

Exact Conic Programming Relaxations for a Class of Convex Polynomial Cone Programs

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Abstract In this paper, under a suitable regularity condition, we establish a broad class of conic convex polynomial optimization problems, called conic sum-of-squares convex polynomial programs, exhibiting exact conic programming relaxation, which can be solved by various numerical methods such as interior point methods. By considering a general convex cone program, we give unified results that apply to many classes of important cone programs, such as the second-order cone programs, semidefinite programs, and polyhedral cone programs. When the cones involved in the programs are polyhedral cones, we present a regularity-free exact semidefinite programming relaxation. We do this by establishing a sum-of-squares polynomial representation of positivity of a real sum-of-squares convex polynomial over a conic sum-of-squares convex system. In many cases, the sum-of-squares representation can be numerically checked via solving a conic programming problem. Consequently, we also show that a convex set, described by a conic sum-of-squares convex polynomial, is (lifted) conic linear representable in the sense that it can be expressed as (a projection of) the set of solutions to some conic linear systems.

Keywords Convex cone programs · Convex polynomial optimization · Conic programming · Semidefinite programming

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1 Introduction

The class of conic convex polynomial optimization problems is an important subclass of conic convex programming problems. It covers broad classes of convex programs such as inequality constrained convex polynomial programs [\[1\]](#page-21-0), the SOS-convex polynomial programs [\[1](#page-21-0)[–3\]](#page-21-1), convex semidefinite programs [\[4,](#page-21-2)[5\]](#page-21-3), second-order cone programs as well as *p*-order cone programs [\[6,](#page-21-4)[7\]](#page-21-5). It frequently appears in robust optimization [\[8\]](#page-21-6). For instance, a robust counterpart of a convex quadratic programming problem with a second-order cone constraint under (unstructured) norm-bounded uncertainty can equivalently be written as a conic convex quadratic optimization problem with the positive semidefinite cone (see Theorem 6.3.2 [\[8\]](#page-21-6)).

Moreover, many basic control problems are modeled as conic polynomial optimization problems. In particular, static output feedback design problems [\[9](#page-21-7)] are formulated as polynomial matrix inequality optimization problems. These conic polynomial problems are solved by approximating them with the Lasserre hierarchy of conic programming relaxation problems such as semidefinite linear programming (SDP) relaxation problems [\[2,](#page-21-8)[9\]](#page-21-7), which can efficiently be solved by interior point methods.

In particular, for an inequality constrained convex polynomial programs, the Lasserre hierarchy of SDP relaxations is known to converge finitely $[10-12]$ $[10-12]$ under suitable assumptions. An exact SDP relaxation has been shown to hold for the special class of SOS-convex programs [\[2](#page-21-8)[,3](#page-21-1)] and for some other classes of robust SOS-convex programs [\[1\]](#page-21-0). Due to the importance of exact SDP relaxations to conic convex programs, it has recently been shown that SOS-convex programs with semidefinite linear constraints [\[4\]](#page-21-2) exhibit exact SDP relaxations under appropriate regularity conditions.

The purpose of this piece of work is to show that a broad class of conic convex polynomial programming problems exhibits exact conic programming relaxations under suitable conditions. As an outcome, we give the following key unified results that apply to many classes of important convex cone programs, such as the second-order cone programs, *p*-order cone programs, semidefinite programs, and polyhedral cone programs.

- (i) By generalizing the notion of sum-of-squares convexity (SOS-convexity) of real polynomials [\[4,](#page-21-2)[13](#page-22-1)[–20\]](#page-22-2) to conic SOS-convexity of vector-valued polynomial mappings, we show that the class of conic SOS-convex polynomials includes various classes of conic convex polynomials such as the matrix SOS-convex polynomials [\[21](#page-22-3)] and conic convex quadratic maps. We establish that the positivity of an SOS-convex polynomial over a conic SOS-convex system can be characterized in terms of a sum-of-squares polynomial representation under a regularity condition. Related recent results for SOS-convex inequality systems with applications to robust optimization can be found in $[1,16]$ $[1,16]$ $[1,16]$. In many cases, the sum-of-squares representation can be numerically checked via solving a conic, such as semidefinite linear program. This representation also paves the way for deriving exact conic programming relaxation results.
- (ii) We prove that the class of SOS-convex cone programs enjoys exact conic programming relaxations under suitable regularity conditions. In the case, where

the cone involved in the program is a convex polyhedral cone, we show that our exact relaxation result collapses to an exact semidefinite linear programming relaxation without any regularity condition. We further show that the relaxation problem attains its maximum under a generalized Slater condition.

(iii) Consequently, we also show that a convex set, described by a conic SOS-convex polynomial, is (lifted) conic linear representable in the sense that it can be expressed as (a projection of) the set of solutions to some conic linear systems.

The outline of the paper is as follows. Section [2](#page-2-0) introduces the notion of conic SOS-convexity and provides various dual characterizations of solvability of conic SOS-convex systems, including a sum-of-squares representation result on positivity of polynomials. Section [3](#page-12-0) presents exact conic programming relaxation results. Section [4](#page-17-0) present results on conic linear representability of some convex sets.

2 Conic SOS-Convexity and Positivstellensatz

In this section we introduce the notion of conic SOS-convexity for polynomial mappings and present various first- and second-order characterizations of conic SOSconvexity. Consequently, we establish characterizations of positivity and nonnegativity of SOS-convex polynomials over conic SOS-convex polynomial systems.

We begin with some preliminaries on polynomials. We say that a real polynomial *f* is sum-of-squares [\[11](#page-21-10)[,22](#page-22-5)] if there exist real polynomials f_j , $j = 1, \ldots, r$, such that $f = \sum_{i=1}^r f_i^2$. The set consisting of all sum-of-squares real polynomials is denoted by $\Sigma^2[x]$. Moreover, the set consisting of all sum-of-squares real polynomials with degree at most *d* is denoted by $\Sigma_d^2[x]$. If each G_i , $i = 1, 2, ..., m$, is a polynomial on \mathbb{R}^n , then the degree of the polynomial map *G*, defined by $G(x) = (G_1(x), \ldots, G_m)$, is denoted by deg *G*. It is defined as the maximum degree of each G_i , $i = 1, \ldots, m$.

In this paper, \mathbb{R}^m denotes the Euclidean space with dimension *m*. For any $x, y \in \mathbb{R}^m$, the inner product between *x*, *y* is denoted by $\langle x, y \rangle$ and the norm induced by the inner product is defined by $||x|| = \sqrt{\langle x, x \rangle}$. We also use $\mathbb B$ (resp. $\mathbb B$) to denote the closed (resp. open) unit ball of \mathbb{R}^m . Throughout the paper, we assume that *K* is a closed and convex cone in \mathbb{R}^m . Its dual cone of *K*, denoted by K^{\oplus} , is given by $K^{\oplus} = \{y \in \mathbb{R}^m : \langle y, k \rangle > 0 \text{ for all } k \in K\}.$

The set $Sⁿ$ denotes the space of symmetric $(n \times n)$ matrices with the trace inner product and \succeq denotes the Löwner partial order of *Sⁿ*, that is, for *M*, $N \in S^n$, $M \succeq N$ if and only if $(M - N)$ is positive semidefinite. Let $S_+^n := \{M \in S_n : M \geq 0\}$ be the closed and convex cone of positive semidefinite $(n \times n)$ matrices. Note that for *M*, $N \in S^n_+$, the inner product, $(M, N) := \text{Tr}[MN]$, where $\text{Tr}[\cdot]$ refers to the trace operation. Note also that $M > 0$ means that *M* is positive definite.

Positive semidefinite cones, second-order cones, and polyhedral cones are some important closed and convex cones *K* that arise in numerous applications of optimization.

Recall that a real polynomial *f* on \mathbb{R}^n is called *SOS-convex* [\[13](#page-22-1)] if for any $\alpha \in$ [0, 1], *h*(*x*, *y*) := α*f* (*x*) + (1 − α) *f* (*y*) − *f* (α*x* + (1 − α)*y*) is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$. An equivalent notion of SOS-convexity was first introduced in [\[15\]](#page-22-6). Clearly, a SOS-convex polynomial is convex. However, the converse is not true, that is, there exists a convex polynomial which is not SOS-convex [\[13](#page-22-1)]. It is known that any convex quadratic function and any convex separable polynomial is a SOS-convex polynomial. Moreover, the class of SOS-convex polynomials is much larger than the convex quadratic functions and convex separable polynomials. For instance, $f(x) = x_1^8 + x_1^2 + x_1x_2 + x_2^2$ is a SOS-convex polynomial (see [\[15\]](#page-22-6)) which is nonquadratic and nonseparable.

We now introduce the new notion of conic SOS-convexity of a polynomial map, extending the notion of the scalar SOS-convex polynomial.

Definition 2.1 (*Conic SOS-convexity*) Let *K* be a closed and convex cone in \mathbb{R}^m . Let $G : \mathbb{R}^n \to \mathbb{R}^m$ be a polynomial vector-valued mapping, that is, $G(x) =$ $(G_1(x), \ldots, G_m(x))$ where each G_i is a polynomial on \mathbb{R}^n . We say $G : \mathbb{R}^n \to \mathbb{R}^m$ is a *K***-SOS-convex polynomial** if and only if, for any $\alpha \in [0, 1]$ and for any $\lambda \in K^{\oplus}$

 $h(x, y) = \langle \lambda, \alpha G(x) + (1 - \alpha)G(y) - G(\alpha x + (1 - \alpha)y) \rangle$

is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$.

The following equivalent characterization for scalar SOS-convex polynomial from [\[13](#page-22-1)] will be useful for us.

Lemma 2.1 [\[13\]](#page-22-1) Let f be a real polynomial on \mathbb{R}^n . Then, the following statements *are equivalent:*

(a) *f is a SOS-convex polynomial on* \mathbb{R}^n ;

(b) $(x, y) \mapsto \frac{1}{2} f(x) + \frac{1}{2} f(y) - f(\frac{x+y}{2})$ *is a sum-of-squares polynomial on* $\mathbb{R}^n \times \mathbb{R}^n$; (c) $(x, y) \mapsto f(x) - f(y) - \nabla f(y)(x - y)$ *is a sum-of-squares polynomial on* $\mathbb{R}^n \times \mathbb{R}^n$; (d) $(x, y) \mapsto y^T \nabla_{xx}^2 f(x) y$ is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$;

It is easy to see that, if $K = \mathbb{R}^m_+$, then *G* is \mathbb{R}^m_+ -SOS-convex if and only if *G_i* is SOSconvex for each $i = 1, 2, \ldots, m$. In particular, if $m = 1$ and $K = \mathbb{R}_+$, then our notion of conic SOS-convexity collapses to the definition of an SOS-convex polynomial [\[13](#page-22-1)].

We now present simple first- and second-order characterizations of conic SOSconvexity.

Lemma 2.2 (First- & second-order characterizations of conic SOS-convexity) *Let K be a closed and convex cone in* \mathbb{R}^m ; *let* G_i : $\mathbb{R}^n \to \mathbb{R}$, $i = 1, \ldots, m$, *be polynomials. Let* $G : \mathbb{R}^n \to \mathbb{R}^m$ *be a polynomial mapping, given by* $G(x) = (G_1(x), \ldots, G_m(x))$ *. Then, the following statements are equivalent:*

- (i) *G is a K -SOS-convex polynomial;*
- (ii) *for any* $\lambda \in K^{\oplus}$ *,*

$$
h(x, y) = \left\langle \lambda, \frac{1}{2}(G(x) + G(y)) - G\left(\frac{x+y}{2}\right) \right\rangle
$$

is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$.

(iii) *for any* $\lambda \in K^{\oplus}$, $(x, y) \mapsto \langle \lambda, G(x) - G(y) - \nabla G(y)(x - y) \rangle$ *is a sum-of-squares polynomial on* $\mathbb{R}^n \times \mathbb{R}^n$;

- (iv) *for any* $\lambda \in K^{\oplus}$, $(x, y) \mapsto y^T \nabla^2_{xx} (\langle \lambda, G(x) \rangle) y$ *is a sum-of-squares polynomial on* $\mathbb{R}^n \times \mathbb{R}^n$;
- (v) *for any* $\lambda \in K^{\oplus}$, $(x, y) \mapsto \langle \lambda, (\nabla G(x) \nabla G(y))(x y) \rangle$ *is a sum-of-squares polynomial on* $\mathbb{R}^n \times \mathbb{R}^n$.

Proof From the definition, we see that *G* is a *K*-SOS-convex polynomial if and only if for any $\alpha \in [0, 1]$ and for any $\lambda \in K^{\oplus}$,

$$
h(x, y) = \langle \lambda, \alpha G(x) + (1 - \alpha)G(y) - G(\alpha x + (1 - \alpha)y) \rangle
$$

is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$, which is further equivalent to $x \mapsto$ $\langle \lambda, G(x) \rangle$ is a scalar SOS-convex polynomial for any $\lambda \in K^{\oplus}$. Thus, the equivalences between (i), (ii) (iii) and (iv) follow immediately from Lemma [2.1.](#page-3-0)

 $[(iii) \Rightarrow (v)]$ Assume that (iii) holds. Fix an arbitrary $\lambda \in K^{\oplus}$. Then, $h_1(x, y) =$ $\langle \lambda, G(x) - G(y) - \nabla G(y)(x - y) \rangle$ is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$. Then,

$$
\langle \lambda, (\nabla G(x) - \nabla G(y))(x - y) \rangle
$$

= $\langle \lambda, G(x) - G(y) - \nabla G(y)(x - y) \rangle + \langle \lambda, G(y) - G(x) - \nabla G(x)(y - x) \rangle$
= $h_1(x, y) + h_1(y, x),$

is a sum-of-squares polynomial. Thus, (v) follows.

 $[(v) \Rightarrow (iv)]$ Assume that (v) holds. Fix an arbitrary $\lambda \in K^{\oplus}$. Let *d* be the degree of *G*. Then, $h_2(x, y) := \langle \lambda, (\nabla G(x) - \nabla G(y))(x - y) \rangle$ is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$ with degree at most *d*. It follows that, for any $\epsilon > 0$

$$
\langle \lambda, (\nabla G(x + \epsilon y) - \nabla G(x))(\epsilon y) \rangle = h_2(x + \epsilon y, x)
$$

is also a sum-of-squares polynomial with degree at most *d* on $\mathbb{R}^n \times \mathbb{R}^n$. Since

$$
y^T \nabla_{xx}^2 (\langle \lambda, G(x) \rangle) y = \lim_{\epsilon \to 0} \frac{\langle \lambda, (\nabla G(x + \epsilon y) - \nabla G(x))(\epsilon y) \rangle}{\epsilon^2}
$$

is a polynomial with degree at most *d*, it follows from the closedness of the cone consisting of all sum-of-squares polynomials with degree at most *d* (see [\[10](#page-21-9)]) that the function $(x, y) \mapsto y^T \nabla_{xx}^2(\langle \lambda, G(x) \rangle) y$ is also a sum-of-squares polynomial with degree at most *d*.

In the special case where K is the positive semidefinite cone in the space of symmetric matrices, by exploiting the structure of the positive semidefinite cone, we obtain a simplified characterizations in terms of vectors (instead of semidefinite matrices). This also allows us to compare the known notion of matrix SOS-convexity introduced in [\[15](#page-22-6)] with our notions in the following corollary. We note that the equivalence between (ii) and (iii) of the following corollary has been established in [\[23](#page-22-7)]. For the sake of convenience of a reader and for completeness, we provide a simple proof here.

Corollary 2.1 (Conic SOS-convexity and matrix SOS-convexity) *Let* $G : \mathbb{R}^n \to S^q$ *be a matrix polynomial, and let S^q* ⁺ *be the positive semidefinite matrix cone. Then, the following statements are equivalent*

- (i) *G* is a S^q_+ -SOS-convex polynomial;
- (ii) *for any* $z \in \mathbb{R}^q$, $(x, y) \mapsto y^T \nabla^2_{xx}(z^T G(x)z)y$ *is a sum-of-squares polynomial on* $\mathbb{R}^n \times \mathbb{R}^n$:
- (iii) *G* is matrix SOS-convex in the sense that, for any $z \in \mathbb{R}^m$, there exist $l \in \mathbb{N}$ and a *matrix polynomial* $F_z : \mathbb{R}^n \to \mathbb{R}^{l \times n}$ *such that, for all* $x \in \mathbb{R}^n$, $\nabla^2_{xx}(z^T G(x)z) =$ $F_z(x)^T F_z(x)$.

Proof As S^q and $\mathbb{R}^{q(q+1)/2}$ have the same dimensions, there exists an invertible linear map $L: S^q \to \mathbb{R}^{q(q+1)/2}$ such that

$$
L(M_1)^T L(M_2) = \text{Tr}(M_1 M_2) \quad \text{for all} \quad M_1, M_2 \in S^q.
$$

We now identify S^q , that is equipped with the trace inner product, as $\mathbb{R}^{q(q+1)/2}$ with the usual Euclidean inner product by associating each symmetric matrix *M* to *L*(*M*). Let $K = S^q_+$. Note that any positive semidefinite matrix can be written as sum of rankone matrices and $Tr((zz^T)G(x)) = z^T(G(x))z$ for any $x \in \mathbb{R}^n$ and $z \in \mathbb{R}^q$. Then, for any $\Lambda \in S^q_+$, $(x, y) \mapsto y^T \nabla^2_{xx} (\text{Tr}(\Lambda G(x))) y$ is a sum-of-squares polynomial is equivalent to the condition that, for any $z \in \mathbb{R}^q$, $(x, y) \mapsto y^T \nabla^2_{xx} (z^T G(x) z) y$ is a sum-of-squares polynomial. Then, the equivalence of (i) and (ii) follows from the preceding lemma with $K = S_+^q$.

To finish the proof, we only need to show the equivalence of (ii) and (iii). Suppose that (iii) holds. Then, for any $y \in \mathbb{R}^n$,

$$
y^T \nabla_{xx}^2 (z^T G(x) z) y = y^T (F_z(x))^T F_z(x) y = ||F_z(x) y||^2 = \sum_{j=1}^l (F_z(x) y)_j^2,
$$

which is a sum-of-squares polynomial. So, (ii) holds. Conversely, if (ii) holds, then for any $z \in \mathbb{R}^m$, there exist $l \in \mathbb{N}$ and polynomials $f_z^j : \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}, j = 1, ..., l$, such that

$$
y^T \nabla_{xx}^2 (z^T G(x) z) y = \sum_{j=1}^l f_z^j (x, y)^2.
$$

As, for each fixed $z \in \mathbb{R}^m$ and $x \in \mathbb{R}^n$, $y^T \nabla^2_{xx}(z^T G(x)z) y$ is a homogeneous quadratic function with respect to the variable *y*, we can assume that $f_z^j(x, y)$ is linear with respect to *y* for each $j = 1, \ldots, l$. Let

$$
(f_z^1(x, y), \ldots, f_z^l(x, y)) = F_z(x)y,
$$

where $F_z : \mathbb{R}^n \to \mathbb{R}^{l \times n}$ is a matrix polynomial. It follows that

$$
y^T \nabla_{xx}^2 (z^T G(x)z) y = y^T (F_z(x))^T F_z(x) y
$$
 for all $y \in \mathbb{R}^n$.

This shows that (iii) holds. \square

Example 2.1 (Classes of conic SOS-convex polynomials) We now provide some simple examples of conic SOS-convex polynomials.

- (a) Let $K = \mathbb{R}^m_+$. Then, $K^{\oplus} = K$, and so a polynomial vector-valued mapping $G = (G_1, \ldots, G_m)$ is a *K*-SOS-convex polynomial if and only if each G_i is a scalar SOS-convex polynomial.
- (b) Let $K = \{x \in \mathbb{R}^m : a_j^T x \geq 0, j = 1, \ldots, s\}$ be a convex polyhedral cone where $a_j = (a_{j1}, \ldots, a_{jm}) \in \mathbb{R}^m$, $j = 1, \ldots, s$. Then, we see that $K^{\oplus} := \{ \sum_{j=1}^{s} \mu_{j} a_{j} : \mu_{j} \geq 0 \}$ and a polynomial vector-valued mapping $G = (G_1, \ldots, G_m)$ is *K*-SOS-convex if and only if $h_j(x) := \sum_{i=1}^m a_{ji} G_i(x)$ is a scalar SOS-convex polynomial for all $j = 1, \ldots, s$.

In particular, if each G_i is a quadratic function of the form

$$
G_i(x) = \alpha_i + v_i^T x + x^T A_i x,
$$

where $\alpha_i \in \mathbb{R}$, $v_i \in \mathbb{R}^n$ and $A_i \in S^n$, then $G = (G_1, \ldots, G_m)$ is *K*-SOS-convex if and only if $\sum_{i=1}^{m} a_{ji} A_i \ge 0$ for all $j = 1, ..., s$.

- (c) Let $K = S_+^q$ be the cone consisting of all the $(q \times q)$ positive semidefinite matri-ces. Then, by corollary [2.1,](#page-4-0) $G : \mathbb{R}^n \to S^q$ is a matrix SOS-convex polynomial [\[21\]](#page-22-3) if and only if *G* is *K*-SOS-convex polynomial.
- (d) Let $K = SOC_m := \{x \in \mathbb{R}^m : x_1 \ge \sqrt{\sum_{i=2}^m |x_i|^2}\}$ be the second-order cone in \mathbb{R}^m . Then, the dual cone $K^{\oplus} = K$. Let $G_i(x)$, for $i = 1, 2, ..., m$, be defined by

$$
G_i(x) = \begin{cases} \alpha_1 + \sum_{j=1}^n a_{1j}x_j + \sum_{j=1}^n \beta_j x_j^d, & \text{if } i = 1, \\ \alpha_i + \sum_{j=1}^n a_{ij}x_j, & \text{if } i \ge 2, \end{cases}
$$

where α_i , $a_{ij} \in \mathbb{R}, \beta_j \geq 0$ and *d* is an even number. Then it can easily be checked that the map *G*, defined by $G(x) = (G_1(x), \ldots, G_m(x))$, is a *K*-SOS-convex polynomial because for any $\lambda = (\lambda_1, \ldots, \lambda_m) \in K^{\oplus} = K$, $\lambda_1 \geq 0$ and

$$
y^T \nabla_{xx}^2 (\langle \lambda, G(x) \rangle) y = \lambda_1 d(d-1) \left(\sum_{j=1}^n \beta_j x_j^{d-2} y_j^2 \right),
$$

which is a sum-of-squares polynomial.

For instance, the map *G*, defined by

$$
G(x) = (1 + 2x_1 - x_2 + x_1^8 + 3x_2^8, x_1 + x_2),
$$

satisfies the condition above and so is an $SOC₂-SOS$ -convex polynomial.

Remark 2.1 (Conic SOS-convexity vs conic convexity) Let G_i be real polynomials with degree *d* on \mathbb{R}^n and let $G = (G_1, \ldots, G_m)$.

- In the case where $K = \mathbb{R}^m_+$, $G = (G_1, \ldots, G_m)$ is a *K*-SOS-convex polynomial if and only if each *Gi* is a scalar SOS-convex polynomial. So, the recent characterization of scalar SOS-convexity $[13]$ $[13]$ implies that \mathbb{R}^m_+ -SOS-convexity is equivalent to \mathbb{R}^m_+ -convexity if and only if either one of the following holds (1) $n = 1$; (2) $d = 2$; (3) $(n, d) = (2, 4)$.
- Let $m = q(q + 1)/2$ and identify S^q with $\mathbb{R}^{q(q+1)/2}$. We now consider the case where $K = S^q_+$. Note that any bivariate quartic nonnegative polynomial is sumof-squares, and *G* is a S^q_+ -SOS-convex polynomial if and only if for any $z \in \mathbb{R}^q$, $(x, y) \mapsto y^T \nabla^2_{xx} (z^T G(x) z) y$ is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$. This shows that, if $(n, d) = (1, 2)$, *G* is S^q_+ -SOS-convex if and only if *G* is S^q_+ -convex.

In the case where the cone K has specific structure, we obtain new tractable class for conic SOS-convexity. To do this, let \circ be a bilinear operation from $\mathbb{R}^m \times \mathbb{R}^m \to \mathbb{R}^m$ satisfying $a \circ b = b \circ a$ for all $a, b \in \mathbb{R}^q$. Let K be a closed and convex cone which is self-dual in the sense that $K^{\oplus} = K$. We say *K* is generated by the bilinear operation ◦ if and only if

$$
K = \{a \circ a : a \in \mathbb{R}^m\}.
$$

We note that many important and widely used closed and convex cones can be generated in this way. For example,

- − The nonnegative orthant $K = \mathbb{R}^m_+ = \{a \circ a : a \in \mathbb{R}^m\}$, where the bilinear operation is defined by $a \circ b = (a_1b_1, \ldots, a_mb_m)^T$ for all $a, b \in \mathbb{R}^m$.
- $-$ In the case where *K* is the second-order cone, that is,

$$
K = \{(r, z) \in \mathbb{R} \times \mathbb{R}^{m-1} : r \geq ||(z_1, \ldots, z_{m-1})||_2\},\
$$

we see that $K = \{a \circ a : a \in \mathbb{R}^m\}$ where bilinear operation is defined by (α, u) $(\beta, v) = (\alpha \beta + u^T v, \alpha v + \beta u) \in \mathbb{R} \times \mathbb{R}^{m-1}$ for all (α, u) and $(\beta, v) \in \mathbb{R} \times \mathbb{R}^{m-1}$. – As S^q and $\mathbb{R}^{q(q+1)/2}$ have the same dimensions, there exists an invertible linear map $L: S^q \to \mathbb{R}^{q(q+1)/2}$ such that

$$
L(M_1)^T L(M_2) = \text{Tr}(M_1 M_2) \text{ for all } M_1, M_2 \in S^q.
$$

We now identify S^q , that is equipped with the trace inner product, as $\mathbb{R}^{q(q+1)/2}$ with the usual Euclidean inner product by associating each symmetric matrix *M* to $L(M)$. Then, the cone of positive semidefinite matrix satisfies $S_+^q = \{a \circ a : a \in L(M)\}$ $a \in \mathbb{R}^{q(q+1)/2}$ where bilinear operation is defined by $a \circ b = \frac{1}{2}(L^{-1}(a)L^{-1}(b) +$ *L*^{−1}(*b*)*L*^{−1}(*a*)) for all *a*, *b* ∈ $\mathbb{R}^{q(q+1)/2}$.

Definition 2.2 (*SDP tractable class of conic SOS-convex mappings: uniformly K - SOS-convex*) Let *K* be a closed and convex self-dual cone which is generated by the bilinear operation \circ , that is, $K = \{a \circ a : a \in \mathbb{R}^m\}$. We say *G* is uniformly *K*-SOS-convex whenever

$$
(x, y, a) \mapsto y^T \nabla_{xx}^2 ((a \circ a, G(x))) y
$$

is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^m$.

Now, it follows from the definition that, uniformly *K*-SOS-convex mappings must be *K*-SOS-convex for any closed and convex self-dual cone *K* generated by a bilinear operation ◦. Moreover, an important feature of the uniformly *K*-SOS-convexity is that checking whether a mapping *G* is uniformly *K*-SOS-convex or not is equivalent to solving a semidefinite programming problem.

Example 2.2 (Classes of uniform conic SOS-convex polynomials)

 $-$ Let $K = \mathbb{R}^m_+$. Then, any *K*-SOS-convex polynomial is uniformly *K*-SOS-convex. To see this, note that $K = \mathbb{R}^m_+ = \{a \circ a : a \in \mathbb{R}^m\}$ where bilinear operation is defined by $a \circ b = (a_1b_1, \ldots, a_mb_m)^T$ for all $a, b \in \mathbb{R}^m$, and any polynomial mapping $G = (G_1, \ldots, G_m)$ is \mathbb{R}^m_+ -SOS-convex if and only if each G_i , $i =$ 1,..., *m*, is a scalar SOS-convex polynomial. Then,

$$
y^T \nabla_{xx}^2 (\langle a \circ a, G(x) \rangle) y = \sum_{i=1}^m a_i^2 (y^T \nabla^2 G_i(x) y)
$$

is a sum-of-squares polynomial.

– Let $K = S_+^q$ be the cone consisting of all the $(q \times q)$ positive semidefinite matrices. By identifying S^q as $\mathbb{R}^{q(q+1)/2}$, we define $G : \mathbb{R}^n \to S^q$ as

$$
G(x) = A_0 + \sum_{i=1}^n x_i A_i + \sum_{l=1}^p f_l(x) B_l,
$$

where for each $l = 1, \ldots, p$, $f_l(x)$ is an SOS-convex scalar polynomial, $A_0, A_1, \ldots, A_n \in S^q$ and $B_l \succeq 0$. Recall that the positive semidefinite cone can be written as $S^q_+ = \{a \circ a : a \in \mathbb{R}^{q(q+1)/2}\}$ where bilinear operation is defined by $a \circ b = \frac{1}{2}(L^{-1}(a)L^{-1}(b) + L^{-1}(b)L^{-1}(a))$ for all $a, b \in \mathbb{R}^{q(q+1)/2}$. Then, *G* is a uniform *K*-SOS-convex polynomial. To see this, we note that, for any $a \in \mathbb{R}^{q(q+1)/2}$,

$$
y^T \nabla_{xx}^2 (\langle a \circ a, G(x) \rangle) y = \sum_{l=1}^p \text{Tr}([L^{-1}(a)]^2 B_l) y^T \nabla_{xx}^2 f_l(x) y.
$$

As each f_l is an SOS-convex scalar polynomial, $(x, y) \mapsto y^T \nabla^2_{xx} f_l(x) y$ is a sum-of-squares polynomial. Moreover, note that $B_l \geq 0$ and $[L^{-1}(a)]^2 \geq 0$. So, $a \mapsto \text{Tr}([L^{-1}(a)]^2 B_l)$ is a nonnegative quadratic function and so is also a sum-of-squares polynomial. This implies that

$$
(x, y, a) \mapsto y^T \nabla_{xx}^2 ((a \circ a, G(x))) y
$$

is also a sum-of-squares polynomial. Thus, *G* is a uniform *K*-SOS-convex polynomial.

As a simple example, consider the map *G*, defined by

$$
G(x) = \begin{pmatrix} 1 + x_2 + x_1^8 + (x_1 + x_2)^4 & -1 + x_1 + x_2 \\ -1 + x_1 + x_2 & x_2 + x_1^8 + (x_1 + x_2)^4 \end{pmatrix}
$$

=
$$
\begin{pmatrix} 1 & -1 \\ -1 & 0 \end{pmatrix} + x_1 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} + x_2 \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} + (x_1^8 + (x_1 + x_2)^4) \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}.
$$

Then, *G* satisfies the condition above and so is a uniform S^2_+ -SOS-convex polynomial.

2.1 Solvability of Conic SOS-Convex Systems

We now establish characterizations of positivity or nonnegativity of SOS-convex polynomials over sets described by conic SOS-convex polynomial systems. To do this, for $\lambda \in \mathbb{R}^m$ and $G : \mathbb{R}^n \to \mathbb{R}^m$, we define the function $\langle \lambda, G \rangle$ by

$$
\langle \lambda, G \rangle(x) := \langle \lambda, G(x) \rangle \text{ for all } x \in \mathbb{R}^n.
$$

Moreover, the following existence result for solutions of a convex polynomial optimization problem will also be useful for our later analysis.

Lemma 2.3 [\[24,](#page-22-8) Theorem 3] *Let* f_0, f_1, \ldots, f_m *be convex polynomials on* \mathbb{R}^n *. Suppose that* $\inf_{x \in C} f_0(x) > -\infty$ *where* $C = \{x \in \mathbb{R}^n : f_i(x) \le 0, i = 1, ..., m\}$ *. Then,* $\argmin_{x \in C} f_0(x) \neq \emptyset$.

Proposition 2.1 (Strict solvability) *Let K be a closed and convex cone in* R*^m with nonempty interior. Let* $G : \mathbb{R}^n \to \mathbb{R}^m$ *be a K-SOS-convex polynomial. Then, exactly one of the following statements holds:*

- (i) *there exists* $x_0 \in \mathbb{R}^n$ *such that* $G(x_0) \in -\text{int}K$;
- (ii) *there exists* $\lambda \in K^{\oplus} \setminus \{0\}$ *such that* $\langle \lambda, G \rangle \in \Sigma^2[x]$ *.*

Proof $[(ii) \Rightarrow Not (i)]$ Suppose (ii) holds. Then,

$$
\langle \lambda, G \rangle \ge 0. \tag{1}
$$

By construction and by the definition of K^{\oplus} , both systems (i) and [\(1\)](#page-9-0) cannot have solutions simultaneously.

[Not(i) \Rightarrow (ii)] Suppose that (i) fails. Then, $0 \notin \text{int}(\lbrace G(x) : x \in \mathbb{R}^n \rbrace + K)$, where $\{G(x) : x \in \mathbb{R}^n\} + K$ is a convex set. Now, it follows by a convex separation theorem that $\lambda \in K^{\oplus} \setminus \{0\}$ and $\langle \lambda, G(x) \rangle \geq 0$ for all $x \in \mathbb{R}^n$. As *G* is a *K*-SOS-convex polynomial, we see that $H(x) := \langle \lambda, G(x) \rangle$ is a convex polynomial. From Lemma [2.3,](#page-9-1) we obtain that argmin_{*x*∈R}_n {*H*(*x*)} $\neq \emptyset$. Now, let $x^* \in \operatorname{argmin}_{x \in \mathbb{R}^n}$ {*H*(*x*)}. Then, $H(x^*) \geq 0$ and $\nabla H(x^*) = 0$. This implies that for any $h \in \mathbb{R}^n$, $\langle \lambda, \nabla G(x^*)h \rangle = 0$. Moreover, from Lemma [2.2,](#page-3-1) we see that $h(x, y) := \langle \lambda, G(x) - G(y) - \nabla G(y)(x - y) \rangle$ is a sum-of-squares polynomial on $\mathbb{R}^n \times \mathbb{R}^n$. It follows that

$$
h(x, x^*) = \langle \lambda, G(x) - G(x^*) - \nabla G(x^*)(x - x^*) \rangle = \langle \lambda, G(x) \rangle - H(x^*)
$$

is an SOS-convex polynomial. It follows that $\langle \lambda, G(x) \rangle = h(x, x^*) + H(x^*)$ is also a sum-of-squares polynomial a sum-of-squares polynomial.

We now provide a characterization of positivity of an SOS-convex polynomial over a set described by conic SOS-convex systems, under the condition that the set

$$
\Omega := (f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K)
$$

= {(s, y) \in \mathbb{R} \times \mathbb{R}^m : \exists x \in \mathbb{R}^n, s \ge f(x), y \in G(x) + K},

is closed. It is easy to check that Ω is a closed set if; for instance, f is coercive in the sense that $\lim_{\|x\| \to \infty} f(x) = +\infty$. The set of the form Ω has played an important role in the study of duality in convex optimization [\[25\]](#page-22-9).

Theorem 2.1 (Positivstellensatz for conic SOS-convex systems) *Let K be a closed* and convex cone in \mathbb{R}^m . Let f be an SOS-convex real polynomial and let G be a K-SOS*convex polynomial. Let* $\{x : G(x) \in -K\} \neq \emptyset$ *. If the convex set* $(f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K)$ *is closed, then the following statements are equivalent*

(i) $G(x) \in -K \implies f(x) > 0;$ (ii) $(\exists \lambda \in K^{\oplus}, \delta_0 > 0)$ $(\exists \sigma_0 \in \Sigma^2[x])$ $(\forall x \in \mathbb{R}^n)$ $f(x) + \langle \lambda, G(x) \rangle = \delta_0 + \sigma_0(x)$ *.*

Proof $[(ii) \Rightarrow (i)]$ Suppose that there exist a sum-of-squares polynomial $\sigma_0, \lambda \in K^{\oplus}$ and $\delta_0 > 0$ such that $f = \delta_0 + \sigma_0 - \langle \lambda, G \rangle$. Then, for each *x* with $G(x) \in -K$, $f(x) = \delta_0 + \sigma_0(x) - \langle \lambda, G(x) \rangle \ge \delta_0 > 0.$

[(i) \Rightarrow (ii)] Suppose that *f* is positive on *F* := {*x* : *G*(*x*) ∈ −*K*}. Then, we see that

$$
0_{m+1} \notin \Omega := (f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K).
$$

Thus, strict separation theorem (cf. [\[26\]](#page-22-10)) shows that there exist $\alpha \in \mathbb{R}, \epsilon > 0$ and $(\mu_0, \mu_1, \ldots, \mu_m) \neq (0, 0, \ldots, 0)$ such that, for all $y = (y_0, y_1, \ldots, y_m) \in \Omega$,

$$
0\leq \alpha < \alpha+\epsilon \leq \langle (y_0,y_1,\ldots,y_m),(\mu_0,\mu_1,\ldots,\mu_m)\rangle.
$$

As $\Omega + (\mathbb{R}_+ \times K) \subseteq \Omega$, we see that $(r, \mu) \in (\mathbb{R}_+ \times K^{\oplus}) \setminus \{0\}$. Since, for each $x \in \mathbb{R}^n$, $(f(x), G(x)) \in \Omega$, it follows that

$$
rf(x) + \langle \mu, G(x) \rangle \ge \alpha + \epsilon \ge \epsilon
$$
, for all $x \in \mathbb{R}^n$.

As $F \neq \emptyset$, we see that $r \neq 0$, and so $r > 0$. So, we obtain that

$$
f(x) + \langle \lambda, G(x) \rangle \ge \delta_0
$$
, for all $x \in \mathbb{R}^n$,

where $\lambda = \frac{\mu}{r} \in K^{\oplus}$ and $\delta_0 = \frac{\epsilon}{r} > 0$. This implies that

$$
r := \inf_{x \in \mathbb{R}^n} \{ f(x) + \langle \lambda, G(x) \rangle \} \ge \delta_0.
$$

Now, by Lemma [2.3,](#page-9-1) we can find x^* such that $f(x^*) + \langle \lambda, G(x^*) \rangle = r$. Let

$$
h(x) := f(x) + \langle \lambda, G(x) \rangle - r, \quad \text{for each} \quad x \in \mathbb{R}^n.
$$

Then *h* is a SOS-convex polynomial with $h(x^*) = 0$ and $\nabla h(x^*) = 0$. The SOSconvexity of *f* and *K*-SOS-convexity of *G* together with Lemma [2.2](#page-3-1) give us that $f(x) - f(x^*) - \nabla f(x^*)(x - x^*) = \theta_0(x)$ for some sum-of-squares polynomial θ₀ and $(λ, G(x) – G(x[*]) – ∇G(x[*])(x – x[*])) = θ_i(x)$ for some sum-of-squares polynomial θ_i . So, we see that, for all $x \in \mathbb{R}^n$

$$
f(x) + \langle \lambda, G(x) \rangle - (f(x^*) + \langle \lambda, G(x^*) \rangle) = \theta_0(x) + \sum_{i=1}^m \lambda_i \theta_i(x).
$$

As $f(x^*) + \langle \lambda, G(x^*) \rangle = h(x^*) + r = r$, we see that

$$
h(x) = f(x) + \langle \lambda, G(x) \rangle - r = \theta_0(x) + \sum_{i=1}^{m} \lambda_i \theta_i(x)
$$

which is a sum-of-squares polynomial. Hence,

$$
f(x) + \langle \lambda, G(x) \rangle = h(x) + r = (h(x) + r - \delta_0) + \delta_0 = \sigma_0(x) + \delta_0,
$$

where $r - \delta_0 > 0$ and $\sigma_0(x) := h(x) + (r - \delta_0)$ is a sum-of-squares polynomial. Thus, the conclusion follows the conclusion follows.

In passing, we note that the sum-of-squares representation (*ii*) of Theorem [2.1](#page-10-0) can, in many cases, be numerically checked by solving a related conic, such as a semidefinite linear program. Now, as a consequence, we obtain a characterization for nonnegativity of an SOS-convex polynomial over a set described by a conic SOS-convex system.

Corollary 2.2 (Nonnegativity characterization) *Let f be an SOS-convex scalar polynomial. Let K be a closed and convex cone, and let G* : $\mathbb{R}^n \to \mathbb{R}^m$ *be a K-SOS-convex polynomial with* $\{x : G(x) \in -K\} \neq \emptyset$. If the convex set $(f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K)$ is *closed, then the following statements are equivalent:*

(i) $G(x) \in -K \implies f(x) > 0;$ (ii) $(\forall \epsilon > 0)$ $(\exists \lambda \in K^{\oplus})$ $(\exists \sigma_0 \in \Sigma^2[x])$ $(\forall x \in \mathbb{R}^n)$ $f(x) + \langle \lambda, G(x) \rangle + \epsilon = \sigma_0(x)$ *.*

Proof $[(ii) \Rightarrow (i)]$. Suppose that for each $\epsilon > 0$ there exist a sum-of-squares polynomial σ_0 , $\lambda \in K^\oplus$ such that $f + \epsilon = \sigma_0 - \langle \lambda, G \rangle$. Let *x* be such that $G(x) \in -K$. Then, $f(x) + \epsilon = \sigma_0(x) - \langle \lambda, G(x) \rangle \ge 0$. Letting $\epsilon \to 0$, we see that $f(x) \ge 0$ for all *x* with $G(x) \in -K$.

 $[(i) \Rightarrow (ii)]$ Assume that (i) holds. Then, for any $\epsilon > 0$, $f + \epsilon$ is positive on *F*. So, Theorem [2.1](#page-10-0) implies that there exist a sum-of-square polynomial σ_0 , $\lambda \in K^{\oplus}$ and $\delta_0 > 0$ such that, for all $x \in \mathbb{R}^n$, $f(x) + \langle \lambda, G(x) \rangle + \epsilon = \delta_0 + \sigma_0$. Hence, the conclusion follows conclusion follows.

3 Exact Conic Programming Relaxations

In this section, we present exact conic relaxations for a class of conic convex nonlinear programming problems. We consider the problem

$$
\inf_{x \in \mathbb{R}^n} \{ f(x) : G(x) \in -K \},\tag{P}
$$

where *f* is a SOS-convex polynomial with degree d_0 on \mathbb{R}^n and $G : \mathbb{R}^n \to \mathbb{R}^m$ is a *K*-SOS-convex polynomial with degree d_1 . Let d be the smallest even number such that $d \ge \max\{d_0, d_1\}$. We assume throughout that the feasible set $F := \{x : G(x) \in$ $-K$ } $\neq \emptyset$.

It is easy to see that the program (P) is equivalent to the program

$$
\sup_{\mu \in \mathbb{R}} \{ \mu : f(x) - \mu \ge 0, \ \forall x \in F \}.
$$

The conic programming relaxation of [\(P\)](#page-12-1) based on certificate of nonnegativity on *F* (see Corollary [2.2\)](#page-11-0) is given by

$$
\sup_{\mu \in \mathbb{R}, \lambda \in \mathbb{R}^m} \{ \mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x], \ \lambda \in K^{\oplus} \},
$$
 (RP)

where $\sum_{d}^{2}[x]$ denotes the sum-of-squares polynomial with degree no larger than *d*. Note that the sum-of-squares condition, $f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x]$, can equivalently be reformulated as a linear matrix inequality [\[10](#page-21-9)]. In particular, if $K = S_+^m$ is the positive semidefinite matrix cone (or more generally, if *K* is a linear matrix inequality representable closed and convex cone), then [\(RP\)](#page-12-2) can equivalently be written as a semidefinite linear programming problem. Moreover, if *K* is the *p*-order cone, then [\(RP\)](#page-12-2) can also be equivalently reformulated as a tractable conic involving semidefinite constraints and norm constraints.

Theorem 3.1 (Exact conic programming relaxation) *For problem* [\(P\)](#page-12-1)*, let f be a SOSconvex polynomial with degree d*⁰ *and let G be a K -SOS-convex polynomial with degree* d_1 *. Let* $F := \{x : G(x) \in -K\} \neq \emptyset$ *. If the set* $(f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K)$ *is closed, then*

$$
\inf(P) = \sup_{\mu \in \mathbb{R}, \lambda \in \mathbb{R}^m} \{ \mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x], \ \lambda \in K^{\oplus} \},
$$

where d is the smallest even number such that $d \geq \max\{d_0, d_1\}$ *.*

Proof Clearly, for any $x \in F$, $\mu \in \mathbb{R}$ and $\lambda \in K^{\oplus}$, with $f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x]$,

$$
f(x) - \mu \ge f(x) + \langle \lambda, G(x) \rangle - \mu \ge 0.
$$

So, $\inf_{x \in F} f(x) \ge \sup_{\mu \in \mathbb{R}, \lambda \in K} \theta \{ \mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x] \}.$

To see the reverse inequality, we assume without loss of generality that inf $\chi \in F$ $f(x)$ $-\infty$; otherwise, the conclusion trivially holds. As $F \neq \emptyset$, we have $r := \inf_{x \in F} f(x) \in$ R. Then, for any $\epsilon > 0$, $f - r + \epsilon$ is positive on *F*. So, our Theorem [3.1](#page-12-3) implies that there exists a sum-of-squares polynomial σ_0 , $\lambda \in K^{\oplus}$ and $\delta_0 > 0$ such that

$$
f + \langle \lambda, G \rangle - (r - \epsilon) = \delta_0 + \sigma_0
$$

is a sum-of-squares polynomial. Note that any sum-of-squares polynomial must have even degree, *f* is of degree d_0 and *G* is of degree d_1 . It follows that $f + \langle \lambda, G \rangle - (r - \epsilon)$ is of degree at most *d*. This shows that, for any $\epsilon > 0$,

$$
\sup_{\mu \in \mathbb{R}, \lambda \in K^{\oplus}} \{\mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x] \} \ge r - \epsilon.
$$

Letting $\epsilon \to 0$, we see that $\sup_{\mu \in \mathbb{R}, \lambda \in K^{\oplus}} {\{\mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x] \}} \ge r$ and so the conclusion follows.

As an easy consequence, we obtain the following exact SDP relaxation for SOSconvex polynomial optimization problem over a SOS-convex polynomial matrix inequality problem under the closedness condition.

Corollary 3.1 *For problem* [\(P\)](#page-12-1), let f be a SOS-convex polynomial with degree d_0 and *let* $G : \mathbb{R}^n \to S^q$ *be a matrix SOS-convex polynomial with degree* d_1 *. Let* $F := \{x : f(x) = 0\}$ *G*(*x*) ∈ $-S_+^q$ } $\neq \emptyset$ *. If* (*f*, *G*)(\mathbb{R}^n) + ($\mathbb{R}_+ \times S_+^{\tilde{q}}$) *is closed, then*

$$
\inf(P) = \sup_{\mu \in \mathbb{R}, \lambda \in S^q} \{ \mu : f + \text{Tr}(\lambda G) - \mu \in \Sigma_d^2[x], \lambda \in S_+^q \},
$$

where, for $\lambda \in S^q$, $\text{Tr}(\lambda G)(x) := \text{Tr}(\lambda G(x))$ *for all* $x \in \mathbb{R}^n$ *and d is the smallest even number such that* $d \geq \max\{d_0, d_1\}$ *.*

Proof By identifying S^q as $\mathbb{R}^{q(q+1)/2}$, Corollary [2.1](#page-4-0) implies that *G* is indeed an S_1^q -SOS-convex polynomial. Thus the conclusion follows from Theorem [3.1](#page-12-3) with $K = S^q_+$. \mathbb{L}

Now, we consider a conic SOS-convex polynomial optimization with quadratic constraints involving the *p*-order cone:

$$
\min_{x \in \mathbb{R}^n} \{ f(x) : G(x) \in -POC_m \},\tag{PP}
$$

where *f* is an SOS-convex polynomial and $G = (G_1, \ldots, G_m)$ where each G_i is a quadratic function given by

$$
G_i(x) = \alpha_i + a_i^T x + x^T A_i x,\tag{2}
$$

with $\alpha_i \in \mathbb{R}, a_i \in \mathbb{R}^n$ and $A_i \in S^n, i = 1, \ldots, m$ and POC_{*m*} is the *p*-order cone given by $\text{POC}_m = \{(x_1, ..., x_m) \in \mathbb{R}^m : x_1 \ge \sqrt[p]{\sum_{i=2}^m |x_i|^p} \}$. We now show that problem [\(PP\)](#page-13-0) admits an exact conic relaxation under easily verifiable conditions.

Recall that, for a symmetric matrix *A*, the minimum eigenvalue of *A* is denoted by $\lambda_{\min}(A)$ and $||A||_2 = \sup{||Ay||_2 : ||y||_2 = 1} = \sqrt{\lambda_{\max}(A^T A)}$ is known as the matrix 2-norm of *A*.

Corollary 3.2 (Exact relaxation with quadratic constraints and *p*-order cone) *For problem* [\(PP\)](#page-13-0), let f be a SOS-convex polynomial with degree d_0 and let G = (G_1, \ldots, G_m) *be a quadratic vector-valued mapping where each* G_i *is given as in* [\(2\)](#page-13-1)*. Let* $F := \{x : G(x) \in \text{--POC}_m\} \neq \emptyset$ *. If* $\nabla^2_{xx} f(x_0) > 0$ for some $x_0 \in \mathbb{R}^n$ and

$$
\lambda_{\min}(A_1) \ge \sqrt[p]{\sum_{i=2}^{m} \|A_i\|_2^p},
$$
\n(3)

then

$$
\inf(PP) = \sup_{\mu \in \mathbb{R}, (\lambda_1, ..., \lambda_m) \in \mathbb{R}^m} \left\{ \mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x], \lambda_1 \geq \sqrt[p*]{\sum_{i=2}^m |\lambda_i|^{p^*}} \right\},
$$

where p^* *is a real number satisfies* $\frac{1}{p} + \frac{1}{p^*} = 1$ *and d is the smallest even number such that* $d \geq \max\{d_0, 2\}$ *.*

Proof Let $K = \text{POC}_m$. Then, the dual cone K^{\oplus} is the *p*[∗]-order cone where *p*[∗] satisfies $\frac{1}{p} + \frac{1}{p^*} = 1$. Let $G(x) = (G_1(x), \ldots, G_m(x))$ be a quadratic vector function where each G_i is defined as in [\(2\)](#page-13-1).

We first verify that G is a POC_m -SOS-convex polynomial. To see this, fix any $\lambda = (\lambda_1, \ldots, \lambda_m) \in K^{\oplus}$. Then, $\lambda_1 \geq \sqrt[p]{\sum_{i=2}^m |\lambda_i|^{p^*}}$ and

$$
y^T \nabla_{xx}^2 (\langle \lambda, G(x) \rangle) y = \left(\sum_{i=1}^m \lambda_i y^T A_i y \right)
$$

which is a quadratic function of *y* only. As any nonnegative quadratic function is a sum-of-squares polynomial, to show $(x, y) \mapsto y^T \nabla^2_{xx} (\langle \lambda, G(x) \rangle) y$ is a sum-ofsquares polynomial, we only need to prove that $\sum_{i=1}^{m} \lambda_i y^T A_i y \ge 0$ for all $y \in \mathbb{R}^n$. To verify this, note from [\(3\)](#page-14-0) that, for all $y \in \mathbb{R}^n$,

$$
y^T A_1 y \geq \lambda_{\min}(A_1) \|y\|^2 \geq \sqrt{\left(\sum_{i=2}^m \|A_i\|_2 \|y\|_2^2\right)^p} \geq \sqrt[p]{\sum_{i=2}^m |y^T A_i y|^p}.
$$

It now follows that

$$
\sum_{i=1}^{m} \lambda_i y^T A_i y = \lambda_1 y^T A_1 y + \sum_{i=2}^{m} \lambda_i y^T A_i y
$$

\n
$$
\geq \lambda_1 y^T A_1 y - \sum_{i=2}^{p^*} \left| \sum_{i=2}^{m} |\lambda_i|^{p^*} \right| \sum_{i=2}^{m} |y^T A_i y|^p
$$

\n
$$
\geq \left(\lambda_1 - \sum_{i=2}^{p^*} |\lambda_i|^{p^*} \right) y^T A_1 y \geq 0,
$$

where the first inequality follows by the Hölder inequality and the last inequality follows from $A_1 \ge 0$ [thanks to [\(3\)](#page-14-0)] and $\lambda_1 \ge \sqrt[n]{\sum_{i=2}^m |\lambda_i|^{p^*}}$. Thus, *G* is a POC_{*m*}-SOS-convex polynomial.

Finally, as *f* is a convex polynomial with $\nabla_{xx}^2 f(x_0) > 0$ for some $x_0 \in \mathbb{R}^n$, [\[12,](#page-22-0) Lemma 3.1] implies that *f* is coercive. This shows that $(f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K)$ is closed. Then, the conclusion follows from Theorem 3.1 with $K = \text{POC}_{\text{max}}$. closed. Then, the conclusion follows from Theorem [3.1](#page-12-3) with $K = \text{POC}_m$.

Corollary 3.3 (Qualification-free exact SDP relaxation and polyhedral cones *K*) *For the problem* [\(P\)](#page-12-1), let $K = \{x \in \mathbb{R}^m : a_j^T x \ge 0, j = 1, ..., s\}$ be the polyhedral *cone where* $a_j \in \mathbb{R}^m$. Let f be a SOS-convex polynomial with degree d_0 and let $G = (G_1, \ldots, G_m)$ *be a K-SOS-convex polynomials with degree d₁. Let* $F = \{x :$ $G(x) \in -K$ $\neq \emptyset$ *. Then,*

$$
\inf_{x \in F} f(x) = \sup_{\mu \in \mathbb{R}, \lambda \in \mathbb{R}^m, \mu_j \in \mathbb{R}} \left\{ \mu : f + \sum_{i=1}^m \sum_{j=1}^s \mu_j a_{ji} G_i - \mu \in \Sigma_d^2[x], \ \mu_j \ge 0 \right\},\
$$

where $a_j = (a_{j1}, \ldots, a_{jm})^T \in \mathbb{R}^m$ *and d is the smallest even number larger than* $max{d_0, d_1}$.

Proof Let $G(x) = (G_1(x), \ldots, G_m(x))$ and denote $a_j = (a_{j1}, \ldots, a_{jm}) \in \mathbb{R}^m$, $j = 1, \ldots, s$. Then, the set $\Omega := (f, G)(\mathbb{R}^n) + (\mathbb{R}_+ \times K)$ becomes

$$
\Omega = \left\{ (y_0, y_1, \dots, y_m) : \exists x \in \mathbb{R}^n \text{ s.t. } f(x) \le y_0, \sum_{i=1}^m a_{ji} (G_i(x) - y_i) \le 0, j = 1, \dots, s \right\}.
$$

We now show that the set Ω is closed. To see this, take any $(y_0^k, y_1^k, \ldots, y_m^k) \in \Omega$ be such that $(y_0^k, y_1^k, \ldots, y_m^k) \to (y_0, y_1, \ldots, y_m)$. By the definition of Ω , for each *k*, there exists $x^k \in \mathbb{R}^n$ such that $f(x^k) \leq y_0^k$ and $\sum_{i=1}^m a_{ji}(G_i(x^k) - y_i) \leq 0$, $j = 1, \ldots, s$. Now, consider the optimization problem

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$$
\min_{x,z_1,\dots,z_m} \left\{ \sum_{i=0}^m (z_i - y_i)^2 : f(x) \le z_0, \sum_{i=1}^m a_{ji} (G_i(x) - z_i) \le 0, j = 1,\dots,s \right\}.
$$
\n
$$
(P_0)
$$

For each $j = 1, \ldots, s$, let $h_j(x, z) := \sum_{i=1}^m a_{ji}(G_i(x) - z_i)$. Note that, for each $j = 1, \ldots, s, a_j \in K^{\oplus}$. As *G* is a *K*-SOS-convex polynomial, it follows that $\sum_{i=1}^{m} a_{ij} G_i(x) = (a_i, G(x))$ is an SOS convex polynomial. This implies that each $\sum_{i=1}^{m} a_{ji} G_i(x) = \langle a_j, G(x) \rangle$ is an SOS-convex polynomial. This implies that each h_i is an SOS-convex polynomial, and so (P_0) is a convex polynomial optimization problem and

$$
0 \le \inf(P_0) \le \sum_{i=0}^m (y_i^k - y_i)^2 \to 0.
$$

So, $\inf(P_0) = 0$. Moreover, Lemma [2.3](#page-9-1) implies that $\inf(P_0)$ is attained, and so there exists x^* ∈ \mathbb{R}^n such that and $f_0(x^*)$ ≤ *y*₀ and $\sum_{i=1}^m a_{ji}(G_i(x^*) - y_i) \le 0$, *j* = $1, \ldots, s$. So, $(y_0, y_1, \ldots, y_m) \in \Omega$. Therefore, Ω is closed.

Then, Theorem [3.1](#page-12-3) gives us that

$$
\inf(P) = \sup_{\mu \in \mathbb{R}, \lambda \in \mathbb{R}^m} \{ \mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x], \ \lambda \in K^{\oplus} \}.
$$

As *K* is a polyhedral cone with the form $\{x : a_j^T x \ge 0\}$, one has

$$
K^{\oplus} = \left\{ \sum_{j=1}^{s} \mu_j a_j : \mu_j \ge 0 \right\}.
$$

So, $\lambda \in K^{\oplus}$ if and only if $\lambda = \sum_{j=1}^{s} \mu_j a_j$ for some $\mu_j \ge 0$, and for all $x \in \mathbb{R}^n$,

$$
\langle \lambda, G(x) \rangle = \sum_{j=1}^{s} \mu_j \sum_{i=1}^{m} a_{ji} G_i(x) = \sum_{i=1}^{m} \sum_{j=1}^{s} \mu_j a_{ji} G_i(x).
$$

Therefore, the conclusion follows.

Corollary 3.4 Let f be a SOS-convex polynomial and let g_i , $i = 1, \ldots, m$, be SOS*convex polynomials with* $F = \{x : g_i(x) \leq 0, i = 1, \ldots, m\} \neq \emptyset$ *. Then,*

$$
\inf_{x\in F} f(x) = \sup_{\mu\in\mathbb{R}, \lambda_i\in\mathbb{R}} \left\{ \mu : f + \sum_{i=1}^m \lambda_i g_i - \mu \in \Sigma_d^2[x], \lambda_i \geq 0, i = 1, 2, \dots, m \right\},\
$$

where d is the smallest even number larger than $\max\{\deg f, \max_{1 \leq i \leq m} \deg g_i\}$ *.*

Proof Let $a_j = e_j$, $j = 1, \ldots, m$, where e_j is the vector whose *j*th coordinate is one, and zero otherwise. Then, $\{x : a_j^T x \ge 0, j = 1, ..., m\} = \mathbb{R}^m_+$ and $G = (g_1, ..., g_m)$

$$
\Box
$$

is \mathbb{R}^m_+ -SOS-convex. Thus, the conclusion follows immediately from the preceding corollary. \Box

We finish this section by establishing an exact conic relaxation with the solution attainment under a strict feasibility condition. In the special case when *G* is an affine function and $K = S_{+}^{m}$, this exact conic relaxation result collapses to the exact SDP relaxation result for nonlinear semidefinite programming problems with SOS-convex polynomial objective functions, established in [\[4\]](#page-21-2).

Theorem 3.2 (Relaxation with attainment) *Let K be a closed and convex cone with nonempty interior. For the problem* [\(P\)](#page-12-1)*, f be a SOS-convex polynomial with degree d*⁰ *and G be a K -SOS-convex polynomial with degree d*1*. Suppose that there exists* $x_0 \in \mathbb{R}^n$ *such that* $G(x_0) \in -\text{int}K$. Then

$$
\inf(P) = \max_{\mu \in \mathbb{R}, \lambda \in \mathbb{R}^m} \{ \mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x], \ \lambda \in K^{\oplus} \},
$$

where d is the smallest even number such that $d \geq \max\{d_0, d_1\}$ *.*

Proof Without loss of generality, let $r := \inf_{x \in F} f(x) \in \mathbb{R}$ be finite. Then, the system

$$
x \in \mathbb{R}^n, G(x) \in -K, f(x) - r < 0
$$

has no solution. So the following system

$$
x \in \mathbb{R}^n, G(x) \in -\mathrm{int}K, f(x) - r < 0
$$

has no solution. Now, by Proposition [2.1,](#page-9-2) there exists $0_{m+1} \neq (\tau, \lambda) \in \mathbb{R}_+ \times K^{\oplus}$ such that

$$
\tau(f - r) + \langle \lambda, G \rangle \in \Sigma_d^2[x].
$$

If $\tau = 0$, then $\lambda \neq 0$ and $\langle \lambda, G \rangle \in \Sigma_d^2[x]$. This gives us a contradiction, by Proposition [2.1,](#page-9-2) to the hypothesis that $G(x_0) \in -\text{int}K$ for some $x_0 \in \mathbb{R}^n$. Hence, $\tau > 0$ and so we get that $f + \langle \overline{\lambda}, G \rangle - r \in \Sigma_d^2[x]$, where $\overline{\lambda} = \frac{\lambda}{\tau}$. Thus, $\max_{\mu \in \mathbb{R}, \lambda \in K^{\oplus}} {\{\mu : f + \langle \lambda, G \rangle - \mu \in \Sigma_d^2[x] \}} \ge r$. The conclusion now holds as the opposite inequality always holds by construction. \Box

4 Conic Representations of Some Convex Sets

In this section, we show that the sets described by conic SOS-convex polynomials admit lifted conic representations. Let $G : \mathbb{R}^n \to \mathbb{R}^m$ be a *K*-SOS-convex polynomial with degree *d*. Its epigraph is given by

$$
epi G = \{(x, y) \in \mathbb{R}^n \times \mathbb{R}^m : y \in G(x) + K\}.
$$

It is easy to verify that epi *G* is a convex set.

Next, we show that the epigraph is lifted conic representable in the sense that it can be rewritten as a projection of a convex set described by a conic linear system. If the degree $d = 1$, then epi *G* can be trivially described by a conic linear system. Thus, without loss of generality, we only consider the case where $d > 2$. To do this, let $\mathbb{R}_d[x_1,\ldots,x_n]$ be the space consisting of all real polynomials on \mathbb{R}^n with degree *d* and let $C(d, n)$ be the dimension of $\mathbb{R}_d[x_1, \ldots, x_n]$. Write the canonical basis of $\mathbb{R}_d[x_1,\ldots,x_n]$ by

$$
x^{(d)} := \left(1, x_1, x_2, \dots, x_n, x_1^2, x_1 x_2, \dots, x_2^2, \dots, x_n^2, \dots, x_1^d, \dots, x_n^d\right)^T
$$

and let $x_{\alpha}^{(d)}$ be the α -th coordinate of $x^{(d)}$, $1 \le \alpha \le C(d, n)$. Then, we can write $f(x) = \sum_{\alpha=1}^{C(d,n)} f_{\alpha} x_{\alpha}^{(d)}$. Let $G : \mathbb{R}^{n} \to \mathbb{R}^{m}$ be a polynomial mapping with degree *d*. Then $G = (G_1, \ldots, G_m)$ and

$$
G(x) = \sum_{1 \leq \alpha \leq C(n,d)} G_{\alpha} x_{\alpha}^{(d)} = \left(\sum_{1 \leq \alpha \leq C(n,d)} (G_1)_{\alpha} x_{\alpha}^{(d)}, \dots, \sum_{1 \leq \alpha \leq C(n,d)} (G_m)_{\alpha} x_{\alpha}^{(d)} \right),
$$

where

$$
G_{\alpha} = ((G_1)_{\alpha}, \dots, (G_m)_{\alpha}).
$$
\n(4)

It is easy to see that $[x^{(d/2)}][x^{(d/2)}]^T$ is a matrix whose entries are real polynomials on \mathbb{R}^n with degree *d*. We then define $A_{\alpha}^{(d)}$, $1 \leq \alpha \leq C(n, d)$, to be symmetric matrices such that

$$
[x^{(d/2)}][x^{(d/2)}]^T = \sum_{1 \le \alpha \le C(n,d)} A_{\alpha}^{(d)} x_{\alpha}^{(d)}.
$$
 (5)

Theorem 4.1 *Let K be a closed and convex cone in* \mathbb{R}^m *. Let* $G : \mathbb{R}^n \to \mathbb{R}^m$ *be a K*-SOS-convex polynomial with degree $d \geq 2$. Suppose that either *K* is a polyhedral *cone or K has a nonempty interior and* $G(\bar{x}_0) \in \text{int}K$ for some $x_0 \in \mathbb{R}^n$. Then,

$$
\operatorname{epi} G = \left\{ (x, y) \in \mathbb{R}^n \times \mathbb{R}^m : \exists z = (z_{\alpha}) \in \mathbb{R}^{C(d,n)}, y - \sum_{1 \le \alpha \le C(d,n)} G_{\alpha} z_{\alpha} \in K, \right\}
$$

$$
\sum_{1 \le \alpha \le C(n,d)} A_{\alpha}^{(d)} z_{\alpha} \ge 0, z_1 = 1, z_2 = x_1, \dots, z_{n+1} = x_n \right\}.
$$

In particular,

 \mathbf{r}

$$
\{x \in \mathbb{R}^n : G(x) \in -K\} = \left\{x \in \mathbb{R}^n : \exists z = (z_\alpha)_{1 \leq \alpha \leq C(d,n)}, \sum_{1 \leq \alpha \leq C(d,n)} G_\alpha z_\alpha \in -K, \sum_{1 \leq \alpha \leq C(n,d)} A_\alpha^{(d)} z_\alpha \geq 0, z_1 = 1, z_2 = x_1, \dots, z_{n+1} = x_n\right\},\
$$

where G_{α} *and* A_{α}^d *are defined as in* [\(4\)](#page-18-0) *and* [\(5\)](#page-18-1) *respectively.*

 ϵ

Proof Denote the degree of *G* be $d \ge 2$. Clearly *d* must be an even number. It follows that

$$
\begin{split} \text{epi } G &= \left\{ (x, y) : y - \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} z_{\alpha} \in K, z_{\alpha} = x_{\alpha}^{(d)} \right\} \\ &= \left\{ (x, y) : y - \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} z_{\alpha} \in K, z_{\alpha} = x_{\alpha}^{(d)}, [x^{(d/2)}][x^{(d/2)}]^T \succeq 0 \right\} \\ &= \left\{ (x, y) : y - \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} z_{\alpha} \in K, z_{\alpha} = x_{\alpha}^{(d)}, \sum_{1 \le \alpha \le C(n,d)} A_{\alpha}^{(d)} z_{\alpha} \succeq 0 \right\}, \end{split}
$$

where the last equality follows from the definitions of $A_{\alpha}^{(d)}$. This shows that epi $G \subseteq S$ where

$$
S = \left\{ (x, y) \in \mathbb{R}^n \times \mathbb{R}^m : \exists z_{\alpha} \text{ such that } y - \sum_{1 \leq \alpha \leq C(n,d)} G_{\alpha} z_{\alpha} \in K, \right\}
$$

$$
\sum_{1 \leq \alpha \leq C(n,d)} A_{\alpha}^{(d)} z_{\alpha} \geq 0, z_1 = 1, z_2 = x_1, \dots, z_{n+1} = x_n \right\}.
$$

To see the converse inclusion, we proceed by the method of contradiction and suppose there exists $(\overline{x}, \overline{y}) \in S$ but $(\overline{x}, \overline{y}) \notin \text{epi } G$. Then, by the strict separation theorem, there exist $(a, b) \in (\mathbb{R}^n \times \mathbb{R}^m) \setminus \{0_{n+m}\}, r \in \mathbb{R}$ and $\delta > 0$ such that

$$
a^T \overline{x} + b^T \overline{y} \le r < r + 2\delta \le a^T x + b^T y \text{ for all } y \in G(x) + K.
$$

Now, consider the following optimization problem

$$
\min_{(x,y)\in\mathbb{R}^n\times\mathbb{R}^m} \{a^T x + b^T y - (r+\delta) : G(x) - y \in -K\},\tag{P}
$$

where val(\overline{P}) := inf_{(*x,y*)∈ $\mathbb{R}^n \times \mathbb{R}^m$ { $a^T x + b^T y - (r + \delta)$: $G(x) - y \in -K$ }. Then,} we see that val $(\overline{P}) \ge \delta > 0$. From our assumption, either *K* is a polyhedral or *K* has a nonempty interior and $G(x_0) \in \text{int } K$ for some $x_0 \in \mathbb{R}^n$. We now claim that under this assumption, one can find $\lambda \in K^{\oplus}$ and $\sigma \in \Sigma^2[x, y]$ such that, for all $(x, y) \in \mathbb{R}^n \times \mathbb{R}^m$,

$$
a^T x + b^T y + \langle \lambda, G(x) - y \rangle - (r + \delta) = \sigma(x, y). \tag{6}
$$

Granting this, note that $a^T x + b^T y + \langle \lambda, G(x) - y \rangle - (r + \delta) + \epsilon$ is affine and any sum-of-squares polynomial must be of even order. This implies that $b = \lambda \in K^{\oplus}$, and so

$$
a^T x + \langle \lambda, G(x) \rangle - (r + \delta) = \sigma(x)
$$
 and $\sigma \in \Sigma^2[x]$.

So, by the positive semidefinite representation of a sum-of-squares polynomial, there exists $W \geq 0$ such that for all $x \in \mathbb{R}^n$

$$
a^T x + \langle \lambda, G(x) \rangle - (r + \delta) = (x^{(d/2)})^T W x^{(d/2)}
$$

=
$$
\sum_{1 \le \alpha \le C(n,d)} x_{\alpha}^{(d)} \text{Tr} \left(W A_{\alpha}^{(d)} \right).
$$
 (7)

Let $a \in \mathbb{R}^{C(n,d)}$. Define a linear map $L_a : \mathbb{R}_d[x] \to \mathbb{R}$ by

$$
L_u(f) = \sum_{\alpha} f_{\alpha} u_{\alpha} \text{ for all } f(x) = \sum_{\alpha} f_{\alpha} x_{\alpha}^{(d)}.
$$

Note that $(\overline{x}, \overline{y}) \in S$. There exists $\overline{z} \in \mathbb{R}^{C(m+n,d)}$ such that

$$
\overline{y} - \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} \overline{z}_{\alpha} \in K, \quad \sum_{1 \le \alpha \le C(n,d)} A_{\alpha}^{(d)} \overline{z}_{\alpha} \succeq 0
$$

and $\overline{z}_1 = 1$, $\overline{z}_2 = \overline{x}_1$, \ldots , $\overline{z}_{n+1} = \overline{x}_n$.

Applying the linear map $L_{\bar{z}}$ to [\(7\)](#page-20-0), we obtain that

$$
a^T \bar{x} + \left\langle \lambda, \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} \bar{z}_{\alpha} \right\rangle - (r+\delta) = \sum_{1 \le \alpha \le C(n,d)} \bar{z}_{\alpha} \text{Tr}\left(W A_{\alpha}^{(d)}\right) \ge 0.
$$

where the last inequality follows from the fact $\sum_{1 \le \alpha \le C(n,d)} A_{\alpha}^{(d)} \overline{z}_{\alpha} \ge 0$ and $W \ge 0$. Recall that $\lambda = b \in K^{\oplus}$. It follows that

$$
a^T \bar{x} + b^T \bar{y} + \left\langle \lambda, \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} \bar{z}_{\alpha} - \bar{y} \right\rangle - (r + \delta) \ge 0.
$$

This together with $a^T \bar{x} + b^T \bar{y} \le r$ and $\bar{y} - \sum_{1 \le \alpha \le C(n,d)} G_{\alpha} \bar{z}_{\alpha} \in K$ gives us that $-\delta = r - (r + \delta) \ge 0$, which is impossible.

To see our claim [\(6\)](#page-19-0), we first consider the case where *K* is polyhedral. Note that val(\bar{P}) $\geq \delta$ > 0. Thus, Corollary [3.3](#page-15-0) implies that

$$
\sup_{\mu \in \mathbb{R}, \lambda \in K^{\oplus}} \{ \mu : a^T x + b^T y + \langle \lambda, G(x) - y \rangle - (r + \delta) - \mu \in \Sigma_d^2 \} \ge \delta > 0.
$$

Thus, there exist $\lambda \in K^{\oplus}$ and $\sigma \in \Sigma^2[x, y]$ such that, for all $(x, y) \in \mathbb{R}^n \times \mathbb{R}^m$, $a^T x + b^T y + \langle \lambda, G(x) - y \rangle - (r + \delta) = \sigma(x, y) + \delta_0$, for some $\delta_0 > 0$. Thus, the claim (6) holds.

In the case where $G(x_0) \in \text{int } K$ for some $x_0 \in \mathbb{R}^n$, then strict feasibility condition holds for problem (\overline{P}) with (x_0 , 0). So, Theorem [3.2](#page-17-1) implies that there exist $\lambda \in K^{\oplus}$ and $\sigma \in \Sigma^2[x, y]$ such that, for all $(x, y) \in \mathbb{R}^n \times \mathbb{R}^m$,

$$
a^T x + b^T y + \langle \lambda, G(x) - y \rangle - (r + \delta) = \sigma(x, y) + \text{val}(\overline{P}).
$$
 (8)

Thus, claim [\(6\)](#page-19-0) also holds.

The conic linear representation of the set $S := \{x \in \mathbb{R}^n : G(x) \in -K\}$ follows as $\in \mathbb{R}^n : G(x) \in -K\} = \{x \in \mathbb{R}^n : (x, 0) \in \text{eni } G\}$ ${x \in \mathbb{R}^n : G(x) \in -K} = {x \in \mathbb{R}^n : (x, 0) \in \text{epi } G}.$

5 Conclusions

In this paper, we extended the notion of SOS-convexity of a real-valued polynomial to a more general notion of conic SOS-convexity for polynomial mappings. We then established that conic sum-of-squares convex (SOS-convex) polynomial programs exhibit exact conic programming relaxations under suitable regularity assumptions, and so they can be solved by various numerical methods such as interior point methods. As conic convexity and conic programming have found many applications in other branches of mathematics, economics and engineering [\[7\]](#page-21-5), our results point a way to obtaining further computational tractable classes of conic convex programs and their associated applications. This is a topic of further research and will be examined in a forthcoming study.

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