original image except the edge of the airplane. (*right*) The target sub-image is the whole airplane, and the background sub-image is the original image except the airplane

Table 2 G_{ϕ} , N (number of pixels of the original image), r (number of pixels of the target sub-image), $\overline{k}(0) = \frac{1}{r} \sum_{i \in A} \theta_i$, $\overline{l}(0) = \frac{1}{N-r} \sum_{i \in B} \theta_i$ and $\sigma_l^2(0) = \frac{1}{N-r} \sum_{i \in B} \theta_i - \overline{l}(0) ^2$ are some priori data	Priori data	No. 1 case	No. 2 case
	G_{ϕ}	577.4167784268036	577.4167784268036
	Ν	262144	262144
	r	1294	29459
	$\overline{k}(0)$	0.001584193144183	0.001589671717044
	$\overline{l}(0)$	0.001896058040768	0.001933113637537
	$\sigma_l^2(0)$	$2.248558073521673 \times 10^{-7}$	$1.896400939105966 \times 10^{-7}$

In (1) step, first, we represent colors of the original image by using angles $\{a_i\}$ (see Eq. (2)). Second, calculating $\theta_i = \frac{a_i}{G_{\phi}}$ where $G_{\phi} = \sqrt{\sum_{y=0}^{N-1} a_y^2}$ and N = 262144, we acquire a quantum state $(|\psi(1)\rangle$ or $|\psi(2)\rangle)$ as the representation of the original image for No. 1 (the edge of an airplane) or No. 2 (the whole airplane) case.

$$|\psi(1)\rangle = \sum_{i=0,i\in B}^{262143} \theta_i |i\rangle \otimes |0\rangle + \sum_{i=0,i\in A}^{262143} \theta_i |i\rangle \otimes |1\rangle$$
(83)

where B and A are defined as two sets containing coordinates of the background sub-image and the edge of the airplane (see the middle of Fig. 5), respectively.

$$|\psi(2)\rangle = \sum_{i=0,i\in B}^{262143} \theta_i |i\rangle \otimes |0\rangle + \sum_{i=0,i\in A}^{262143} \theta_i |i\rangle \otimes |1\rangle$$
(84)

where *B* and *A* are defined as two sets contain coordinates of the background subimage and the whole airplane (see the right of Fig. 5), respectively. Third, we create the quantum state $(|\psi(1)\rangle \text{ or } |\psi(2)\rangle)$ by the quantum circuit in Fig. 1. Finally, we store some priori data in quantum systems shown in Table 2.

Table 3	Some calculated data	Calculated data	No. 1 case	No. 2 case	
		α	π	2.545259554463482	
		t	11	2	
		P _{max}	1,294	29,459	
		$\overline{k}(0)$	0.9413	0.9559	



Fig. 6 The quantum circuit of the improved Grover's algorithm. **a** for No. 1 case and **b** for No. 2 case, *dashed box* is an oracle (i.e., I_r in Eq. (48)) and *solid box* is $-H^{\otimes 18}(I^{\otimes 18} - (1 - e^{i\alpha})|0\rangle \langle 0|)H^{\otimes 18}$ where $\alpha = \pi$ and $\alpha = 2.545259554463482$, respectively. *Note* $|\chi_i\rangle$ is a state $|0\rangle$ or $|1\rangle$, $R(\alpha) = \begin{bmatrix} e^{i\alpha} & 0 \\ 0 & e^{i\alpha} \end{bmatrix}$ in Fig. 6

In (2) step, α (a rotation angle of marked and unmarked states corresponding to the target sub-image and the background sub-image), *t* (the number of iterations of Grover), P_{max} (the max probability of success after t iterations to search out marked states) are calculated by Eqs. (71–75), (66) and shown in Table 3. And then, we design the quantum circuit of the improved Grover's algorithm shown in Fig. 6.

In (3) step, for No. 1 case, we retrieve a coordinate of the edge of the airplane with probability 94.13 % by measuring the first 19 qubits in (a) of Fig. 6 every time. For No. 2 case, we know that $|\psi(t)\rangle = \sum_{i \in B} l_i(t) |i\rangle |0\rangle + \sum_{i \in A} k_i(t) |i\rangle |1\rangle$ after *t* Grover iterations where $k_i(t)$ and $l_i(t)$ are amplitudes of the marked and unmarked states, and obtain $|k_i(t)|^2$ by measuring the first 19 qubits in (b) of Fig. 6. Then, we get $k_i(0)$ (i.e., θ_i) by Eqs. (77–82) and calculate $a_i = \theta_i \times G_{\phi}$ which represents a color. We can retrieve a coordinate of the airplane with probability 95.59 % by measuring (b) of Fig. 6 every time, but we retrieve a coordinate of the airplane with probability 8.63 % $(\sum_{i \in A} (\theta_i)^2 = 0.0863)$ by measuring directly the original image in Fig. 5.

If the number of sub-images $2 < m \le 8$, the NAQSS state $|\psi(r)\rangle$ in (43) is small changed into



Fig. 7 A $4 \times 4 \times 2$ image which is divided into four sub-images (numbered as 0, 1, 2, 3), and on the *right* for *lower* sub-image

$$|\psi_m\rangle = \sum_{i=0}^{N-1} \theta_i |i\rangle \otimes |j\rangle, \, j \in \left\{0, 1, \cdots (2^{\lceil \log(m) \rceil} - 1)\right\}$$
(85)

For example, suppose m = 4, let us consider a $4 \times 4 \times 2$ image shown in Fig. 7 as an example. Grover search algorithm for state $|\psi_m\rangle$ in (85) just needs to redefine the quantum Oracles as following: $f(i_1i_2\cdots i_n00) = 1$, $f(i_1i_2\cdots i_n01) =$ 1, $f(i_1i_2\cdots i_n10) = 1$ and $f(i_1i_2\cdots i_n11) = 1$ correspond, respectively, to No. 0, 1, 2, 3 sub-images in Fig. 7. For example, we retrieve No. 0 sub-image using Grover search algorithm by redefined A and B

$$A = \{i | f(i_1 i_2 \cdots i_n j_1 j_2) = 1\}, B = \{i | f(i_1 i_2 \cdots i_n j_1 j_2) = 0\}$$
(86)

where $f(i_1i_2\cdots i_n00) = 1$, $f(i_1i_2\cdots i_nj_1j_2) = 0$ for $j_1j_2 \neq 00$, $i = i_1i_2\cdots i_n$ and $i_1, i_2, \ldots, i_n, j_1, j_2 \in \{0, 1\}$.

5 Conclusion

In this paper, a (n + 1)-qubit NAQSS state can represent a k-dimensional color image, where *n* qubits represent colors and coordinates of 2^n pixels and 1 qubit represents an image segmentation information to improve the accuracy of image segmentation, while only the image (not including additional segmentation information) is represented by (n + 1)-qubits in paper [7]. And we design a general quantum circuit to store an image in quantum systems and study the different strategies to retrieve images from quantum systems. So, the NAQSS provides a foundation not only to represent images but also to explore theoretical and practical aspects of image processing on quantum computers.

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