

On some constacyclic codes over $\mathbb{Z}_4[u] / \langle u^2 - 1 \rangle$, their \mathbb{Z}_4 images, and new codes

Nuh Aydin¹ · Yasemin Cengellenmis² · Abdullah Dertli³

Received: 14 February 2017 / Revised: 8 July 2017 / Accepted: 11 July 2017 / Published online: 19 July 2017 © Springer Science+Business Media, LLC 2017

Abstract In this paper, we study λ -constacyclic codes over the ring $R = \mathbb{Z}_4 + u\mathbb{Z}_4$ where $u^2 = 1$, for $\lambda = 3+2u$ and 2+3u. Two new Gray maps from R to \mathbb{Z}_4^3 are defined with the goal of obtaining new linear codes over \mathbb{Z}_4 . The Gray images of λ -constacyclic codes over R are determined. We then conducted a computer search and obtained many λ -constacyclic codes over R whose \mathbb{Z}_4 -images have better parameters than currently best-known linear codes over \mathbb{Z}_4 .

Keywords Constacyclic codes \cdot Codes over rings \cdot Quaternary codes \cdot Gray map \cdot Best known codes

Mathematics Subject Classification 94B15 · 94B60

1 Introduction

Codes over finite rings are a topic of much research in algebraic coding theory for over two decades [1,3-6,9,10,12,14,15,17,20,21]. Among all finite rings, the ring \mathbb{Z}_4 of integers modulo 4 has a special place in coding theory. In addition to the original discovery in [10]

Communicated by T. Helleseth.

Nuh Aydin aydinn@kenyon.edu

> Yasemin Cengellenmis ycengellenmis@gmail.com

Abdullah Dertli abdullah.dertli@gmail.com

¹ Department of Mathematics and Statistics, Kenyon College, Gambier, OH 43022, USA

² Mathematics Department, Faculty of Arts and Sciences, Trakya University, Edirne, Turkey

³ Mathematics Department, Faculty of Arts and Sciences, Ondokuz Mayis University, Samsun, Turkey

that showed that certain binary non-linear codes can be obtained as Gray images of \mathbb{Z}_4 -linear codes, codes over \mathbb{Z}_4 are known to be related to lattices [19], combinatorial designs [7], and low correlation sequences [18]. As a result of increased research activity on codes over \mathbb{Z}_4 , a database of \mathbb{Z}_4 -codes was introduced in [2], which is available online [8]. The database is updated as researchers discover new codes with better parameters. When a new code is added to the database, the previous codes are not removed. Several recent papers contributed new codes to the database such as [14,20].

More recently, the rings of order 16 that are extensions of \mathbb{Z}_4 have received special attention. The interest in these rings increased after the introduction of linear and cyclic codes over the ring $\mathbb{Z}_4 + u\mathbb{Z}_4$, $u^2 = 0$ in [20,21]. In [1], Ashraf and Mohammed studied (1 + 2u)-constacyclic codes over this ring.

In [3], cyclic, quasi cyclic (QC), and constacyclic codes over $\mathbb{Z}_4 + w\mathbb{Z}_4$, $w^2 = 2$ have been studied. Two Gray maps are defined and Gray images of codes over this ring have been described. Moreover, new linear codes over \mathbb{Z}_4 are obtained using these Gray maps.

Due to the increased interest in codes over rings of order 16, these rings are classified in [11]. One of the special properties of rings of order 16 is the fact that the smallest local, finite, Frobenius, commutative, non-chain ring has order 16 [11]. In this work, we consider one of the rings of order 16, $\mathbb{Z}_4[u] / \langle u^2 - 1 \rangle$ which can also be described as $R = \mathbb{Z}_4 + u\mathbb{Z}_4$ with $u^2 = 1$. This ring is isomorphic to the ring $\mathbb{Z}_4[u] / \langle u^2 + 2u \rangle$ via $u \mapsto u + 1$.

In [14], Ozen et al. studied cyclic codes and (2 + u)-constacyclic codes over R. They determined the form of generators of cyclic codes over this ring. They proved that the \mathbb{Z}_4 -image of a (2 + u)-constacyclic code of odd length over R is a cyclic code over \mathbb{Z}_4 . They presented many examples of good codes over \mathbb{Z}_4 obtained as images of cyclic codes over R.

In [17], Shi et al. defined a new Gray map from $\mathbb{Z}_4 + u\mathbb{Z}_4$ to \mathbb{Z}_4^2 and studied (1 + 2u)constacyclic codes over R. They investigated properties of (1 + 2u)-constacyclic codes over R. They proved that the Gray images of (1 + 2u)-constacyclic codes of length n over R are
cyclic codes of length 2n over \mathbb{Z}_4 . They presented many examples of (1 + 2u)-constacyclic
codes R whose \mathbb{Z}_4 -images have good parameters.

Inspired by these works, we introduce two new Gray maps from R to \mathbb{Z}_4^3 . We are particularly interested in finding new codes with better parameters than the currently best known linear codes over \mathbb{Z}_4 . In fact, this was our main purpose and motivation for this work. We have been able to obtain such codes using each map. Moreover, we give a description of Gray images of cyclic and constacyclic codes over R.

2 Preliminaries

In [4,14,17], the authors considered the finite commutative ring $R = \mathbb{Z}_4 + u\mathbb{Z}_4$, where $u^2 = 1$. Clearly, $R \cong \mathbb{Z}_4[u]/\langle u^2 - 1 \rangle$. Its units are given by $\{1, 3, u, 3u, 2 + u, 1 + 2u, 3 + 2u, 2 + 3u\}$. There are 7 ideals in this ring of characteristic 4 given by $\{\langle 0 \rangle, \langle 2u \rangle, \langle 1 + u \rangle, \langle 3 + 3u \rangle, \langle 2 + 2u \rangle, \langle 2u, 1 + u \rangle, R\}$. It is a local ring with the maximal ideal $\langle 2u, 1 + u \rangle$.

In [14,17], constacyclic codes over *R* and their Gray images are studied for the constants 2 + u and 1 + 2u. In this paper, we consider constacyclic codes for the units 3 + 2u and 2 + 3u, and their Gray images under the new maps we introduced. Whereas the Gray maps on this ring considered in earlier works double the code length, our Gray maps triple it. Hence, we are able to find codes with new parameters not obtained in previous works.

A code of length *n* over *R* is a subset of R^n . *C* is a linear code if *C* is a *R*-submodule of R^n . Let λ be a unit in *R*. A linear code *C* over *R* of length *n* is called λ -constacyclic if it is

invariant under the constacyclic shift operator ρ ($r_0, r_1, \ldots, r_{n-1}$) = ($\lambda r_{n-1}, r_0, \ldots, r_{n-2}$), where ($r_0, r_1, \ldots, r_{n-1}$) $\in C$. The cyclic codes are a special case of constacyclic codes with $\lambda = 1$. If σ is the cyclic shift operator, that is σ ($r_0, r_1, \ldots, r_{n-1}$) = ($r_{n-1}, r_0, \ldots, r_{n-2}$), then *C* is said to be cyclic if σ (*C*) = *C*, and *C* is said to be λ -constacyclic if ρ (*C*) = *C*.

As usual, a λ -constacyclic code of length *n* over *R* can be identified as an ideal of the quotient ring $R[x] / \langle x^n - \lambda \rangle$ via the *R*-module isomorphism.

$$\psi : \mathbb{R}^{n} \longrightarrow \mathbb{R}[x] / \langle x^{n} - \lambda \rangle$$

$$(a_{0}, \dots, a_{n-1}) \longmapsto a_{0} + a_{1}x + a_{2}x^{2} + \dots + a_{n-1}x^{n-1} (mod \langle x^{n} - \lambda \rangle)$$

Hence, we use polynomial and vector representations of codes interchangeably.

Next, we define two Gray maps from R to \mathbb{Z}_4^3 . The Gray map on R used in previous works doubles the length of a code. Our Gray maps triple the code length, hence we obtain codes of different lengths over \mathbb{Z}_4 from a given code over R. Our first Gray map is

$$\phi_1 : R \longrightarrow \mathbb{Z}_4^3$$
$$a + ub \longmapsto (b, 2a + 3b, a + 3b)$$

This map is extended component-wise to

$$\begin{split} \phi_1 &: R^n \longrightarrow \mathbb{Z}_4^{3n} \\ \alpha &= (\alpha_0, \dots, \alpha_{n-1}) \longmapsto \phi(\alpha) \\ &= (b_0, \dots, b_{n-1}, 2a_0 + 3b_0, \dots, 2a_{n-1} + 3b_{n-1}, a_0 + 3b_0, \dots, a_{n-1} + 3b_{n-1}) \end{split}$$

where $\alpha_i = a_i + ub_i$, where i = 0, 1, ..., n - 1. ϕ_1 is a \mathbb{Z}_4 -module isomorphism.

We define another Gray map as follows,

$$\phi_2 : R \longrightarrow \mathbb{Z}_4^3$$

 $a + ub \longmapsto (a, 2a + 3b, 3a + 3b)$

This map is extended component-wise to

$$\begin{split} \phi_2 &: R^n \longrightarrow \mathbb{Z}_4^{3n} \\ \alpha &= (\alpha_0, \dots, \alpha_{n-1}) \longmapsto \phi_2(\alpha) \\ &= (a_0, \dots, a_{n-1}, 2a_0 + 3b_0, \dots, 2a_{n-1} + 3b_{n-1}, 3a_0 + 3b_0, \dots, 3a_{n-1} + 3b_{n-1}) \end{split}$$

where $\alpha_i = a_i + ub_i$, where i = 0, 1, ..., n - 1. ϕ_2 is a \mathbb{Z}_4 -module isomorphism. The polynomial correspondence of the Gray maps can be described as

$$\begin{split} \Phi &: R[x] / \langle x^n - \lambda \rangle \longrightarrow \mathbb{Z}_4[x] / \langle x^{3n} - 1 \rangle \\ a(x) + ub(x) \longmapsto \Phi(a(x) + ub(x)) = b(x) + x^n (2a(x) + 3b(x)) + x^{2n} (a(x) + 3b(x))) \\ (\Phi(a(x) + ub(x)) = a(x) + x^n (2a(x) + 3b(x)) + x^{2n} (3a(x) + 3b(x))) \end{split}$$

In the rest of the paper λ will denote either the constant 3 + 2u or 2 + 3u, unless otherwise stated. We note that $\lambda^n = \begin{cases} \lambda, & \text{if } n \text{ is odd.} \\ 1, & \text{if } n \text{ is even.} \end{cases}$

Definition 1 Let $a \in \mathbb{Z}_4^{3n}$ with $a = (a^{(0)} | a^{(1)} | a^{(2)})$, $a^{(i)} \in \mathbb{Z}_4^n$ for i = 0, 1, 2. Let φ be a map from \mathbb{Z}_4^{3n} to \mathbb{Z}_4^{3n} given by $\varphi(a) = (\sigma(a^{(0)}) | \sigma(a^{(1)}) | \sigma(a^{(2)}))$ where σ is the cyclic shift from \mathbb{Z}_4^n to \mathbb{Z}_4^n given by $\sigma(a^{(i)}) = (a^{(i,n-1)}, a^{(i,0)}, \dots, a^{(i,n-2)})$ for every $a^{(i)} = (a^{(i,0)}, \dots, a^{(i,n-1)})$ where $a^{(i,j)} \in \mathbb{Z}_4$, $j = 0, 1, \dots, n-1$. A code of length 3n over \mathbb{Z}_4 is said to be quasi cyclic (QC) code of index 3 if $\varphi(C) = C$.

🖄 Springer

3 The Gray images of some constacyclic codes over *R*

Proposition 1 Let ρ be the (3 + 2u)-constacyclic shift on \mathbb{R}^n , let ϕ_1 be the Gray map from \mathbb{R}^n to \mathbb{Z}_4^{3n} , and φ be the map defined above. Then we have $\phi \rho = \psi \varphi \phi_1$, where the permutation ψ of \mathbb{Z}_4^{3n} is defined by $\psi(x_1, x_2, \dots, x_n, x_{n+1}, \dots, x_{2n}, x_{2n+1}, \dots, x_{3n}) = (x_{\zeta(1)}, x_{\zeta(2)}, \dots, x_{\zeta(3n)})$ with the permutation $\zeta = (1, n + 1)$ of $\{1, 2, \dots, 3n\}$.

Proof Let $r = (r_0, ..., r_{n-1}) \in \mathbb{R}^n$, where $r_i = a_i + ub_i$, i = 0, 1, ..., n - 1. Then we have

$$\varrho(r) = ((3+2u)r_{n-1}, r_0, \dots, r_{n-2})$$

= ((3a_{n-1}+2b_{n-1}) + u(3b_{n-1}+2a_{n-1}), a_0 + ub_0, \dots, a_{n-2} + ub_{n-2})

By applying the Gray map ϕ_1 , we have

$$\phi_1(\varrho(r)) = (2a_{n-1} + 3b_{n-1}, b_0, \dots, b_{n-2}, b_{n-1}, 2a_0 + 3b_0, \dots, 2a_{n-2} + 3b_{n-2}, a_{n-1} + 3b_{n-1}, a_0 + 3b_0, \dots, a_{n-2} + 3b_{n-2})$$

On the other hand,

$$\varphi(\phi_1(r)) = (b_{n-1}, b_0, \dots, b_{n-2}, 2a_{n-1} + 3b_{n-1}, 2a_0 + 3b_0, \dots, 2a_{n-2} + 3b_{n-2}, a_{n-1} + 3b_{n-1}, a_0 + 3b_0, \dots, a_{n-2} + 3b_{n-2})$$

Finally, applying the permutation ψ to $\varphi(\phi_1(r))$ we obtain $\phi_1 \varrho = \psi \varphi \phi_1$.

Theorem 1 The ϕ_1 -image of a (3+2u)-constacyclic code over R of length n is permutation equivalent to a QC code of index 3 over \mathbb{Z}_4 of length 3n.

Proof Let *C* be a (3 + 2u)-constacyclic code over *R*. Then $\rho(C) = C$. Applying ϕ_1 , we have $\phi_1(\rho(C)) = \phi_1(C)$. By Proposition 2, we get $\phi_1(\rho(C)) = \psi(\varphi(\phi_1(C))) = \phi_1(C)$. Therefore, $\phi_1(C)$ is permutation equivalent to a QC code of index 3 over \mathbb{Z}_4 with length 3*n*.

We also have similar results for the map ϕ_2 with a different shift constant. The proofs of the following proposition and the theorem are very similar to the corresponding results for ϕ_1 , hence they are omitted here.

Proposition 2 Let ρ be the (2 + 3u)-constacyclic shift of \mathbb{R}^n , ϕ_2 and φ be defined as in the previous section. Then we have $\phi_2 \rho = \psi \varphi \phi_2$, where the permutation ψ of \mathbb{Z}_4^{3n} is defined by $\psi(x_1, x_2, \ldots, x_n, x_{n+1}, \ldots, x_{2n}, x_{2n+1}, \ldots, x_{3n}) = (x_{\zeta(1)}, x_{\zeta(2)}, \ldots, x_{\zeta(3n)})$ with the permutation $\zeta = (1, n + 1)$ of $\{1, 2, \ldots, 3n\}$.

Theorem 2 *The Gray image of a* (2 + 3u)*-constacyclic codes over R of length n is permutation equivalent to a QC code of index 3 over* \mathbb{Z}_4 *with length 3n.*

We note that for odd lengths, cyclic codes and constacyclic codes are equivalent.

Proposition 3 Let *n* be an odd integer and let λ be one of the units of *R* specified above. Then

$$\Psi : R[x] / \langle x^n - 1 \rangle \longrightarrow R[x] / \langle x^n - \lambda \rangle$$

$$c(x) \longmapsto \Psi(c(x)) = c(\lambda x)$$

is a ring isomorphism.

Springer

Proof The proof similar to the proof Proposition 3.6 in [17].

It follows that for an odd length *n* there is a one-to-one correspondence between cyclic codes and constacyclic codes. Next two corollaries formally state this fact.

Corollary 1 Let *n* be an odd integer and let $\lambda \in R$. Then *I* is an ideal of $R[x] / \langle x^n - 1 \rangle$ if and only if $\Psi(I)$ is an ideal of $R[x] / \langle x^n - \lambda \rangle$. Hence there is a one-to-one correspondence between the ideals of $R[x] / \langle x^n - 1 \rangle$ and the ideals of $R[x] / \langle x^n - \lambda \rangle$ such that the cyclic code *I* is equivalent to the constacyclic code $\Psi(I)$.

Corollary 2 Let $\overline{\zeta}$ be the map from \mathbb{R}^n to \mathbb{R}^n defined as

$$\overline{\zeta} : \mathbb{R}^n \longrightarrow \mathbb{R}^n$$
$$z = (z_0, \dots, z_{n-1}) \longmapsto (z_0, \lambda z_1, \lambda^2 z_2, \dots, \lambda^{n-1} z_{n-1})$$

C is a cyclic code over *R* of odd length *n* if and only if $\overline{\zeta}(C)$ is a λ -constacyclic code of length *n* over *R*.

As in [17], we characterize λ -constacyclic codes over *R* of odd length as follows.

Theorem 3 Let *n* be an odd number and *C* be a λ -constacyclic code of length *n* over *R*. Then *C* is an ideal in $R[x]/\langle x^n - \lambda \rangle$ generated by

$$C = \langle t_1(\widetilde{x})(p_1(\widetilde{x}) + 2) + (1+u)t_2(x)(p_2(\widetilde{x}) + 2), (1+u)t_3(\widetilde{x})(p_3(\widetilde{x}) + 2) \rangle$$

where $\tilde{x} = \lambda x$, $t_i(x)$, $p_i(x)$, $g_i(x)$ are monic pairwise coprime polynomials in $\mathbb{Z}_4[x]$ such that $x^n - 1 = t_i(x)p_i(x)g_i(x)$, where i = 1, 2, 3.

By using a special generator polynomial of a cyclic code and by the relationship between cyclic codes and λ -constacyclic codes given in Proposition 3, we can express a generator of a λ -constacyclic code *C* as

$$C = \langle g_1(\widetilde{x}), ug_2(\widetilde{x}) \rangle$$

where $\widetilde{x} = \lambda x$ as in [17].

Proposition 4 Let $\overline{\zeta}$ be defined as above. Then $\phi_1\overline{\zeta} = \Pi\phi_1$ (and $\phi_2\overline{\zeta} = \Pi\phi_2$) where the permutation Π of \mathbb{Z}_4^{3n} is defined by

$$\Pi(x_1, x_2, \dots, x_n, x_{n+1}, \dots, x_{2n}, x_{2n+1}, \dots, x_{3n}) = (x_{\tau(1)}, x_{\tau(2)}, \dots, x_{\tau(3n)})$$

with the permutation $\tau = (2, n+2)(4, n+4) \dots (n-1, 2n-1)$ of $\{1, 2, \dots, 3n\}$.

Proof Let $r = (r_0, ..., r_{n-1}) \in \mathbb{R}^n$, where $r_i = a_i + ub_i$, i = 0, 1, ..., n-1. Since $\overline{\zeta}(r) = (r_0, (3+2u)r_1, (3+2u)^2r_2, ..., (3+2u)^{n-1}r_{n-1}) = (r_0, (3+2u)r_1, r_2, (3+2u)r_3, ..., (3+2u)r_{n-2}, r_{n-1})$, we have $\phi_1(\overline{\zeta}(r)) = (b_0, 2a_1 + 3b_1, b_2, ..., b_{n-1}, 2a_0 + 3b_0, b_1, ..., a_0 + 3b_0, a_1 + 3b_1, ..., a_{n-1} + 3b_{n-1})$.

On the other hand, $\phi_1(r) = (b_0, \dots, b_{n-1}, 2a_0 + 3b_0, \dots, 2a_{n-1} + 3b_{n-1}, a_0 + 3b_0, \dots, a_{n-1} + 3b_{n-1})$. By applying Π , we have $\phi_1\overline{\zeta} = \Pi\phi_1$. Similarly, we have $\phi_2\overline{\zeta} = \Pi\phi_2$.

Corollary 3 Let n be an odd integer. If Γ is the Gray image (under ϕ_1 or ϕ_2) of a linear cyclic code over R, then Γ is permutation equivalent to a QC code over \mathbb{Z}_4 of index 3 of length 3n, via the permutation Π defined above.

1253

Element of <i>R</i>	Encoding	
0	0	
1	1	
2	2	
3	3	
u	4	
1 + u	5	
2+u	6	
3 + u	7	
2 <i>u</i>	8	
1 + 2u	9	
2 + 2u	А	
3 + 2u	В	
3и	С	
1 + 3u	D	
2 + 3u	Е	
3 + 3u	F	

Table 1 Representations of theelements of R

Table 2 New linear codes over \mathbb{Z}_4 obtained from constacyclic codes over *R*

n	λ	Map	Generator polynomial	Parameters of \mathbb{Z}_4 image
15	2 + 3u	ϕ_1	1E3E260E3806	$[45, 4^8 2^0, 24]$
15	2 + 3u	ϕ_1	181806303 <i>E</i> 1	$[45, 4^{10}2^0, 20]$
15	2 + 3u	ϕ_1	162 <i>E</i> 26	$[45, 4^{20}2^0, 8]$
21	2 + 3u	ϕ_1	162E201E22320626	$[63, 4^{12}2^0, 24]$
21	2 + 3u	ϕ_1	160026162E16201	$[63, 4^{14}2^0, 24]$
21	2 + 3u	ϕ_1	1 <i>E</i> 1 <i>E</i> 3806	$[63, 4^{28}2^0, 8]$
21	2 + 3u	ϕ_1	18303 <i>E</i> 162 <i>E</i> 0 <i>E</i> 36	$[63, 4^{16}2^0, 16]$
21	2 + 3u	ϕ_1	1628181	$[63, 4^{30}2^0, 8]$
10	3 + 2u	ϕ_1	13FA5B81	$[30, 4^8 2^0, 16]$
15	2 + 3u	ϕ_2	16262806	$[45, 4^{16}2^0, 12]$
21	2 + 3u	ϕ_2	1E161038100E30201	$[63, 4^{10}2^0, 28]$
21	2 + 3u	ϕ_2	102016	$[63, 4^{32}2^0, 8]$
10	3 + 2u	ϕ_2	13DA7B81	$[30, 4^8 2^0, 14]$
20	3 + 2u	ϕ_2	12C070A2D2E29	$[60, 4^{16}2^0, 16]$
30	3 + 2u	ϕ_2	$\frac{x^{30} - \lambda}{x^2 + x + u}$	$[90, 4^42^0, 60]$

Proof Let Γ be such that $\Gamma = \phi_1(D)$ (or $\Gamma = \phi_2(D)$), where D is a linear cyclic code over R. From Proposition 10, $\phi_1(\overline{\zeta}(D)) = \Pi(\phi_1(D)) = \Pi(\Gamma)$ (and $\phi_2(\overline{\zeta}(D)) = \Pi(\phi_2(D)) = \Pi(\Gamma)$). We know that $\overline{\zeta}(D)$ is a linear λ -constacyclic code C. So $\phi_1(\overline{\zeta}(D)) = \phi_1(C)$ (and $\phi_2(\overline{\zeta}(D)) = \phi_2(C)$). By Proposition 4, $\phi_1(C)$ (and $\phi_2(C)$) is permutation equivalent to a QC code of index 3 over \mathbb{Z}_4 with length 3n.

4 Computational results

Finding new linear codes over \mathbb{Z}_4 was our main motivation for this work. We conducted a computer search to look for new linear codes over \mathbb{Z}_4 by generating constacyclic codes over R (for the two values of λ discussed in this paper) and computing their \mathbb{Z}_4 -images using each of the Gray map we introduced. We have been able to find 15 new linear codes over \mathbb{Z}_4 with better parameters than the currently best-known codes in the database [8]. We present the details of these codes in Table 2 below. The generator polynomial g(x) of each code is a monic divisor of $x^n - \lambda$ for $\lambda = 2+3u$ or 3+2u. For the sake of saving space, we use a compact representation of the polynomials where coefficients are listed in descending order and elements of R are encoded according to Table 1 below. Each element of R is encoded by a single character in the set $\{0, 1, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$. Thus the generator polynomial in the third row of Table 2, for example, is $x^5 + (2 + u)x^4 + 2x^3 + (2 + 3u)x^2 + 2x + u + 2$. For the last entry in the same table, we gave an alternative description of the generator polynomial because it has a large degree. We only computed the Lee weights of the codes over \mathbb{Z}_4 .

References

- 1. Ashraf M., Mohammed G.: (1 + 2u)-constacyclic codes over $\mathbb{Z}_4 + u\mathbb{Z}_4$. arXiv:1504.03445v1.
- 2. Aydin N., Asamov T.: A database of \mathbb{Z}_4 codes. J. Comb. Inf. Syst. Sci. **34**, 1–12 (2009).
- 3. Aydin N., Dertli A., Cengellenmis Y.: Cyclic and constacyclic codes over $\mathbb{Z}_4 + w\mathbb{Z}_4$. Pre-print.
- 4. Bandi R., Bhaintwal M.: Self dual codes over $\mathbb{Z}_4 + w\mathbb{Z}_4$. Discret. Math. Alg. Appl. 7, 1550014 (2015).
- 5. Bandi R., Bhaintwal M.: Negacyclic codes over $\mathbb{Z}_4 + u\mathbb{Z}_4$. arXiv:1412.3751v1.
- 6. Bandi R., Bhaintwal M.: Cyclic codes over $\mathbb{Z}_4 + u\mathbb{Z}_4$. arXiv:1501.01327.
- Bonnecaze A., Rains A.M., Solé P.: 3-Colored 5-designs and Z₄-codes. J. Stat. Plan. Inf. 86, 349–368 (2000).
- 8. Database of \mathbb{Z}_4 codes [online]. Z4Codes.info (Accessed February 2017).
- 9. Gao J., Gao Y., Fu F.: On linear codes over $\mathbb{Z}_4 + v\mathbb{Z}_4$. arXiv:1402.6771v2.
- Hammons Jr. A.R., Kumar P.V., Calderbank A.R., Solane N.J.A., Solé P.: The Z₄ linearity of Kerdock, Preparata, Goethals and related codes. IEEE Trans. Inf. Theory 40, 301–319 (1994).
- Martinez-Moro E., Szabo S.: On codes over local Frobenius non-chain rings of order 16. Contemp. Math. 634, 227–241 (2015).
- Martinez-Moro E., Szabo S., Yildiz B.: Linear codes over Z₄[x]/⟨x² + 2x⟩. Int. J. Inf. Coding Theory 2, 78–96 (2015).
- 13. Ozen M., Ozzaim T., Aydin N.: Cyclic codes over $\mathbb{Z}_4 + u\mathbb{Z}_4 + u^2\mathbb{Z}_4$. Turk. J. Math. (to appear)
- Ozen M., Uzekmek F.Z., Aydin N., Özzaim N.T.: Cyclic and some constacyclic codes over the ring Z₄[u]/⟨u² − 1⟩. Finite Fields Appl. 38, 27–39 (2016).
- 15. Pattanayak S., Singh A.: A class of cyclic codes over the ring $Z_4[u]/\langle u^2 \rangle$ and its Gray image. arXiv:1507.04938v1.
- Shi M., Qian L., Sok L., Aydin N., Solé P.: On constacyclic codes over Z₄[u]/⟨u² − 1⟩. Finite Fields Appl. 45, 86–95 (2017).
- 17. Shi M., Qian L., Sok L.: On constacyclic codes over $\mathbb{Z}_4[u]/\langle u^2 1 \rangle$ and their Gray images. arXiv:1608.00820.
- Solé P.: A quaternary cyclic code, and a family of quadriphase sequences with low correlation properties. In: Springer Lecture Notes in Computer Science, vol. 388, pp. 193–201 (1998).
- 19. Wan Z.X.: Quaternary Codes. World Scientific, Singapore (1997).
- Yildiz B., Aydin N.: On cyclic codes over Z₄ + uZ₄ and their Z₄ images. Int. J. Inf. Coding Theory 2, 226–237 (2014).
- Yildiz B., Karadeniz S.: Linear codes over Z₄ + uZ₄, Macwilliams identities, projections and formally self dual codes. Finite Fields Appl. 27, 24–40 (2014).