Space-filling Latin hypercube designs for computer experiments

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Abstract In the area of computer simulation, Latin hypercube designs play an important role. In this paper the classes of maximin and Audze-Eglais Latin hypercube designs are considered. Up to now only several two-dimensional designs and a few higher dimensional designs for these classes have been published. Using periodic designs and the Enhanced Stochastic Evolutionary algorithm of Jin et al. (J. Stat. Plan. Interference $134(1)$:268–687, [2005\)](#page-18-0), we obtain new results which we compare to existing results. We thus construct a database of approximate maximin and Audze-Eglais Latin hypercube designs for up to ten dimensions and for up to 300 design points. All these designs can be downloaded from the website <http://www.spacefillingdesigns.nl>.

Keywords Audze-Eglais · Computer experiment · Enhanced stochastic evolutionary algorithm · Latin hypercube design · Maximin · Non-collapsing · Packing problem · Simulated annealing · Space-filling

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1 Introduction

A *k*-dimensional Latin hypercube design (LHD) of *n* points, is a set of *n* points $x_i = (x_{i1}, x_{i2}, \ldots, x_{ik}) \in \{0, \ldots, n-1\}^k$ such that for each dimension *j* all x_i are distinct. In this definition, we assume that our design space is equal to the $[0; n-1]^k$ hypercube. However by scaling, we can use LHDs for any rectangular design space. Alternative definitions of LHDs also occur in the literature. One alternative definition is to divide each axis into *n* equally sized bins and randomly select points such that each bin contains exactly one point. However, we refer to this technique as Latin hypercube sampling (LHS). In this paper the term 'LHD' thus only refers to the first definition.

An LHD is called maximin when the separation distance $\min_{i \neq j} d(x_i, x_j)$ is maximal among all LHDs of given size *n*, where *d* is a certain distance measure. In this paper, we concentrate on the Euclidean (or ℓ^2) distance measure, i.e.,

$$
d(x_i, x_j) = \sqrt{\sum_{l=1}^{k} (x_{il} - x_{jl})^2},
$$
 (1)

since this measure is often the first choice in practice.

Besides maximin LHDs, we also treat Audze-Eglais LHDs. These LHDs minimize the following objective:

$$
\sum_{i=1}^{n} \sum_{j=i+1}^{n} \frac{1}{d(x_i, x_j)^2},\tag{2}
$$

where $d(x_i, x_j)$ is again the Euclidean distance between points x_i and x_j . By minimizing this objective, we can also obtain LHDs with "evenly spread" points (Bates et al. [2004\)](#page-18-2).

For both classes of LHDs, we aim to construct a database of the best designs known in literature. We do this by generating new designs and comparing them with existing results. These designs are often approximate maximin or Audze-Eglais designs in the sense that optimality of the objective is not guaranteed. The reason for this is that optimization over the total set of LHDs can be very time-consuming for larger values of *k* and *n*. Therefore, in order to find good designs, optimization is often done over a certain class of LHDs or heuristics are used which do not guarantee optimality. The periodic LHDs described in this paper are a good example of the first case. Examples of the second case are simulated annealing used by Morris and Mitchell [\(1995](#page-19-0)), the permutation genetic algorithm of Bates et al. ([2004\)](#page-18-2) and the Enhanced Stochastic Evolutionary (ESE) algorithm of Jin et al. [\(2005](#page-18-0)).

The designs which are best according to the comparison in this paper are added to the website <http://www.spacefillingdesigns.nl> where they can be downloaded for free. As far as we know this is the first extensive online catalogue of maximin and Audze-Eglais LHDs, although there are several catalogues for classical design of experiments, see e.g., the WebDOE™ website of Crary ([2008\)](#page-18-3). Crary et al. [\(2000](#page-18-4)) developed I-OPT™ to generate designs with minimal integrated mean squared error (IMSE). They found that IMSE-optimal designs can have proximate design points, which they call "twin points"; see also Crary ([2002\)](#page-18-5).

Our main motivation for investigating this subject is that maximin and Audze-Eglais Latin hypercube designs are very useful in the area of computer simulation. One important area where computer simulation is used a lot is engineering. Engineers are confronted with the task of designing products and processes. Since physical experimentation is often expensive and difficult, computer models are frequently used for simulating physical characteristics. The engineer often needs to optimize the product or process design, i.e., to find the best settings for a number of design parameters that influence the critical quality characteristics of the product or process. A computer simulation run is usually time-consuming and there is a great variety of possible input combinations. For these reasons, meta-models that model the quality characteristics as explicit functions of the design parameters are constructed. Such a meta-model, also called a (global) approximation model or surrogate model, is obtained by simulating a number of design points. Well-known meta-model types are polynomials and Kriging models. Since a meta-model evaluation is much faster than a simulation run, in practice such a meta-model is used, instead of the simulation model, to gain insight into the characteristics of the product or process and to optimize it. A review of meta-modeling applications in structural optimization can be found in Barthelemy and Haftka [\(1993](#page-18-6)), and in multidisciplinary design optimization in Sobieszczanski-Sobieski and Haftka ([1997\)](#page-19-1).

As observed by many researchers, there is an important distinction between designs for computer experiments and designs for the more traditional response surface methods. Physical experiments exhibit random errors and computer experiments are often deterministic (cf. Simpson et al. [2004](#page-19-2)). This distinction is crucial and much research is therefore aimed at obtaining efficient designs for computer experiments.

As is recognized by several authors, such a design for computer experiments should at least satisfy the following two criteria (see Johnson et al. [1990](#page-18-7), Morris and Mitchell [1995,](#page-19-0) and Simpson et al. [2001\)](#page-19-3). First of all, the design should be *spacefilling* in some sense. When no details on the functional behavior of the response parameters are available, it is important to be able to obtain information from the entire design space. Therefore, design points should be "evenly spread" over the entire region. One of the measures often used to obtain space-filling designs is the maximin measure (see p. 148 of Santner et al. [2003](#page-19-4) and p. 17 of Forrester et al. [2006\)](#page-18-8). The Audze-Eglais measure is another measure used for this purpose. Secondly, the design should be *non-collapsing*. When one of the design parameters has (almost) no influence on the function value, two design points that differ only in this parameter will "collapse", i.e., they can be considered as the same point that is evaluated twice. For deterministic simulation models this is not a desirable situation. Therefore, two design points should not share any coordinate values when it is not known a priori which dimensions are important. Note that in other fields of research such designs are referred to as *low discrepancy* designs. To obtain non-collapsing designs the Latin hypercube structure is often enforced. It can be shown that if the function of interest is independent of one or more of the *k* parameters then, after removal of the irrelevant parameters, the projection of the LHD onto the reduced design space retains good spatial properties; see Koehler and Owen ([1996\)](#page-18-9). Maximin LHDs are frequently used in practical applications, see e.g., the examples given in Driessen et al. [\(2002](#page-18-10)), den Hertog and Stehouwer [\(2002](#page-18-11)), Alam et al. [\(2004](#page-17-0)), and Rikards and Auzins ([2004\)](#page-19-5).

Only a few authors consider the construction of maximin LHDs. For example, Morris and Mitchell ([1995\)](#page-19-0) used simulated annealing to find approximate maximin LHDs for up to five dimensions and up to 12 design points, and a few larger values, with respect to the ℓ^1 - and ℓ^2 -distance measure. van Dam et al. ([2007](#page-18-12)) derived general formulas for two-dimensional maximin LHDs, when the distance measure is ℓ^{∞} or ℓ^1 , while for the ℓ^2 -distance measure (approximate) maximin LHDs up to 1000 design points were obtained by using a branch-and-bound algorithm and constructing (adapted) periodic designs. Ye et al. ([2000\)](#page-19-6) proposed an exchange algorithm for finding approximate maximin symmetric LHDs. The symmetry property is used as a compromise between computing effort and design optimality. Jin et al. ([2005\)](#page-18-0) described an enhanced stochastic evolutionary (ESE) algorithm for finding approximate maximin LHDs. They also apply their method to other space-filling criteria. The Statistics Toolbox of Matlab also contains a function lhsdesign to generate approximate maximin LHDs. This function randomly generates a number of LHDs and picks the one with the largest separation distance. Although this method is very fast, other methods generally result in much better space-filling LHDs. To asses the quality of approximate maximin LHDs, van Dam et al. [\(2009](#page-18-13)) generated upper bounds on the separation distance for certain classes of maximin LHDs. By comparing the separation distances of LHDs to these bounds, we can get an indication of their quality.

There is much more literature related to maximin designs that are not restricted to LHDs. Note that a maximin design is certainly space-filling, but not necessarily non-collapsing.

First of all, the problem of finding the maximal common radius of *n* circles which can be packed into a square is equivalent to the maximin design problem in two dimensions. Melissen ([1997\)](#page-19-7) gives a comprehensive overview of the historical developments and state-of-the-art research in this field. For the ℓ^2 -distance measure in the two-dimensional case, optimal solutions are known for $n \leq 30$ and $n = 36$, see e.g., Kirchner and Wengerodt [\(1987\)](#page-18-14), Peikert et al. ([1991\)](#page-19-8), Nurmela and Östergård [\(1999](#page-19-9)), and Markót and Csendes ([2005\)](#page-19-10). Furthermore, many good approximating solutions have been found for $n \geq 31$; see the Packomania website of Specht ([2008\)](#page-19-11). Baer [\(1992](#page-18-15)) solved the maximum ℓ^{∞} -circle packing problem in a *k*-dimensional unit cube. The ℓ^1 -circle packing problem in a square has been solved for many values of *n*; see Fejes Tóth ([1971\)](#page-18-16) and Florian [\(1989\)](#page-18-17).

Secondly, the maximin design problem has been studied in location theory. In this area of research, the problem is usually referred to as the *max-min facility dispersion problem* (see Erkut [1990](#page-18-18)). Facilities are placed such that the minimal distance to any other facility is maximal. Again, the resulting solution is certainly space-filling, but not necessarily non-collapsing. A few papers consider maximin designs in higher dimensions, e.g., Trosset [\(1999](#page-19-12)), Locatelli and Raber ([2002\)](#page-18-19), Stinstra et al. ([2003\)](#page-19-13), and Dimnaku et al. ([2005\)](#page-18-20). These papers describe nonlinear programming heuristics to find approximate maximin designs. In most papers, a rectangular design space is assumed, but Trosset ([1999\)](#page-19-12), Stinstra et al. ([2003\)](#page-19-13) and Dimnaku et al. ([2005\)](#page-18-20) also specifically consider design spaces with different shapes.

Audze-Eglais LHDs are also constructed by only a few authors. The criterion was first introduced by Audze and Eglais [\(1977](#page-17-1)) and is based on the analogy of minimizing forces between charged particles. In Bates et al. [\(2004](#page-18-2)), the problem of finding Audze-Eglais LHDs is formulated and a permutation genetic algorithm is used to generate them. Liefvendahl and Stocki ([2006\)](#page-18-21) compared maximin and Audze-Eglais LHDs and recommend the Audze-Eglais criterion over the maximin criterion. Examples of practical applications of Audze-Eglais LHDs can be found in Rikards et al. [\(2001](#page-19-14)), Bulik et al. [\(2004](#page-18-22)), Stocki ([2005\)](#page-19-15), and Hino et al. [\(2006](#page-18-23)).

There are several other measures proposed in the literature besides maximin and Audze-Eglais, e.g., maximum entropy, minimax, IMSE, and discrepancy. For a good overview, we refer to Koehler and Owen ([1996\)](#page-18-9). In statistical environments, Latin hypercube sampling (LHS) is often used. In such an approach, points on the grid are sampled without replacement, thereby deriving a random permutation for each dimension; see McKay et al. ([1979\)](#page-19-16). Giunta et al. ([2003\)](#page-18-24) give an overview of pseudoand quasi-Monte Carlo sampling, LHS, orthogonal array sampling, and Hammersley sequence sampling. They notice that the basic LHS technique can lead to designs with poor space-filling properties. Extensions to the basic LHS technique are therefore necessary to obtain better designs but these are unfortunately not standard yet in all software packages. Bates et al. [\(1996](#page-18-25)) obtained designs for computer experiments by exploring so-called lattice points and using results from number theory.

Several papers combine space-filling criteria with the Latin hypercube structure. Jin et al. [\(2005](#page-18-0)) described an enhanced stochastic evolutionary algorithm for finding maximum entropy and uniform designs. van Dam ([2008\)](#page-18-26) derived interesting results for two-dimensional minimax LHDs. Rennen et al. [\(2010](#page-19-17)) consider nested maximin LHDs which consist of two separate designs, one being a subset of the other.

This paper is organized as follows. Section [2](#page-4-0) describes how periodic designs can be used to obtain good approximate maximin and Audze-Eglais LHDs. In Sect. [3](#page-7-0), we shortly describe some heuristics found in literature used for this same purpose. The ESE-algorithm of Jin et al. ([2005](#page-18-0)) described in this section and periodic designs are used to generate new approximate maximin and Audze-Eglais LHDs. Computational results for up to ten dimensions and for up to 300 design points, as well as a comparison of the new and existing results, are provided in Sect. [4.](#page-8-0) Finally, Sect. [5](#page-11-0) contains conclusions.

2 Periodic designs

Van Dam et al. ([2007\)](#page-18-12) showed that two-dimensional maximin Latin hypercube designs often have a nice, periodic structure. By constructing (adapted) periodic designs, many maximin LHDs and, otherwise, good LHDs, are found for up to 1000 points. Therefore, extending this idea to higher dimensions seems natural.

Let a *k*-dimensional Latin hypercube design of *n* points be represented by the sequences y_1, \ldots, y_k , with every y_i a permutation of the set $\{0, \ldots, n-1\}$. As in the two-dimensional case, a design is constructed by fixing the first dimension, without loss of generality, to the sequence $y_1 = (0, \ldots, n-1)$ and assigning (adapted) periodic sequences to all other dimensions. Two types of periodic sequences are considered. The first one is the sequence (v_0, \ldots, v_{n-1}) , where

$$
v_i = (i + 1)p \mod (n + 1) - 1, \quad \text{for } i = 0, \dots, n - 1.
$$
 (3)

Here, *p* is the period of the sequence, which is chosen such that $n + 1$ and *p* have no common divisor, i.e., $gcd(n + 1, p) = 1$, resulting in a permutation of the set $\{0, \ldots, n-1\}.$

Note that the periodic designs obtained in this way resemble *lattices*; see e.g., Bates et al. ([1996\)](#page-18-25). The main difference is that lattices are infinite sets of points, which may collapse, and, hence, to construct a (finite) Latin hypercube design a proper subset of non-collapsing lattice points should be chosen. For given *n*, the structure of the lattice will, however, not always lead to a Latin hypercube design with a sufficient number of points. This is in contrast to periodic designs, for which the modulooperator insures that for every combination of periods p_j , with $gcd(n + 1, p_j) = 1$, $j = 2, \ldots, k$, a feasible Latin hypercube design is obtained.

The second type of sequence that is considered is the more general sequence (w_0, \ldots, w_{n-1}) , where $w_i = (s + ip) \mod n$ (note that we changed the modulus), for $i = 0, \ldots, n - 1$. In this case, all starting points $s = 0, \ldots, p$ and all periods $p = 1, \ldots, \lfloor \frac{n}{2} \rfloor$ will be considered. Note, however, that the resulting sequence *w* may no longer be one-to-one, i.e., some values may occur more than once, and, hence, the resulting design may no longer be an LHD. Now, let $r > 0$ be the smallest value for which $w_r = w_0$; it then follows that $r = \frac{n}{\gcd(n, p)}$. When $r < n$ a way to construct a one-to-one sequence of length n is by shifting parts of the sequence by, say, q , and repeating this when necessary. To formulate this more explicitly, for the updated sequence *w* it now holds that

$$
w_i = (s + ip + jq) \mod n,
$$

for $i = jr, ..., (j + 1)r - 1$, and $j = 0, ..., \gcd(n, p) - 1$. (4)

Let *m* represent the modulus and, hence, the type of sequence used, i.e., $m = n + 1$ corresponds to the first type and $m = n$ to the second. For given *n*, we now have to set the parameters (p, q, s, m) for every sequence y_2, \ldots, y_k .

To find the best settings for the parameters it would be best to test all values. However, when the dimension and the number of points increase the number of possibilities increases rapidly. Hence, computing all possibilities gets very time-consuming or even impossible. Therefore, three classes of parameter settings (named A, B, and C) are distinguished. The largest one, class A, consists of checking the following parameter values: $p = 1, ..., \lfloor \frac{n}{2} \rfloor, q = 1 - p, ..., p - 1, s = 0, ..., p$, and $m \in \{n, n + 1\}.$ Testing in three and four dimensions indicated that almost all adapted periodic maximin designs are based on a shift of $1 - p$, -1 , or 1 (as was the case for two dimensions; see van Dam et al. ([2007\)](#page-18-12)). Furthermore, most maximin designs are found to have a starting point equal to either $p - 1$ or p . Class B is therefore set up to be a subset of class A with the aforementioned restrictions on the parameters *q* and *s*. Finally, for the dimensions 5 to 7 the number of possibilities has to be reduced even further, leading to parameter class C, which (based on some more test results) restricts class B to the values $q = 1$ $q = 1$ and $s = p$, leaving the other parameters unchanged. Table 1 shows the different classes used in the computations of the approximate maximin LHDs for each dimension. For the approximate Audze-Eglais LHDs only class C is used.

As an example, consider a three-dimensional adapted periodic LHD of 22 points. For the maximin criterion, a best parameter setting (class A) is found to

Fig. 1 Two-dimensional projection of the three-dimensional LHD *(y*1*,y*2*,y*3*)* of 22 points

 b **e** $(p_2, q_2, s_2, m_2) = (8, -7, 7, 22)$ and $(p_3, q_3, s_3, m_3) = (3, 0, 2, 23)$ and, hence, the corresponding maximin LHD, with separation distance 69, is defined by the sequences

 $y_1 = (0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21)$ $y_2 = (7, 15, 1, 9, 17, 3, 11, 19, 5, 13, 21, 0, 8, 16, 2, 10, 18, 4, 12, 20, 6, 14)$ $y_3 = (2, 5, 8, 11, 14, 17, 20, 0, 3, 6, 9, 12, 15, 18, 21, 1, 4, 7, 10, 13, 16, 19).$ (5)

Thus, y_3 is a periodic sequence, with $m = n + 1$, and y_2 is an adapted periodic sequence, with $m = n$ and $q_2 = -7$. Note that to obtain a one-to-one sequence, the second part of y_2 , i.e., $(0, 8, \ldots, 14)$, is formed by shifting the first part of y_2 , i.e., *(*7*,* 15*,...,* 21*)*, by −7. The periods and shift are clearly visible in the twodimensional projection of the LHD in Fig. [1](#page-6-1). In this figure the *y*3-values are depicted at the design points.

Like in the two-dimensional case, it may happen that for a given *n* the corresponding maximin LHD has a separation distance that is smaller than the distance of a design of *n* − 1 points. For these *n*, however, better designs can usually be derived

by adding an extra "corner point" to the LHD of *n*−1 points. In this way, a monotone nondecreasing sequence of separation distances was found for all dimensions.

3 Other methods

3.1 Enhanced stochastic evolutionary algorithm

Besides restricting ourselves to a certain class of LHDs, we can also generate good maximin or Audze-Eglais LHDs using heuristics. The ESE-algorithm of Jin et al. [\(2005](#page-18-0)) is one of the methods developed for this purpose and is used in this paper to generate new approximate maximin and Audze-Eglais LHDs.

This method starts with an initial design and tries to find better designs by iteratively changing the current design. To determine if a new design is accepted, a threshold-based acceptance criterion is used. This criterion is controlled in the outer loop of the algorithm. In the inner loop of the algorithm new designs are explored.

The inner loop explores the design space as follows. At each iteration, the algorithm creates a fixed number of new designs by exchanging two randomly chosen elements. The new design with the largest separation distance or with the smallest Audze-Eglais objective value is then compared to the current design with a threshold criterion. The criterion is such that it ensures that better designs are always accepted and that worse designs can also be accepted with a certain probability. If the new design is accepted, it replaces the current design. This process is repeated for a user defined number of iterations.

The outer loop controls the threshold value. After the inner loop is completed, the outer loop determines how much improvement is made in the inner loop. If the amount of improvement is above a certain level, the algorithm starts an improving process in which it tries to rapidly find a local optimum. It does this by lowering the threshold value and thus accepting less deteriorations in the inner loop. If too little improvement is made, an exploration process is started which is intended to escape from a local optimum. The threshold value is first rapidly increased to move away from a local optimum and later slowly decreased to find better designs after moving away. The final design of the algorithm is the best design found during all iterations of the inner loop.

For a more detailed description of the algorithm, we refer to the original paper of Jin et al. ([2005\)](#page-18-0). To find maximin and Audze-Eglais LHDs, we implemented the ESE-algorithm in Matlab. The parameters of the algorithm were set to the values suggested in Jin et al. [\(2005](#page-18-0)). The only adjustment we made to the original algorithm is in the choice of stopping criterion. Instead of stopping after a fixed number of runs of the outer loop, our criterion is to stop when in the last 1000 runs of the outer loop no improvement is made.

3.2 Simulated annealing

Another heuristic used to find maximin LHDs is simulated annealing. Morris and Mitchell [\(1995](#page-19-0)) were the first to apply simulated annealing for this purpose. The simulated annealing method tries to find good designs by iteratively changing a random starting design. These changes are chosen randomly from a predefined class of possible changes. For each design, these possible changes define a set of designs which are called the neighborhood of the design. Before a change is accepted, the new neighbor design obtained by applying the selected change is evaluated. If a change improves the current design, the change is always accepted. A key characteristic of simulated annealing is however that changes which result in a worse design can also be accepted. This enables simulated annealing to escape from local optima. A worse design is accepted with a probability which depends on two factors. Firstly, designs which are only slightly worse are accepted with a higher probability than design which are much worse. Secondly, the acceptance probability is changed during the course of the algorithm. Generally, worse designs are accepted with a higher probability at the beginning of the algorithm than at the end.

Besides Morris and Mitchell ([1995\)](#page-19-0), also Husslage ([2006\)](#page-18-27) used simulated annealing for finding maximin LHDs. One of the main differences between the two methods is the used objective function. Husslage [\(2006\)](#page-18-27) directly used the separation distance of a design, whereas Morris and Mitchell ([1995\)](#page-19-0) used a surrogate measure ϕ_p . This measure also takes into account the number of pairs of points with a certain distance between them. By including this information, it is easier to decide which design is best if they have the same separation distance. This surrogate measure is also used by other authors like Jin et al. [\(2005](#page-18-0)) and Palmer and Tsui ([2001\)](#page-19-18).

Simulated annealing and ESE are similar in many respects. Both algorithms create new designs by changing a current design. Furthermore, both algorithms accept worse designs with a positive probability. The change of this acceptance probability in simulated annealing is similar to the change of the threshold value in the outer loop of ESE. The main difference between the two methods is that the ESE-algorithm creates several new designs and compares the best of these designs to the current design, whereas simulated annealing only creates one new design. The ESE-algorithm can thus be regarded as an enhancement of simulated annealing.

3.3 Permutation genetic algorithm

To obtain Audze-Eglais LHDs, Bates et al. ([2004\)](#page-18-2) used a permutation genetic algorithm. The genetic algorithm starts with generating a set of LHDs called a "population". The Audze-Eglais distance of each design in this population is then calculated. Based on these distances, a subset of designs is selected using so-called elitist and tournament selection. A new population of designs is created by applying mutation and crossover operations to the selected designs. By repeatedly selecting and creating designs, the Audze-Eglais distances of the LHDs in the population gradually increase. Results of this algorithm were reported by Bates et al. [\(2004](#page-18-2)) for eight different combinations of *n* and *k*. In Sect. [4,](#page-8-0) we make a comparison between these results, the designs obtained with periodic designs, and the designs obtained with ESE.

4 Computational results

Using (adapted) periodic designs and the ESE-algorithm, approximate maximin and Audze-Eglais LHDs have been obtained for the cases described in Table [2.](#page-9-0) All computations have been performed on PCs with a 2*.*8-GHz Pentium D processor. For the cases with $n > 100$, a limit of 6 hours was imposed on the calculation time.

Fig. 2 Ratio between separation distance of ESE and periodic designs

Table [5](#page-11-1) shows the squared ℓ^2 -separation distance of the (approximate) maximin LHDs that were obtained by applying periodic designs and of those obtained by the ESE-algorithm. The column for two-dimensional periodic designs contains the results obtained in van Dam et al. [\(2007](#page-18-12)). From this table it can be seen that (adapted) periodic designs work particularly well for larger values of *n* and small *k*.

For dimension 2 to 4, Fig. [2](#page-9-1) shows the ratio between the squared separation distance of ESE and periodic designs. A ratio larger than one indicates that the ESE design is better than the (adapted) periodic design. A break-even point, i.e., a point (or, better, an interval) where the preference shifts from the designs found by ESE to (adapted) periodic designs, is clearly visible in this figure. Furthermore, these breakeven points seem to increase with the dimension of the design and it is to be expected that break-even points for *k*-dimensional designs, with $k \geq 5$, will occur for larger values of *n*, i.e., *n >* 250. Because all six- and seven-dimensional (adapted) periodic designs, of 3 to 100 points, are dominated by the designs found by ESE, the former are not computed for larger dimensions.

| \boldsymbol{n} | 3 dim | | 4 dim | | 5 dim | | 6 dim | | 7 dim | | 8 dim | | 9 dim | |
|------------------|-------|----|-------|----|-------|----|---|-----|-------|-----|-------|----|-------|-----|
| | | | | | | | M&M ESE | | | | | | | |
| 3 | 6 | 6 | 7 | 7 | 8 | 8 | | | | | | | | |
| 4 | 6 | 6 | 12 | 12 | 14 | 14 | | | | | | | | |
| 5 | 11 | 11 | 15 | 15 | 24 | 24 | | | | | | | | |
| 6 | 14 | 14 | 22 | 22 | 32 | 32 | 40 | 40 | | | | | | |
| 7 | 17 | 17 | 28 | 28 | 40 | 40 | | | 61 | 61 | | | | |
| 8 | 21 | 21 | 42 | 42 | 50 | 50 | | | | | 91 | 89 | | |
| 9 | 22 | 22 | 42 | 42 | 61 | 61 | | | | | | | 126 | 126 |
| 10 | 27 | 27 | 50 | 47 | 82 | 82 | | | | | | | | |
| 11 | 29 | 30 | 55 | 55 | 80 | 80 | | | | | | | | |
| 12 | 36 | 36 | 63 | 63 | 91 | 91 | 139 | 136 | | | | | | |
| 13 | | | | | | | | | | | | | | |
| 14 | | | | | | | | | 219 | 215 | | | | |

Table 3 Squared ℓ^2 -separation distance of designs found by Morris and Mitchell ([1995\)](#page-19-0) vs. the ESEalgorithm

In Table [3](#page-10-0), we compare the LHDs found by Morris and Mitchell ([1995\)](#page-19-0) and the ESE-algorithm. The ESE-algorithm is able to match the results of Morris and Mitchell ([1995\)](#page-19-0) for most combinations of *k* and *n*. Only for the cases $(k, n) = (4, 10)$, *(*6*,* 12*)*, *(*7*,* 14*)*, and *(*8*,* 8*)* are slightly worse designs obtained. Three of these four design satisfy the property that $n = k$ or $n = 2k$. According to Morris and Mitchell [\(1995](#page-19-0)), these designs exhibit special symmetric properties; they refer to them as *foldover designs*. These special properties are probably the main explanation for the better results in these cases. For the case $(k, n) = (3, 11)$, we obtained an improved (and optimal) design. Furthermore, using a branch-and-bound algorithm, the threedimensional designs of up to 15 points have been verified to be optimal (van Dam et al. [2009\)](#page-18-13). From the above results, we can conclude that performances of the ESEalgorithm and the simulated annealing algorithm of Morris and Mitchell [\(1995](#page-19-0)) are closely matched. However, the numerical results of Morris and Mitchell [\(1995\)](#page-19-0) are probably too limited to be useful in most practical applications.

We also compared the ESE results with the SA results in Husslage [\(2006](#page-18-27)) and saw that the ESE-algorithm gives better or equally good results for most combination of *k* and *n*. For only nine combinations the results are better of the SA algorithm and for 7 percent of the compared combinations the results are equally good. However, especially for larger values of *n*, the ESE algorithm found many designs with a more than 15 percent higher separation distance.

The results obtained for the Audze-Eglais measure are given in Table [6](#page-15-0). We can easily see that the results of the ESE-algorithm are better for almost all cases. It is likely that by running ESE for some more starting solutions, better or equally good designs can be found for all cases. The ESE algorithm thus outperforms the periodic designs for the Audze-Eglais measure.

When we compare the results with those found by Bates in Table [4,](#page-11-2) we see that the ESE-algorithm gives better or equally good results. This shows that the ESEalgorithm is quite successful in finding LHDs with a good Audze-Eglais value.

5 Conclusions

This paper discusses existing and new results in the field of maximin and Audze-Eglais Latin hypercube designs. Such designs play an important role in the area of computer simulation. The new results are obtained using two heuristics. The first heuristic is based on the observation that many optimal LHDs, and two-dimensional LHDs in particular, exhibit a periodic structure. By considering periodic and adapted periodic designs, approximate maximin LHDs for up to seven dimensions and for up to 300 design points are constructed. The second heuristic uses the ESE-algorithm of Jin et al. [\(2005](#page-18-0)) to find approximate maximin LHDs for up to ten dimensions. These new results are compared to each other and to existing results obtained with simulated annealing and permutation genetic algorithms. In most cases, the ESE-algorithm resulted in the best maximin and Audze-Eglais LHDs. However when the number of design points is large with respect to the dimension of the design, the periodic designs tend to work better. [Appendix](#page-11-3) gives the squared ℓ^2 -separation distances and Audze-Eglais values of the designs described in this paper. All corresponding designs can be downloaded from the website <http://www.spacefillingdesigns.nl>.

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Appendix: Tables of numerical results

| \mathbf{n} | | 2 dim | | | | | | 3 dim 4 dim 5 dim 6 dim 7 dim 8 dim 9 dim 10 dim | | | | | | | |
|--------------|------------|---------------|---------|-------------------------|------------------------|--|--|--|-------|-------|----------------------------|----|---|--------------------------|------|
| | | | | | | | | | | | | | ESE Per ESE Per ESE Per ESE Per ESE Per ESE Per ESE | ESE | ESE |
| | $2\quad 2$ | $\frac{2}{2}$ | | $3 \quad 3$ | | | | | | | 4 4 5 5 6 6 7 7 8 | | | $\overline{}$ | - 10 |
| | $3\quad 2$ | $\mathbf{2}$ | 6 | $\overline{\mathbf{3}}$ | 7 | | | | | | | | 4 8 5 12 6 13 7 14 | 18 | -19 |
| | $4\quad 5$ | 5. | 6 | 6 | | | | | | | 12 12 14 11 20 15 21 16 26 | | | 28 | 33 |
| | 5 5 | | 5 11 6 | | | | | | | | 15 12 24 11 27 15 32 16 40 | | | 43 | 50 |
| | 6 5 | | | | 5 14 14 22 16 32 23 40 | | | | | 28 47 | | 29 | -53 | 61 | 68 |
| | 7 8 | | 8 17 14 | | 28 16 40 | | | | 23 52 | 28 61 | | 31 | 70 | 80 | 89 |

Table 5 Squared ℓ^2 -separation distance found using periodic designs (PD) vs. the ESE-algorithm (ESE)

Table 5 (*Continued*)

| ESE Per ESE Per ESE ESE ESE ESE ESE Per Per Per Per 203 464 803 1592 1989 37 50 196 415 634 1167 893 1074 49 | ESE 2397 2492 2566 | ESE 2828 2893 |
|---|-----------------------------|---------------------|
| | | |
| | | |
| 830 1203 1639 2041 37 52 206 213 478 415 663 893 1113 50 | | |
| 37 52 206 213 490 421 850 692 1230 1662 2132 51 917 1161 | | 3006 |
| 37 58 217 213 504 883 709 1274 1003 1734 1231 2203 52 455 | 2686 | 3134 |
| 219 2234 53 37 58 216 515 455 894 1340 1003 1808 1241 716 | 2713 | 3261 |
| 233 2356 37 58 209 534 932 760 1359 1019 1856 1288 54 477 | 2805 | 3339 |
| 2429 55 40 58 230 243 546 483 956 760 1421 1082 1896 1325 | 2935 | 3452 |
| 58 230 243 558 515 982 1431 1104 2003 1358 2444 56 41 784 | 3021 | 3551 |
| 41 58 249 261 574 515 1007 846 1488 1136 2024 1479 2554 57 | 3119 | 3651 |
| 1554 58 41 61 245 261 594 539 1035 846 1166 2043 1479 2650 | 3187 | 3795 |
| 59 41 61 254 266 609 544 1063 849 1564 1223 2136 1509 2733 | 3297 | 3889 |
| 2232 41 65 261 273 568 1094 904 1631 1242 2796 60 618 1577 | 3420 | 4090 |
| 41 65 274 630 620 2266 2868 61 266 1128 904 1667 1258 1615 | 3525 | 4158 |
| 41 65 269 283 657 620 1150 1715 2345 2977 62 934 1306 1680 | 3636 | 4313 |
| 63 50 65 281 297 670 620 1178 1781 2376 3056 967 1380 1680 | 3690 | 4355 |
| 50 65 278 297 684 625 1206 985 1804 2452 3097 64 1430 1769 | 3820 | 4514 |
| 50 290 694 630 1216 997 1868 2492 3219 65 68 314 1430 1786 | 3932 | 4581 |
| 50 68 299 314 718 666 1261 1874 1476 2543 3279 66 1050 1857 | 4004 | 4769 |
| 50 74 294 314 735 666 1299 1072 1954 2638 3399 67 1482 1868 | 4081 | 4942 |
| 68 50 74 306 314 746 685 1330 1087 1983 1538 2693 1940 3453 | 4212 | 4995 |
| 69 50 306 324 765 698 1351 1112 2028 2746 3520 74 1588 1965 | 4317 | 5127 |
| 50 74 314 325 779 1378 2094 2838 3588 70 716 1150 1633 2130 | 4464 | 5276 |
| 50 74 314 325 793 716 1413 2141 2871 2130 3749 71 1150 1644 | 4548 | 5437 |
| 72 50 74 314 341 810 750 1430 1203 2136 2960 3810 1768 2177 | 4666 | 5556 |
| 73 50 74 329 350 834 759 1462 1229 2197 3042 2206 3932 1768 | 4776 | 5661 |
| 50 74 341 1512 1229 2291 3120 2244 3941 74 350 842 767 1774 | 4915 | 5817 |
| 50 80 341 350 867 1530 1274 2303 1862 3157 2295 4073 75 771 | 5006 | 5937 |
| 50 341 363 882 813 1569 1300 2387 1935 3218 2375 4178 76 85 | 5179 | 6111 |
| 77 50 85 341 363 894 823 1591 1308 2433 1947 3323 2403 4266 | 5222 | 6272 |
| 50 85 371 387 910 844 1621 1382 2479 2014 3387 2505 4390 78 | 5385 | 6384 |
| 50 374 387 927 848 1639 1382 2498 2037 3474 2525 4465 79 85 | 5535 | 6466 |
| 50 374 949 873 1691 1395 2554 2037 80 85 403 3550 2590 4565 | 5577 | 6653 |
| 50 406 1406 81 85 381 963 916 1730 2648 2064 3619 2642 4679 | 5748 | 6780 |
| 85 989 82 50 374 406 938 1742 1475 2680 2141 3669 2753 4719 | 5859 | 6935 |
| 417 1002 1762 1501 2696 3723 2767 4848 83 50 90 374 940 2141 | 5976 | 7094 |
| 1021 50 406 426 967 1818 1534 2790 2229 3870 2838 4920 84 90 | 6119 | 7256 |
| 1043 50 413 426 967 1866 1552 2819 2232 3919 2874 5032 85 90 | 6212 | 7357 |
| 1053 50 413 428 967 1882 1573 2875 2375 3958 3103 5164 86 97 | 6346 | 7532 |
| 1073 1934 50 97 413 428 976 1598 2913 2375 4095 3103 5225 87 | 6469 | 7639 |
| 437 1086 1050 1954 2975 50 97 434 1685 2398 4166 3183 5340 88 | 6660 | 7877 |
| 50 443 1102 1050 1990 89 97 426 1690 3067 2400 4176 3183 5450 | 6750 | 7950 |

| n | 2 dim | 3 dim | | 4 dim | | 5 dim | | 6 dim | | 7 dim | | | | 8 dim 9 dim 10 dim |
|-----|-----------------|-------|--|---------------------|--|---------|--|---------|--|---------|--|-----|-----|--------------------|
| | ESE Per ESE Per | | | ESE Per | | ESE Per | | ESE Per | | ESE Per | | ESE | ESE | ESE |
| 255 | | | | 1417 1923 4100 5122 | | | | | | | | | | |
| 260 | | | | 1445 1971 4164 5236 | | | | | | | | | | |
| 265 | | | | 1449 2021 4182 5519 | | | | | | | | | | |
| 270 | | | | 1464 2144 4361 5656 | | | | | | | | | | |
| 275 | | | | 1478 2150 4487 5746 | | | | | | | | | | |
| 280 | | | | 1493 2184 4388 6023 | | | | | | | | | | |
| 285 | | | | 1501 2209 4607 6094 | | | | | | | | | | |
| 290 | | | | 1476 2269 4722 6380 | | | | | | | | | | |
| 295 | | | | 1526 2354 4726 6590 | | | | | | | | | | |
| 300 | | | | 1542 2409 4898 6604 | | | | | | | | | | |

Table 5 (*Continued*)

Table 6 Audze-Eglais values found using periodic designs (PD) vs. the ESE-algorithm (ESE)

| \boldsymbol{n} | 2 dim | | 3 dim | | 4 dim | | 5 dim | | 6 dim | 7 dim 8 dim | | 9 dim | $10 \dim$ |
|------------------|------------|-------|------------|-------|------------|-------|------------|-------|------------|-------------|------------|------------|------------|
| | ESE | Per | ESE | Per | ESE | Per | ESE | Per | ESE | ESE | ESE | ESE | ESE |
| 2 | 0.500 | 0.500 | 0.333 | 0.333 | 0.250 | 0.250 | 0.200 | 0.200 | 0.167 | 0.143 | 0.125 | 0.111 | 0.100 |
| 3 | 0.900 | 0.900 | 0.611 | 0.611 | 0.386 | 0.450 | 0.321 | 0.362 | 0.250 | 0.230 | 0.193 | 0.200 | 0.151 |
| 4 | 1.000 | 1.000 | 0.642 | 0.642 | 0.454 | 0.489 | 0.367 | 0.382 | 0.300 | 0.260 | 0.225 | 0.201 | 0.180 |
| 5 | 1.298 | 1.390 | 0.727 | 0.891 | 0.509 | 0.658 | 0.401 | 0.527 | 0.336 | 0.287 | 0.250 | 0.222 | 0.200 |
| 6 | 1.521 | 1.521 | 0.794 | 0.800 | 0.561 | 0.594 | 0.431 | 0.476 | 0.358 | 0.307 | 0.268 | 0.238 | 0.215 |
| 7 | 1.598 | 1.598 | 0.867 | 0.975 | 0.599 | 0.694 | 0.464 | 0.532 | 0.376 | 0.322 | 0.282 | 0.250 | 0.225 |
| 8 | 1.804 | 1.879 | 0.921 | 0.960 | 0.619 | 0.696 | 0.488 | 0.538 | 0.398 | 0.334 | 0.292 | 0.260 | 0.234 |
| 9 | 1.935 | 1.935 | 0.971 | 1.052 | 0.660 | 0.742 | 0.504 | 0.567 | 0.414 | 0.349 | 0.301 | 0.267 | 0.240 |
| 10 | 2.066 | 2.066 | 1.020 | 1.085 | 0.686 | 0.744 | 0.515 | 0.556 | 0.425 | 0.360 | 0.311 | 0.273 | 0.246 |
| 11 | 2.196 | 2.279 | 1.069 | 1.137 | 0.709 | 0.785 | 0.536 | 0.612 | 0.434 | 0.369 | 0.319 | 0.281 | 0.250 |
| 12 | 2.273 | 2.273 | 1.095 | 1.163 | 0.724 | 0.785 | 0.551 | 0.589 | 0.441 | 0.375 | 0.326 | 0.287 | 0.256 |
| 13 | 2.401 | 2.487 | 1.128 | 1.191 | 0.746 | 0.825 | 0.563 | 0.632 | 0.453 | 0.381 | 0.331 | 0.292 | 0.261 |
| 14 | 2.476 | 2.476 | 1.167 | 1.252 | 0.762 | 0.829 | 0.575 | 0.635 | 0.462 | 0.385 | 0.335 | 0.296 | 0.265 |
| 15 | 2.578 | 2.643 | 1.194 | 1.255 | 0.775 | 0.818 | 0.583 | 0.636 | 0.470 | 0.393 | 0.339 | 0.299 | 0.268 |
| 16 | 2.666 | 2.683 | 1.221 | 1.290 | 0.791 | 0.848 | 0.589 | 0.642 | 0.477 | 0.398 | 0.341 | 0.302 | 0.271 |
| 17 | 2.721 | 2.721 | 1.246 | 1.340 | 0.805 | 0.866 | 0.600 | 0.656 | 0.483 | 0.404 | 0.347 | 0.305 | 0.273 |
| 18 | 2.819 | 2.848 | 1.271 | 1.337 | 0.816 | 0.875 | 0.609 | 0.655 | 0.488 | 0.408 | 0.350 | 0.307 | 0.275 |
| 19 | 2.890 | 2.984 | 1.292 | 1.374 | 0.827 | 0.895 | 0.615 | 0.667 | 0.492 | 0.413 | 0.354 | 0.310 | 0.277 |
| 20 | 2.959 | 2.962 | 1.318 | 1.394 | 0.835 | 0.907 | 0.620 | 0.681 | 0.496 | 0.416 | 0.358 | 0.313 | 0.278 |
| 21 | 3.025 | 3.033 | 1.339 | 1.408 | 0.847 | 0.914 | 0.625 | 0.671 | 0.501 | 0.419 | 0.361 | 0.316 | 0.281 |
| 22 | 3.070 | 3.070 | 1.357 | 1.426 | 0.856 | 0.922 | 0.632 | 0.687 | 0.505 | 0.422 | 0.363 | 0.318 | 0.283 |
| 23 | 3.138 | 3.159 | 1.377 | 1.454 | 0.868 | 0.925 | 0.638 | 0.693 | 0.510 | 0.425 | 0.366 | 0.321 | 0.285 |
| 24 | 3.197 | 3.201 | 1.396 | 1.458 | 0.875 | 0.931 | 0.644 | 0.677 | 0.513 | 0.427 | 0.368 | 0.323 | 0.287 |
| 25 | 3.254 | 3.293 | 1.412 | 1.485 | 0.884 | 0.940 | 0.648 | 0.701 | 0.516 | 0.430 | 0.370 | 0.324 | 0.289 |

Table 6 (*Continued*)

| n | 2 dim | | 3 dim | | 4 dim | | 5 dim | | 6 dim | 7 dim | 8 dim | 9 dim | 10 dim |
|-----|------------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|------------|
| | ESE | Per | ESE | Per | ESE | Per | ESE | Per | ESE | ESE | ESE | ESE | ESE |
| 67 | 4.636 | 4.642 | 1.761 | 1.812 | 1.047 | 1.100 | 0.744 | 0.787 | 0.581 | 0.478 | 0.407 | 0.355 | 0.315 |
| 68 | 4.661 | 4.681 | 1.766 | 1.819 | 1.049 | 1.096 | 0.745 | 0.790 | 0.581 | 0.478 | 0.407 | 0.356 | 0.316 |
| 69 | 4.683 | 4.713 | 1.771 | 1.818 | 1.052 | 1.104 | 0.746 | 0.791 | 0.582 | 0.479 | 0.408 | 0.356 | 0.316 |
| 70 | 4.703 | 4.710 | 1.775 | 1.818 | 1.053 | 1.100 | 0.747 | 0.790 | 0.583 | 0.480 | 0.408 | 0.356 | 0.316 |
| 71 | 4.727 | 4.742 | 1.780 | 1.831 | 1.055 | 1.108 | 0.747 | 0.795 | 0.584 | 0.480 | 0.409 | 0.357 | 0.317 |
| 72 | 4.743 | 4.746 | 1.784 | 1.829 | 1.057 | 1.103 | 0.749 | 0.791 | 0.584 | 0.481 | 0.409 | 0.357 | 0.317 |
| 73 | 4.763 | 4.781 | 1.789 | 1.836 | 1.059 | 1.106 | 0.749 | 0.794 | 0.585 | 0.481 | 0.409 | 0.357 | 0.317 |
| 74 | 4.781 | 4.820 | 1.793 | 1.842 | 1.061 | 1.112 | 0.751 | 0.795 | 0.586 | 0.482 | 0.410 | 0.358 | 0.317 |
| 75 | 4.803 | 4.817 | 1.796 | 1.847 | 1.063 | 1.112 | 0.752 | 0.797 | 0.586 | 0.482 | 0.410 | 0.358 | 0.318 |
| 76 | 4.823 | 4.828 | 1.801 | 1.850 | 1.064 | 1.110 | 0.753 | 0.796 | 0.587 | 0.483 | 0.410 | 0.358 | 0.318 |
| 77 | 4.838 | 4.853 | 1.805 | 1.851 | 1.066 | 1.112 | 0.755 | 0.799 | 0.587 | 0.483 | 0.411 | 0.359 | 0.318 |
| 78 | 4.863 | 4.883 | 1.809 | 1.847 | 1.068 | 1.114 | 0.755 | 0.795 | 0.588 | 0.484 | 0.411 | 0.359 | 0.318 |
| 79 | 4.882 | 4.934 | 1.812 | 1.863 | 1.070 | 1.119 | 0.756 | 0.801 | 0.589 | 0.484 | 0.411 | 0.359 | 0.318 |
| 80 | 4.895 | 4.922 | 1.816 | 1.864 | 1.071 | 1.117 | 0.757 | 0.799 | 0.589 | 0.484 | 0.412 | 0.359 | 0.319 |
| 81 | 4.920 | 4.942 | 1.820 | 1.869 | 1.072 | 1.120 | 0.758 | 0.802 | 0.590 | 0.485 | 0.412 | 0.360 | 0.319 |
| 82 | 4.936 | 4.944 | 1.824 | 1.862 | 1.074 | 1.120 | 0.759 | 0.801 | 0.590 | 0.485 | 0.413 | 0.360 | 0.319 |
| 83 | 4.949 | 4.949 | 1.827 | 1.879 | 1.076 | 1.126 | 0.760 | 0.805 | 0.591 | 0.486 | 0.413 | 0.360 | 0.319 |
| 84 | 4.968 | 4.992 | 1.831 | 1.876 | 1.077 | 1.122 | 0.761 | 0.802 | 0.591 | 0.486 | 0.413 | 0.360 | 0.320 |
| 85 | 4.985 | 5.014 | 1.834 | 1.879 | 1.079 | 1.124 | 0.761 | 0.804 | 0.592 | 0.486 | 0.414 | 0.360 | 0.320 |
| 86 | 5.003 | 5.014 | 1.838 | 1.882 | 1.081 | 1.125 | 0.762 | 0.804 | 0.592 | 0.487 | 0.414 | 0.361 | 0.320 |
| 87 | 5.019 | 5.060 | 1.842 | 1.891 | 1.082 | 1.130 | 0.763 | 0.808 | 0.593 | 0.487 | 0.414 | 0.361 | 0.320 |
| 88 | 5.034 | 5.047 | 1.845 | 1.885 | 1.083 | 1.130 | 0.764 | 0.805 | 0.594 | 0.487 | 0.414 | 0.361 | 0.320 |
| 89 | 5.056 | 5.096 | 1.848 | 1.895 | 1.085 | 1.133 | 0.765 | 0.810 | 0.594 | 0.488 | 0.415 | 0.361 | 0.321 |
| 90 | 5.070 | 5.063 | 1.852 | 1.885 | 1.086 | 1.131 | 0.766 | 0.807 | 0.594 | 0.488 | 0.415 | 0.361 | 0.321 |
| 91 | 5.086 | 5.113 | 1.854 | 1.890 | 1.088 | 1.134 | 0.766 | 0.809 | 0.595 | 0.489 | 0.415 | 0.362 | 0.321 |
| 92 | 5.104 | 5.114 | 1.858 | 1.902 | 1.089 | 1.135 | 0.767 | 0.809 | 0.595 | 0.489 | 0.416 | 0.362 | 0.321 |
| 93 | 5.119 | 5.122 | 1.861 | 1.903 | 1.090 | 1.136 | 0.768 | 0.810 | 0.596 | 0.489 | 0.416 | 0.362 | 0.321 |
| 94 | 5.130 | 5.143 | 1.864 | 1.900 | 1.092 | 1.138 | 0.769 | 0.810 | 0.596 | 0.490 | 0.416 | 0.362 | 0.321 |
| 95 | 5.151 | 5.177 | 1.867 | 1.909 | 1.093 | 1.138 | 0.769 | 0.813 | 0.597 | 0.490 | 0.416 | 0.362 | 0.322 |
| 96 | 5.163 | 5.183 | 1.870 | 1.910 | 1.094 | 1.139 | 0.770 | 0.811 | 0.597 | 0.490 | 0.417 | 0.363 | 0.322 |
| 97 | 5.177 | 5.179 | 1.872 | 1.915 | 1.096 | 1.138 | 0.771 | 0.814 | 0.598 | 0.490 | 0.417 | 0.363 | 0.322 |
| 98 | 5.198 | 5.223 | 1.876 | 1.915 | 1.097 | 1.142 | 0.771 | 0.812 | 0.598 | 0.491 | 0.417 | 0.363 | 0.322 |
| 99 | 5.211 | 5.244 | 1.879 | 1.923 | 1.098 | 1.143 | 0.772 | 0.815 | 0.599 | 0.491 | 0.417 | 0.363 | 0.322 |
| 100 | 5.223 | 5.221 | 1.882 | 1.921 | 1.099 | 1.143 | 0.773 | 0.812 | 0.599 | 0.491 | 0.418 | 0.363 | 0.322 |

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