

Noncoherent spectral optical code division multiple access system using one-dimensional added length codes

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Abstract This work proposes a new family of 1-D added length codes for use in spectral amplitude coding optical code division multiple access (OCDMA) networks. The proposed structure uses the proposed optical line terminal (OLT) and optical network units (ONUs) to produce a simple architecture for ease of operation. The OLT transfers the information bits into the code sequences of 1-D added length codes in OCDMA. The ONUs use the code sequences of the 1-D added length codes to perform a modified cross-correlation, which produces the recovered “1” and “0” bits at the photodiodes. The photodiodes remove the interference caused by other simultaneous users, which is called multi-user interference. Additionally, the proposed modified cross-correlation process suppresses the phase-induced intensity noise. The numerical results demonstrate the bit error rate of the proposed system using 1-D added length codes to be better than that of other comparable systems using 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. The data transmission rate using the 1-D added length codes can reach 2.5 Gbps.

Keywords Optical code division multiple access · Optical line terminal · Optical network units · Multi-user interference · Phase-induced intensity noise · Photodiode

1 Introduction

The maturity of optical code division multiple access (OCDMA) technology had the use of optical components to produce multiple access networks (Kataoka et al. 2011; Gharaei et al. 2010; Lalmahomed et al. 2010; Chanclou et al. 2012; Kazovsky et al. 2007). Passive

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optical networks (PONs) used passive optical components to produce network services. The OCDMA PON was based on use of various multiplexing technologies, including time division multiplexing (TDM), wavelength division multiplexing (WDM), hybrid TDM/WDM, code division multiplexing (CDM), and optical code division multiplexing (OCDM) for access networks (Ayotte and Rusch 2005; Salehi and Brackett 1989; Ahmed et al. 2013; Shalaby 2002). The OCDMA-based PON produced a highly promising solution for fiber-to-the-home (FTTH) applications. The PON technology used for OCDMA has seen some improvements in recent decades. The establishment of PON technology (Yoshima et al. 2010; Yang et al. 2004; Ghafouri-Shiraz et al. 2007; Mohamed et al. 2013; Prucnal et al. 1986; Ugale and Mishra 2010) adopted this promising technique to generate broadband access networks with the following advantages: multiple access capability, accurate timing, ability to support variable bit rates, and security against unauthorized users.

Tainta et al. (2011) developed a coding/decoding setup for a spectral phase encoding OCDMA (SPE-OCDMA) system. The use of an easily tunable electro-optic phase modulator based on the temporal self-imaging effect achieved line-by-line coding of the transmitted signal and assured its compatibility with WDM techniques. Modulation of the code sequences was performed at the same rate that was used for the data. The modulation of the code sequences in this manner avoided the need to use high-bandwidth electro-optic modulators. Based on the concept for this technique, an experimental back-to-back coder/decoder setup was produced that transmitted an unmodulated optical pulse train within an optical window. The proposed system avoided the problems inherent in similar schemes related to temperature control or requirements for high-bandwidth electro-optic modulators.

Fathallah et al. (2014) used OCDMA to produce a promising candidate for next generation passive optical networks (NG-PON). The OCDMA-PON can potentially allow all customers to use a Gb/s-class bandwidth upstream with inherent flexibility. OCDMA was subject to multi-user interference (MUI) and various detection noises. The development of an efficient interference-cancelling detection method was based on mixed Poisson–Gaussian noise. A maximum likelihood framework was developed for the system detection process using Gaussian, Poisson and Poisson–Gaussian noise. Then, the interference-cancelling detector was developed for the Poisson–Gaussian noise case after derivation of a conventional expectation–maximization (EM) detector for the Poisson noise case. Finally, the detector was simulated and compared with the projected parallel interference-cancelling EM and other types of detectors to validate its superiority and enhanced performance.

In this paper, we propose a family of new codes called 1-D added length codes for use in OCDMA systems to achieve the proposed structure using an optical line terminal (OLT) and optical network units (ONUs). The proposed structure has a simple architecture for ease of operation, which is produced using the proposed OLT and the ONUs. The OLT transfers the information bits into the code sequences of the 1-D added length codes in OCDMA. The ONUs receive the OCDMA signals and perform a modified cross-correlation, which produces the recovered “1” and “0” bits at the photodiodes. The proposed system suffers from the interference from other simultaneous users, which is called MUI, but the modified cross-correlation of the photodiode signals removes the MUI. The proposed modified cross-correlation process creates a modified photocurrent that suppresses the phase-induced intensity noise (PIIN). In the numerical results, we use the bit error rate (BER) to observe the corresponding number of simultaneous users, data transmission rate, and effective source power. The number of simultaneous users for the proposed system

using the 1-D added length codes is 23 users and is a larger number of simultaneous users than that for the other systems using 1-D M-sequence codes (Yang et al. 2004), 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes (Wei and Ghafouri-Shiraz 2002). When the BER = 10⁻⁹, the data transmission rate for the proposed system using the 1-D added length codes is 2.5 Gbps and is higher than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. The effective source power for the proposed system when using the 1-D added length codes is -2 dBm, and is less power than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes.

This paper is organized as follows. In Sect. 2, the code sequences of the 1-D added length codes are developed. In Sect. 3, we use the OLT and ONUs to produce a description of the proposed system. In Sect. 4, a performance analysis is performed to determine the signal-to-noise ratio (SNR) and the BER. In Sect. 5, the numerical results are discussed in terms of the number of simultaneous users, data transmission rate, and effective source power. In Sect. 6, conclusions are drawn.

2 1-D added length codes

We propose the 1-D added length codes for use in spectral amplitude coding (SAC) OCDMA networks. The proposed codes have code structures that enable construction of the new code families for the 1-D added length codes. The proposed codes use elements of the number sequences of the 1-D added length codes that consist of groups. We have a series of new code families for the prime number α used in each group. The 1-D added length codes use the prime number α to determine the number sequences. The proposed codes use the 1-D added length codes to produce number sequences for these 1-D added length codes. The new code families of the 1-D added length codes have the code characteristics required to create the two steps for the 1-D added length codes as follows.

The first step uses the code weight position to derive the following expression:

$$\sigma_{c,\varepsilon}(\delta) = \begin{cases} 2[(\delta * (\alpha + c) - \varepsilon - \kappa)] + \nu - 1, & \delta = 0, 1, \dots, \alpha - 1, \\ [\alpha + cc] \bmod \alpha, & \delta = \alpha, \end{cases} \tag{1}$$

where δ is [0, 1, ..., α], α is the prime number, c is [0, 1, 2, ..., $\alpha - 1$], ε is [0, 1, 2, ..., $\alpha - 1$], κ is a selection from the sequence [0, 1, 2, ..., $\alpha - 1$], cc is a selection from the sequence [0, 1, 2, ..., $\alpha - 1$], and ν is taken from [1, 2]. We construct the code weight positions in the elements of the number sequences, which are from the sequence [0, 1, ..., $2\alpha - 1$] in the finite field $GF(\alpha)$ over the prime number α . The number sequence $\sigma_{c,\varepsilon}(\delta)$ is used to create the code weight positions. The upper part of Eq. (1) in the 1-D added length codes is formed by the multiplication of δ and $(\alpha + c)$ and subsequent subtraction of $\delta * (\alpha + c)$ and $(\varepsilon + \kappa)$. The lower part of Eq. (1) in the 1-D added length codes is formed by addition of α and cc . The proposed codes use different number sequences for each pair of fixed parameters κ and cc by changing the parameters c and ε . When both c and ε change, the different number sequences constitute the code family of the 1-D added length codes.

The second step uses the number sequences $\sigma_{c,\varepsilon}(\delta)$ to map the code sequences of the 1-D added length codes. Then, based on the number sequences $\sigma_{c,\varepsilon}(\delta)$, the proposed codes

have binary number code sequences $Q_{\mu,\varphi}$ that are found using the elements of the following mapping method:

$$q_{\mu,\varphi} = \begin{cases} 1, & \varphi = \sigma_{c,e}(\delta) + 2 * \delta * \alpha, \\ 0, & \text{otherwise,} \end{cases} \tag{2}$$

where $\mu = [0, 1, 2, \dots, 2\alpha^2 - 1]$ and $\varphi = [0, 1, 2, \dots, 2\alpha^2 + 2\alpha]$. The proposed codes use the elements of the following mapping method to produce the binary number code sequences $Q_{\mu,\varphi}$. The 1-D added length codes use the number sequences $\sigma_{c,e}(\delta)$ to produce the required code sequences for the 1-D added length codes. Table 1 shows the code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes using the prime number 3. Each code sequence has $2\alpha^2 + 2\alpha$ elements, which are to be divided into $2(\alpha + 1)$ groups. Each group has one “1” and $(\alpha - 1)$ “0”s. Therefore, the proposed codes transfer the information bits into the code sequences for the 1-D added length codes.

We use the code sequence $Q_{\mu,\varphi}$ and the code sequence $Q_{\mu',\varphi}$ to produce a cross-correlation between the two code sequences as follows:

$$R(\mu, \mu') = \sum_{\mu=0}^{2\alpha^2+2\alpha-1} q_{\mu,\varphi} q_{\mu',\varphi}, \tag{3}$$

where $q_{\mu,\varphi}$ and $q_{\mu',\varphi}$ are the elements of code sequence $Q_{\mu,\varphi}$ and code sequence $Q_{\mu',\varphi}$, respectively. The proposed codes based on the 1-D added length codes use the code sequence $Q_{\mu,\varphi}$ and code sequence $Q_{\mu',\varphi}$ to produce the cross-correlation $R(\mu, \mu')$. The cross correlation $R(\mu, \mu')$ of the 1-D added length codes is derived as follows:

Table 1 Code sequence $Q_{\mu,\varphi}$ of 1-D added length codes using the prime number 3

μ	Code sequences $Q_{\mu,\varphi}$
0	0,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,1,0,0,0
1	0,0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0
2	1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,0,0,1,0,0,0
3	0,0,0,0,1,0,1,0,0,0,0,0,0,0,1,0,0,0,0,0,0,0,1,0
4	0,0,1,0,0,0,0,0,0,0,1,0,1,0,0,0,0,0,0,0,0,0,0,1,0
5	1,0,0,0,0,0,0,0,1,0,0,0,0,0,0,0,1,0,0,0,0,0,1,0
6	0,0,0,0,1,0,0,0,1,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0
7	0,0,1,0,0,0,1,0,0,0,0,0,0,0,0,0,1,0,1,0,0,0,0,0
8	1,0,0,0,0,0,0,0,0,1,0,0,0,1,0,0,0,1,0,0,0,0,0,0
9	0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,1,0,0
10	0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0
11	0,1,0,0,0,0,0,1,0,0,0,0,0,1,0,0,0,0,0,0,0,1,0,0
12	0,0,0,0,0,1,0,1,0,0,0,0,0,0,0,1,0,0,0,0,0,0,0,1
13	0,0,0,1,0,0,0,0,0,0,0,1,0,1,0,0,0,0,0,0,0,0,0,1
14	0,1,0,0,0,0,0,0,1,0,0,0,0,0,0,0,1,0,0,0,0,0,1
15	0,0,0,0,0,1,0,0,0,1,0,0,0,1,0,0,0,0,0,1,0,0,0,0
16	0,0,0,1,0,0,0,1,0,0,0,0,0,0,0,0,0,1,0,1,0,0,0,0
17	0,1,0,0,0,0,0,0,0,0,0,1,0,0,0,1,0,0,0,1,0,0,0,0

$$R(\mu, \mu') = \begin{cases} \alpha + 1, & \mu = \mu', \\ 0, & \mu \neq \mu', v \neq v', \\ 1, & \text{otherwise.} \end{cases} \tag{4}$$

The cross-correlation $R(\mu, \mu')$ of the 1-D added length codes has the classifications $\mu = \mu'$, $\mu \neq \mu'$, $v \neq v'$, and *otherwise*. If we find that $\mu = \mu'$, then the cross-correlation $R(\mu, \mu')$ of the 1-D added length codes produces the result $\alpha + 1$. If we find that $\mu \neq \mu'$ and $v \neq v'$, then the cross correlation $R(\mu, \mu')$ of the 1-D added length codes produces the result 0. If we obtain *otherwise*, then the cross-correlation $R(\mu, \mu')$ of the 1-D added length codes produces the result 1. Therefore, the proposed codes use the code sequence $Q_{\mu,\varphi}$ and the code sequence $Q_{\mu',\varphi}$ to produce the cross-correlation $R(\mu, \mu')$ of the 1-D added length codes. The cross-correlation between the code sequence $Q_{\mu,\varphi}$ and the code sequence $Q_{\mu',\varphi}$ produces the interference originating from the other simultaneous users called MUI.

The proposed codes use the code sequence $Q_{\mu',\varphi}$ to operate the modified complementary code sequence $Q_{\mu,\varphi}$. The cross-correlation between the code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes and the modified complementary code sequence $Q_{\mu',\varphi}$ of the 1-D added length codes is follows:

$$R'(\mu, \mu') = \sum_{\mu=0}^{2\alpha^2+2\alpha-1} q_{\mu,\varphi} \overline{q_{\mu',\varphi}}, \tag{5}$$

where $R'(\mu, \mu')$ is the cross-correlation of the 1-D added length codes, and $\overline{q_{\mu',\varphi}}$ represents the elements of the modified complementary code sequence $Q_{\mu',\varphi}$ of the 1-D added length codes. The cross-correlation between the code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes and the modified complementary code sequence $Q_{\mu',\varphi}$ of the 1-D added length codes has the classifications $\mu = \mu'$, $\mu \neq \mu'$, $v \neq v'$, and *otherwise*, which are given as follows:

$$R'(\mu, \mu') = \begin{cases} 0, & \mu = \mu', \\ 0, & \mu \neq \mu', v \neq v', \\ \alpha, & \text{otherwise.} \end{cases} \tag{6}$$

If we obtain $\mu = \mu'$, then the cross-correlation $R'(\mu, \mu')$ of the 1-D added length codes has the value 0. If we obtain $\mu \neq \mu'$ and $v \neq v'$, then the cross correlation $R'(\mu, \mu')$ of the 1-D added length codes has the value 0. If we obtain *otherwise*, then the cross-correlation $R'(\mu, \mu')$ of the 1-D added length codes has the value α . We use the code sequence $Q_{\mu,\varphi}$ and the modified complementary code sequence $Q_{\mu',\varphi}$ to produce the cross-correlations $R(\mu, \mu')$ and the cross-correlations $R'(\mu, \mu')$ for the 1-D added length codes. We use the cross-correlation $R(\mu, \mu')$ between the code sequence $Q_{\mu,\varphi}$ and code sequence $Q_{\mu',\varphi}$ and cross-correlation $R'(\mu, \mu')$ between the code sequence $Q_{\mu,\varphi}$ and modified complementary code sequence $Q_{\mu',\varphi}$ as the equations for Eqs. (4) and (6), respectively. The cross-correlations $R(\mu, \mu')$ and $R'(\mu, \mu')$ of the 1-D added length codes thus derive the relationships required to produce the MUI.

The modified cross-correlation creates a subtraction between the cross-correlation $R(\mu, \mu')$ and cross-correlation $R'(\mu, \mu')/\alpha$. We then use the cross-correlation $R(\mu, \mu')$ and cross-correlation $R'(\mu, \mu')/\alpha$ of the 1-D added length codes to produce the modified cross-correlation as follows:

$$\begin{aligned}
 R_M(\mu, \mu') &= R(\mu, \mu') - \frac{R'(\mu, \mu')}{\alpha} \\
 &= \begin{cases} \alpha + 1, & \mu = \mu', \\ 0, & \text{otherwise,} \end{cases}
 \end{aligned}
 \tag{7}$$

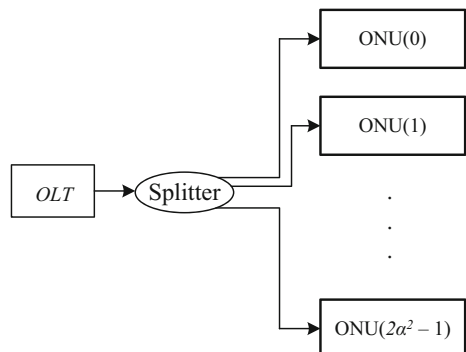
where $R_M(\mu, \mu')$ is the modified cross-correlation. $R_M(\mu, \mu)$ is the auto-cross-correlation = “ $\alpha + 1$ ” for $\mu = \mu'$ or the cross-correlation = “0” for *otherwise*. The modified cross-correlation provides the ideal cross-correlation required to produce $\mu = \mu'$ as the specific recovered bit from the 1-D added length codes. We thus perform the modified cross-correlation to produce the recovered bits. The proposed codes obtain the information bits required to output the code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes, which produces the modified cross-correlation required to recover the recovered bits. The modified cross-correlation $R_M(\mu, \mu')$ of the 1-D added length codes then eliminates the MUI using the interference contributions of the other simultaneous users. Thus, we obtain the modified cross-correlation $R_M(\mu, \mu')$ of the 1-D added length codes that is required to perform the ideal cross-correlation, which removes the MUI. We use the photodiodes to perform the modified cross-correlation of the 1-D added length codes to suppress the PIIN. Therefore, the proposed codes create a modified cross-correlation to recover the recovered bits that eliminates the MUI and suppresses the PIIN.

3 System description

Figure 1 shows a schematic block diagram of the 1-D added length codes, where the proposed architecture comprises the OLT, the $1 \times 2\alpha^2$ optical splitter, and $2\alpha^2$ of the ONU(μ), where $\mu = 0, 1, \dots, 2\alpha^2 - 1$. The OLT connects to the $1 \times 2\alpha^2$ optical splitter. The $1 \times 2\alpha^2$ optical splitter then connects to the $2\alpha^2$ of the ONU(μ), where $\mu = 0, 1, \dots, 2\alpha^2 - 1$. The system operation is as follows. The OLT transmits the code sequences of the 1-D added length codes to be sent to the $1 \times 2\alpha^2$ optical splitter. The $1 \times 2\alpha^2$ optical splitter separates the code sequences of the 1-D added length codes into the $2\alpha^2$ numbers corresponding to the ONUs. The $2\alpha^2$ ONUs then receive the code sequences of the 1-D added length codes to decode the recovered bits. Therefore, we use the OLT, $1 \times 2\alpha^2$ optical splitter, and $2\alpha^2$ ONUs to create the optical fiber connections.

Figure 2 shows the proposed architecture of the OLT. The proposed OLT architecture comprises an unpolarized broadband laser source (BLS), a $1 \times 2\alpha^2$ optical splitter, $2\alpha^2$ electro-optic modulators (EOMs), $2\alpha^2$ optical circulators, $2\alpha^2$ fiber Bragg gratings

Fig. 1 Schematic block diagram of the 1-D added length codes



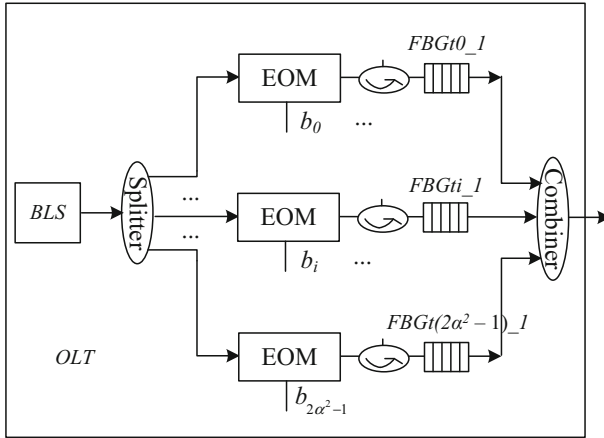


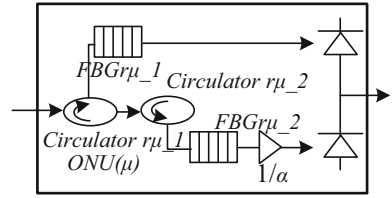
Fig. 2 Proposed architecture of the OLT

$FBGt_{\mu-1}$, where $\mu = 0, 1, 2, \dots, 2^{\alpha^2} - 1$, and a $2^{\alpha^2} \times 1$ optical combiner. The unpolarized BLS connects to the $1 \times 2^{\alpha^2}$ optical splitter. The $1 \times 2^{\alpha^2}$ optical splitter then connects to the 2^{α^2} EOMs. The 2^{α^2} EOMs then connect to the 2^{α^2} optical circulators. The 2^{α^2} optical circulators then connect to the 2^{α^2} fiber Bragg gratings $FBGt_{\mu-1}$, where $\mu = 0, 1, 2, \dots, 2^{\alpha^2} - 1$. The 2^{α^2} fiber Bragg gratings $FBGt_{\mu-1}$ then connect to the $2^{\alpha^2} \times 1$ optical combiner.

The OLT system operation is as follows. The unpolarized BLS emits optical light that is delivered to the $1 \times 2^{\alpha^2}$ optical splitter, which then creates the 2^{α^2} optical light beams. The 2^{α^2} optical light beams are delivered to the 2^{α^2} EOMs with the controlled values b_{μ} , where $\mu = 0, 1, 2, \dots, 2^{\alpha^2} - 1$, to produce the information bits. The 2^{α^2} EOMs then deliver the resulting beams to the 2^{α^2} optical circulators to be sent to the 2^{α^2} fiber Bragg gratings $FBGt_{\mu-1}$, where $\mu = 0, 1, 2, \dots, 2^{\alpha^2} - 1$. The gratings of fiber Bragg gratings $FBGt_{\mu-1}$ is set to match the complementary code sequence $Q_{\mu,\phi}$ of the 1-D added length codes. The fiber Bragg grating $FBGt_{\mu-1}$ then filters out the complementary code sequence $Q_{\mu,\phi}$ of the 1-D added length codes and filters the code sequence $Q_{\mu,\phi}$ of the 1-D added length codes. The reverse path of the optical circulator has the complementary code sequence $Q_{\mu,\phi}$ of the 1-D added length codes to be filtered out. We use the 2^{α^2} fiber Bragg gratings $FBGt_{\mu-1}$, where $\mu = 0, 1, 2, \dots, 2^{\alpha^2} - 1$, to provide the inputs of the code sequence $Q_{\mu,\phi}$ of the 1-D added length codes to the $2^{\alpha^2} \times 1$ optical combiner. Therefore, the OLT uses the unpolarized BLS to produce the code sequences of the 1-D added length codes, which are sent into the $1 \times 2^{\alpha^2}$ optical splitter, and 2^{α^2} ONU(μ), where $\mu = 0, 1, \dots, 2^{\alpha^2} - 1$, in Fig. 2.

Figure 3 shows the proposed architecture of the ONU(μ), where $\mu = [0, 1, \dots, 2^{\alpha^2} - 1]$. The proposed architecture of the ONU(μ) is composed of a left optical circulator and a right optical circulator, fiber Bragg grating $FBGr_{\mu-1}$, fiber Bragg grating $FBGr_{\mu-2}$, optical attenuator $1/\alpha$, and balanced detector using the upper photodiode and the lower photodiode. The code sequences of the 1-D added length codes connect to the left optical circulator of the ONU(μ). The left optical circulator connects to the fiber Bragg grating $FBGr_{\mu-1}$ and the right optical circulator. The right optical circulator then connects to the fiber Bragg grating $FBGr_{\mu-2}$. The fiber Bragg grating $FBGr_{\mu-1}$ and fiber Bragg grating $FBGr_{\mu-2}$ with optical attenuator $1/\alpha$ then connect to the balanced detector using the upper photodiode and the lower photodiode, respectively.

Fig. 3 Proposed architecture of the ONU(μ)



ONU(μ) system operation is as follows. The code sequences of the 1-D added length codes are received by the left optical circulator. The passed path of the left optical circulator then transmits the code sequences of the 1-D added length codes to operate fiber Bragg grating $FBGr_{\mu_1}$. The fiber Bragg grating $FBGr_{\mu_1}$ has the required number of the gratings of the complementary code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes. $FBGr_{\mu_1}$ thus filters the code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes and filters out the complementary code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes, which represents the reverse path of the left optical circulator. The reverse path of the left optical circulator is the passed path of the right optical circulator. The passed path of the right optical circulator is then used to transmit the complementary code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes to operate fiber Bragg grating $FBGr_{\mu_2}$. The fiber Bragg grating $FBGr_{\mu_2}$ has the number of gratings to be equal to the modified complementary code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes. $FBGr_{\mu_2}$ thus filters the modified complementary code sequence $Q_{\mu,\varphi}$ of the 1-D added length codes. The outputs of $FBGr_{\mu_1}$ and $FBGr_{\mu_2}$ with the optical attenuator $1/\alpha$ then deliver to the upper photodiode and lower photodiode in the balanced detector, respectively.

The output photocurrent of photodiode PD0 is proportional to $R^{(0)}(\mu, \mu')$. The output photocurrent of photodiode PD1 is proportional to $R^{(1)}(\mu, \mu')$. The proposed ONU(μ) structure suffers the interference from the ONU(μ) using unmatched codewords in the MUI. According to Eq. (7), the modified photocurrent of the photodiodes is proportional to $R^{(0)}(\mu, \mu') - R^{(1)}(\mu, \mu')/\alpha$, which is $(a + 1)$ for $\mu = \mu'$ or zero otherwise. The MUI is thus completely eliminated by the modified photocurrents of the photodiodes. Additionally, because the PIIN comes from the simultaneous users, the PIIN can also be suppressed using the modified photocurrents of the photodiodes. Therefore, the modified photocurrents of the photodiodes combine to remove the MUI and suppress the PIIN.

4 Performance analysis

The performance analysis provides the number of the simultaneous users, data transmission rate, and effective source power, which are required to measure the BER. A Gaussian approximation is used to calculate the BER. The BER is used to calculate the SNR, which determines the square of the photocurrent from the modified photocurrent and the photocurrent noise variance. The photocurrent noise variance contains the shot noise, PIIN, and thermal noise. Based on a simplified analysis, some assumptions are defined as follows. Each of the spectral components has an identical spectral width. Each simultaneous user uses equal amounts of power at the ONUs. The broadband light source is ideally unpolarized and has a flat spectrum over $[f_o - \Delta f/2, f_o + \Delta f/2]$. f_o and Δf are the central frequency and bandwidth of the source, respectively. The code sequences of the 1-D added

length codes arrive at the ONUs. Therefore, the performance analysis uses the BER to measure the proposed system overall.

The power spectral density (PSD) of the received optical laser signals in the ONU(μ) is described as follows:

$$\Lambda(f) = \frac{P_{sr}}{\Delta f} \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} \Psi(f, \varphi), \tag{8}$$

where P_{sr} is the effective source power at the ONU(μ), Δf is the source bandwidth, $Z(\theta)$ is the data bit of the θ -th user (which can be “1” or “0”), O is the number of simultaneous users, α is the code weight of the code sequence in the 1-D added length codes, and $q_{\mu,\varphi}$ is an element of the θ -th user’s in the code sequence $Q_{\mu,\varphi}$. For convenience in the analysis, $\Psi(f, \varphi)$ produces the following equation:

$$\Psi(f, \varphi) = \left\{ v \left[f - f_0 - \frac{\Delta f}{2(2\alpha^2 + 2\alpha)} (- (2\alpha^2 + 2\alpha) + 2\varphi) \right] - v \left[f - f_0 - \frac{\Delta f}{2(2\alpha^2 + 2\alpha)} (- (2\alpha^2 + 2\alpha) + 2\varphi + 2) \right] \right\}, \tag{9}$$

$v(f)$ is the unit step function, as follows:

$$v(f) = \begin{cases} 1, & f \geq 0, \\ 0, & f < 0. \end{cases} \tag{10}$$

Therefore, the PSD of the received optical laser signals in the ONU(μ) is given by Eq. (8).

The cross-correlation between the codewords $Q_{\mu,\varphi}$ and $Q_{0,\varphi}$ is given by the PSDs at the photodiodes of the ONUs. Therefore, the cross-correlation during a single bit period using the PSDs at the photodiodes of the ONUs can be written as follows:

$$H_0(f) = \frac{P_{sr}}{\Delta f} \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} q_{0,\varphi} \Psi(f, \varphi), \tag{11}$$

$$H_1(f) = \frac{P_{sr}}{\Delta f \alpha} \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} \overline{q_{0,\varphi}} \Psi(f, \varphi), \tag{12}$$

where α is the code weight of the code sequence in the 1-D added length codes. Therefore, $H_0(f)$ and $H_1(f)$ correspond to the cross-correlations. Because the photodiodes are PDO and PDI , the photocurrents I_{PD0} and I_{PD1} are developed as follows:

$$\begin{aligned} I_{PD,i,0} &= R \int_0^\infty H_0(f) df \\ &= \frac{RP_{sr}}{\Delta f} \int_0^\infty \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} q_{0,\varphi} \Psi(f, \varphi) df \\ &= \frac{RP_{sr}}{2\alpha^2 + 2\alpha} \left((\alpha + 1) + \frac{(O-1)(\alpha^2-1)}{(2\alpha^2-1)} \right), \end{aligned} \tag{13}$$

$$\begin{aligned}
 I_{PD,i,1} &= R \int_0^\infty H_1(f)df \\
 &= \frac{RP_{sr}}{\Delta f} \int_0^\infty \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} \overline{q_{0,0,\varphi}} \Psi(f, \varphi) df \\
 &= \frac{RP_{sr}}{(2\alpha^2 + 2\alpha)} \left(\frac{(O-1)(\alpha^2-1)}{(2\alpha^2-1)} \right),
 \end{aligned} \tag{14}$$

where R is the photodiode responsivity given by $R = \eta e \lambda_o / hc$, η is the quantum efficiency of the photodiode, e is the electron charge, λ_o is the central wavelength, h is Planck's constant, c is the speed of light, and P_{sr} is the effective source power of each ONU. The proposed system uses $Z(\theta) = 1$ and $\theta = 1, 2, \dots, O$ in the worst case scenario. Therefore, the photocurrents I_{PD0} and I_{PD1} in the photodiodes $PD0$ and $PD1$, respectively, are based on the cross-correlations.

The cross-correlations operate $R^{(0)}(\mu, \mu')$ and $R^{(1)}(\mu, \mu')$. The proposed system uses these cross-correlations to produce a modified cross-correlation. The modified cross-correlation thus produces the modified photocurrent. The modified photocurrent can be written as:

$$\begin{aligned}
 I_a &= R \int_0^\infty [H_0(f) - H_1(f)]df \\
 &= I_{PD0} - I_{PD1} \\
 &= \frac{RP_{sr}(\alpha + 1)}{2\alpha^2 + 2\alpha}.
 \end{aligned} \tag{15}$$

Because the proposed system uses $Z(\theta) = 1$ and $\theta = 1, 2, \dots, O$ in the worst case scenario, the worst case increases the maximum MUI. Therefore, the modified photocurrent is produced to eliminate the MUI.

The photodiodes have photocurrent noise variances that are described as follows:

$$\langle i_{noise}^2 \rangle = \langle i_{PIIN}^2 \rangle + \langle i_{shot}^2 \rangle + \langle i_{thermal}^2 \rangle = I_a^2 B_r \mu_a + 2eI_t B_r + 4K_a T_n B_r / R_o, \tag{16}$$

where the PIIN, shot noise, and thermal noise are considered in the photocurrent analysis, effect of the dark current is neglected, I_a is the modified photocurrent, I_t is the total photocurrent, B_r is the electrical bandwidth, μ_a is the coherence time of the light incident on the photodiode, e is the electron charge, K_a is Boltzmann's constant, T_n is the absolute noise temperature, and R_o is the load resistance. Therein, μ_a is derived as follows:

$$\mu_a = \int_0^\infty H^2(f)df \bigg/ \left[\int_0^\infty H(f)df \right]^2, \tag{17}$$

where $H(f)$ is the single sideband PSD of the light incident on the photodiode. Therefore, the proposed system suffers from the variances of the photocurrent noise, which provide the PIIN, shot noise, and thermal noise.

The variances of the photocurrent noise are based on statistically independent noise characteristics to produce the PIIN, shot noise, and thermal noise. The variance of the PIIN is derived as follows:

$$\langle i_{PIIN}^2 \rangle = B_r I_a^2 \mu_a. \tag{18}$$

The PIIN has a variance in the ONU that can be written as follows:

$$\begin{aligned} \langle i_{PIIN}^2 \rangle &= B_r I_a^2 \int_0^\infty H^2(f) df / \left[\int_0^\infty H(f) df \right]^2 \\ &= B_r I_a^2 \frac{\int_0^\infty [H_0(f) - H_1(f)]^2 df}{\left\{ \int_0^\infty [H_0(f) - H_1(f)] df \right\}^2} \\ &= B_r R^2 \int_0^\infty [H_0(f) - H_1(f)]^2 df \\ &= B_r R^2 \int_0^\infty [H_0^2(f) + H_1^2(f) - 2H_0(f)H_1(f)] df. \end{aligned} \tag{19}$$

The proposed system uses the spectrum of $H_0(f)$ and $H_1(f)$ such that it does not overlap the spectrum of $H_0(f)$ and $H_1(f)$ given by Eqs. (11) and (12). When the proposed system uses Eqs. (11) and (12), we have $H_i^2(f)$, where $i = [0, 1]$, as follows:

$$\begin{aligned} \int_0^\infty H_0^2(f) df &= \int_0^\infty \left[\frac{P_{sr}}{\Delta f} \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} q_{0,\varphi} \Psi(f, \varphi) \right]^2 df \\ &= \frac{P_{sr}^2}{\Delta f (2\alpha^2 + 2\alpha)(\alpha + 1)} \left[(\alpha + 1) + \frac{(O - 1)(\alpha^2 - 1)}{2\alpha^2 - 1} \right]^2, \end{aligned} \tag{20}$$

$$\begin{aligned} \int_0^\infty H_1^2(f) df &= \int_0^\infty \left[\frac{P_{sr}}{\Delta f} \sum_{\theta=1}^O Z(\theta) \sum_{\mu=0}^{(2\alpha^2-1)} \sum_{\varphi=0}^{(2\alpha^2+2\alpha-1)} q_{\mu,\varphi} q_{\overline{0},\varphi} \Psi(f, \varphi) \right]^2 df \\ &= \frac{P_{sr}^2 (\alpha + 1)(O - 1)^2 (\alpha^2 - \alpha)^2}{\Delta f \alpha^2 (2\alpha^2 + 2\alpha)(2\alpha^2 - 1)^2}. \end{aligned} \tag{21}$$

All simultaneous users use the information bit “1” from the ONU, according to Eqs. (20) and (21). The proposed system suffers from PIIN. $\langle i_{PIIN}^2 \rangle$ is described as follows:

$$\begin{aligned} \langle i_{PIIN}^2 \rangle &= \frac{B_r R^2 P_{sr}^2}{\Delta f (2\alpha^2 + 2\alpha)(\alpha + 1)} \left[(\alpha + 1) + \frac{(O - 1)(\alpha^2 - 1)}{2\alpha^2 - 1} \right]^2 \\ &+ \frac{P_{sr}^2 (\alpha + 1)(O - 1)^2 (\alpha^2 - \alpha)^2}{\Delta f \alpha^2 (2\alpha^2 + 2\alpha)(2\alpha^2 - 1)^2}. \end{aligned} \tag{22}$$

Because every simultaneous user uses the information bit “1” and information bit “0” to be sent with equal probability, we revise Eq. (22) as follows:

$$\begin{aligned} \langle i_{PIIN}^2 \rangle &= \frac{B_r R^2 P_{sr}^2}{2\Delta f (2\alpha^2 + 2\alpha)(\alpha + 1)} \left[(\alpha + 1) + \frac{(O - 1)(\alpha^2 - 1)}{2\alpha^2 - 1} \right]^2 \\ &+ \frac{P_{sr}^2 (\alpha + 1)(O - 1)^2 (\alpha^2 - \alpha)^2}{2\Delta f \alpha^2 (2\alpha^2 + 2\alpha)(2\alpha^2 - 1)^2} \end{aligned} \tag{23}$$

The proposed system considers the shot noise to be independent of the PIIN noise and thermal noise. The variance of the shot noise can thus be written as follows:

$$\begin{aligned} \langle i_{shot}^2 \rangle &= 2eB(I_{total}) \\ &= \frac{2eBRP_{sr}}{2\alpha^2 + 2\alpha} \left((\alpha + 1) + \frac{2(O-1)(\alpha^2 - 1)}{(2\alpha^2 - 1)} \right), \end{aligned} \tag{24}$$

where I_t is the total photocurrent given by $I_t = I_0 + I_1$. The proposed system allows the information bit “1” and the information bit “0” to be sent with equal probability in every simultaneous user. Thus, the variance of the shot noise can be written as:

$$\langle i_{shot}^2 \rangle = \frac{eBRP_{sr}}{2\alpha^2 + 2\alpha} \left((\alpha + 1) + \frac{2(O - 1)(\alpha^2 - 1)}{(2\alpha^2 - 1)} \right), \tag{25}$$

The proposed system also considers the thermal noise to be independent of the PIIN noise and the shot noise. The variance of the thermal noise is described as follows:

$$\langle i_{thermal}^2 \rangle = 4K_a T_n B_r / R_L, \tag{26}$$

where K_a is Boltzmann’s constant, T_n is the absolute temperature, and R_L is the load resistance. Therefore, the proposed system provides the square of the modified photocurrent and photocurrent noise variances using the PIIN, shot noise, and thermal noise. The SNR is thus depicted as follows:

$$\begin{aligned} SNR &= I_a^2 / \langle i_{noise}^2 \rangle \\ &= I_a^2 / (\langle i_{PIIN}^2 \rangle + \langle i_{shot}^2 \rangle + \langle i_{thermal}^2 \rangle). \end{aligned} \tag{27}$$

Because the BER uses a Gaussian approximation, the BER is used to calculate the SNR. The BER is thus expressed as follows:

$$BER = \operatorname{erfc}(\sqrt{SNR/8}) / 2, \tag{28}$$

where $\operatorname{erfc}(\cdot)$ is the complementary error function, which is given as follows:

$$\operatorname{erfc}(y) = \frac{2}{\sqrt{\pi}} \int_y^\infty \exp(-z^2) dz. \tag{29}$$

We also use the BER to observe the numbers of simultaneous users, data transmission rate, and effective source power.

5 Numerical results

The results of the numerical calculations are based on use of the following parameters: PD quantum efficiency $\eta = 0.6$, data transmission rate = 2.5 Gbps, receiver noise temperature $T_n = 300$ K, and receiver load resistance $R_L = 1030 \Omega$. The spectral efficiency is used by $SE = \text{aggregate information rate}/\text{total spectral bandwidth} = \theta_{BER=10^{-9}} Rb / \Delta f$, where $\theta_{BER=10^{-9}}$ is the number of simultaneous user at $BER = 10^{-9}$, and $Rb = 1/\text{data transmission rate}$ is the data transmission rate per ONU (Chang and Sargent 2003). Figure 4 shows the code length versus the spectral efficiency for SAC-OCDMA systems using the 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. The spectral efficiency using the 1-D added length codes is better than that using the 1-D M-sequence codes, 1-D composite

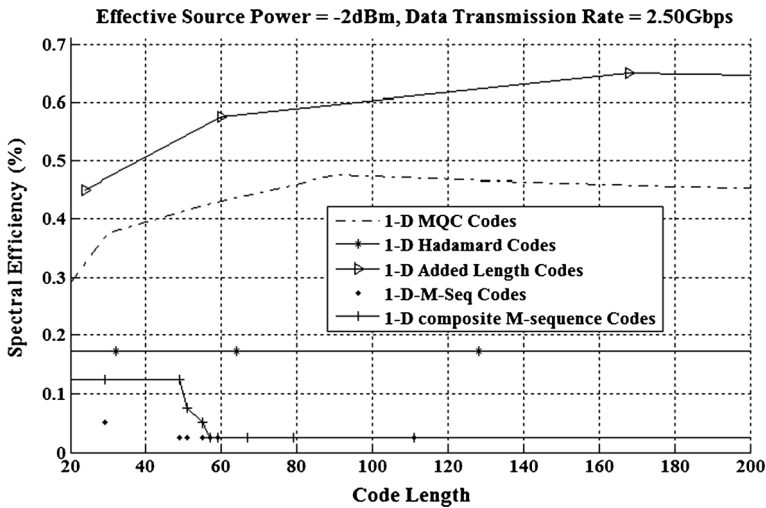


Fig. 4 Code length versus the spectral efficiency for SAC-OCDMA systems using the 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes

M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. Figure 5 shows the number of simultaneous user versus BER characteristics for the SAC-OCDMA systems using two code lengths of the 1-D added length codes. We use the effect of code set's properties to achieve the code length. The proposed codes use the larger code length to get the larger number of simultaneous user. Hence, the 1-D added length codes can operate the modified cross-correlation to achieve the larger number of simultaneous user using the larger code length. Figure 6 shows the number of simultaneous users versus the BER for SAC-OCDMA systems using the 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. These systems use similar code lengths. Because these systems using the 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes produce the effective source power of -2 dBm, a data transmission rate of 2.5 Gbps, and a BER of 10^{-9} , the numbers of simultaneous users can be compared. The proposed system using the 1-D added length codes achieves the 23 users, system using the 1-D M-sequence codes achieves only one user, system using the 1-D composite M-sequence codes achieves three users, system using the 1-D Hadamard codes achieves seven users, and system using the 1-D MQC codes achieves seventeen users. The number of simultaneous users for the proposed system using the 1-D added length codes is thus much larger than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. Based on the number of the simultaneous users 23 users, the BER for the proposed system using the 1-D added length codes is 10^{-9} , that for the system using the 1-D M-sequence codes is $10^{-0.5}$, that for the system using the 1-D composite M-sequence codes is $10^{-0.9}$, that for the system using the 1-D Hadamard codes is $10^{-1.8}$, that for the system using the 1-D MQC codes is $10^{-5.9}$. The BER for the proposed system using the 1-D added length codes is thus much lower than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes.

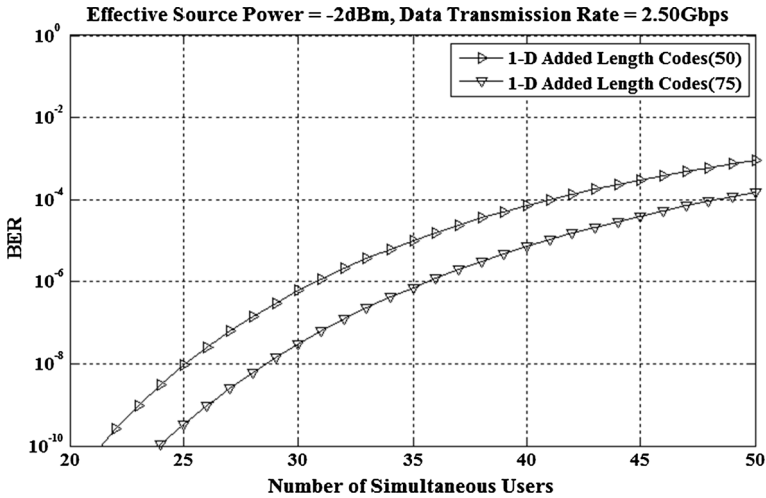


Fig. 5 Number of simultaneous users versus BER in the systems using two code lengths of the 1-D added length codes

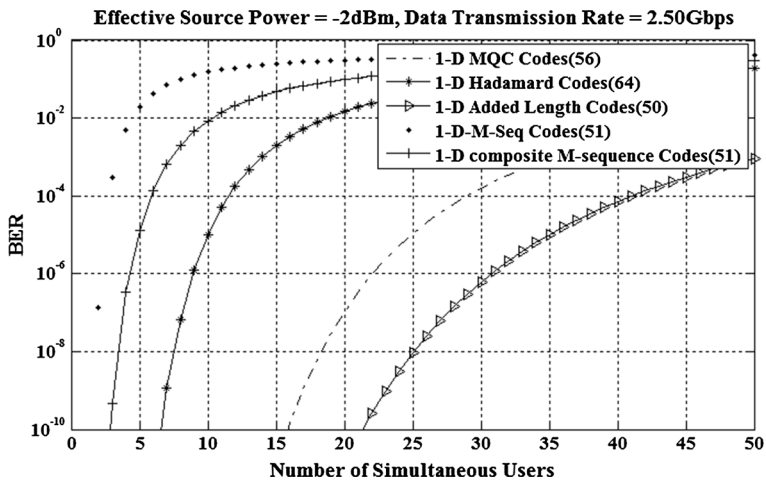


Fig. 6 Number of simultaneous users versus BER in the systems using 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes

Figure 7 shows the data transmission rate versus BER characteristics for SAC-OCDMA systems using the 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. The data transmission rate describes the information bit required to achieve the desired communication speed. Because Fig. 7 describes the number of simultaneous users to be 23 users, effective source power to be -2 dBm, and BER to be 10^{-9} , the data transmission rate for the proposed system using the 1-D added length codes is 2.5 Gbps; in contrast, that for the system using the 1-D M-sequence codes is 0.1 Gbps, that for the system using the 1-D composite M-sequence codes is 0.2 Gbps, that for the system using the 1-D Hadamard codes is 0.3 Gbps, and that for the system using the 1-D MQC codes is 1.3 Gbps. The observed data

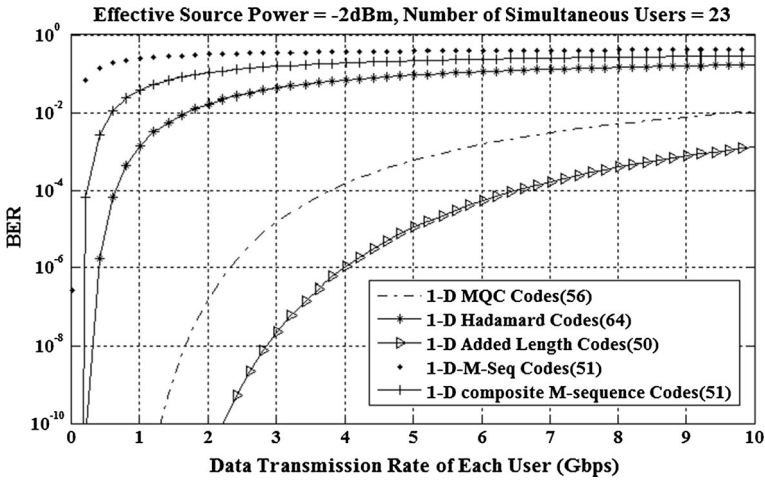


Fig. 7 Data transmission rate versus BER in the systems using 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes

transmission rate for the proposed system using the 1-D added length codes is thus much higher than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. At the same time, when the number of simultaneous users is 23 users, effective source power is -2 dBm, and data transmission rate is 2.5 Gbps, the BER for the proposed system using the 1-D added length codes is 10^{-9} , that for the system using the 1-D M-sequence codes is $10^{-0.5}$, that for the system using the 1-D composite M-sequence codes is $10^{-0.9}$, that for the system using the 1-D Hadamard codes is $10^{-1.7}$, and that for the system using the 1-D MQC codes is $10^{-5.8}$. The BER for the proposed system using the 1-D added length codes is thus much lower than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes.

Figure 8 shows the effective source power versus BER characteristics for the SAC-OCDMA systems using the 1-D added length codes, 1-D M-sequence codes, and 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. Again using the number of simultaneous users of 23, the data transmission rate of 2.5 Gbps, and the BER of 10^{-9} , the effective source power for the proposed system using the 1-D added length codes is -2 dBm, that for the systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes fail to show cross points. The effective source power for the proposed system using the 1-D added length codes is less power than that for the other systems using the 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. Based on the effective source power of -2 dBm, the BER for the proposed system using the 1-D added length codes is 10^{-9} , that for the system using the 1-D M-sequence codes is $10^{-0.5}$, and that for the system using the 1-D composite M-sequence codes is $10^{-0.9}$, that for the system using the 1-D Hadamard codes is $10^{-1.8}$, that for the system using the 1-D MQC codes is $10^{-5.8}$. The BER for the proposed system using the 1-D added length codes is thus obtained at lower power than that for the other systems using the 1-D M-sequence codes and 1-D composite M-sequence codes.

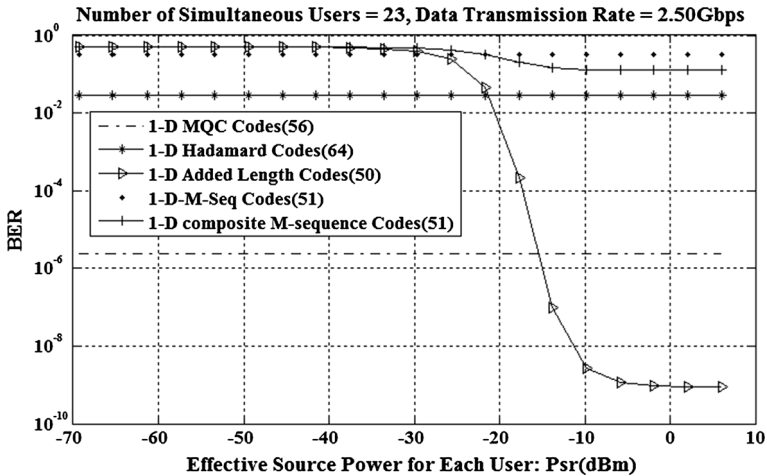


Fig. 8 Effective source power versus BER in the systems using 1-D added length codes, 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes

6 Conclusions

In this paper, we propose a family of new codes called 1-D added length codes for OCDMA systems. The proposed structure uses the 1-D added length codes to construct the OLT and ONUs. The proposed structure has a simple architecture for ease of operation that produces the proposed OLT and ONUs. The OLT transfers the information bits into the code sequences of the 1-D added length codes in OCDMA. The ONUs receive the OCDMA signals to create a modified cross-correlation, which produces the recovered “0” and “1” bits of the photodiodes. The modified photocurrent removes the interference from other simultaneous users that constitutes the MUI. The proposed modified cross-correlation suppresses the PIIN. In the numerical results, the number of simultaneous users for the proposed system using the 1-D added length codes is 23 users, which is much larger than that for the other systems using 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes based on a BER of 10^{-9} . Using this BER of 10^{-9} , the data transmission rate for the proposed system using the 1-D added length codes is 2.5 Gbps, and is much higher than that for the other systems using 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes. For the same BER, the effective source power for the proposed system using 1-D added length codes is -2 dBm, which is lower than that for the other systems using 1-D M-sequence codes, 1-D composite M-sequence codes, 1-D Hadamard codes, and 1-D MQC codes.

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