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Analyzing the power consumption behavior of a large scale data center

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Abstract

The aim of this paper is to illustrate the use of application and system level logs to better understand scientific data center behavior and energy-spending. Analyzing a data center log of 900 nodes (Sandy Bridge and Haswell), we study node power consumption and describe approaches to estimate and forecast it. Our results include methods to cluster nodes based on different vmstat and RAPL measurements as well as Gaussian and GAM models for estimating the plug power consumption. We also analyze failed jobs and find that non-successfully terminated jobs consume around 40% of computing time. While the actual numbers are likely to vary in different data centers at different times, the purpose of the paper is to share ideas of what can be found by statistical and machine learning analysis of large amount of log data.

Keywords RAPL · Energy modeling · Energy efficiency · Data center log analysis

1 Introduction

According to a recent report by Lawrence Berkeley National Laboratory [16] the data centers in United States consumed 70 billion kWh of electricity in 2014. The consumption is predicted to grow even higher although the growth has been more moderate than expected earlier. One reason for the moderate growth of power consumption while the computing needs have drastically increased, has been the attention of both high performance computing (HPC) industry and researchers to improve the energy efficiency. Reduced consumption results both in smaller electricity bill and reduced environmental load.

Möbius et al. [13] provide a comprehensive survey of electricity consumption estimation in HPC systems. The techniques can be broadly categorized as direct measurements and power modeling. Direct measurement techniques involve power measuring devices or sensors to monitor the current draw [14] whereas power modeling techniques estimate the power draw with system utilization metrics such as hardware counters or Operating System (OS) counters [5].

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Intel's Running Average Power Limit (RAPL) is one such power measurement tool, which has been useful in power measurement and modeling research [8,11,17]. RAPL reports the real time power consumption of the CPU package, cores, DRAM and attached GPUs using Model Specific Registers (MSRs). Since its introduction in Sandy Bridge it has evolved and in newer architectures, Haswell and Skylake, RAPL works as a reliable and handy power measurement tool [8].

In this paper, we study and analyze the energy consumption of a computing cluster named *Taito*, which is a part of the CSC - IT Center for Science in Finland. In *Taito*, most of the jobs come from universities and research institutes. They are typically simulation or data analysis jobs and run parallel on multiple cores and nodes. We utilize a dataset of 900 nodes (Sandy Bridge and Haswell) which includes OS counter logs from vmstat tool (see Table 1), CPU package power consumption values from RAPL and plug power consumption value sampled at a frequency of approximately 0.5 Hz over a period of 42 h (more details in Sect. 3).

The aim of this study is to show examples of information that can be extracted from data center logs. In particular we

- 1. Investigate how OS counters and RAPL measurements can be used to explain and estimate the total power consumption of a computing node (Sects. 4, 5 and 7).
- 2. Analyse failed jobs and their influence in energy spending (Sect. 6).

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 Table 1
 Vmstat output variables

 used: description and min and
 max values in CSC dataset

Description	Min.	Max.	
# of processes waiting CPU time		200	
# of processes waiting on I/O	0	97	
# of virtual memory blocks	0	9,775,548	
# of blocks of idle memory	393,316	876,866,240	
# of memory blocks used as cache	15,656	622,179,392	
# of blocks per sec swapped in	0	27	
# of blocks per sec swapped out	0	27	
# of blocks received from HD	0	1247	
# of blocks sent to from HD	0	3461	
# of interrupts per sec	0	74	
# of context switches per sec	0	73	
User time % of CPU time	1	97	
System (kernel) % of CPU time	0	20	
Idle % of CPU time	2	92	
% of CPU time waiting for IO	0	29	
	 # of processes waiting CPU time # of processes waiting on I/O # of virtual memory blocks # of blocks of idle memory # of blocks of idle memory # of blocks per sec swapped in # of blocks per sec swapped out # of blocks received from HD # of blocks sent to from HD # of of context switches per sec User time % of CPU time System (kernel) % of CPU time Idle % of CPU time waiting for IO 	BeschiptionInit.# of processes waiting CPU time0# of processes waiting on I/O0# of processes waiting on I/O0# of virtual memory blocks0# of blocks of idle memory393,316# of blocks of idle memory393,316# of blocks of idle memory15,656# of blocks per sec swapped in0# of blocks per sec swapped out0# of blocks received from HD0# of blocks sent to from HD0# of of context switches per sec0User time % of CPU time1System (kernel) % of CPU time0Idle % of CPU time2% of CPU time waiting for IO0	

- 3. Cluster the nodes based on the OS counter and RAPL values. This gives an indication of the opportunities to combine different workload in a way which uses the resources in a balanced way (Sect. 5).
- 4. Use machine learning to map power consumption to OS counter values (Sects. 7 and 8).

2 Related works

Power measurement is one of the key inputs in any energy efficient system design. As such, it has been quite extensively studied in the energy efficiency literature for HPC systems and data centers. As described in Sect. 1, the measurement techniques can be categorized as direct measurements and power modeling. Direct measurements using external power meters provide accurate measurements and can give real time power consumption of different components of the system depending on the type of hardware and software instrumentation [6,7]. However, direct measurement techniques often require physical system access and custom and complex instrumentations. Sometimes such techniques may hinder the normal operation of the data center [5].

Modern day data centers also make use of sensors and/or Power Distribution Units (PDUs) that monitor and report useful runtime information about the system such as power or temperature. Such tools also show good accuracy. However, PDUs and sensors can be costly to deploy and may not scale well as the demand increases. These devices are not yet commonly deployed and they might have usability issues as reported in [5].

Power modeling using performance counters are quite useful with regards to cost, usability and scaling. There are mainly two types of such counters which can be used in power modeling of computing systems, namely hardware performance counters (often referred as Performance monitoring counters(PMC)) and OS provided utilization counters or metrics. PMCs have been used quite extensively in monitoring the system behavior and finding correlation with power expenditure of systems thus providing a useful input for power modeling approaches [2,9]. However, such models often suffer from problems like limited number of events that can be monitored and then PMCs are often architecture dependent and so the models may not be transferable from one architecture to the other [13]. The accuracies of such models are also often workload dependent and as such may not be reliable at times [5,13].

Intel introduced the RAPL interface [10] to limit and monitor the energy usage on its Sandy Bridge processor architectures. It is designed as a power limiting infrastructure which allows users to set a power cap and as a part of this process it also exposes the power consumption readings of different domains. RAPL is implemented as Model-Specific Registers (MSRs) which are updated roughly every millisecond. RAPL provides energy measurements for processor package (PKG), power plane 0 (PP0), power plane 1 (PP1), DRAM, and PSys which concerns entire System on Chip (SoC). PKG includes the processor die that contains all the cores, on-chip devices, and other uncore components, PP0 reports the consumption of CPU cores only, PP1 holds the consumption of on-chip graphics processing units (GPU) and DRAM plane gives the energy consumption of dual in-line memory modules (DIMMs) installed in the system. From Intel's Skylake architecture onwards RAPL also reports the consumption of entire SoC in PSys domain (it may not be available on all Skylake versions). In Sandy Bridge, RAPL domain values were modeled (not measured) and thus it had some deviations from the actual measurements [8]. With the

introduction of Fully Integrated Voltage Regulators (FIVRs) in Haswell, RAPL readings have promisingly improved and it has proved its usefulness in power modeling also [11].

There has also been interesting works regarding the job power consumption and estimation for data centers [3,15]. Borghesi et al. [3] proposed machine learning technique to predict the consumption of HPC system using real production data from Eurora supercomputer. Their prediction technique show an average error of approximately 9%. In our analysis, we show a different analysis of data center power consumption since we use system utilization metrics from OS counters and RAPL. Our results confirm a few of the observations already seen in literature. However, our approach is different since we make use of tools like vmstat and RAPL from a real life production dataset. We show the power consumption predictability of such tools and we pinpoint metrics which tend to correlate more with the power readings than the other as we cluster nodes based on vmstat and RAPL values. This paper also demonstrates different modeling techniques (leveraging machine learning) to model the plug power from OS counter and RAPL values and pinpoints essential parameters that influence the accuracy of such techniques.

3 Dataset description

The CSC dataset consists of around 900 nodes which are all part of *Taito* computing cluster. Among the 900 nodes, there are approximately 460 Sandy Bridge compute nodes, 397 Haswell nodes and a smaller number of more specialized nodes with GPUs, large amounts of memory or fast local disks for I/O intensive workloads. Since there are different hardwares and hence performance differences between the two types of nodes, their power consumption exhibit different patterns (see Fig. 1).

 Table 2
 Hardware configurations—*Taito* compute nodes

Туре	Haswell	Sandy bridge
Number of nodes	397	496
Node model	HP XL230a G9	HP SL230s G8
Cores/node	24	16
Memory/node	128 GB	64 GB

The dataset, captured in June 2016, consists of vmstat output (Table 1), RAPL package power readings, plug power obtained from Intelligent Platform Management Interface (IPMI) and job ids. All of these are sampled at a frequency of approximately 0.5 Hz over a period of 42 h. The hardware configurations of *Taito*'s compute nodes are given in Table 2 [1].

vmstat (Virtual memory statistics) is a Linux tool, which reports the usage summary of memory, interrupts, processes, CPU usage and block I/O. The vmstat variables that we have used are presented in Table 1. The CSC dataset reports the energy consumption of two RAPL PKG domains for the dual socket based server systems in *Taito*. The metrics collection for this dataset was done manually. In order to continuously collect and analyze this type of data, better high-resolution energy measurement tools are needed which should ideally work in a cross-platform basis across different hardware and batch job schedulers.

4 Power consumption of computing nodes

We start by inspecting how the variable of interest: power consumption (measured directly at the plug) changes over time at different nodes. First observation is that there are



Fig. 1 Power consumption differences on Haswell and Sandy Bridge nodes. **a** Distributions of average values per node, **b** whisker diagrams with all the values



Fig. 2 Power consumption of nodes running mostly a single job. a Node C581, b node C836, c node C749



Fig. 3 Power consumption of nodes running a highly variable number of jobs. a Node C585, b node C626, c node C819

considerable variations in the measured power consumption between different nodes (see Fig. 1), and even at a single node, at different time intervals during the observed period. This is not surprising, as the node power consumption at any point is dependent on the type of computing jobs running on that node. In order to illustrate this variability, we show the power consumption plots of several nodes with rather diverse patterns in Figs. 2 and 3.

From Fig. 2 we observe that single running jobs also exhibit different patterns and variability in how they consume power. While the influence of the number of jobs running on a node on its power consumption is evident from Fig. 3, it is also clear that this dependency is very subtle and not straight forward to express.

5 Vmstat and RAPL variables statistics

After the observations on the power consumption in relation to the number of jobs running on a node, we turn to the observation of power consumption in relation to the vmstat output values. Namely, vmstat output informs us about the consumption of different computing resources on a node and hence captures more subtle properties of the jobs running on the node. The description of the vmstat output variables in CSC dataset is presented in Table 1.



Fig. 4 Power consumption and number of user and kernel processes running—node C775 (color figure online)

Taking the same set of nodes introduced earlier (Figs. 2 and 3), we investigate visually the interplay of vmstat and RAPL variables and power consumption. We observe that vmstat values **r,b** (see Table 1 for explanation) change even on a node running no jobs. Looking at similar analysis for the nodes running several jobs in Fig. 4, the relationship between vmstat values **r,b** and power consumption values is evident. Similarly, Fig. 5 illustrates the interplay between memory



Fig. 5 Power (in blue) and two types of memory consumption (see legend). a Node C581, b node C836, c node C749 (color figure online)



Fig. 6 Power (in blue) and two types of CPU consumption (see legend). a Node C585, b node C626, c node C819

RAPL values (DRAM) and power consumption, and Fig. 6 between CPU RAPL values and power consumption.

Figure 7 presents Self-Organizing Maps (SOM) model [12] classification output on the CSC dataset. SOM is a unsupervised classification technique to visualize high dimensional data in low dimensional space. In this figure, we cluster all the nodes in 9 clusters based on the similarity in Node data. Node count per class shows the number of nodes in different clusters as a heat map. Clusters represented in 'white' color contain around 200+ nodes whereas clusters represented in 'red' color contain around 50 or less nodes with the other colors falling in between. If we now see the same clusters in the Node data (left sub-figure of Fig. 7), we observe which variables dominate the similarities in that cluster. For example, Node data for the 'white' colored cluster in the top-right corner shows that the variables us, CPU1, CPU2 and plug dominate the cluster (CPU1, CPU2 correspond to the RAPL package power).

6 Analysis of unsuccessful jobs

Table 3 presents statistics of the jobs executed on the *Taito* cluster. We focus on the job exit status, number of jobs which have the same status, elapsed time per job (in hours) and total CPU Time used (user time plus system time). The dataset from *Taito* contains four types of job status: completed, failed, cancelled and timeout. Completed jobs are successful

jobs that ran to completion. Failed jobs are jobs that failed to complete successfully and did not produce desirable outputs. Cancelled jobs are cancelled by their users. These are often failures but sometimes cancellation is done on purpose after the job has produced the desirable results. Timeout jobs did not run to successful completion within a given time limit. Timeouts are not necessarily failures, they are done occasionally on purpose and can produce useful outputs.

From Table 3 we can see that approximately 84% of the jobs are completed jobs and they consume 56.95% of the total CPU time. Failed jobs on the other hand constitute of 12.5% of the total jobs and they consume around 14.75% of the total CPU time. Interestingly, only 0.5% of the total jobs are timed out but they consume around 19.34% of the total CPU time. Timeout jobs also have elapsed time of 25 h per job which is by far the maximum.

If we have a pessimistic assumption that all the noncompleted jobs are unsuccessful it turns out that 16% of such jobs consumed around 43% of total CPU time. This shows that the wasted resources and energy in terms of unsuccessful jobs can be as much as 43% in typical data centers. This is approximately 280.000 days of CPU time in numbers. If these failures are identified in relatively early stage of a job lifetime, the potential CPU time and energy save can be significant. It can be a potential target for energy efficiency in data center workload management.





Table 3Job statistics—total of809,178jobs

Job status	Nr. of jobs (%)	Elapsed time/job (h)	CPU time (%)	
Completed	84.0	1.0	56.95	
Failed	12.5	0.7	14.75	
Cancelled	3.0	8.0	8.96	
Timeout	0.5	25	19.34	

7 Estimation results

In this section we present results of power consumption estimation based on historical power consumption, vmstat and RAPL data (input to build the model) and current vmstat and RAPL values (intervention variables). We take first twothirds of the time period (around 1 day) as historical data and we build the model on it. Afterwards we test the accuracy of prediction of such a model on the last third of the data (around half a day).

At first we tested building a model on data from a single node and predicting power at the same node. We do not report these results, as on some nodes this approach has worked rather well, but on some other nodes the results were under an acceptable limit. However, such an exercise taught us that the 'problematic' nodes on which prediction performance was poor, featured a sudden change in the patterns of power consumption and job execution during the period we were trying to predict. Since ML algorithms are designed to learn from 'seen' values, and they do not perform well on 'unseen' ones, which result in poor performance in such cases.

After such an understanding, we try building ML models on a random sample of shuffled data coming from **all** the nodes (of type Haswell) in our dataset. Precisely, we sample 2% of data from all the nodes (251, 244 data samples) and evaluate performance of different ML algorithms on it using standard 10-fold cross validation approach. The best result is achieved using Random Forest [4] as shown in Table 4.

In addition to a high correlation coefficient, the regression model makes mean absolute error (MAE) of 3.12, which is measured in the units of target variable (power consumption). If we remind ourselves of the power consumption values on Haswell nodes in Fig. 1b, we see that such an error compared to average values around 300 yields a good result. Root mean squared error (RMSE) is more sensitive to sudden changes in the target variable, which are present in our data. Relative errors measure how well our estimation compares to a null model that would always predict the average value. The value larger than 100% would mean that our model is performing worse, while smaller values are better (Table 5).

 Table 4
 Power estimation: 10-fold cross validation results on a 2% sample from all nodes

Correlation coefficient (corrcoef)	0.97
Mean absolute error (MAE)	3.12
Root mean squared error (RMSE)	9.11
Relative absolute error (RAE)	12.25%
Root relative squared error (RRSE)	21.83%
Total number of instances	251,244

Table 5 Power estimation results per node					
Node id	Corr. coef.	RMSE	MAE	# Instances	
C832	0.92	13.92	10.84	15,838	
C907	0.93	16.98	10.68	26,149	
C836	0.79	1.37	1.94	27,962	
C775	0.99	6.68	12.09	26,756	
C581	0.96	1.60	2.00	28,136	
C585	0.91	6.42	13.17	28,174	
C626	0.74	6.60	13.80	28,594	
C742	0.99	10.64	14.01	19,727	
C749	0.68	3.34	3.73	28,197	
C819	0.73	10.62	16.14	27,505	



Fig. 8 Original correlation matrix

8 Modeling plug power

We take a sample of 30,000 measurements focusing on the 'Haswell' type computing nodes. 80% of this is used for the training set and 20% for the test set.

We aim at modelling the plug power using both OS counters and RAPL measurements. The variables and their linear correlations are shown in Fig. 8.

The distribution of the plug variable is shown in Fig. 9. The distribution does not match very well with any common theoretical distribution. However, using the normal distribution gives the best results when using regression models. We also tested whether there is any lag between the RAPL values and the plug power values and found out the best results are received when using the plug values 10 s after the RAPL measurements. The variable is named 'lag5', since we used 0.5 Hz sampling frequency.



Fig. 9 Distribution of plug variable

We first fitted a linear model for estimating the plug power consumption using the RAPL parameters.

$$f(x) = a_0 + a_2 CPU1 + a_3 CPU2 + a_4 DRAM1 + a_5 DRAM2 + e$$
(1)

Fitting the model to our training set gave the following result:

When testing the accuracy using the test sample, the linear model gave 2.10% mean absolute percentage error. Next, we applied generalized additive models (GAM):

$$g(u) = \beta_0 + f_1(x_1) + f_2(x_2) + \dots + f_n(x_n) + e$$

Where x_i are covariates, β_0 the intercept, f_i smooth functions, e_i the error terms, and g() the link function. This makes it possible to model non-linear relationships in a regression model. We use the same covariants as above and no link function. The mean absolute percentage error slightly decreased to 1.97%. Figure 10 shows the smooth functions of each independent variable in the GAM model. As we can see, the effect of the DRAM is much smaller than the effect of CPU. The curves are not totally linear meaning that the effect of RAPL values to the plug power is not exactly linear.

Finally, we include possible interactions among the RAPL variables into the model meaning that 2 or 3 variables can have a common effect. For example, CPU1 and DRAM2 together could increase the plug power more than both of them as separate components do. This is not included in the previous models.

Figure 11 illustrates the accuracy of the model. Large values match very well but the model has difficulties to estimate very small values. The mean absolute percentage error was slightly smaller again, 1.87%.



Fig. 10 Smooth functions of the GAM model



Fig. 11 Testing the GAM model with interactions against the test set. (Black asterisk = real value, red circle = estimated value, grey circles are 95% CI for the estimation) (color figure online)

In Figs. 12 and 13, we see plots illustrating the combined effects among variables. The total effect to the power consumption is shown in z-axis (upwards) while x- and y-axis represent the values of the variables. For example, in Fig. 12, we have an example of combined effect of CPU1 and CPU2 to the total power consumption. We see that the effect of CPU1 to the total power consumption decreases when its value increases, and when both the CPUs run at medium power, the total effect is slightly higher. In any case the combined effects are relative small compared to direct effects (e.g. Fig. 10).



Fig. 12 The combined effect of CPU1 and CPU2 measurements. In the middle range of the both values the actual effect seems to increase



Fig. 13 The combined effect of CPU1 and DRAM1 measurements. When little DRAM power is used, the CPU power has large effect

9 Conclusion

In this paper we have presented different approaches for analyzing data center power and OS counter based utilization logs. We have shown that estimating plug power from utilization metrics is promising and the logs can be used in different ways for producing effective power models for data centers. Tools such as RAPL add to the accuracy of the models by providing real time power consumption data. For example, the GAM model shows that RAPL values can predict the plug power with mean absolute error rate of 1.97%. If we consider interactions among RAPL variables the error reduces to 1.87%. Apart from modeling, our analysis also shows that unsuccessful jobs can consume significant resources and power. If the problems can be identified early in job life cycle, resource and energy waste can be reduced. In the future, we aim to utilize such data center logs to produce job specific power consumption models and identify power consumption anomalies within data center workload management.

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