

On the Number of Group‑Separable Preference Profles

Alexander Karpov1,[2](http://orcid.org/0000-0002-2693-2389)

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Abstract

The paper studies group-separable preference profles. Such a profle is group-separable if for each subset of alternatives there is a partition in two parts such that each voter prefers each alternative in one part to each alternative in the other part. We develop a parenthesization representation of group-separable domain. The precise formula for the number of group-separable preference profles is obtained. The recursive formula for the number of narcissistic group-separable preference profles is obtained. Such a profle is narcissistic group-separable if it is group-separable and each alternative is preferred the most by exactly one voter.

Keywords Schröder paths · Schröder numbers · Separable permutations · Permutation patterns · Narcissistic preferences

1 Introduction

We consider classical (Arrovian) social choice theory framework. We address *n* preference orders of an *m*-element set $X = \{1, ..., m\}$ that satisfy specific properties. Each preference order represents a binary relation of being better for a voter. The elements of set *X* are called alternatives. A multi-set of *n* preference orders is called a preference profle. We do not assume anonymity. A nontrivial

 \boxtimes Alexander Karpov akarpov@hse.ru

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¹ National Research University Higher School of Economics, 20 Myasnitskaya str., Moscow, Russia 10100

² Institute of Control Science, Russian Academy of Science, 65 Profsoyuznaya Str., Moscow, Russia 117997

permutation of diferent preference orders leads to a nonequivalent preference profle. A permutation of equal preference orders does not change preference profle.

Within the unrestricted set of preference profles, it is impossible to fnd an aggregation rule (social welfare function) that satisfes a set of desirable properties (Arrow [1963;](#page-15-0) Campbell and Kelly [2002](#page-15-1)). In particular, the simple majority rule leads to intransitive social ordering.

One of the possible solutions of this problem is a restricted domain of preference profles. Inada ([1964,](#page-16-0) [1969](#page-16-1)) and Sen [\(1966](#page-16-2)) introduced several types of Condorcet domains (if a society consists of voters whose preferences belongs to a Condorcet domain, then the simple majority rule implies transitive social ordering). The aim of this paper is to analyze the group-separable preferences^{[1](#page-1-0)}. Inada [\(1964](#page-16-0)) defined separability property in the following way. Let A and B be the two groups of alternatives. For each possible preference order, any alternative in A is always preferred to any alternative in B, or any one in B is always preferred to any one in A. Then, we call the set of alternatives separable into two groups. The group-separability property implies a partition of each set of alternatives in two parts such that each voter prefers each alternative in one part to each alternative in the other part. Despite the common partition, a better part difers from voter to voter.

Structured preferences have theoretical importance. There are many examples of computational social choice problems that become easy when there are structured preferences (Elkind et al. [2017;](#page-15-2) Faliszewski et al. [2011](#page-16-3)). Brandt et al. [\(2015](#page-15-3)) presented a detailed study about the computational complexity of social choice problems for single-peaked preferences. In further research papers e.g. Bredereck et al. [\(2016](#page-15-4)), Faliszewski et al. ([2014\)](#page-16-4) a weaker notions of structured preferences (near single-peaked preferences, etc.) were introduced. Elkind et al. ([2012\)](#page-15-5) explored clone structures in preferences as useful tool for studying structured preferences. Our study applies clone structures for group-separable preferences analysis.

Ballester and Haeringer ([2011\)](#page-15-6) demonstrated that group-separable preferences can be defned via forbidden sequences in preference orders. There are two types of restrictions for group-separable preferences. The frst condition is medium restrictedness; that is, for each triple of alternatives, the set of alternatives, which are in the middle of at least one suborder, has cardinality 1 or 2. This restriction was frst discovered by Sen [\(1966](#page-16-2)). The second restriction is the separability of each pair of preference orders. Separability means that one preference order can be obtained from another preference order by means of a separable permutation. A permutation is separable if and only if it contains neither 2413 nor 3142 as a pattern. Separable permutations are counted by the large Schröder numbers (see the diferent proofs in West [1995;](#page-16-5) Stankova [1994;](#page-16-6) Ehrenfeucht et al. [1998\)](#page-15-7). Separable permutations have diferent tree representations (West [1996;](#page-16-7) Bose et al. [1998](#page-15-8); Albert et al. [2015;](#page-15-9) Kitaev [2011](#page-16-8)), which are applicable for related combinatorial problems.

Enumerative combinatorics of preference profles is aimed at fnding the number of preference profles with the desired properties. It is applied for calculating

 1 Group separable preferences are also known as severe disagreement preferences (Van Deemen 2014).

corresponding probabilities. The impartial culture assumption states every possible preference order has the same probability of occurrence. The number of all possible preference profles equals (*m*!)*ⁿ* .

Gehrlein and Fishburn [\(1976](#page-16-10), [1979\)](#page-16-11) obtained recursive and precise formulas for the number of preference profles with a Condorcet winner. Lackner and Lackner [\(2017](#page-16-12)) obtained an asymptotic formula of the number of single-peaked preference profles. Chen and Finnendahl ([2018\)](#page-15-10) found the number of single-peaked narcissistic and single-crossing narcissistic preference profles. Following Chen and Finnendahl [\(2018](#page-15-10)), the narcissistic preference profle is a preference profle with *n* orders and n alternatives such that order i has alternative i in the first place. Our paper defines group-separable narcissistic preference profles and counts them.

The only known result regarding the number of group-separable preference pro-files is Lackner and Lackner's [\(2017](#page-16-12), Corollary 10) inequality:

$$
\#^{GSPP}(n,m) \le m! \left(3 + 2\sqrt{2}\right)^{m(n-1)}
$$

This paper has two main results. First, the precise formula for the number of group-separable preference profles is obtained (Theorem [1\)](#page-9-0). The combinatorial proof is based on the bijection between group-separable domains and Schröder paths without the peaks at level one counted by the small Schröder numbers. Second, the recursive formula for the number of narcissistic group-separable preference profles is obtained (Theorem [2](#page-10-0)).

The structure of the paper is as follows. Section [2](#page-2-0) describes the properties of preference profles and preference domains. Section [3](#page-3-0) presents tree representations of a preference profle. Section [4](#page-5-0) describes the group-separable domains and contains the enumeration result for group-separable preference profles. Section [5](#page-10-1) contains the enumeration result for narcissistic group-separable preference profles. Section [6](#page-10-2) concludes. [Appendix](#page-11-0) contains all proofs.

2 Framework

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Let a finite set $X = \{1, ..., m\}$ be the set of alternatives and a finite set $\mathcal{N} = \{1, ..., n\}$ be the set of voters. Each voter $i \in \mathcal{N}$ has a linear preference order P_i over *X*. The first element is the best alternative. The last element is the worst. Let $\mathcal{L}(X)$ be the set of all possible linear orders on *X*. An *n*-tuple of the preference orders generates the preference profile $P = (P_1, ..., P_n) \in \mathcal{L}(X)^n$. Function $pos(P_i, j) =$ | $\left\{ x \in X | x P_j \right\} \right| + 1$ | indicates the position of alternative *j* in preference profile P_i .

A preference profile P is *group-separable* if for every $A \subseteq X$, $|A| \ge 2$ there exists a partition on two nonempty sets A' , $A \setminus A'$ such that for each $i \in \mathcal{N}$, we have either $aP_i b$ for each $a \in A', b \in A \backslash A'$ or $bP_i a$ for each $a \in A, b \in A \backslash A'$.

A preference profile P is *medium restricted* if for any triple of alternatives *A* \subseteq *X*, $|A| = 3$ we have $\left| \{ y \in A | \exists i \in \mathcal{N} : \max_{j \in A} pos(P_i, j) > pos(P_i, y) > min_{j \in A} pos(P_i, j) \} \right| \le 2$. $\{y \in A | \exists i \in \mathcal{N} : \max_{j \in A} pos(P_i, j) > pos(P_i, y) > min_{j \in A} pos(P_i, j) \} \leq 2.$

A preference profile P is *β-restricted* if there is no $i, j \in \mathcal{N}$ and four alternatives $a, b, c, d \in X$ such that $aP_i bP_i cP_i d$ and $bP_j dP_j aP_j c$. This pattern leads to separable pairs of permutations. The number of separable permutations with *m* elements is the *m* −1*th* large Schröder number (it is the A006318 sequence in the On-line Encyclopedia of Integer Sequences, published electronically at [http://](http://oeis.org) [oeis.org;](http://oeis.org) henceforth, OEIS). Proposition [1](#page-3-1) characterizes group-separable preference profles via forbidden subprofles.

Proposition 1 (Ballester and Haeringer [2011\)](#page-15-6) *A preference profle is group*-*separable if and only if it is medium restricted and β*-*restricted.*

A *domain* $D \in L(X)^d$ is a set of *d* distinct preference orders. We say that a preference profle belongs to a domain if all of the preference orders from this preference profle belong to the domain.

A domain of preferences is a *maximal Condorcet domain* if each preference profle with an odd number of voters from this domain does not yield cycles in the simple majority relation. Condorcet domains are extensively studied in social choice theory (see Monjardet [2009](#page-16-13) for a review of this feld).

A domain of preferences is *symmetric* if it contains the reversed order of each preference order in the domain. A domain, which is maximal Condorcet domain and also a symmetric domain, is a *maximal symmetric Condorcet domain* (for an extensive study of this property, see Danilov and Koshevoy [2013](#page-15-11)).

A domain of preferences is *minimally rich* if for any alternative $x \in X$ there is an order $P \in \mathcal{D}$ such that *P* has *x* as the first alternative. The minimal richness plays an important role in the characterization of the single-peaked domain (Puppe [2018](#page-16-14)).

A domain of preferences is *normal* if it contains the order *123…m* and the opposite to it (Danilov and Koshevoy [2013](#page-15-11)).

3 Clone Sets

A nonempty set $A \subseteq X$ is *a clone set* if and only if when $\forall z \in X \setminus A$, $\forall x, y \in A$, and $∀i ∈ N$, we have either $xP_i z$, $yP_i z$ or $zP_i x$, $z P_i y$. This definition is common in the social choice literature and was considered in Tideman [\(1987\)](#page-16-15), Zavist and Tideman ([1989](#page-16-16)) and Elkind et al. ([2012](#page-15-5)). A clone set is *proper* if it is neither singleton nor set *X*.

A partition $D(X, P)$ of set X into disjointed clone sets is called a *clone partition*. A clone partition $D(X, P)$ is *finer* than another clone partition $D'(X, P)$ if for any $D \in \mathcal{D}(X)$ there exists $D' \in \mathcal{D}'(X, \mathcal{P})$ such that $D = D'$ or $D \subseteq D'$. The inverse relation is called being *coarser*. A clone partition $D(X, P)$ is *coarser* than another clone partition $D'(X, P)$ if for any $D \in D(X, P)$ there exists $D' \in D'(X, P)$ such that $D = D'$ or $D' \subseteq D$.

A clone partition $D(X, P)$ is *trivial* if all its parts are singletons. A clone partition $D(X, P)$ is a *null partition* if the only part is the set *X*. A clone partition

 $D(X, P)$ is *proper* if it is neither trivial nor a null partition. A clone partition $D(X, P)$ is *minimal* if its only coarser clone partition is the null partition.

If there is a two-part clone partition of set *X* with respect to preference profle , then this set and corresponding preference profle is called *reducible*, and otherwise, it is *irreducible*. A clone partition $D(X, P)$ is *ordered* if there is an order of parts such that the union of any number of subsequent parts is a clone set.

We diferentiate *reversible and nonreversible* clone sets. If for each preference order, a clone set reversal does not break other clone sets induced by the preference profle, then this clone set is reversible. Otherwise, it is nonreversible. For example, in the preference profle

(12345, 12354),

 $\{4, 5\}$ is a reversible clone set and $\{1, 2, 3\}$ is a nonreversible clone set. Clone set {3, 4, 5} would be broken after 123 reversal in any preference order.

3.1 Tree Representations of the Preference Profle

Elkind et al. [\(2012](#page-15-5)) proposed a PQ-tree representation of a family of clone sets [PQtrees are formally introduced in Booth and Lueker ([1976\)](#page-15-12)]. Here, we defne this representation in our terms. A *PQ*-*tree* is an ordered tree with two types of nodes. If a node is of type P, then its children can be permuted arbitrarily. If a node is of type Q, then the order of its children can be reversed.

A *clone decomposition tree* $G(P)$ is a PO-tree. Set *X* is the root of the tree. If *X* is irreducible, then the node *X* is of type P and parts of the minimal clone partition are the children. If X is reducible, then the node X is of type Q and parts of the finest ordered clone partition are the children. Each non-singleton child becomes a root vertex and generates children from its own minimal clone partition (for irreducible set) or its own fnest ordered clone partition. Tree leaves are singletons.

Proposition 2 *Each preference profle has a unique clone decomposition tree.*

The proof for Proposition [2](#page-4-0) and subsequent propositions are given in the [Appen](#page-11-0)[dix.](#page-11-0) Similar result is known for tournament decomposition trees (Laslier [1997,](#page-16-17) Brandt et al. [2011](#page-15-13)). Tournament decomposition trees also exploit reducibility property.

For a node with two children, the type of node does not matter. In our framework, all such nodes are Q-nodes. All P-nodes have at least 3 children.

If a Q-node has more than two children, then for any three parts *A*, *B*, and *C*, one part is between the other parts according to the decomposition ordering. If *B* is *between* A and C, then for any $a \in A, b \in B$, and $c \in C$ and for any voter $i \in \mathcal{N}$, we have either $a_i^b P_i^c$, or $c_i^b P_i^d$. The betweenness relation reflects the structure of alternatives even without introducing any geometric space. We call the clone decomposition tree with only Q-nodes the *Q*-*tree*. From the reducibility of each clone set, we obtain Proposition [3](#page-5-1).

Proposition 3 *A preference profle is group*-*separable if and only if its clone decomposition tree is a Q*-*tree.*

Each node in a Q-tree represents reversible clone set and each nontrivial reversible clone set corresponds to node in a Q-tree representation.

It is convenient to represent Q-trees and corresponding reversible clone sets as a parenthesization for set *X*. Each pair of parentheses corresponds to a reversible clone set. For example, parenthesization $((12)3)$ means that $\{1, 2\}$ are $\{1, 2, 3\}$ are reversible clone sets. $\{1, 2\}, \{3\}$ is the finest ordered partition of set $\{1, 2, 3\}$ with respect to preference profile P. $\{1\}$, $\{2\}$ is the finest ordered partition of set $\{1, 2\}$ with respect to preference profile P . The next section utilizes parenthesization representation for group-separable domains.

4 Group‑Separable Domains

A domain of preferences is group-separable if all of the preference profles from this domain are group-separable. After having a labeled clone decomposition *Q*-tree and corresponding parenthesization, we define a *parenthesization domain* $D(G)$ such that a preference profle containing all of the preference orders from this domain generates a parenthesization \mathcal{G} : $\mathcal{G} = \mathcal{G}(D(\mathcal{G}))$, and if we add any other preference order, then a parenthesization will be different. Any parenthesization $\mathcal G$ induces a parenthesization domain.

Each parenthesization domain is group-separable. Each group-separable preference profle belongs to corresponding parenthesization domain.

A parenthesization with one parentheses generates a domain with two reversed orders. Any additional parentheses doubles the number of preference orders in a domain. A parenthesization domain has 2^f preference orders, where f is the number of parentheses.

A maximal group-*separable domain* is a group-separable domain such that any other domain of a higher cardinality, which includes this domain, is not a groupseparable domain. Parenthesizations that generate maximal group-separable domain are called *binary parenthesizations.* These parenthesizations have *m*−1 parentheses. Thus, each maximal group-separable domain contains 2*^m*−1 orders. Each group-separable preference profle belongs to a maximal group-separable domain.

The number of labeled binary parenthesizations of a word with m letters equals the number of maximal group-separable domains. From Murtagh ([1984\)](#page-16-18), the number of such parenthesizations and number of maximal group-separable domains is equal to

$$
\#^{MGSD}(m) = \prod_{i=2}^{m} (2i - 3),
$$

with $\#^{MGSD}(1) = 1$.

Theorem 3 from Danilov and Koshevoy ([2013\)](#page-15-11) states that each normal symmetric Condorcet domain of size 2*^m*−1 has the form of some parenthesization of 12…n (*normal parenthesization*). Thus, we state the following characterization a maximal group-separable domain.

Proposition 4 (Danilov and Koshevoy [2013](#page-15-11)) *A domain is a maximal symmetric Condorcet domain of size* 2*^m*−1 *if and only if it is a maximal group*-*separable domain.*

All maximal group-separable domains also satisfy minimal richness. Maximal symmetric Condorcet domains of a size smaller than 2*^m*−1 may not be group-separable domains. For example, the domain constructed in Proposition 5 of Danilov and Koshevoy ([2013\)](#page-15-11)

(12345, 54321, 24153, 35142)

is a maximal symmetric Condorcet domain, but it is neither group-separable nor minimally rich. Dittrich ([2018\)](#page-15-14) found minimally rich maximal symmetric Condorcet domain, which is not group-separable:

(54321, 45321, 25413, 24513, 31542, 31452, 12354, 12345).

4.1 Number of Group‑Separable Preference Profles

Here, we construct two bijections. The first one is a bijection μ between the normal parenthesization domains and Schröder paths of semilength *m*−1 (i.e., the lattice paths from (0, 0) to $(2m-2, 0)$ with steps H = (2, 0), U = (1, 1) and D = (1, −1) that do not go below the *x*-axis) with no peaks at level one. The second one is a bijection *φ* between the normal parenthesization domains and the sets of all proper reversible clone sets (reversible clone sets are defned below in this section). All of these objects are counted by the small Schröder numbers (it is the A001003 sequence in OEIS).

Schröder path UD has peak at level one, Schröder path UUDD has peak at level two. By counting the number of peaks at diferent levels it is convenient to categorize Schröder paths.

Let us define a function μ from normal parenthesization domains to Schröder paths of semilength m−1 without peaks at level 1.

- 1. *m*=1. Normal parenthesization domain generates an empty Schröder path.
- 2. *m*=2. Normal parenthesization domain generates the Schröder path H.
- 3. For $m \ge 3$, function μ is defined recursively. There are three types of normal parenthesizations.

Type A

$$
(1(2...m)) \leftrightarrow U\mu(2...m)D.
$$

Type A parenthesization is a unique parenthesization type, which has "(1", where 1 is the frst element of parenthesized word. Type A Schröder path is a unique Schröder path type, which has no peaks at level 2 and has D at the end.

Type B

$$
((1\ldots m-1)m)\leftrightarrow \mu(1\ldots m-1)H.
$$

Type B parenthesization is a unique parenthesization type, which has "m)", where m is the last element of parenthesized word. Type B Schröder path is a unique Schröder path type, which has H at the end.

Type C

$$
((1...k_1)(k_1+1...k_2)...(k_{p-2}+1...k_{p-1})(k_{p-1}+1...m)) \leftrightarrow \mu(1...k_1)
$$

$$
U\mu(k_1+1...k_2)UD... \mu(k_{p-2}+1...k_{p-1})UD\mu(k_{p-1}+1...m)D,
$$

where $(k_{i-1} + 1 \ldots k_i)$ is the *ith* reversible clone (it can be a singleton).

Type C parenthesization is a unique parenthesization type, which has neither (1) , where 1 is the first element of parenthesized word, nor (m) , where m is the last element of parenthesized word. Type C Schröder path is a unique Schröder path type, which has an inner common point with the x-axis and has no *H* at the end.

Because all types have unique features, function μ is bijective. The image of function µ is the set of all Schröder paths of semilength m−1 without peaks at level 1. Examples [1](#page-7-0) and 2 clarify the bijection.

Example 1 $\mu((123)4(5(67))) = \mu(123)U\mu(4)UD\mu(5(67))D$ (Type C). Having $\mu(123) = UUDD$ (Type C), $\mu(5(67)) = U\mu(67)D = UHD$ (Type A), we obtain

$$
\mu((123)4(5(67))) = UUDDUUDUHDD.
$$

Inverse problem. Because *UUDD*UUD*UHD*D has an inner common point with the x-axis and has no H at the end, it is type C. Thus, we have

 $\mu^{-1}(UUDDUUDDUHDD) = \mu^{-1}(UUDD)U\mu^{-1}(UDD)\mu^{-1}(UHD)D = (123)4(5(67)).$

Let Ω be the set of all proper reversible clone sets of a preference profile with all of the preference orders from a normal parenthesization domain. Each set between a pair of parentheses is a proper reversible clone set. One proper reversible clone set can be a subset of another proper reversible clone set, but it is impossible to have nonempty intersection without inclusion. By introducing parentheses, we obtain a correct parenthesization. Thus, we construct a bijection φ between normal parenthesization domains and sets of all proper reversible clone sets of normal parenthesization domains. Example 2 in Table [1](#page-8-0) clarifes the constructed bijections.

Proposition 5 *For m*≥*2, the sum of the number of proper reversible clone sets and number of peaks in the corresponding Schröder path equals m*−2.

Each maximal parenthesization domain contains *m*−2 proper reversible clone sets. There are $\binom{n-2}{i}$ \setminus parenthesization domains with *i* proper reversible clone sets (*m*−*i*−2 peaks in the corresponding Schröder path) that belong to the maximal parenthesization domain.

Table 1 Example 2

The case of $m=4$. There are 11 normal parenthesizations that generate normal domains

Chen et al. [\(2006](#page-15-15), Theorem 3.3) found that the number of Schröder paths with no peaks at level one contained exactly *k* peaks at other levels. From Proposition [5,](#page-7-1) bijections μ we know, that these Schröder paths correspond normal parenthesization domains.

From the bijection μ , Chen et al.'s ([2006](#page-15-15)) Theorem 3.3 and the sequence A126216 from OEIS, we directly have Proposition [6.](#page-8-1)

Proposition 6 *The number of normal parenthesization domains with 2m*–*k*−1 *preference orders (k is the number of peaks in the corresponding Schröder path) is equal to*

$$
\#^{NPD}(k,m) = \frac{1}{m-1} \binom{m-1}{k} \binom{2m-2-k}{m}.
$$

The total number of normal parenthesization domains can be obtained by summation $\#^{NPD}(k,m)$ over k from 0 to $m-2$. Chen et al. [\(2006](#page-15-15)) showed, that this number is the *m*−1*th* small Schröder number (A001003 in OEIS).

Proposition 7 *The number of normal parenthesization domains equals the m*−1*th small Schröder number (A001003 in OEIS).*

The small Schröder numbers have no brief explicit formula. Stanley [\(1997](#page-16-19)) and Deutsch [\(2001](#page-15-16)) discussed in details several interpretations and related bijections of these numbers. Some paths, trees, parenthesizations, triangulations, etc. are counted by the small Schröder numbers.

Substituting $k=0$ in formula from proposition [6](#page-8-1) we obtain the number of normal maximal parenthesization domains.

Proposition 8 *The number of normal maximal parenthesization domains equals the m*−1*th Catalan number (A000108 in OEIS)*

$$
\#^{NMPD}(m) = \frac{1}{m-1} \binom{2m-2}{m}.
$$

Proposition [8](#page-9-1) also immediately follows from Danilov and Koshevoy ([2013,](#page-15-11) p 191). Proposition [6](#page-8-1) also leads to the main result.

Theorem 1 *The number of group*-*separable preference profles is equal to*

$$
\#^{GSPP}(n,m) = m! \sum_{k=0}^{m-2} \frac{(-1)^k}{m-1} {m-1 \choose k} {2m-2-k \choose m} 2^{(m-k-1)(n-1)}.
$$

For $m=3$, $m=4$, and $m=5$, we have the following formulas:

$$
#^{GSPP}(n,3) = 6(2 \cdot 4^{n-1} - 2^{n-1});
$$

\n
$$
#^{GSPP}(n,4) = 24 \cdot (5 \cdot 8^{n-1} - 5 \cdot 4^{n-1} + 2^{n-1});
$$

\n
$$
#^{GSPP}(n,5) = 120 \cdot (14 \cdot 16^{n-1} - 21 \cdot 8^{n-1} + 9 \cdot 4^{n-1} - 2^{n-1}).
$$

These formulas have a polynomial form with coefficients from the sequence A126216 from OEIS.

For $m=3$, the number of group-separable preference profiles is equal to the number of single-peaked preference profles. By using the asymptotic formula for the number of single-peaked preference profles from Lackner and Lackner [\(2017](#page-16-12)) for *m*≥2 and *n*≥2, we obtain the limits

$$
\lim_{n\to\infty}\frac{\#^{GSPP}(n,m)}{\#^{SPPP}(n,m)}=2\frac{(2m-3)!!}{m!};\quad \lim_{m\to\infty}\frac{\#^{GSPP}(n,m)}{\#^{SPPP}(n,m)}=\infty.
$$

Both the maximal group-separable and maximal single-peaked domains contain 2^{*m*−1} preference orders. The number of group-separable preference profiles is higher than the number of single-peaked preference profles, but for small numbers of alternatives, this diference is not too large.

5 Number of Narcissistic Group‑Separable Preference Profles

Following Chen and Finnendahl ([2018\)](#page-15-10), the *narcissistic preference profle* is a preference profle with *n* orders and *n* alternatives such that order *i* has alternative *i* in the frst place.

Theorem 2 *The following recursive formula for the number of narcissistic groupseparable preference profles holds:*

$$
\#^{NGSPP}(n) = \sum_{k=1}^{n-1} {n-1 \choose k-1} 2^{2nk-2k^2-n} \#^{NGSPP}(k) \#^{NGSPP}(n-k),
$$

with $\#^{NGSPP}(1) = 1$.

The frst terms of #*NGSPP*(*n*) are 1, 1, 6, 144, 13,440, 4,976,640, 7,390,494,720, and so on (for $n=1, 2, 3,...$). Starting from $n=3$, the number of narcissistic group separable preference profles is higher than the number of narcissistic single peaked preference profles and narcissistic single crossing preference profles.

6 Conclusion

This paper shows that group-separable domains are a subclass of maximal symmetric Condorcet domains and that all maximal symmetric Condorcet domains of the largest cardinality are group-separable. All of the presented examples of the maximal symmetric Condorcet domain have a cardinality of 2^k for some $k \in \mathbb{N}$. Whether this is true for all maximal symmetric Condorcet domains is an open question.

This paper connects group-separable preference domains, ordered trees, and Schröder paths and thus provides new bijections. These bijections lead to a new combinatorial result of the enumeration of group-separable preference profles. This result fundamentally solves the enumeration problem frst studied in Lackner and Lackner [\(2017](#page-16-12)).

Appendix

Proof of Proposition 2 From Theorem 3.12 in Elkind et al. ([2012\)](#page-15-5) we have that the family of all clone sets satisfy axioms $A1-A5$. Let $\mathcal F$ be a family of clone sets. Axioms A1-A5 are the following.

A1. $\{x\} \in \mathcal{F}$ for any $x \in X$, $\emptyset \notin \mathcal{F}$, and $X \in \mathcal{F}$. A2. If $A, B \in \mathcal{F}$ and $A \cap B \neq \emptyset$, then $A \cup B \in \mathcal{F}$ and $A \cap B \in \mathcal{F}$. A3. If $A, B \in \mathcal{F}, A \cap B \neq \emptyset, A \setminus B \neq \emptyset, B \setminus A \neq \emptyset$, then $A \setminus B \in \mathcal{F}$ and $B \setminus A \in \mathcal{F}$. A4. For each $C \in \mathcal{F}$ there are at most two sets (minimal supersets) $Z \in \mathcal{F}$ with the following properties: $C \subseteq Z$, $C \neq Z$ and there is no set $Y \in \mathcal{F}$ such that $Y \neq Z$ and $C \subseteq Y \subseteq Z$.

A5. There is no set family A_0, \ldots, A_{k-1} with $k \geq 3$ such that for all $i=0, \ldots, k-1$

- (1) *Ai* ∩ *Ai*+¹ ≠ �, *Ai* �*Ai*+¹ ≠ �, *Ai*+¹�*Ai* ≠ �;
- (2) $A_{i-1} \cap A_i \cap A_{i+1} = \emptyset;$
- (3) *Ai* ∈ *Ai*−¹ ∪ *Ai*+1;

where all indices are computed modulo *k*.

(i) Suppose that the reducible preference profle has two diferent fnest ordered partitions, $\mathcal{D}^{A}(X,\mathcal{P})$ and $\mathcal{D}^{B}(X,\mathcal{P})$, with ordered clone partitions $A_1, ..., A_a$ and $B_1, ...,$ *B_b*, respectively. Thus, there is A_i and $l \in \{1, ..., b\}$ such that $A_i \cap (\bigcup_{j=1}^l B_j) \neq \emptyset$ and $A_i \cap (\bigcup_{j=l+1}^b B_j) \neq \emptyset$. If $(\bigcup_{j=1}^{i-1} A_j) \cap (\bigcup_{j=1}^l B_j) \neq \emptyset$ and $\left(\bigcup_{j=1}^{i-1} A_j\right) \bigcap \left(\bigcup_{j=l+1}^{b} B_j\right) \neq \emptyset$, then sets A_i , $\left(\bigcup_{j=1}^{l} B_j\right)$, $\left(\bigcup_{j=l+1}^{b} B_j\right)$, and $\left(\bigcup_{j=1}^{i-1} A_j\right)$ contradict with axiom A Thus, we have either $(\bigcup_{j=1}^{i-1} A_j) \cap (\bigcup_{j=1}^{l} B_j) = \emptyset$ or $(\bigcup_{j=1}^{i-1} A_j) \cap (\bigcup_{j=i+1}^{b} B_j) = \emptyset$. By applying the same argument for $(\bigcup_{j=i+1}^{a} A_j)$, we have either $\left(\bigcup_{j=i+1}^a A_j\right) \cap \left(\bigcup_{j=1}^l B_j\right) = \emptyset$ or $\left(\bigcup_{j=i+1}^a A_j\right) \cap \left(\bigcup_{j=l+1}^b B_j\right) = \emptyset$.

Case *i*=1. We have either $\left(\bigcup_{j=1}^{l} B_j\right) \bigcap A_2 \neq \emptyset$ or $\left(\bigcup_{j=l+1}^{b} B_j\right) \bigcap A_2 \neq \emptyset$. If $\left(\bigcup_{j=1}^{l} B_j\right) \bigcap A_2 \neq \emptyset$; then, $\left(A_1 \bigcup A_2\right) \cap \left(\left(\bigcup_{j=1}^{l} B_j\right) \bigcup A_2\right)$ is a clone set and an ordered partition C_1, \ldots, C_{a+1} such that $C_1 = A_1 \cap \left(\bigcup_{j=l+1}^b B_j\right), C_2 = A_1 \cap \left(\bigcup_{j=1}^l B_j\right)$ and $C_{j+1} = A_j$ for $j \in \{2, ..., a\}$ is a clone ordered partition. If $\left(\bigcup_{j=l+1}^b B_j\right) \cap A_2 \neq \emptyset$, then $(A_1 \cup A_2) \cap \left(\left(\bigcup_{j=l+1}^b B_j \right) \cup A_2 \right)$ is a clone set and an ordered partition C_1, \ldots, C_{a+1} such that $C_1 = A_1 \cap \left(\bigcup_{j=1}^l B_j \right), C_2 = A_1 \cap \left(\bigcup_{j=l+1}^b B_j \right)$, and $C_{j+1} = A_j$ for $j \in \{2, ..., a\}$ is a clone ordered partition. These partitions are finer than $\mathcal{D}^{A}(X,\mathcal{P}).$

Case $i = a$ is similar.

 $\text{Case } a > i > 1. \text{ If } \left(\bigcup_{j=1}^{i-1} A_j \right) \bigcap \left(\bigcup_{j=1}^{i} B_j \right) = \emptyset \text{ and } \left(\bigcup_{j=i+1}^{a} A_j \right) \bigcap \left(\bigcup_{j=1}^{i} B_j \right) = \emptyset \text{, then }$ clone set $A_i \cap (\bigcup_{j=l+1}^b B_i)$ has three minimal supersets: the subsets of $(\bigcup_{j=1}^{i} A_j) \cap (\bigcup_{j=l+1}^{b} B_j)$, $(\bigcup_{j=l+1}^{a} A_j) \cap (\bigcup_{j=l+1}^{b} B_j)$, and A_i . It contradicts with axiom A4. The same argument exists for the case of $(\bigcup_{j=1}^{i-1} A_j) \cap (\bigcup_{j=l+1}^{b} B_j) = \emptyset$ and $\left(\bigcup_{j=i+1}^a A_j\right) \cap \left(\bigcup_{j=l+1}^b B_j\right) = \emptyset$. If $\left(\bigcup_{j=1}^{i-1} A_j\right) \cap \left(\bigcup_{j=1}^l B_j\right) = \emptyset$. and $\left(\bigcup_{j=i+1}^a A_j\right) \cap \left(\bigcup_{j=l+1}^b B_j\right) = \emptyset$, then an ordered partition C_1, \ldots, C_{a+1} such that $C_j = A_j$ for $j \in \{1, ..., i-1\}$, $C_i = A_i \cap \left(\bigcup_{j=l+1}^b B_j\right)$, $C_{i+1} = A_i \cap \left(\bigcup_{j=1}^l B_j\right)$, and $C_{j+1} = A_j$ for $j \in \{i+1,\ldots,a\}$ is a clone ordered partition. If $\left(\bigcup_{j=1}^{i-1} A_j\right) \cap \left(\bigcup_{j=i+1}^{b} B_j\right) = \emptyset$ and $\left(\bigcup_{j=i+1}^a A_j\right) \cap \left(\bigcup_{j=1}^l B_j\right) = \emptyset$, then an ordered partition C_1, \ldots, C_{a+1} such that $C_j = A_j$ for $j \in \{1, ..., i-1\}$, $C_i = A_i \cap \left(\bigcup_{j=1}^l B_j\right)$, and $C_{j+1} = A_j$ for $j \in \{i+1, ..., a\}$ is a clone ordered partition. These partitions are finer than $D^A(X, \mathcal{P})$.

(ii) Suppose that an irreducible preference profle has two diferent minimal partitions $\mathcal{D}^A(X, \mathcal{P})$ and $\mathcal{D}^B(X, \mathcal{P})$ with clone partitions A_1, \ldots, A_a and B_1, \ldots, B_b , where $a \leq b$.

Let us consider the bipartite graph with vertices A_1 , …, A_a and B_1 , …, B_b . If $A_i \cap B_j \neq \emptyset$, then we have an edge between A_i and B_j . Because of Axiom A5, this graph is a bipartite graph without cycles, Because of Axiom A4, each vertex is connected with, at most, two vertices with a degree of two or more. If this graph is not connected, then clone partitions A_1, \ldots, A_a and B_1, \ldots, B_b are not minimal. Thus, there is a chain, that connects vertices A_1 , …, A_a and a chain that connects $a-1$ *A* type vertices. By axiom A2, this chain corresponds to a clone set. Thus, $\mathcal{D}^{A}(X, \mathcal{P})$ is reducible. \Box

Proof of Proposition 3 By proposition [2](#page-4-0) each preference profile has unique clone decomposition tree. Suppose that a group-separable preference profle has P-node in a clone decomposition tree. The set associated with this node is irreducible. It contradicts with defnition of group-separability. Thus, there is no P-node.

If a clone decomposition tree of a preference profle is a Q-tree, then for each subset of alternatives $A \subseteq X$, $|A| \ge 2$, there is a minimal reversible clone superset (it is a clone set which contains set *A*, but there is no other smaller reversible clone set, which also contains set *A*). Tree representation presents the fnest ordered partition of this set. Because we fnd minimal reversible clone superset, the set *A* belongs to at least two parts of this partition. Thus, if for every $A \subseteq X$, $|A| \ge 2$ there exists a partition on two nonempty sets *A'*, *A**A'* such that for each $i \in \mathcal{N}$, we have either $aP_i b$ for each $a \in A'$, $b \in A \backslash A'$ or $bP_i a$ for each $a \in A, b \in A \backslash A'$. \square

Proof of Proposition 5 Let *a* be a function, which possess the value of the number of proper reversible clone sets, and *b* be a function, which possess the value of the number of peaks in the corresponding Schröder path.

According to bijection μ for $m=2$, we have Schröder path H and the sum equals to 0. Suppose that for all \hat{m} < *m* we have $a+b=m-2$. Let us prove for *m*.

For $m \geq 3$, function μ is defined recursively. For type A, we have

$$
a(1(2...m)) = a(2...m) + 1;
$$

$$
b(U\mu(2...m)D) = b(\mu(2...m)).
$$

From induction assumption we have $a(2 \dots m) + b(\mu(2 \dots m)) = m - 3$. Thus, we $obtain a(1(2...m)) + b(U\mu(2...m)D) = m - 2.$

For type B, we have

$$
a((1 \dots m - 1)m) = a(1 \dots m - 1) + 1;
$$

$$
b(\mu(1 \dots m - 1)H) = b(\mu(1 \dots m - 1));
$$

From induction assumption we have $a(1 \dots m-1) + b(\mu(1 \dots m-1)) = m-3$. Thus, we obtain $a((1 \dots m-1)m) + b(\mu(1 \dots m-1)H) = m-2$.

For type C, we have

$$
a((1...k_1)(k_1+1...k_2)...(k_{p-2}+1...k_{p-1})(k_{p-1}+1...m))
$$

= $a(1...k_1) + a(k_1+1...k_2) + ... + a(k_{p-1}+1...m) + p$ - #singletons;
 $b(\mu(1...k_1)U\mu(k_1+1...k_2)UD... \mu(k_{p-2}+1...k_{p-1})UD\mu(k_{p-1}+1...m)D)$
= $b(\mu(1...k_1)) + b(\mu(k_1+1...k_2)) + ... + b((k_{p-1}+1...m)) + p - 2;$

From induction assumption we have for nonsingletons $a(k_i + 1 ... k_{i+1}) + b(\mu(k_i + 1 ... k_{i+1})) = k_{i+1} - k_i - 2$, and for singletons $a(k_i + 1 ... k_{i+1}) + b(\mu(k_i + 1 ... k_{i+1})) = k_{i+1} - k_i - 1 = 0$. Thus, we obtain

$$
a((1...k_1)(k_1+1...k_2)...(k_{p-2}+1...k_{p-1})(k_{p-1}+1...m))
$$

+ $b(\mu(1...k_1)U\mu(k_1+1...k_2)UD... \mu(k_{p-2}+1...k_{p-1})UD\mu(k_{p-1}+1...m)D)$
= $m-2p+p+p-2=m-2$.

□

Proof of Theorem 1 Each group-separable preference profle belongs to a parenthesization domain. Each parenthesization domain is group-separable. The number of group-separable preference profles with the frst order 1, 2 … *m* is *m*! times less than $\#^{GSPP}(n, m)$. We call this preference profile a normal preference profile.

The number of preference orders in a normal parenthesization domain with *r* proper reversible clone sets is equal to $2^{(r+1)}$. The number of normal preference profiles in a normal parenthesization domain with *r* proper reversible clone sets is equal to $2^{(r+1)}$ (*ⁿ*−1). From inclusion–exclusion principle the number of normal preference profles in two normal parenthesization domains with $r_1 = |\Omega_1|, r_2 = |\Omega_2|, r_3 = |\Omega_1 \cap \Omega_2|$ is equal to $2^{(r_1+1)(n-1)} + 2^{(r_2+1)(n-1)} - 2^{(r_3+1)(n-1)}$ equal to $2^{(r_1+1)(n-1)} + 2^{(r_2+1)(n-1)} - 2^{(r_3+1)(n-1)}$.

Applying inclusion–exclusion principle we obtain the number of normal preference profles that belong to a parenthesization domain with *k* peaks in the corresponding

Schröder path and *m*−2−*k* proper reversible clone sets, which do not belong to any other normal parenthesization domains with the lower number preference orders

$$
2^{(m-k-1)(n-1)} + \sum_{j=1}^{m-2-k} (-1)^j {m-2-k \choose j} 2^{(m-k-j-1)(n-1)}.
$$

The parenthesization domain without proper reversible clone sets does not contain any other parenthesization domain. The number of normal preference profles in this domain equals $2^{(n-1)}$. By knowing the number of normal parenthesization domains from proposition 6 , we obtain

$$
\begin{aligned} \#^{GSPP}(n,m) &= m!2^{(n-1)} + m! \sum_{k=0}^{m-3} \frac{1}{m-1} \binom{m-1}{k} \binom{2m-2-k}{m} \\ &\times \left[2^{(m-k-1)(n-1)} + \sum_{j=1}^{m-2-k} (-1)^j \binom{m-2-k}{j} 2^{(m-k-j-1)(n-1)} \right]. \end{aligned}
$$

By introducing $a = k + j$ and modifying, we have

$$
\begin{split} \#^{GSPP}(n,m) &= m! \left(\sum_{k=0}^{m-2} \frac{1}{m-1} \binom{m-1}{k} \binom{2m-2-k}{m} 2^{(m-k-1)(n-1)} \right. \\ &\quad \left. + \sum_{a=1}^{m-2} \frac{2^{(m-a-1)(n-1)}(m-2)!}{m(m-2-a)!a!} \sum_{k=0}^{a-1} (-1)^{(a-k)} \binom{a}{k} \binom{2m-k-2}{m-k-1} \right). \end{split}
$$

By using the following combinatorial identity

$$
\sum_{k=0}^{a} (-1)^{k} {a \choose k} {2(m-1)-k \choose (m-1)-k} = {2(m-1)-a \choose (m-1)-a},
$$

(both sides are the number of ways of selecting *m*−1 elements from the 2(*m*−1) element set such that *a* particular elements are not included in the selection), we obtain the result

$$
\#^{GSPP}(n,m) = m! \sum_{k=0}^{m-2} \frac{(-1)^k}{m-1} {m-1 \choose k} {2m-2-k \choose m} 2^{(m-k-1)(n-1)}.
$$

□

Proof of Theorem 2 Each narcissistic group-separable preference profile generates a parenthesization, which corresponds to a maximal parenthesization domain. Each maximal parenthesization domain has 2*ⁿ*−1 orders.

Each maximal parenthesization partitions the set of alternatives into two parts (two clone sets). Let, a part with alternative 1 be a set A , $|A| = k$. The number of ways to choose *A* with *k* elements equals to $\binom{n-1}{k-1}$ $\left(\frac{1}{k^{NGSPP}}(k) \right)$ is the number of ways to defne narcissistic group-separable preference subprofle with *k* voters and the frst alternatives from *A*. Each such preference profle belongs to a maximal parenthesization domain with 2^{$k-1$} preference orders. $\#^{NGSPP}(n-k)$ is the number of ways to defne narcissistic group-separable preference subprofle with *n*-*k* voters and alternatives from *X**A X*∖*A*. Each such preference profle belongs to a maximal parenthesization domain with 2*ⁿ*−*^k*−1 preference orders. For each pair of narcissistic group-separable preference subprofiles there are $(2^{n-k-1})^k (2^{k-1})^{n-k}$ ways to define remainder preference profle. Thus, we have

$$
\#^{NGSPP}(n) = \sum_{k=1}^{n-1} {n-1 \choose k-1} (2^{n-k-1})^k (2^{k-1})^{n-k} \#^{NGSPP}(k) \#^{NGSPP}(n-k)
$$

Through simplification we obtain the result. \square

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