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# **Entanglement-Assisted Quantum Negacyclic BCH Codes**

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Received: 7 June 2018 / Accepted: 30 January 2019 / Published online: 12 February 2019 © Springer Science+Business Media, LLC, part of Springer Nature 2019

#### **Abstract**

The entanglement-assisted quantum error correcting codes (EAQECCs) are a simple and important class of quantum codes. The entanglement-assisted formalism can transform arbitrary classical linear codes into EAQECCs by using pre-shared entanglement between the sender and the receiver. In this paper, by decomposing the defining set of negacyclic BCH codes, we construct a class of new EAQECCs with length  $n = \frac{q^{4m}-1}{q^2-1}$ .

**Keywords** Negacyclic codes · BCH codes · EAQECCs

# **1 Introduction**

Quantum error-correcting codes (QECCs) play an important role in quantum information and computation. As we all know, constructing good QECCs is a crucial subject of research [\[1–](#page-13-0)[8\]](#page-13-1) all the time. Recently, such theory has been extended to EAQECCs. Customarily, an entanglement-assisted quantum error correcting code (EAQECC) can be denoted as  $[[n, k, d; c]]_q$ , which encodes *k* information qubits into *n* channel qubits with the help of *c* pairs of maximally entangled states and corrects up to  $\lfloor \frac{d-1}{2} \rfloor$  errors, where *d* is the minimum distance of the code. If  $c = 0$ , then it is called a *q*-ary standard [[*n*, *k*, *d*]] quantum code. The performance of an EAQECC is measured by its rate  $\frac{k}{n}$  and net rate  $\frac{k-c}{n}$ .

Brun et al. [\[9\]](#page-13-2) proposed an entanglement-assisted stabilized formalism, which overcame the barrier of the dual-containing condition in constructing standard quantum codes from classical codes. They proved that if shared entanglement is available between the sender

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This research is supported by the National Natural Science Foundation of China (No.61772168; No.61572168) and the Natural Science Foundation of Anhui Province (No.1808085MA15).

and the receiver in advance, non-dual-containing classical quaternary codes can be used to construct EAQECCs. Since then, more and more scholars begin to study EAQECCs [\[10–](#page-13-3)[14\]](#page-13-4).

Hsieh et al. [\[15\]](#page-13-5) constructed some EAQECCs with good parameters from quasicyclic low-density parity-check codes. Fujiwara et al. [\[16\]](#page-13-6) used low-density parity-check codes to construct some good parameters' EAQECCs with different lengths soon afterwards. In Refs. [\[17\]](#page-13-7) and [\[18\]](#page-13-8), Li et al. proposed the concept about decomposing the defining set of BCH cyclic codes, transformed the problem of calculating the number of share pairs into determining a special subset of the defining set of a BCH code, and constructed some EAQECCs with good parameters. Afterwards, Lü and Li made a further study on constructing of EAQECCs by using primitive quaternary BCH codes with length  $n = 4^m - 1$  in Ref. [\[19\]](#page-14-0). Recently, Chen et al. [\[20\]](#page-14-1) generalized their method to apply in negacyclic codes, and obtained four classes of optimal EAQECCs. Lü et al.  $[21]$  constructed six classes of  $q$ -ary entanglement-assisted quantum MDS codes based on classical negacyclic MDS codes.

Most of them committed themselves to the construction of entanglement-assisted quantum MDS codes, while the larger length case has received less attention. This reality inspires us to construct EAQECCs with the larger length. In this paper, we obtain a class of new EAQECCs by negacyclic BCH codes with length  $n = \frac{q^{4m}-1}{q^2-1}$ , where *q* is odd and  $m \ge 2$ . Speaking specifically, we construct a class of EAQECCs with parameters as follows:

$$
(1) \quad (i) \quad \text{If } \delta \leq \frac{q^{2m}+3}{2},
$$

$$
\begin{cases} \left[ [n, n - 4m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + 4m\varepsilon - mq^2 + 7m, \ge \delta; 4m] \right]_q, if \varepsilon = 2or3, \\ \left[ [n, n - 4m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + mq^2 + m, \ge \delta; 4m] \right]_q, if \varepsilon = 4. \end{cases}
$$
  
(ii) If  $\frac{q^{2m} + 5}{2} \le \delta \le q^{2m}$ ,

$$
[[n, n-4m](\delta - \frac{3}{2})(1-q^{-2})] + mq^2 + m, \ge \delta; 4m]]_q,
$$
  
where  $q = 3$  and  $m \ge 2$ .

(2)

$$
\begin{cases} \n\quad & [[n, n - 4m\lceil(\delta - \frac{3}{2})(1 - q^{-2})] + 4m, \ge \delta; 4m\lceil \lg, if \varepsilon < 6, \\
\quad & [[n, n - 4m\lceil(\delta - \frac{3}{2})(1 - q^{-2})] + 4\varepsilon m - mq^2 + 7m, \ge \delta; 4m\lceil \lg, if \delta \le \varepsilon \le 11,\n\end{cases}
$$

where  $q = 5$  and  $m > 2$ .

(3)  $[[n, n - 4m](\delta - \frac{3}{2})(1 - q^{-2})] + 4m$ ,  $\geq \delta$ ;  $4m$ ]]<sub>*q*</sub>, where  $q \geq 7$  is a power of an odd prime *p* and  $m \geq 2$ .

This paper is organized as follows. In Section [2,](#page-2-0) some basic background and results about negacyclic codes and BCH codes are reviewed. In Section [3,](#page-3-0) we briefly review some basic definitions and results of EAQECCs. In Section [4,](#page-3-1) we construct a class of EAQECCs with new parameters. In Section [5,](#page-13-9) we give an example to illustrate the significance of results in this paper. Section [6](#page-13-10) concludes the paper.

#### <span id="page-2-0"></span>**2 Preliminaries**

Let *q* be a power of an odd prime *p* and  $\mathbb{F}_{q^2}$  be a finite field with  $q^2$  elements. For any element  $a \in \mathbb{F}_{q^2}$ , we denote the conjugate  $a^q$  of *a* by  $\overline{a}$ . Given two vectors **a** =  $(a_0, a_1, \ldots, a_{n-1})$  and **b** =  $(b_0, b_1, \ldots, b_{n-1}) \in \mathbb{F}_{q^2}^n$ , their Hermitian inner product is defined as

$$
\langle \mathbf{a}, \mathbf{b} \rangle = a_0 \overline{b}_0 + a_1 \overline{b}_1 + \cdots + a_{n-1} \overline{b}_{n-1} \in \mathbb{F}_{q^2}.
$$

The vectors **a** and **b** are called orthogonal with respect to the Hermitian inner product if  $\langle \mathbf{a}, \mathbf{b} \rangle = 0$ . A *q*<sup>2</sup>-ary linear code C of length *n* is a nonempty subspace of the vector space  $\mathbb{F}_{q^2}^n$ . For a  $q^2$ -ary linear code C, the Hermitian dual code of C is defined as

$$
\mathcal{C}^{\perp_h} = \{ \mathbf{a} \in \mathbb{F}_{q^2}^n | \langle \mathbf{a}, \mathbf{b} \rangle = 0 \text{ for all } \mathbf{b} \in \mathcal{C} \}.
$$

A  $q^2$ -ary linear code C of length *n* is called Hermitian self-orthogonal if  $C \subseteq C^{\perp_h}$ , and it is called Hermitian self-dual if  $C = C^{\perp_h}$ . If a  $q^2$ -ary linear code C of length *n* satisfies the property that

$$
(-c_{n-1}, c_0, \ldots, c_{n-2}) \in \mathcal{C}
$$
, forall $(c_0, c_1, \ldots, c_{n-1}) \in \mathcal{C}$ ,

then C is said to be a negacyclic code of length *n* over  $\mathbb{F}_{q^2}$ . Customarily, a codeword  $\mathbf{c} =$  $(c_0, c_1, \ldots, c_{n-1})$  in C is identified with its polynomial representation  $c(x) = c_0 + c_1x + c_2$  $\cdots + c_{n-1}x^{n-1}$ . It is well known that a  $q^2$ -ary negacyclic code of length *n* is precisely an ideal of the quotient ring  $\mathbb{F}_{q^2}[x]/\langle x^n + 1 \rangle$  and C can be generated by a monic divisor  $g(x)$ of  $x^n + 1$ . The polynomial  $g(x)$  is called the generator polynomial of the code C and the dimension of C is  $n - k$ , where  $k = \deg(g(x))$ .

In the following, we assume *q* is a power of an odd prime *p* with  $gcd(n, p) = 1$ , where *n* is a positive integer. Let *β* be a primitive 2*n*-th root of unity in some extension field of  $\mathbb{F}_{q^2}$  and  $\eta = \beta^2$ . Then  $\eta$  is a primitive *n*-th root of unity. Hence,

$$
x^{n} + 1 = \prod_{j=0}^{n-1} (x - \beta \eta^{j}) = \prod_{j=0}^{n-1} (x - \beta^{1+2j}).
$$

Let  $\Omega = \{1 + 2j | 0 \le j \le n - 1\}$ . For each  $i \in \Omega$ , let  $C_i$  be the  $q^2$ -cyclotomic coset modulo 2*n* containing *i*,

$$
C_i = \{i, iq^2, iq^4, \ldots, iq^{2(m_i-1)}\},
$$

where  $m_i$  is the smallest positive integer such that  $iq^{2m_i} \equiv i \mod 2n$ . Each  $C_i$  corresponds to an irreducible divisor of  $x^n + 1$  over  $\mathbb{F}_{q^2}$ . Let C be a negacyclic code of length *n* over  $\mathbb{F}_{q^2}$ with generator polynomial *g(x)*. Then the set  $Z = {i \in \Omega | g(\delta^i) = 0}$  is called the defining set of C. Obviously, the defining set of C must be a union of some  $q^2$ -cyclotomic cosets modulo 2*n* and dim( $C$ ) =  $n - |Z|$ .

For a negacyclic code of length *n* over  $\mathbb{F}_{q^2}$ , it is easy to verify that its Hermitian dual code is still a negacyclic code. Therefore, the Hermitian dual code  $C^{\perp_h}$  of C is still an ideal of  $\mathbb{F}_{q^2}[x]/\langle x^n + 1 \rangle$ . Hence, if *Z* is the defining set of *C*, then its Hermitian dual code  $C^{\perp_h}$ has defining set  $Z^{\perp_h} = \{z \in \Omega | - qz \mod 2n \notin Z\}$ . Note that  $Z^{-q} = \{-qz \mod 2n | z \in Z\}$ . Then C contains its Hermitian dual code if and only if  $Z \cap Z^{-q} = \emptyset$  from Lemma 2.2 in Ref. [\[22\]](#page-14-3).

Let *q* be a power of an odd prime *p* with  $gcd(n, p) = 1$  and  $\beta$  be a primitive 2*n*-th root of unity. A negacyclic BCH code of length *n* over  $\mathbb{F}_{q^2}$  with designed distance  $\delta$  is a negacyclic code with generator polynomial

$$
g(x) = \prod_{j \in Z} (x - \beta^j), where Z = \bigcup_{j=b}^{b+\delta-2} C_{1+2j} and b is some integer.
$$

Let  $\mathcal{C}_{(n,q^2,b,\delta)}$  denote the negacyclic BCH codes of length *n* with generator polynomial *g*(*x*). If  $b = 0$ , then we abbreviate  $C_{(n,q^2,b,\delta)}$  as  $C_{(n,q^2,\delta)}$ . Similarly to BCH codes, negacyclic BCH codes have the following property.

**Theorem 1** [\[24\]](#page-14-4) *(The BCH bound for negacyclic codes) Assume that*  $gcd(n, q) = 1$ *. Let* C be a negacyclic code of length *n* over  $\mathbb{F}_{q^2}$ , and let its generator polynomial  $g(x)$  have *elements*  $\{\beta^{1+2j} | 0 \le j \le d-2\}$  *as the roots, where*  $\beta$  *is a primitive* 2*n-th root of unity. Then the minimum distance of*  $C$  *is at least d.* 

### <span id="page-3-0"></span>**3 Review of EAQECCs**

In this section, we give some basic definitions and results of EAQECCs. For more details about EAQECCs theory, please refer to Refs. [\[9–](#page-13-2)[21\]](#page-14-2) therein.

Suppose that *H* is an  $(n - k) \times n$  parity check matrix of C over  $\mathbb{F}_{q^2}$ . Then  $C^{\perp_h}$  has an  $n \times (n-k)$  generator matrix *H*<sup>†</sup>, where *H*<sup>†</sup> is the conjugate transpose matrix of *H* over  $\mathbb{F}_{q^2}$ .

Similarly to the CSS construction of stabilizer quantum codes, there is the following construction method for EAQECCs in Refs. [\[9\]](#page-13-2) and [\[10\]](#page-13-3).

**Theorem 2** [\[9,](#page-13-2) [10\]](#page-13-3) *If*  $C = [n, k, d]_q^2$  *is a classical code over*  $\mathbb{F}_{q^2}$  *and H is its parity check matrix, then*  $C^{\perp_h}$  *stabilizes an entanglement-assisted code with parameters* [[*n,* 2*k* − *n* +  $c, d; c]$ <sub>*d*</sub>, where  $c = rank(HH^{\dagger})$  *is the number of maximally entangled states required and*  $H^{\dagger}$  *is the conjugate matrix of*  $H$  *over*  $\mathbb{F}_{q^2}$ *.* 

#### <span id="page-3-1"></span>**4 Construction of Entanglement-Assisted Quantum BCH Codes**

In Ref. [\[20\]](#page-14-1), the authors gave the following definition and lemma which can determine the number of entangled states by decomposing the defining set of negacyclic codes.

**Definition 1** [\[20\]](#page-14-1) Let C be a negacyclic code of length *n* with defining set Z. Assume that *Z*<sub>1</sub> = *Z* ∩ (−*qZ*) and *Z*<sub>2</sub> = *Z* \ *Z*<sub>1</sub>, where −*qZ* = {*n* − *qx*|*x* ∈ *Z*}. Then *Z* = *Z*<sub>1</sub> ∪ *Z*<sub>2</sub> is called a decomposition of the defining set of  $C$ .

**Lemma 1** [\[20\]](#page-14-1) *Let* C *be a negacyclic code of length <i>n over*  $\mathbb{F}_{q^2}$ *, where*  $gcd(n, q) = 1$ *. Suppose that Z is the defining set of the negacyclic code*  $C$  *and*  $Z = Z_1 \cup Z_2$  *is a decomposition of Z. Then the number of entangled states required is*  $c = |Z_1|$ *. In order to construct the entanglement-assisted quantum BCH codes, we firstly give two lemmas below.*

**Lemma 2** [23] 
$$
\mathcal{C}^{\perp_h}_{(n,q^2,\delta)} \subseteq \mathcal{C}_{(n,q^2,\delta)}
$$
 if and only if  $2 \leq \delta \leq \delta_{max}$ , where  $\delta_{max} = \frac{q^{2m+1}-q}{q^2-1}$ .

**Lemma 3** [\[23\]](#page-14-5) *Let*  $n = \frac{q^{4m}-1}{q^2-1}$  *and*  $m \ge 2$ *. Let i be an integer such that*  $0 \le i \le q^{2m}$  *and*  $i \neq \frac{q^2-1}{2} \mod q^2$ . If  $q^2 \equiv 1 \mod 4$ , then

$$
|C_{1+2i}| = \begin{cases} m, & i = \frac{q^{2m}-1}{4}, \frac{3q^{2m}+1}{4}, \\ 2m, & otherwise. \end{cases}
$$

In addition,  $1 + 2i$  is not a coset leader in the following cases:

$$
i\in\left\{b\cdot\frac{q^{2m}-1}{q^2-1}:b\in\left[\frac{q^2+3}{4},\frac{q^2-3}{2}\right]\right\}\bigcup\left\{\frac{q^{2m}+1}{2}+b\cdot\frac{q^{2m}-1}{q^2-1}:b\in\left[\frac{q^2+3}{4},\frac{q^2-1}{2}\right]\right\}.
$$

#### **4.1 The Number of Entangled States**

**Theorem 3** *Let*  $n = \frac{q^{4m}-1}{q^2-1}$ , where q is a power of an odd prime p and  $m \ge 2$ . Then we have  $|C_{1+2(\frac{q^{2m+1}-q}{q^2-1}-1)}| = |C_{\frac{2q^{2m+1}-q^2-2q+1}{q^2-1}}| = 2m$  and  $|C_{\frac{2q^{4m}-2q^{2m+2}+q^3+2q^2-q-2}{q^2-1}}| = 2m$ .

*Proof* On the one hand, since

 $(q^{2m} - 1)(q^2 - 1) - 4(q^{2m+1} - q^2 - q + 1) = q^{2m+2} - 4q^{2m+1} - q^{2m} + 3q^2 + 4q - 3$ if *q* = 3, then  $(q^{2m}-1)(q^2-1)$  <  $4(q^{2m+1}-q^2-q+1)$ ; if *q* ≥ 5, then  $(q^{2m}-1)(q^2-1)$  >  $4(q^{2m+1} - q^2 - q + 1)$ . Therefore,  $\frac{q^{2m+1}-q}{q^2-1} - 1 \neq \frac{q^{2m}-1}{4}$ .

On the other hand,

$$
\frac{3q^{2m}+1}{4} - (\frac{q^{2m+1}-q}{q^2-1}-1) = \frac{3q^{2m+2}-4q^{2m+1}-3q^{2m}+5q^2+4q-5}{4(q^2-1)},
$$

if *q* = 3, then  $\frac{3q^{2m}+1}{4} < \frac{q^{2m+1}-q}{q^2-1} - 1$ ; if *q* ≥ 5, then  $\frac{3q^{2m}+1}{4} > \frac{q^{2m+1}-q}{q^2-1} - 1$ . Therefore,  $q^{2m+1}-q$  *q*<sup>2−1</sup> − 1 ≠  $\frac{3q^{2m}+1}{4}$ . From Lemma 3, we have  $|C_{1+2(q^{2m+1}-q-1)}| = 2m$  immediately. From  $\frac{2q^{4m}-2q^{2m+2}+q^3+2q^2-q-2}{q^2-1} = 2n - q(1+2(\frac{q^{2m+1}-q}{q^2-1}-1))$  and  $|C_{1+2(\frac{q^{2m+1}-q}{q^2-1}-1)}| =$ 2*m*, we have  $|C_{\frac{2q^{4m}-2q^{2m+2}+q^3+2q^2-q-2}{q^2-1}}| = 2m$  immediately.

**4.2** The Dimension of EAQECCs with  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ 

**Lemma 4** [\[23\]](#page-14-5) *Let*  $n = \frac{q^{4m}-1}{q^2-1}$ *, where q is odd and*  $m \ge 2$ *. (i) If*  $\delta \leq \frac{q^{2m}+3}{2}$ , then we define  $\varepsilon = \left(\frac{(\delta-2)(q^2-1)}{q^{2m}-1}\right)$ *q*2*m*−1  $\int$  *and*  $C_{(n,q^2,\delta)}$  *has dimension*  $k =$  $\sqrt{ }$  $\int$  $\overline{\mathcal{N}}$  $n-2m\left[ (\delta - \frac{3}{2})(1-q^{-2}) \right],$  *if*  $\varepsilon < \left[ \frac{q^2-1}{4} \right],$  $n-2m\left[(\delta-\frac{3}{2})(1-q^{-2})\right]+m(2\varepsilon-\frac{q^2-3}{2}), if \left|\frac{q^2-1}{4}\right| \leq \varepsilon \leq \frac{q^2-3}{2},$  $n-2m\left[(\delta-\frac{3}{2})(1-q^{-2})\right]+m\frac{q^2-3}{2}, if \varepsilon > \frac{q^2-3}{2}.$ 

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(*ii*) If  $\frac{q^{2m}+5}{2} \le \delta \le q^{2m}$ , then we define  $\varepsilon = \left( \frac{(\delta - \frac{q^{2m}+5}{2})(q^2-1)}{q^{2m}-1} \right)$ *q*2*m*−1 *and*  $\mathcal{C}_{(n,q^2,\delta)}$  *has dimension*

$$
k = \begin{cases} n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m \frac{q^2 - 3}{2}, if \varepsilon < \left[ \frac{q^2 - 1}{4} \right], \\ n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + 2m\varepsilon, if \left[ \frac{q^2 - 1}{4} \right] \le \varepsilon \le \frac{q^2 - 3}{2}, \\ n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m(q^2 - 3), if \varepsilon > \frac{q^2 - 3}{2}. \end{cases}
$$

*In order to calculate the dimension of EAQECCs with*  $n = \frac{q^{4m}-1}{q^2-1}$  *and*  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ *,* where q is a power of an odd prime p,  $m \ge 2$  and  $1 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , we need to determine the range of  $\delta$  for distinct q from Lemma 4. So we give the following theorem *firstly.*

**Theorem 4** *Let*  $n = \frac{q^{4m}-1}{q^2-1}$  *and*  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ *, where q is a power of an odd prime p,*  $m \geq 2$  *and*  $1 \leq t \leq \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ .

(i) If 
$$
q = 3
$$
, then  $\delta_{\min} < \frac{q^{2m} + 3}{2}$  and  $\frac{q^{2m} + 5}{2} < \delta_{\max} < q^{2m}$ .

$$
(ii) \tIf q \ge 5, then \delta_{max} < \frac{q^{2m}+3}{2}.
$$

Proof Since 
$$
\delta = \frac{q^{2m+1}-q}{q^2-1} + t
$$
 and  $1 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , we have  

$$
\frac{q^{2m+1}-q}{q^2-1} + 1 \le \delta \le \frac{3q^{2m+1}+2q^{2m}-3q^{2m-1}+q^2-3}{2q^2-2}.
$$

From

$$
\frac{q^{2m}+3}{2}-\delta_{max}=\frac{q^{2m}+3}{2}-\frac{3q^{2m+1}+2q^{2m}-3q^{2m-1}+q^2-3}{2q^2-2}=\frac{q^{2m-1}(q^3-3q^2-3q+3)+2q^2}{2q^2-2},
$$

if  $q = 3$ , then  $\frac{q^{2m}+3}{2} - \delta_{max} < 0$ ; if  $q \ge 5$ , then  $\frac{q^{2m}+3}{2} - \delta_{max} > 0$ . When  $q = 3$ , we have

$$
q^{2m} - \delta_{max} = q^{2m} - \frac{3q^{2m+1} + 2q^{2m} - 3q^{2m-1} + q^2 - 3}{2q^2 - 2} = \frac{q^{2m-1}(2q^3 - 3q^2 - 4q + 3) - q^2 + 3}{2q^2 - 2} > 0,
$$

and

$$
\frac{q^{2m}+3}{2} - \delta_{min} = \frac{q^{2m}+3}{2} - (\frac{q^{2m+1}-q}{q^2-1}+1) = \frac{q^{2m+2}-2q^{2m+1}-q^{2m}+q^2+2q-1}{2(q^2-1)} > 0,
$$
  
i.e.,  $\delta_{min} < \frac{q^{2m}+3}{2}$  and  $\frac{q^{2m}+5}{2} < \delta_{max} < q^{2m}$ .

From the discussion above, we can determine the dimension of the negacyclic BCH codes with length  $n = \frac{q^{4m}-1}{q^2-1}$ , where *q* is a power of an odd prime *p* and  $m \ge 2$ .

**Theorem 5** *Let*  $n = \frac{q^{4m}-1}{q^2-1}$  *and*  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ *, where*  $q = 3$ *,*  $m \ge 2$  *and*  $1 \le t \le$  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ .  $f(i)$  *If*  $\delta \leq \frac{q^{2m}+3}{2}$ , then  $\varepsilon = \left(\frac{(\delta-2)(q^2-1)}{q^{2m}-1}\right)$ *q*2*m*−1 *and*  $\mathcal{C}_{(n,q^2,\delta)}$  *has dimension*  $k =$  $\sqrt{ }$ ⎨  $\mathbf{I}$  $n-2m\left[(\delta-\frac{3}{2})(1-q^{-2})\right]+m(2\varepsilon-\frac{q^2-3}{2}),$  *if*  $\varepsilon=2$  *or* 3*,*  $n-2m\left[(\delta-\frac{3}{2})(1-q^{-2})\right]+m\frac{q^2-3}{2}, if \varepsilon=4.$ *(ii) If*  $\frac{q^{2m}+5}{2} \le \delta \le q^{2m}$ , then  $\varepsilon = \left| \frac{(\delta-\frac{q^{2m}+5}{2})(q^2-1)}{(q^{2m}-1)} \right|$ *q*2*m*−1  $\left| \right.$  *and*  $C_{(n,q^2,\delta)}$  *has dimension*  $k = n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m \frac{q^2 - 3}{2}.$ 

*Proof* From Lemma 4 and Theorem 4, we have the following results:

(i) If 
$$
\frac{q^{2m+1}-q}{q^2-1} + 1 \le \delta \le \frac{q^{2m}+3}{2}
$$
, then  $\varepsilon = \left[ \frac{(\delta-2)(q^2-1)}{q^{2m}-1} \right]$ . Therefore, we have  
\n
$$
\frac{q^{2m+1}-q^2-q+1}{q^{2m}-1} \le \varepsilon \le 4.
$$
\nSince  $\frac{q^2-1}{4} - \varepsilon_{min} = \frac{q^{2m+2}-4q^{2m+1}-q^{2m}+3q^2+4q-3}{4(q^{2m}-1)} < 0$ , then  $\varepsilon_{min} > \left[ \frac{q^2-1}{4} \right]$ . From  $\frac{q^2-3}{2} - \varepsilon_{min} = \frac{q^2-3}{2} - \frac{q^{2m+1}-q^2-q+1}{q^{2m}-1} = \frac{q^2+2q+1}{2(q^{2m}-1)} > 0$ , we have  $\varepsilon_{min} < \frac{q^2-3}{2}$ . Besides,  
\n $\varepsilon_{max} = 4 > \frac{q^2-3}{2} = 3$ . Therefore,  $C_{(n,q^2,\delta)}$  has dimension as below:

$$
k = \begin{cases} n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m(2\varepsilon - \frac{q^2 - 3}{2}), if \varepsilon = 2 \text{ or } 3, \\ n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m\frac{q^2 - 3}{2}, if \varepsilon = 4. \end{cases}
$$

(ii) If 
$$
\frac{q^{2m}+5}{2} \le \delta \le \frac{q^{2m+2}+q^{2m}+6}{2(q^2-1)}
$$
, then  $\varepsilon = \left\lfloor \frac{(\delta - \frac{q^{2m}+5}{2}) (q^2-1)}{q^{2m}-1} \right\rfloor$ . Therefore, we have  
\n
$$
0 \le \varepsilon \le \frac{2q^{2m} - 5q^2 + 11}{2(q^{2m} - 1)}.
$$
\nFrom  $\frac{q^2-1}{4} - \varepsilon_{max} = \frac{q^2-1}{4} - \frac{2q^{2m}-5q^2+11}{2(q^{2m}-1)} = \frac{q^{2m+1}+q^{2m}+q^4-21}{4(q^{2m}-1)} > 0$ , we have  $\varepsilon_{max}$ 

 $\frac{q^2-1}{4}$ . Therefore,  $C_{(n,q^2,\delta)}$  has dimension  $k = n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m \frac{q^2 - 3}{2}$ .

**Theorem 6** *Let*  $n = \frac{q^{4m}-1}{q^2-1}$ ,  $\varepsilon = \left[ \frac{(\delta-2)(q^2-1)}{q^{2m}-1} \right]$ *q*2*m*−1  $\int$  *and*  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ *, where*  $q \ge 5$  *is odd,*  $m \geq 2$  *and*  $1 \leq t \leq \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ .

*(i)* If  $q = 5$ , then  $\mathcal{C}_{(n,q^2,\delta)}$  has dimension

$$
k = \begin{cases} n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right], \text{if } \varepsilon < 6, \\ n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m(2\varepsilon - \frac{q^2 - 3}{2}), \text{if } 6 \le \varepsilon \le 11. \end{cases}
$$

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(*ii*) If 
$$
q \ge 7
$$
, then  $C_{(n,q^2,\delta)}$  has dimension  $k = n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right]$ .

*Proof* From Lemma 4 and Theorem 4, we can define  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$  and  $1 \le t \le$  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ . Then we have

$$
\frac{q^{2m+1}-q}{q^2-1}+1 \le \delta \le \frac{3q^{2m+1}+2q^{2m}-3q^{2m-1}+q^2-3}{2q^2-2}.
$$

From  $\varepsilon = \left| \frac{(\delta - 2)(q^2 - 1)}{q^{2m} - 1} \right|$ *q*2*m*−1 , we have

$$
\frac{q^{2m+1}-q^2-q+1}{q^{2m}-1}\leq \varepsilon\leq \frac{3q^{2m+1}+2q^{2m}-3q^{2m-1}-3q^2+1}{2(q^{2m}-1)}.
$$

On the one hand,  $\frac{q^2-1}{4} - \varepsilon_{min} = \frac{q^{2m+2}-4q^{2m+1}-q^{2m}+3q^2+4q-3}{4(q^{2m}-1)} > 0$ , then  $\varepsilon_{min} < \left\lfloor \frac{q^2-1}{4} \right\rfloor$ ; on the other hand,

$$
\varepsilon_{max} - \frac{q^2 - 1}{4} = \frac{-q^{2m-1}(q^3 - 6q^2 - 5q + 6) - 5q^2 + 1}{4(q^{2m} - 1)}.
$$

If  $q = 5$ , then  $\varepsilon_{max} > \frac{q^2-1}{4}$ ; if  $q \ge 7$ , then  $\varepsilon_{max} < \frac{q^2-1}{4}$ . Therefore, when  $q \ge 7$ , we have  $\varepsilon < \frac{q^2-1}{4}$ , then  $C_{(n,q^2,\delta)}$  has dimension  $k = n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right]$ .

As for  $q = 5$ , from  $\frac{q^2-3}{2} - \varepsilon_{max} = \frac{q^{2m-1}(q^3-3q^2-5q+3)+2q^2+2}{2(q^{2m}-1)} > 0$ , we have  $\varepsilon_{max} <$  $\frac{q^2-3}{2}$ . Then  $\mathcal{C}_{(n,q^2,\delta)}$  has dimension:

$$
k = \begin{cases} n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right], \text{if } \varepsilon < 6, \\ n - 2m \left[ (\delta - \frac{3}{2})(1 - q^{-2}) \right] + m(2\varepsilon - \frac{q^2 - 3}{2}), \text{if } 6 \le \varepsilon \le 11. \end{cases}
$$

#### **4.3 The Construction of EAQECCs**

Based on the discussions above, we give a theorem below.

**Theorem 7** *Let*  $n = \frac{q^{4m}-1}{q^2-1}$ ,  $q \ge 7$  *is a power of an odd prime*  $p, m \ge 2$  *and*  $\delta =$  $\frac{q^{2m+1}-q}{q^2-1}$  + t, where  $q^2 \equiv 1 \mod 4$  and  $1 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ . If C is a  $q^2$ -ary negacyclic  $\int_{c}^{2m+1} - q^{2m+1} - q^{2}}{C_{1+2i}}$ , then there exist EAQECCs *with parameters*  $[[n, n - 4m[(\delta - \frac{3}{2})(1 - q^{-2})] + 4m, \ge \delta; 4m]]_q$ .

*Proof* From Lemma 2, we can assume that the defining set of the negacyclic code C is  $Z = \bigcup_{i=0}^{\delta-2} C_{1+2i}$ , where  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ ,  $q \ge 7$ ,  $m \ge 2$  and  $1 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ . Then C is a negacyclic code with parameters  $[n, n - 2m[(\delta - \frac{3}{2})(1 - q^{-2})], \ge \delta]_{q^2}$  from Theorems 1 and 6. Therefore, we have the following result:

<span id="page-8-0"></span>
$$
Z_{1} = Z \cap (-qZ)
$$
\n
$$
= ((\bigcup_{i=0}^{q^{2m+1}-q} - 1 C_{1+2i}) \cup (\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i})
$$
\n
$$
= ((\bigcup_{i=0}^{q^{2m+1}-q} - 1 C_{1+2i}) \cup -q(\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i})
$$
\n
$$
\cap (-q(\bigcup_{i=0}^{q^{2m+1}-q} - 1 C_{1+2i}) \cup -q(\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i})
$$
\n
$$
= ((\bigcup_{i=0}^{q^{2m+1}-q} - 1 C_{1+2i}) \cap -q(\bigcup_{i=0}^{q^{2m+1}-q} - 1 C_{1+2i}) )
$$
\n
$$
= \bigcup ((\bigcup_{i=0}^{q^{2m+1}-q} - 1 C_{1+2i}) \cap -q(\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i})
$$
\n
$$
\bigcup ((\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i}) \cap -q(\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 1 C_{1+2i})
$$
\n
$$
\bigcup ((\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i}) \cap -q(\bigcup_{i=0}^{q^{2m+1}-q} - 2 + i C_{1+2i}) )
$$
\n
$$
= \bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i} \cap -q(\bigcup_{i=\frac{q^{2m+1}-q}{q^{2-1}} - 2 + i C_{1+2i})
$$
\n
$$
= C_{\frac{2q^{2m+1}-q^{2}-2q+1}{q^{2-1}}} \cup C_{\frac{2q^{4m}-2q^{2m+2}+q^{3}+2q^{2}-q-2}{q^{2-1}}}.
$$
\n(1)

In order to get the result of  $(1)$ , we have to show that

$$
\begin{array}{l} q^{2m+1}-q \cr (\cup_{i=0}^{q^{2-1}}-1 \cr \cr (-q^{2-1}-1 \cr \cr (-q^{2m+1}-q \cr )-q \cr (\cup_{i=0}^{q^{2m+1}-q}q^{2} \cr ) \cr (1+2i) \cap -q \cr (\cup_{i=0}^{q^{2m+1}-q}2+i \cr (\cup_{i=0}^{q^{2m+1}-q}1 \cr (-q^{2m+1}-q \cr \cr (-q^{2m
$$

and

$$
(\cup_{\substack{q^{2m+1}-q \\ i \equiv \frac{q^{2m+1}-q}{q^2-1}}}^{ \frac{q^{2m+1}-q}{q^2-1}-2+t} C_{1+2i}) \cap -q (\cup_{\substack{q^{2m+1}-q \\ i \equiv \frac{q^{2m+1}-q}{q^2-1}}}^{ \frac{q^{2m+1}-q}{q^2-1}-2+t} C_{1+2i}) = \emptyset.
$$

Firstly, we demonstrate that

$$
(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-1}C_{1+2i})\cap-q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-1}C_{1+2i})=C_{\frac{2q^{2m+1}-q^2-2q+1}{q^2-1}}\cup C_{\frac{2q^{4m}-2q^{2m+2}+q^3+2q^2-q-2}{q^2-1}}.
$$

From Lemma 2,  $\mathcal{C}^{\perp_h}_{(n,q^2,\delta_{max})} \subseteq \mathcal{C}_{(n,q^2,\delta_{max})}$  means that

$$
(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-2}C_{1+2i})\cap-q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-2}C_{1+2i})=\emptyset.
$$

<sup>2</sup> Springer

And from  $C^{\perp_h}_{(n,q^2,\delta_{max}+1)} \nsubseteq C_{(n,q^2,\delta_{max}+1)}$  in the proof of Lemma 2 in Ref. [\[23\]](#page-14-5), we can get *C*1+2*(δmax*−1*)* ∩ −*q(*∪  $q^{2m+1-q}_{q^2-1}$  *C*<sub>1+2*i*</sub>) = *C*<sub>1+2</sub>( $\delta_{max-1}$ ) immediately, where  $C_{1+2(\delta_{max}-1)}$  =  $C_{\frac{2q^{2m+1}-q^{2}-2q+1}{q^{2}-1}}$ , i.e.,

$$
C_{\frac{2q^{2m+1}-q^{2}-2q+1}{q^{2}-1}}\cap -q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^{2}-1}-1}C_{1+2i})=C_{\frac{2q^{2m+1}-q^{2}-2q+1}{q^{2}-1}}.
$$

From the equation above, we can get

$$
-qC_{\frac{2q^{2m+1}-q^2-2q+1}{q^2-1}}\cap (\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}}C_{1+2i})=-qC_{\frac{2q^{2m+1}-q^2-2q+1}{q^2-1}}=C_{\frac{2q^{4m}-2q^{2m+2}+q^3+2q^2-q-2}{q^2-1}}.
$$
 Thus, the desired result 
$$
(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}}C_{1+2i})\cap -q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}}C_{1+2i})=C_{\frac{2q^{2m+1}-q^2-2q+1}{q^2-1}}\cup C_{\frac{2q^{4m}-2q^{2m+2}+q^3+2q^2-q-2}{q^2-1}} \text{ follows.}
$$
 Secondly, we testify

$$
(\cup_{\substack{q^{2m+1}-q \\ i \equiv q^{2m+1}-q}}^{\frac{q^{2m+1}-q}{q^2-1}-2+t} C_{1+2i}) \cap -q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}} C_{1+2i}) = \emptyset,
$$

for 
$$
2 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}
$$
.  
\nIf  $\left(\bigcup_{\substack{q^{2m+1}-q\\i=\frac{2m+1-q}{q^2-1}}} \frac{q^{2m+1}-q}{C_{1+2i}}\right) \cap -q\left(\bigcup_{\substack{q^{2-1} \\i=0}} \frac{q^{2m+1}-q}{C_{1+2i}}\right) \ne \emptyset$ , for  $2 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , i.e.,

$$
(\cup_{i=2}^t C_{1+2(i+\frac{q^{2m+1}-q}{q^2-1}-2)})\cap-q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-1}C_{1+2i})\neq\emptyset,
$$

for  $2 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , then there exist two integers *l* and *j*, where  $2 \le l \le$  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , 0 ≤  $j \leq \frac{q^{2m+1}-q}{q^2-1} - 1$ , such that

$$
1 + 2(\frac{q^{2m+1} - q}{q^2 - 1} - 2 + l) \equiv -q(1 + 2j)q^{2k} \mod 2n,
$$

for some  $k \in \{0, 1\}$ . We can seek contradictions as follows.

(i) When  $k = 0$ , we have  $1 + 2(\frac{q^{2m+1}-q}{q^2-1} - 2 + l) \equiv -q(1+2j) \mod 2n$ , which is equivalent to

<span id="page-9-0"></span>
$$
\frac{2q^{2m+1} - 2q}{q^2 - 1} + q - 3 + 2l + 2jq \equiv 0 \mod 2n,
$$
 (2)

where 
$$
n = \frac{q^{4m}-1}{q^2-1}
$$
. From  $2 \le l \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}, 0 \le j \le \frac{q^{2m+1}-q}{q^2-1} - 1$ , then  $\frac{2q^{2m+1}-2q}{q^2-1} + q + 1 \le \frac{2q^{2m+1}-2q}{q^2-1} + q - 3 + 2l + 2jq \le \frac{2q^{2m+2}+3q^{2m+1}+2q^{2m}-3q^{2m-1}-q^3-4q^2+q}{q^2-1}$ . Because of  $\frac{2q^{2m+2}+3q^{2m+1}+2q^{2m}-3q^{2m-1}-q^3-4q^2+q}{q^2-1} < 2n$ , (2) is not established.

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(ii) When  $k = 1$ , we have  $1 + 2(\frac{q^{2m+1}-q}{q^2-1} - 2 + l) \equiv -q^3(1+2j) \mod 2n$ , which is equivalent to

<span id="page-10-0"></span>
$$
\frac{2q^{2m+1} - 2q}{q^2 - 1} + q^3 - 3 + 2l + 2jq^3 \equiv 0 \mod 2n,
$$
 (3)

where 
$$
n = \frac{q^{4m}-1}{q^2-1}
$$
. From  $2 \le l \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}, 0 \le j \le \frac{q^{2m+1}-q}{q^2-1} - 1$ , we have:  $\frac{2q^{2m+1}-2q}{q^2-1} + q^3 + 1 \le \frac{2q^{2m+1}-2q}{q^2-1} + q^3 - 3 + 2l + 2jq^3 \le \frac{2q^{2m+4}+3q^{2m+1}+2q^{2m}-3q^{2m-1}-q^5-2q^4+q^3-2q^2}{q^2-1}$ . Because of  $\frac{2q^{2m+4}+3q^{2m+1}+2q^{2m}-3q^{2m-1}-q^5-2q^4+q^3-2q^2}{q^2-1} < 2n$ , (3) is not established.

From the discussions above, we can see

$$
(\cup_{\substack{q^{2m+1}-q\\i=q^{2m+1}-q\\ q^2-1}}^{q^{2m+1}-q-2+t}C_{1+2i})\cap-q(\cup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}}C_{1+2i})=\emptyset,
$$

for 2 ≤ *t* ≤  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ .<br>Next, we show that

$$
(\cup_{i=0}^{q^{2m+1}-q}C_{1+2i})\cap-q(\cup_{\substack{q^{2m+1}-q\\i=\frac{q^{2m+1}-q}{q^{2}-1}}}^{q^{2m+1}-q-2+t}C_{1+2i})=\emptyset.
$$

Since

$$
-q((\bigcup_{\substack{q^{2m+1}-q\\i=\frac{q^{2m+1}-q}{q^2-1}}}^{\frac{q^{2m+1}-q}{q^2-1}-2+t}C_{1+2i})\cap-q(\bigcup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-1}C_{1+2i}))=-q\emptyset=\emptyset,
$$

it follows that

$$
(\bigcup_{i=0}^{q^{2m+1}-q}C_{1+2i})\cap-q(\bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}}^{q^{2m+1}-q-2+t}C_{1+2i})=\emptyset.
$$

Finally, we show that

$$
\sum_{i=\frac{q^{2m+1}-q}{q^2-1}} \frac{q^{2m+1}-q}{C_{1+2i}} - 2 + t
$$
\n
$$
(\bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}} \frac{1}{C_{1+2i}}) \cap -q \bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}} \frac{1}{C_{1+2i}} - 2 + t
$$
\n
$$
\text{for } 2 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}. \text{ If } (\bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}} \frac{1}{C_{1+2i}}) \cap -q \bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}} \frac{1}{C_{1+2i}} - 2 + t
$$
\n
$$
(\bigcup_{i=2}^t C_{1+2(i+\frac{q^{2m+1}-q}{q^2-1}-2)}) \cap -q \bigcup_{i=2}^t C_{1+2(i+\frac{q^{2m+1}-q}{q^2-1}-2)}) \ne \emptyset,
$$

for  $2 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , then there exist two integers *l* and *j*, where  $2 \le l, j \le$  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , such that

$$
1 + 2(l + \frac{q^{2m+1} - q}{q^2 - 1} - 2) \equiv -q[1 + 2(j + \frac{q^{2m+1} - q}{q^2 - 1} - 2)]q^{2k} \mod 2n,
$$

for some  $k \in \{0, 1\}$ . We can seek contradictions as follows.

(i) When  $k = 0$ , we have  $1 + 2(l + \frac{q^{2m+1}-q}{q^2-1} - 2) \equiv -q[1+2(j + \frac{q^{2m+1}-q}{q^2-1} - 2)] \mod 2n$ , which is equivalent to

<span id="page-11-0"></span>
$$
\frac{2q^{2m+2} + 2q^{2m+1} - 2q^2 - 2q}{q^2 - 1} - 3q - 3 + 2l + 2jq \equiv 0 \mod 2n,\tag{4}
$$

where 
$$
n = \frac{q^{4m}-1}{q^2-1}
$$
. From  $2 \le l, j \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , we have:  $\frac{2q^{2m+2}+2q^{2m+1}+q^3-q^2-3q-1}{q^2-1} \le \frac{2q^{2m+2}+2q^{2m+1}-2q^2-2q}{q^2-1} - 3q -$   
\n $3 + 2l + 2jq \le \frac{3q^{2m+2}+5q^{2m+1}-q^{2m}-3q^{2m-1}-2q^3-2q^2}{q^2-1}$ . Because of  $\frac{3q^{2m+2}+5q^{2m+1}-q^{2m}-3q^{2m-1}-2q^3-2q^2}{q^2-1} < 2n$ , (4) is not established.

(ii) When  $k = 1$ , we have  $1+2(l+\frac{q^{2m+1}-q}{q^2-1}-2) \equiv -q^3[1+2(j+\frac{q^{2m+1}-q}{q^2-1}-2)] \mod 2n$ , which is equivalent to

<span id="page-11-1"></span>
$$
2l + 2jq^3 \equiv \frac{2q^{4m} - 2q^{2m+4} - 2q^{2m+1} + 3q^5 + 2q^4 - 3q^3 + 3q^2 + 2q - 5}{q^2 - 1}
$$
 mod 2n. (5)

From  $2 \le l$ ,  $j \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ , we have

$$
4+4q^3\le 2l+2jq^3\le \frac{q^{2m+4}+2q^{2m+3}-3q^{2m+2}+q^{2m+1}+2q^{2m}-3q^{2m-1}+q^5+2q^4-3q^3+q^2+2q-3}{q^2-1}.
$$

Then we have the following results:

- (a) when  $m = 2$ , since  $(4 + 4q^3)(q^2 1) = 4q^5 4q^3 + 4q^2 4 > q^5 + 2q^4$  $3q^{3} + 3q^{2} + 2q - 5$ , [\(5\)](#page-11-1) is not established;
- (b) when  $m > 3$ , we have

$$
\frac{2q^{4m} - 2q^{2m+4} - 2q^{2m+1} + 3q^5 + 2q^4 - 3q^3 + 3q^2 + 2q - 5}{q^2 - 1} > (2l + 2jq^3)_{max}.
$$

Therefore, [\(5\)](#page-11-1) is not established.

From the discussions above, we can see

$$
\bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}}^{\frac{q^{2m+1}-q}{q^2-1}-2+t}C_{1+2i}\big) \cap -q\bigcup_{i=\frac{q^{2m+1}-q}{q^2-1}}^{\frac{q^{2m+1}-q}{q^2-1}-2+t}C_{1+2i}\big) = \emptyset,
$$

for  $2 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ .

From Lemma 1 and Theorem 3, we have  $c = 4m$ . From Theorem 2, there exist entanglement assisted quantum codes with parameters

$$
[[n, n - 4m[(\delta - \frac{3}{2})(1 - q^{-2})] + 4m, \ge \delta; 4m]]_q,
$$
  

$$
\le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)}.
$$

When  $q = 3$  or 5, we can construct EAQECCs in the following two theorems. The proof is similar to Theorem 7, so we omit it here.

 $where 1$ 

**Theorem 8** *Let*  $q = 3$ *,*  $m \ge 2$ *,*  $n = \frac{q^{4m}-1}{q^2-1}$  *and*  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ *, where*  $q^2 \equiv 1 \mod 4$ *and* 1 ≤ *t* ≤  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ . *If* C *is a q*<sup>2</sup>*-ary negacyclic code of length n with defining set*  $Z = \bigcup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-2+t} C_{1+2i}$ , then there exsit EAQECCs with parameters  $(f)$  *If*  $\delta \leq \frac{q^{2m}+3}{2}$ ,

$$
\begin{cases} \left[ [n, n-4m \left[ (\delta - \frac{3}{2})(1-q^{-2}) \right] + 4m\varepsilon - mq^2 + 7m, \ge \delta; 4m] \right]_q, if \varepsilon = 2 \text{ or } 3, \\ \left[ [n, n-4m \left( (\delta - \frac{3}{2})(1-q^{-2}) \right] + mq^2 + m, \ge \delta; 4m] \right]_q, if \varepsilon = 4. \end{cases}
$$
  
(ii) If  $\frac{q^{2m}+5}{2} \le \delta \le q^{2m}$ ,  $\left[ [n, n-4m \left( (\delta - \frac{3}{2})(1-q^{-2}) \right] + mq^2 + m, \ge \delta; 4m] \right]_q$ .

**Theorem 9** *Let*  $q = 5$ *,*  $m \ge 2$ *,*  $n = \frac{q^{4m}-1}{q^2-1}$  *and*  $\delta = \frac{q^{2m+1}-q}{q^2-1} + t$ *, where*  $q^2 \equiv 1 \mod 4$ *and* 1 ≤ *t* ≤  $\frac{(q^{2m-1}+1)(q+3)}{2(q+1)}$ . *If* C *is a q*<sup>2</sup>*-ary negacyclic code of length n with defining set*  $Z = \bigcup_{i=0}^{\frac{q^{2m+1}-q}{q^2-1}-2+t} C_{1+2i}$ , then there exsit EAQECCs with parameters  $\left\{ \left[ (n, n - 4m\left[(\delta - \frac{3}{2})(1 - q^{-2})\right] + 4m, \geq \delta; 4m] \right]_q, \text{if } \varepsilon < 6, \right\}$ 

$$
[[n, n-4m\lceil(\delta-\frac{3}{2})(1-q^{-2})]+4\varepsilon m-mq^2+7m, \geq \delta; 4m]]_q, if \ 6 \leq \varepsilon \leq 11.
$$

<span id="page-12-0"></span>**Table 1** new EAQECCs from



#### <span id="page-13-9"></span>**5 Example**

Let  $q = 3$  and  $m = 2$ . Then  $n = \frac{q^{4m}-1}{q^2-1} = 820$ ,  $1 \le t \le \frac{(q^{2m-1}+1)(q+3)}{2(q+1)} = 21$  and  $31 \leq \delta = \frac{q^{2m+1}-q}{q^2-1} + t \leq 51$ . Then from Theorem 8, we can construct some EAQECCs with new parameters in Table [1.](#page-12-0)

# <span id="page-13-10"></span>**6 Conclusion**

In this work, we have constructed a class of EAQECCs from negacyclic BCH codes over the finite fields  $\mathbb{F}_{q^2}$  of length  $n = \frac{q^{4m}-1}{q^2-1}$ , where  $q \ge 3$  is some odd prime power and  $m \ge 2$ . The construction is through cyclotomic cosets and ideal theory. It would be interesting to construct EAQECCs with different lengths from other types of linear codes.

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