

# **DAGWO based secure task scheduling in Multi‑Cloud environment with risk probability**

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### **Abstract**

The Cloud with pay-per-use functioning has attained a great attraction towards on-demand applications. However, the availability of such services by a single data centre is limited, particularly, during the peak demand season. This is due to the fact that it has restricted resource availability. Therefore, the multi-cloud framework has been implemented. In this framework, more clouds get incorporated in a shared manner. All-private, public, or a blend of both may be a multi-cloud environment. The count of virtual machines is higher in the public cloud; however, security is not ensured. So far, most of the works have considered only the metrics like makespan, execution time, and execution time while allocating the tasks. However, the assurances of security while tasks' execution is still an issue in many complex environments. This research work intends to propose a secure task scheduling scheme in the multi-cloud environment with the assessment of risk probability. The suggested study focuses on using the optimization idea to allocate the tasks in the best way possible. As a result, the suggested optimal task allocation includes four objectives like "makespan, execution time, utilization cost, and security constraints (risk evaluation)". Also, a unique hybrid technique called Dragon Aided Grey Wolf optimization (DAGWO) is presented to address this optimization problem. Lastly, the performance of the suggested scheme is compared with theextantapproachesin terms of convergence, energy, makespan, etc. Especially, the risk probability of the proposed model while scheduling 100 tasks is 5.99%, 49.93%, 50.10%, 21.48%, and 31.557% better than existing PSO, WOA, DA, GWO, and MGWO methods respectively.

**Keywords** Multi-Cloud · Resource Allocation · Resource Scheduling · 4 fold Objectives · DAGWO

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#### **Nomenclature**

### **1 Introduction**

Cloud computing is a network-based evolving technology that aims to deliver diferent IT services for a wide range of enterprises and customers on a pay-for-use basis [\[26–](#page-22-0)[28](#page-22-1)]. Cloud computing ofers many advantages with respect to low cost and data usability. Especially, cloud computing promotes improvedresponsiveness to supply chain interruptions [[41](#page-23-0)]. Furthermore, the cloud model is used in review-based recommender systems [\[3\]](#page-21-0), and biometric authentication [[19](#page-22-2)]. In the cloud, data values from data owners are handled by the cloud provider [[6](#page-22-3), [13,](#page-22-4) [55\]](#page-23-1). Moreover, the risks of service degradation and the malicious insider are the issues in the single cloud. However, the multi-cloud environment monitors several cloud infrastructures and prevents dependence on other clouds [\[8](#page-22-5), [9](#page-22-6), [34](#page-23-2), [46](#page-23-3)]. The storage and device partitioning mechanism [\[37\]](#page-23-4) is used as the multi-cloud resource sharing scheme. For the delivery of services to various cloud vendors, four types of architectures are used. They are application duplication, application device partition, application logic partition in fragments, and application data partition in fragments [\[2,](#page-21-1) [3](#page-21-0)].

In the cloud, the data centre handles huge requests for applications from every corner of the world [\[17,](#page-22-7) [38\]](#page-23-5). In each application, there have many contingent and autonomous activities. To execute such requirements as well as to decide the execution order of the tasks, each cloud needs a scheduling strategy. Diferent clouds may get their scheduling strategies [[5,](#page-21-2) [44,](#page-23-6) [51\]](#page-23-7). Task scheduling objectives [[1,](#page-21-3) [11\]](#page-22-8) specifcally involve reducing the time and energy consumption of task execution and optimizing the usage of resources and the ability to balance the loads. Further, reducing work completion time is benefcial for enhancing the customer experience with the drastic rise of the number of cloud users [\[21\]](#page-22-9) [[7](#page-22-10)]. In order to avoid the performance degradation due to the enlargement of resources or waste

obligatedvia unnecessary idle resources, load balancing capability gets improved, which further leads to maximum utilization of VM. The TS algorithm has therefore been found to be NP-complete, and it is not feasible to achieve the optimum solution in a fnite amount of time [\[23,](#page-22-11) [49\]](#page-23-8).

Till date, evolutionary algorithms such as GA and distribution estimation algorithms have been introduced to solve multiple scheduling and mapping problems [\[42,](#page-23-9) [48](#page-23-10), [53,](#page-23-11) [54](#page-23-12)]. With the single-objective approach, schemes like "Min-Min, Max-Min, and Sufrage models" are implemented. However, they are not very extensible or adaptable. In most circumstances, the model merely takes into account the shortest possible time to complete the task, ignoring safety precautions.

The maincontribution of this research study is listed beneath:

- Asecure task scheduling scheme in the multi-cloud environment is proposed with the assessment of risk probability.
- For scheduling the tasks, a four-fold multi-objective constraint-based optimization model is introduced.
- The defined 4 fold-objective constraints are "makespan, execution time, utilization cost, and risk probability". In addition, the weight in the objective function is calculated by the fuzzy triangular membership function.
- A unique hybrid DAGWO algorithm is proposed for optimal task scheduling. The proposed DAGWO model is a combination of traditional DA and GWO algorithms.

**Paper organization** Section [2](#page-2-0) shows therecent studies in the same field. Section [3](#page-5-0) discusses theoverview and problem statement of thesuggested secure task scheduling. Section [4](#page-5-1) shows the cloud setup as well as defned 4 foldobjectives. Section [5](#page-13-0) explaines the obtained outcomes. Section [6](#page-21-4) ends the work.

### <span id="page-2-0"></span>**2 Literature review**

In this section, the latest task scheduling researches are catagorized based on their approaches.

#### **2.1 Related works**

#### a) *Optimization based approaches*

In 2019, Pang *et al*. [[37\]](#page-23-4) have developed an EDA-GA-based hybrid algorithm for task scheduling. At frst, EDA's sampling & probability concepts have beendeployed to givespecifc workable solutions. The ideal scheduling model for assigning the work to VM's was fnally realized.

In 2019, Sanaj and Joe [\[45\]](#page-23-13) have explored CSSA for optimal "multitask scheduling in an IaaS cloud environment". The method continuously generates work schedules, which increases cost-efectiveness. The prior network was prouced using chaotic optimization for task allocation in order to assure superior global convergence.

In 2019, Yusuf [\[30\]](#page-22-12) has introduced a new TS model, where the major intention relied on optimizing the TS by reducing the energy consumption and MS. The MGWO was able to solve the scheduling difficulties, which enhanced system efficiency. Encircling  $\&$  hunting in MGWO have been modifed utilizing the mean value to raise GWO's efectiveness.

In 2020, Neelima and Reddy [[31](#page-22-13)] have developed a new load balancing task scheduling model in the cloud by means of ADA, which was an amalgamation of DA and FF algorithms. The multi-objective function was defned on the basis of"load, processing costs, & completion time" constraints toattain the enhanced performance.

In 2020, Ismayilov and Haluk [[14](#page-22-14)] have investigated a prediction-oriented approach named as "NN-DNSGA-II algorithm" that incorporated NSGA-II with ANN framework. To address the job scheduling issue, the top fve non-prediction-oriented dynamic techniques have beenutilized. The six objectives of the work that is being provided are "improvement of utilization and reliability, decrease in energy cost, makespan, & imbalance level".

#### b) *other approaches*

In 2017, Liu *et al*. [[22](#page-22-15)] have established a multi-task scheduling model, which incorporated a task workload model by considering service quantity and service coefficient. The efects of several workload-oriented task scheduling strategies have been then examined with respect to utilization  $\&$  overall completion time.

In 2019, Panda *et al*. [[36](#page-23-14)] have developed the multi-cloud network, where several clouds were integrated mutually for providing a combined service. In this method, an "allocationaware task scheduling algorithm" was proposedand itdepend on the conventional "Min-Min and Max-Min algorithm".

In 2020, Lavanya *et al*. [\[20\]](#page-22-16) have explored two allocation models termed TBTS and SLA-LB models. Task scheduling in VM with various setups is made easier by TBTS, which scheduled the tasks in batches. Tasks were dynamically scheduled by SLA-LB based on user requirements such as budget as well as deadline.

#### **2.2 Research Gaps**

Table [1](#page-4-0) represents the review of recent works. The following are the limitations of the extant models.

- (i) The SLA-LB model provides minimal complexity. But the main disadvantage was that the energy utilization factor was not evaluated [\[20\]](#page-22-16). Because the cost of energy utilization in the cloud is an important factor in determining the system's efectiveness.
- (ii) In task scheduling, the continuous arrival of tasks is considered for efective system performance, but this factor is not considered in [[22](#page-22-15)].
- (iii) In the ADA model, VM scheduling was not clearly considered and it needs more consideration [\[31\]](#page-22-13). Furthermore, the MGWO model requires attention to task priority [[30\]](#page-22-12).
- (iv) Security is one of the major issues in cloud computing. The related works failed to analyze the security constraints.

To overcome these limitations, a novel DAGWO based task scheduling is proposed in this study. Compared with the related works, this research introduces 4-fold objectives like"makespan, execution time, and utilization cost, along with security constraints". Moreover, a novel DAGWO model is utilized for optimal task scheduling.

<span id="page-4-0"></span>



## <span id="page-5-0"></span>**3 Proposed secure task scheduling model: an overview and problem statement**

### **3.1 An Overview**

The major components of the suggested work are listed below:

- Customer: This is the cloud's service client. The CU can produce their demands with the assistance of the cloud manager.
- Cloud Manager: This is a central entity that handles customers' service requests and receives the status of cloud providers' VMs.
- Cloud Service Provider: This is a cloud distributor that offers on-demand services by setting up virtual machines on actual servers. Each cloud contains a management server that contacts with other manager servers to send customer requests to manage the peak demands. Customers' requests are frequently scattered over several clouds as well as are handled using a distributed model.

Between the components of the cloud model, the following connections occur: "(a) CUcloud manager, (b) cloud manager-CSP, (c) CSP-CSP, (d) CSP-cloud manager, and (e) cloud manager-CU". Furthermore, the overhead will happen to decide on the individual part. Nevertheless, we believe that the fexibility of the cloud model makes these overheads marginal.

Figure [1](#page-6-0) represents a multi-cloud task scheduling scheme. The proposed model aims to produce satisfaction to the customer via efective task scheduling. Initially, users at various felds send a request tothe cloud. The cloud manager manages the service request and receives the cloud provider's VM status. If a task is submitted to the cloud manager, the manager adds this task to a waiting queue as well as locates an active VM to set the task to based on the specifed 4-fold objectives. " Execution time, makespan, utilization cost, and risk probability" are the established 4-fold objectives in this research work. The work that has been presented deals with applying the optimization notion to task allocation. In the VMs, the tasks are continuously assigned, and the scheduling happens concurrently as well. A new hybrid method called DAGWO is presented for optimal scheduling. Furthermore, taking risk probability into account is the main contribution because this parameter guarantees the need for security while scheduling the activity. Each task's VM allocation is done based on the risk probability (and with other metrics). Finally, the efficient task is allocated with minimum execution time, minimum makespan, minimum utilization cost, and high security (low-risk probability).

### <span id="page-5-1"></span>**4 Cloud setup and defned 4 fold‑objective: makespan, execution time, utilization cost, risk probability**

### **4.1 Cloud Setup**

A set of clouds  $c = c_1, c_2, c_3, \ldots, c_M$  is considered in this work. The cloud service provider *CSP<sub>I</sub>* is denoted as  $c_I$ ,  $1 \le I \le M$ . Here, the count of clouds,  $M = 3$ . The number of VMs ( $|v_I|$ ), as well as the scheduling strategy,  $(S<sub>I</sub>)$  will be assigned by the  $CSP<sub>I</sub>$ . The CSP provides the



<span id="page-6-0"></span>**Fig. 1** The framework of the suggested task scheduling approach

clouds or processing certain requests from the user. Within the cloud, resides the PM with a huge number of VMs. The 3-tuple fashion is utilized to represent *cI* as *cI*= <*CSPI* , *SI* ,  $|v_I| >$ ,  $1 \le I \le M$ .

The request from the users is considered as the task. The task set  $t = t_1, t_2, t_3, \ldots, t_L$  is of independent tasks in which  $t_I$ ,  $1 \le I \le L$  are  $I^{th}$  tasks with an instruction set  $ins_I$  (in MI). To process these tasks, a set of VM is considered  $v = v_1, v_2, v_3, \ldots, v_M$ , in which the VMs set

is  $v_I$ ,  $1 \le I \le M$  as well as it is deployed under  $c_I$ . In addition, a virtual machine  $VM_J \in v_I$ ; 1≤*I*≤*M*; 1≤*J*≤|*vI* | has the processing speed in MIPS *PJ*. The total VM in all the clouds is depicted by  $M' = \sum_{l=1}^{m} |v_l|$ .

The execution time of a task ( $t_I$ ,  $1 \le I \le L$ ) on a VM ( $VM_J$ ,  $1 \le J \le M'$ ) is given in the form of the matrix  $\&$  it is reffered as ETC matrix. Mathematically, the ETC matrix is shown in Eq.  $(1)$  [[36](#page-23-14)].

$$
ETC = \begin{cases}\nV M_1 & ETC_{11} & tr_2 & \cdots & tr_L \\
& ETC_{11} & ETC_{21} & \cdots & ETC_{11} \\
& \vdots & \vdots & \vdots & \vdots \\
& VM_{|v_1|} & ETC_{1|v_1|} & ETC_{11|v_1|} & \cdots & ETC_{L|v_1|} \\
& \vdots & \vdots & \vdots & \vdots & \vdots \\
& VM_{\gamma+1} & ETC_{1(\gamma+1)} & ETC_{2(\gamma+1)} & \cdots & ETC_{1(\gamma+1)} \\
& \vdots & \vdots & \vdots & \vdots & \vdots \\
& VM_{\gamma+1} | v_M | ETC_{1(\gamma+|v_M|)} & ETC_{2(\gamma+|v_M|)} & \cdots & ETC_{L(\gamma+|v_M|)}\n\end{cases}
$$
\n(1)

In which,  $\gamma = \sum_{k=1}^{M-1} |\nu_k|$ . Moreover,  $ETC_{IJ}$ ,  $1 \leq I \leq L$ ;  $1 \leq J \leq M'$  is the ratio of the instruction set (in MI) to the processing speed. Mathematically,  $ETC<sub>II</sub>$  is given in Eq. ([2](#page-7-1)).

<span id="page-7-1"></span><span id="page-7-0"></span>
$$
ETC_{IJ} = \frac{ins_I}{P_J} \tag{2}
$$

"The mapping is done in terms of allocating, matching, and scheduling of tasks." The mapping function *F*gets the request set  $r = {r_1, r_2, r_3, ... r_R}$  from the CUset,  $CU = \{CU_1, CU_2, CU_3, \ldots, CU_R\}$ . The request  $r_i = \{t_1, t_2, t_3, \ldots, t_L\}$ ;  $1 \leq I \leq R$ ;  $L' < L$  is assigned to the set of VM's  $v = \{v_1, v_2, v_3, ..., v_M\}$ . Moreover,  $|t| = \sum_{k=1}^{R} |r_k|$ 

The secure mapping of the task onto the cloud, which has the VM set,  $v(F:t\rightarrow v)$ . This work intends to schedule the task optimally concerning the described 4-fold objective "(a) minimized makespan *S*, (b) lower cloud utilization cost of tasks *U*, (c) lower execution time of VM*E*, (d) lower security risk probability *G*." The upcoming section comprehensively portrays these defned 4-fold objectives. The fow chart representation is shown in Fig. [2](#page-8-0).

#### **4.2 Defned 4‑fold objectives**

The objective of this study is evaluated as per Eq.  $(3)$ . This Eq.  $(3)$  is also known as a ftness function and it is defned as: "A ftness value is a form of the objective function which summarises, in a single measure, how close a given design solution is to attaining the defned goals".

<span id="page-7-2"></span>
$$
obj = \min\{(W_1.S) + (W_2.U) + (W_3.E^{task}) + (W_4.G)\}\
$$
 (3)

In which, *S*,*U*,*Etask*, and *G* denotes makespan, utilization cost of task, execution time for all tasks, and security risk probability. The weights  $W_1$ ,  $W_2$ ,  $W_3$ ,  $\& W_4$  are calculated using fuzzy triangular membership function. The weight calculation is described in Eq. [\(4](#page-9-0)). Here *p*, *q*, *r* is the vertices of triangular membership function T(f). Lower boundaryis *p*, medium



<span id="page-8-0"></span>**Fig. 2** Workfow model

boundary with membership value 1 is *q* as well as *r* is the upper boundary with membership value 0.

<span id="page-9-0"></span>
$$
W = \begin{cases} 0 & \text{if } r < f \\ \frac{r - f}{p - f} & \text{if } f \le r \le p \\ \frac{q - r}{q - p} & \text{if } p \le r \le q \\ 0 & \text{if } r \ge q \end{cases} \tag{4}
$$

**Execution time of VM (** $E^{VM}$ **)** "Execution time of VM is the amount of time taken required by task to complete its executionon VM". The mathematical formula  $E<sup>VM</sup>$  is given as per Eq. ([5](#page-9-1)). In which,  $H_{length(I,J)}$  is the task required to execute the instruction length,  $V_{CPU(I,J)}$ is the count of CPU virtual machine *J*, and  $V_{MIPS(I,I)}$  is the VM processing capability as VM *J* [\[10\]](#page-22-17).

<span id="page-9-1"></span>
$$
E^{VM} = \frac{H_{length (I,J)}}{V_{CPU(I,J)} \times V_{MIPS (I,J)}}
$$
(5)

**Utilization cost of tasks,** *U* "Utilization cost is cost or total amount of payment from a cloud user to cloud provider against the utilization of resources to execute tasks". It is rep-resented in Eq. ([6](#page-9-2)). In which, *V<sub>I</sub>* denotes the count of VM  $\&ct_I$  indicates the completion time of VM. The expression for  $ct<sub>I</sub>$  is given in Eq. [\(7](#page-9-3)), where, *pesnumber* is the count of processing element for running a task on a suitable VM, *MIPS* is the execution speed per processing element of a VM, and *R* is the count of tasks [[24](#page-22-18)].

<span id="page-9-3"></span><span id="page-9-2"></span>
$$
U_{task} = \sum_{I=1}^{V_I} PRICE_I * ct_I
$$
 (6)

$$
ct_I = \sum_{J=1}^{R} \frac{H_J.LENGTH}{V_I.pesnumber \times V_I.MIPS}; I \in \{1, 2, \dots, V_I\} \text{ and } J \in \{1, 2, \dots, R\}
$$
 (7)

**Execution time for all tasks** *Etask* **"**The amount of time needed to complete the execution of all tasks is referred as execution time for all tasks" and it is evaluated as per Eq. [\(8\)](#page-9-4). Here,  $R$  is the count of tasks  $[32]$ .

$$
E^{task} = \frac{1}{\text{MAX (execution time)} \times R} \sum_{l=1}^{R} \text{(Execution time of respective } \text{VM} \times \text{Size of the task)} \tag{8}
$$

**Makespan** *S* "Makespan is the cumulative time that the resources are required to complete the execution of all tasks". In general, VM usage is characterized by how well the resources in the cloud are used [[35](#page-23-15)]. The scheduler is delivering efective and good task planning to resources if the makespan is low. The mathematical formula for *S* is given as per Eq. [\(9](#page-9-5)). Here,  $V_I$  denotes the number of VM and  $ct_I$  indicates the completion time of VM.

<span id="page-9-5"></span><span id="page-9-4"></span>
$$
S = \underset{1 \leq l \leq V_I}{\text{MAX}} \quad \left\{ ct_I \right\} \tag{9}
$$

**Security risk probability** "The determination of the chance of a risk occurring is known as risk probability". It is used to measure the security risk in execution that is to evaluate the risk of the scheduling tasks. For a particular task, the scheduling risk probability of tasks  $t_I V M_j$  are computed using Eq. [\(10\)](#page-10-0).

<span id="page-10-0"></span>
$$
\Pr{ob^q(t_I^q, VM_I^q)} = \left\{ \begin{array}{ll} 0 & \text{if } SD_I^q \leq SS_J^q\\ 1 - e^{-(SD_I^q \leq SS_J^q)} & otherwise \end{array} \right. \tag{10}
$$

Here, the notation  $SD_i^q$ ,  $SS_j^q$  denotes the security demand & security services  $t_i$ . "The *J* risk probability of the task is the overall risk probability of the task corresponding to the security service." Fort<sub>I</sub>, if  $SD_j^q < SS_j^q$ , then the risk probability of scheduling  $t_I$  on  $VM_j$  is zero. In addition,  $Prob(t_1, VM_J)$  denotes the probability of  $t_I$  that is attacked during the execution and it is shown in Eq.  $(11)$  $(11)$ .

$$
\Pr ob(t_I, VM_J) = 1 - \prod_q \left(1 - \Pr ob^q\left(t_I^q, VM_J^q\right)\right) \tag{11}
$$

In addition, the risk probability of workfow is an average of the probabilities of all tasks. The composed task's average probability *Prob*(*W*) is being assaulted during the workflow in Eq.  $(12)$ .

<span id="page-10-2"></span><span id="page-10-1"></span>
$$
\Pr ob(W) = \frac{\sum_{t_i \in t} \Pr ob(t_i, VM_j)}{M}
$$
\n(12)

The VM, as well as tasks, are the input solution to DAGWO for optimal scheduling with the consideration of the above-defned objectives. Since 3 PM with 30 sets of VM in each PM is utilized in this research study, the solution to the proposed model looks as in Fig. [3](#page-11-0). In Fig. [3](#page-11-0), the solution encoding of cloud1, cloud 2, & cloud 3 is manifested.

#### **4.3 Proposed DAGWO for optimal task scheduling**

**Dragonfy Algorithm (DA)** The DA [\[15,](#page-22-20) [16](#page-22-21)] was developed based on the inspiration of dragonfies. The three elementary principles like separation, alignment, and cohesion are observed by the swarming behavior of the dragonfies. In addition, each dragonfy should obey the separation operation, alignment operation, food attraction operation, cohesion operation, and enemy distraction operation.

**Grey Wolf Optimization (GWO)** Grey wolves are highly sociable animals with a very rigid hierarchy. They spend the majority of their lives hunting, seeking for prey frst, encircling prey, as well as eventually attacking prey. Grey wolves have four primary societal structures: alpha, beta, delta, & omega, each of which plays a particular role in the group. The GWO algorithm is a new meta-heuristic based on these grey wolf characteristics [\[50\]](#page-23-16) [[47](#page-23-17)] [[12](#page-22-22), [52\]](#page-23-18).

The DA has the beneft ofa higher convergence rate, which includes the ability to solve continuous problems. Similarly, the hunting habits of grey wolves served as an infuence for the development of GWO. Since it has a straightforward structure, it is simple to build, requires less data than other techniques, converges quickly, and avoids local optima when

<span id="page-11-0"></span>

used with composite functions. This research work intends to hybridize the concept of DA and GWO to make the model even stronger for optimal solutions with better convergence. Generally, hybridization of optimization algorithms can solve diverse classifcation and optimization problems [[2,](#page-21-1) [47](#page-23-17)] [\[16,](#page-22-21) [25](#page-22-23), [33,](#page-23-19) [40\]](#page-23-20) [[4](#page-21-5), [39](#page-23-21), [43\]](#page-23-22). The proposed model is referred to as DAGWO, and its step-by-step procedure is described below:

In the beginning, the population of the search agent is initialized as  $X_i = X_1, X_2, \ldots$  $X_W$ . The step vector  $\Delta X_{iter+1}$  in the DA algorithm is deployed to evaluate the individual's movement direction and the mathematical formula is shown in Eq. [\(13\)](#page-11-1). Here, the current iteration is denoted as *iter*. Also, *p*, *b*, *a*, *f*, *and e* represents the separation weight, alignment weight, cohesion weight, food factor, as well as enemy factor. The inertia weight of the step vector is indicated by *w*.

<span id="page-11-1"></span>
$$
\Delta X_{iter+1} = (pO_i + bB_i + aA_i + f, food_i + e. enemy_i) + w\Delta X_{iter}
$$
\n(13)

Compute the value of separation criteria *O*, alignmen t*B*, cohesion criteria *C*, attraction to food resource (*food*) as well asdistractionaway froman enemy (*enemy*) by Eq. ([14](#page-12-0))- Eq. ([18](#page-12-1)), respectievly. In Eq. [\(17\)](#page-12-2) and Eq. [\(18\)](#page-12-1), *X*+,*X*−&*X*signifes the food sourceposition, enemy source, &position of the present individual.

$$
O_i = \sum_{j=1}^{D} X - X_j \tag{14}
$$

<span id="page-12-0"></span>
$$
B_i = \frac{\sum_{j=1}^{D} X_j}{D} \tag{15}
$$

$$
A_i = \frac{\sum_{j=1}^{D} X_j}{D} - X
$$
 (16)

<span id="page-12-2"></span>
$$
food_i = X^+ - X \tag{17}
$$

<span id="page-12-1"></span>
$$
enemy_i = X^- + X \tag{18}
$$

The adaptive knowledge rate of  $i<sup>th</sup>$  dragonfliesat the current iteration is computed by using Eq. [\(19\)](#page-12-3).

<span id="page-12-4"></span><span id="page-12-3"></span>
$$
d_i^{iter} = \frac{1}{1 + e^{-V}}\tag{19}
$$

 $\text{Here,} V = \frac{|fit(X_i^{iter}) - fit(X_{best}^{iter})|}{fit(X_{best}^{iter}) + \epsilon}$  $f_i(x_i) = \frac{f_i(x_i)}{f_i(x_i)} + \frac{f$ and  $\text{fit}(X_{\text{best}}^{\text{iter}})$  is the best fitness value. Also, *k*denotes the constant value, which is employed to avoid the zero division error. Then update the radius of the neighbour. If there is only one neighbour to the present search agent, then update the search agent position by Eq. ([20](#page-12-4)).

$$
X_{iter+1}^i = d_{iter}^i X_{iter}^i + \Delta X X_{iter+1}^i
$$
\n(20)

**Based on the hybrid proposed contribution,** initialize a random variable *rand*. On the basis of this *rand*, the position update gets varied. The threshold value fxed here is 0.5, and if the assigned *rand* is less than the defned threshold value (*rand*<0.5), then update the search agent position utilizing standard GWO, else keep the existing solution as it is. The position update of GWO is evaluated as per Eq. ([21](#page-12-5)).

<span id="page-12-5"></span>
$$
X(\text{iter} + 1) = \frac{X_1 + X_2 + X_3}{3} \tag{21}
$$

In which,  $X_1 = X_\alpha - Y_1 \cdot (s_\alpha),$  $X_2 = X_\beta - Y_2 \cdot (s_\beta),$  $X_3 = X_{\delta} - Y_3 \cdot (s_{\delta}).$ 

 $X_{\alpha}$ ,  $X_{\beta}$  &  $X_{\delta}$  points to the position of the 1st, 2nd and 3rd best solutions.

$$
s_{\alpha} = | h_1.X_{\alpha} - X |,
$$
  
(Here, 
$$
s_{\beta} = | h_2.X_{\beta} - X |,
$$

$$
s_{\delta} = | h_3.X_{\delta} - X |
$$

In which,  $s$ ,  $Y$  are the coefficient vectors and are mathematically defined in Eq. ([22](#page-13-1)) and Eq. [\(23\)](#page-13-2), respectively.

$$
TY = 2e.u_1 - \vec{e}
$$
\n<sup>(22)</sup>

<span id="page-13-2"></span><span id="page-13-1"></span>
$$
h = 2.u_2 \tag{23}
$$

Here, *e* is a constant, which is gradually reduced from 2 to 0 over the count of iterations. In addition,  $u_1$  and  $u_2$  is a random vector within 0 to 1.

Algorithm [1](#page-13-3) manifests the pseudo-code of the DAGWO algorithm. Figure [4](#page-14-0) represents the block diagram of the proposed approach.

### <span id="page-13-0"></span>**5 Results & Discussions**

### **5.1 Simulation procedure**

The suggestedapproach was executed in **Python**. Three clouds as well as three sets of PM were used to make up the simulation scenario. The number of VMs in each group of PM is 30. The whole number of tasks to be completed ranges from 100 to 150 to 175 to 200, correspondingly. DAGWO outperforms other models in terms of optimal scheduling in terms of "migration cost, total cost, energy consumption, response time, and security analysis". The number of VMs used in this evaluation varies. Two scenarios are considered for evaluation. In scenario 1: 30 counts of VM are considered for scheduling the varying count of tasks (i.e.10 VM in each PM). In Scenario 2: 60 counts of VM are considered (i.e. 20

<span id="page-13-3"></span>**Algorithm 1** Pseudo code of DAGWO Algorithm

Algorithm 1: Pseudo code of DAGWO Algorithm
Initialize the population $X_i = X_1, X_2, , X_w$
step vector $\Delta X_{\text{iter+1}}$ is computed using Eq. (12)
While <i>iter</i> $\lt$ <i>iter</i> <sup><i>Max</i></sup>
Evaluate the search agent fitnessby Eq. (3).
Update food and enemyusing Eq. (17) and Eq. (18), respectively.
Update the value of p, b, a, f, e and w of the $i^{th}$ individual.
Compute the value of O, B, C, food and enemy using Eq. (14) Eq. (18), respectively
Compute the adaptive knowledge rate of $ith$ dragonfly <i>iter</i> by Eq. (18)
Update the neighbour radius
Positional update via Eq. (20)
Initialize a random value rand
If rand $\leq 0.5$
Positional update via Eq. (21), the GWO update
else
Keep the existing solution
End if
end while
Terminate



<span id="page-14-0"></span>**Fig. 4** Flow chart of the suggested DAGWO model

counts of VM in each PM). The DAGWO is evaluated over the existing works like PSO [[29](#page-22-24)], WOA [\[45\]](#page-23-13), DA [\[16\]](#page-22-21), GWO [\[47\]](#page-23-17), MGWO [[30](#page-22-12)] and GA [[18](#page-22-25)].

#### **5.2 Convergence analysis**

The convergence of the DAGWO&existing scheme is determined for scenario1 and scenario 2 throughchanging the iterations number from 20, 40, 60, 80, and 100, respectively. The results acquired with scenario1 and scenario 2 are shown in Fig. [5](#page-15-0). As per the defned 4-fold objectives in Eq. ([3\)](#page-7-2), the schemewhich attains the least cost function is said to be the most appropriate method. At the least count of iteration, the cost function achievedvia the DAGWO &extant work is greater. When the number of iterations increases, the cost function of DAGWO, as well as existing work, gets minimized. In the case of scenario1, the cost function of DAGWO had attained a steep fall in between the range 0 to 20 count of iterations. In addition, when analyzing Fig.[4](#page-14-0) b, the cost function of the suggested&extant technique seem to be greater at the 3rd iteration. Till the 65th iterations, the cost function of the proposed seems was little worse than the existing one, Beyond the 65th iteration, the DAGWO had attained the least value as  $(\sim)0.125$ . Thus, The DAGWO clearly showed the small cost function that suggests that the proposed model may schedule the tasks more efectively.

#### **5.3 Evaluation on energy consumption**

Energy management is becoming more crucial in cloud storage as a result of rising energy prices and greater consumption of cloud computing resources. The call of energy-efficient solutions is the reduced the total energy usage of computing, storage, and communication devices. In data centres, optimum energy consumption is increasingly necessary. The current research tends on energy-efficient resource allocation. The energy required for executing a task must be lower, which leads to the expansion of network lifetime. The con-sumed energy by VM is represented in Fig. [6](#page-16-0). In Fig. 6 a, the DAGWO seems to exhibit the least energy value while evaluating the tasks. For the task count =200, the DAGWO is 16.6%, 96.8%, 3.8%, 3.25%, 2.85%and 1.08%superior to the extant PSO, WOA, DA, GWO,MGWO & GA schemesfor optimal scheduling. On the other hand, the energy consumed by the DAGWO in scenario2 is fuctuating while the task execution. However, the overall performance shows that the proposed algorithm is more efective in solving the task scheduling issue with minimal energy consumption.



<span id="page-15-0"></span>**Fig. 5** Convergence analysis: DAGWO &extantschemes for (**a**) scenario 1 (VM=30) & (**b**) scenario 2  $(VM = 60)$ 



<span id="page-16-0"></span>**Fig.** 6 Energy consumption evaluation: DAGWO &extantmodels for (a) scenario1 (VM=30) & (b) scenario2 ( $VM = 60$ )

### **5.4 Evaluation on execution time**

The computation time recorded via the suggested as well as an existing model is repre-sented in Fig. [7](#page-16-1). Further analyzing Fig. 7 a for task count=200, the DAGWO is  $11.1\%$ . 6.9%, 12.2%, 42.85%, 23.07%, and 0.40%better than the extant PSO, WOA, DA, GWO, MGWO,&GA works with minimum execution time. Furthermore, on analyzingscenario2 at task count = 100, the DAGWO model is attained the least execution time ( $\sim 0.28$ ) than the extant PSO, WOA, DA, GWO,& MGWO models. Even though the execution time of the DAGWO seems to have fuctuated for certain task counts, the overall performance of the suggestedscheme is promising as well as it reveals the adopted scheme betterment.

#### **5.5 Evaluation on makespan**

Figure [8](#page-17-0) depicts themakespan evaluation of the suggestedas well asextantschmes. On analyzing the outcomes, at task=150, the DAGWO scheme is  $25\%$ ,  $52\%$ ,  $33.3\%$ ,  $45.45\%$ ,



<span id="page-16-1"></span>**Fig. 7** Execution time evaluation: DAGWO &extantmodels for (**a**) scenario 1 (VM=30) & (**b**) scenario2  $(VM=60)$ 



<span id="page-17-0"></span>**Fig. 8** Makespan evaluation: DAGWO &extant schemes for (**a**) scenario1 (VM=30) & (**b**) scenario2  $(VM=60)$ 

20% and 64.51% superior to the extant PSO, WOA, DA, GWO, MGWO,and GAapproaches respectively with minimal makespan. Then, forscenario2, the suggested modelis attained the small value at scheduling 200 numbers of tasks and it is smaller thanthe extantapproaches.

### **5.6 Evaluation on resource utilization**

Minimal resource utilization is required during the scheduling  $\&$  task execution in order to meet the given aim in Eq. [\(3\)](#page-7-2). The outcomes of the suggestedas well as the extant approachesis depicted in Fig. [9](#page-17-1). The DAGWO has scheduled and completed 175 counts of jobs with the least amount of resource use, according to the scenario 1 analysis. In this instance, the given work outperforms existing methodologies including PSO, WOA, DA, GWO, MGWO, and GA by 50%, 20%, 33.3%, 3.86%, and 3.21%, respectively. The given work, has recorded the lowest value of 0.05 in the scenario 2 condition, which is determined to be a better outcome when compared to the conventional schemes. The analysis so



<span id="page-17-1"></span>**Fig. 9** Resource utilization evaluation: DAGWO &extant models for (**a**) scenario1 (VM=30) & (**b**) scenario2 ( $VM = 60$ )

Count of Tasks	<b>PSO[10]</b>	WOA [45]		DA [16] GWO [47]	MGWO [30]	GA [18]	DAGWO
100	38.264	33.635	33.746	21.447	24.610	55.849	16.839
150	21.996	19.613	18.366	21.577	19.800	31.633	23.174
175	35.997	37.352	31.179	39.325	24.666	38.613	27.438
200	52.141	55.639	48.461	44.231	55.119	54.519	45.906

<span id="page-18-0"></span>**Table 2** Risk probabilityevaluation: DAGWO &extant schemesfor scenario1

demonstrated that the suggested model was superior for task scheduling while utilizing the fewest resources.

#### **5.7 Evaluation on security or risk probability**

The security analysis for adopted as well as extant approaches for scenarios is repre-sented in Table [2](#page-18-0) and [3](#page-18-1). The findings demonstrate that the DAGWO balances excellent security with reduced risk. On analyzing scenario 1, the DAGWO attask=100 is 16.83, and the extant model values are  $PSO = 38.2$ ,  $WOA = 33.63$ ,  $DA = 33.7$ ,  $GWO = 21.4$ ,  $MGWO = 24.6$ , and  $GA = 55.84$ . On observing all other tasks too, the risk probability of the suggestedshows to be lower than the conventional models. Further, in the case of scenario2, the risk probability of the DAGWO is  $63.046$  at task $=200$  and this is the less value. Altogether, the analysis proves that the suggested model has the ability to process the scheduling as well as execution of tasks even in a secured manner.

#### **5.8 Evaluation on throughput**

The throughput is the greatest rate at which tasks may be completed in a given amount of time. It assesses the efectiveness of the scheduling method. Low response time and large execution rate result from a high throughput rate. Table [4](#page-19-0) and [5](#page-19-1) represents the throughput of the adoptedas well asextant schems at various tasks. In both scenarios, the throughput value of the suggested approach is increased with increasing the number of tasks. In Table [4,](#page-19-0) the throughput of the proposed work at scheduling tasks 200 is 36.11%, 33.03%, 29.26%, 44.34%, 0.34%, and 36.11% better than the existing PSO, WOA, DA, GWO, GA and MGWO approaches respectively. For scenario 2, the throughput of the adopted research at scheduling tasks 200 is2.72%, 11.21%, 7.44%, 12.96%, 10.57%, and 13.61% superior to the extantPSO, WOA, DA, GWO, GA and MGWO approaches respectively. Thus, the proposed model guarantees good throughput at various tasks.

Count of Tasks	<b>PSO[10]</b>	WOA [45]	DA [16]	GWO [47]	MGWO [30]	GA [18]	DAGWO
100	24.517	23.993	25.594	26.481	26.027	36.182	23.692
150	30.597	31.897	25.924	28.701	22.909	39.735	21.567
175	39.534	50.870	37.291	49.644	48.394	66.901	43.298
200	43.198	51.077	48.291	61.009	57.095	86.832	63.047

<span id="page-18-1"></span>**Table 3** Risk probability evaluation: DAGWO &extant schemesfor scenario2

No of Tasks	PSO <sub>[10]</sub>	WOA [45]	DA [16]	GWO [47]	MGWO [30]	GA [18]	<b>DAGWO</b>
100	623.283	562.671	603.216	696.414	595.819	701.865	557.279
150	663.166	684.119	631.255	711.986	665.493	714.844	621.717
175	999.058	1080.647	1053.837	1082.442	1182.263	995.408	995.576
200	1073.911	1095.957	1021.057	965.206	1091.058	1076.208	956.885

<span id="page-19-0"></span>**Table 4** Throughput (pps) evaluation: DAGWO &extant schemes for scenario1

### **5.9 Time complexity**

The time consumed through the adopted as well as traditional schemes is depicted in Table [6](#page-20-0) and [7](#page-20-1) for scenario 1as well as scenario2. For scenario 1, the adopted DAGWOfor task 175 is 15.492%, 19.659%, 11.674%, 38.054%,41.12% and 9.186% superior to the extant PSO, WOA, DA, GWO, GA & MGWO approaches. In the case of Table [7](#page-20-1), the computation time obtained by the developed DAGWO while scheduling tasks 200 is 0.195 s and it is 36.11%, 33.03%, 29.26%, 44.43%, 0.34%, and 36.11% superior to the extant PSO, WOA, DA, GWO, GA & MGWO methods. Hence, the superiority of the adopted scheme is verifed with respect to computation time.

# **5.10 Rank analysis**

The rank analysis of the adopted as well as existing schemes is shown in Table [8.](#page-20-2) This table shows the rank of the 4-objectives like "energy consumption, makespan, resource utilization cost and execution time" and also the overall rank of the methods are represented. On analyzing the results, the proposed model does not attain the frst rank. But considering the overall analysis, the rank of the methods like PSO, WOA, DA, GWO, GA, MGWO, and proposed DAGWO are 7, 3, 2, 6, 5, 4 and 1. Therefore, the efectiveness of the proposed approach is proved successfully.

# **5.11 Discussions**

The suggestedscheme considers the objectives along with the security constraints. The resultants show the efficiency of the proposed work and it is well suited for a multi-cloud environment. Our method attains high makespan, good resource utilization, less risk probability, low energy consumption, and computation time. Furthermore, the proposed model achieves high security than other existing methods with low risk. Sometimes the proposed model lacks its

No of Tasks	<b>PSO[10]</b>	WOA [45]	DA [16]	GWO [47]	<b>MGWO [30]</b>	GA [18]	<b>DAGWO</b>
100	544.738	596.841	572.542	608.823	592.565	613.443	529.929
150	839.786	861.599	895.672	845.055	888.883	819.852	828.403
175	876.909	884.807	847.006	790.269	923.281	926.231	790.921
200	998.548	1058.586	974.607	1006.813	978.763	1079.772	979.557

<span id="page-19-1"></span>**Table 5** Throughput (pps) evaluation: DAGWO &extant schemesfor scenario2

<span id="page-20-0"></span>**Table 6** Computation time (s) evaluation: DAGWO &extant schemes for scenario1

<span id="page-20-1"></span>**Table 7** Computation time (s)evaluation: DAGWO &extant schemesfor scenario2

Number of task PSO [29] WOA [45] DA [16] GWO [47] MGWO [30] GA [18] DAGWO						
100		0.097172 0.100926 0.147046 0.165688		0.071388	0.097172 0.111682	
150		0.216606 0.251551	0.266266 0.273753	0.152369	0.216606 0.163678	
175		0.238304 0.324933	0.25915 0.318829	0.201988	0.238304 0.200971	
200	0.30624	0.292166	0.276618 0.351549	0.196334	0.306240 0.195666	

200 0.220585 0.238271 0.203186 0.231051 0.183171 0.207329 0.176934

<span id="page-20-2"></span>**Table 8** Rank analysis of the DAGWO and extant approaches

Method	Energy con- sumption	Makespan	Resource utiliza- tion cost	Execution time	Overall rank
<b>PSO</b> [29]	<sub>(</sub>		h		
WOA [45]		6			
DA [16]					
GWO [47]					h
GA [18]				h	
MGWO $[30]$	4				
<b>DAGWO</b>	∍	4			

performance when compared to MGWO and this model is restricted by dynamic scheduling. In the future, we enhanced the suggested work to achieve improved results. Also, we consider some task features like deadlines and the data transmission between the workfow tasks.

# <span id="page-21-4"></span>**6 Conclusion**

Data transfer across application tasks is unavoidable in a virtualized cloud environment. If a virtual machine's security mechanism isn't strong enough, malicious activity can disrupt other tasks via altering intermediate data. This research introducesa secure task scheduling scheme for a multi-cloud environment. The study that was given dealt with using the optimization concept to allocate the tasks in the best possible way. As a result, the objectives are taken into account by the suggested optimal task scheduling frameworklike"makespan, execution time, and utilization cost along with security constraints". A novel risk probability evaluation was taken into account for ensuring security. For optimal scheduling, a uniqueapproach referred to as DAGWO that hybridizes the concepts of GWO & DA was introduced. The DAGWO is determined over the existing works in terms of convergence, energy, makespan as well. The experimental fndings reveal that, when compared to existing methods, our method efectively reduces the security risk while keeping a respectable completion time. However, the execution time, makespan, utilization cost of the DAGWO model seems to oscillate in terms of performance, the overall objective is found to be achieved.

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### **Declarations**

**Confict of Interest** The authors declare that they have no confict of interest.

#### **Informed consent** not applicable

**Ethical approval** This article does not contain any studies with human or animal subjects performed by any of the authors.

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