

ORIGINAL PAPER

## **Spanning Tree Decompositions of Complete Graphs Orthogonal to Rotational 1-Factorizations**

**John Caughman[1](http://orcid.org/0000-0002-2692-9198) · John Krussel<sup>2</sup> · James Mahoney1**

Received: 9 June 2016 / Revised: 18 December 2016 / Published online: 10 January 2017 © Springer Japan 2017

**Abstract** In Krussel et al. (ARS Comb 57:77–82, [2000\)](#page-12-0), Krussel, Marshall, and Verall proved that whenever  $2n - 1$  is a prime of the form  $8m + 7$ , there exists a spanning tree decomposition of  $K_{2n}$  orthogonal to the 1-factorization  $GK_{2n}$ . In this paper, we develop a technique for constructing spanning tree decompositions that are orthogonal to rotational 1-factorizations of  $K_{2n}$ . We apply our results to show that, for *every n* > 2, there exists a spanning tree decomposition orthogonal to  $GK_{2n}$ . We include similar applications to other rotational families of 1-factorizations, and provide directions for further research.

**Keywords** Spanning tree · 1-factorization · Complete graph · Matching

**Mathematics Subject Classification** 05C05 · 05C15 · 05C70

## **1 Introduction**

A graph  $G = (V, E)$  consists of a finite set V of **vertices** and a set E of 2-element subsets of *V* called **edges**. A **1-factor** in *G* is a set of edges that partition *V*, and a **1-factorization** is a partition of *E* into 1-factors. The 1-factorizations of the complete graph  $K_{2n}$  are well-studied and we refer to  $[5,6]$  $[5,6]$  for an introduction to the topic and a survey of many known results.

The number of non-isomorphic 1-factorizations of  $K_{2n}$  grows very rapidly with *n* (see [\[2](#page-12-3)]), but one well-known family is denoted  $GK_{2n}$ , and can be easily described as

B John Caughman caughman@pdx.edu

<sup>&</sup>lt;sup>1</sup> Portland State University, Portland, OR 97207, USA

<sup>2</sup> Lewis and Clark College, Portland, OR 97219, USA

follows (see Example 1.1 below). Place vertices  $1, 2, \ldots, 2n - 1$  in a circle around a central vertex labeled 2*n*. Add the edge  $\{1, 2n\}$  and the edges  $\{i + 1, 2n - i\}$  for  $0 < i < n$  to form a 1-factor. Rotating this 1-factor around vertex 2*n* gives a partition of the edges of  $K_{2n}$  into exactly  $2n - 1$  different 1-factors. This partition is the 1factorization known as  $GK_{2n}$ .

<span id="page-1-1"></span>*Example 1* A 1-factor (left) of the complete graph  $K_8$  (right) can be rotated to partition the edge set, giving the 1-factorization  $G K_8$ .



A 1-factorization of  $K_{2n}$  is said to be **rotational** (also called *1-rotational*) if it is stabilized by a permutation of the vertices that fixes one vertex and cyclically permutes the rest. The 1-factorization  $G K_8$  illustrated above is an example, since the permutation  $\rho = (1, 2, 3, \ldots, 7)$  fixes vertex 8 and cyclically permutes the other vertices. In doing so,  $\rho$  also cyclically permutes the seven 1-factors of this 1-factorization, thereby stabilizing the set of them. So  $GK_8$  is rotational, and analogously, we see that  $GK_{2n}$ is rotational for any *n*. For rotational 1-factorizations (and the more general notions of *pyramidal* and *k-pyramidal* 1-factorizations), we refer the reader to [\[4](#page-12-4)[,5](#page-12-1)].

Given any 1-factorization  $F$  of  $K_{2n}$ , we say that a subgraph *G* is **orthogonal** to  $F$  if each 1-factor of  $F$  shares at most one edge with  $G$ . For example, the star graph, in which vertex 8 is adjacent to vertices 1 through 7, is a spanning tree for  $K_8$  that is orthogonal to the 1-factorization  $G K_8$  described above.

To avoid trivialities, let us assume that  $n > 2$ . In [\[1\]](#page-12-5), Brualdi and Hollingsworth conjectured that for any 1-factorization  $F$  of  $K_{2n}$ , there exists a partition of  $K_{2n}$  into *n* disjoint spanning trees that are orthogonal to  $\mathcal F$ . In that same paper, they were able to prove that for any 1-factorization  $F$  of  $K_{2n}$ , there exist at least two disjoint spanning trees that are orthogonal to  $\mathcal{F}$ . In [\[3](#page-12-0)], Krussel et al. proved that there were at least three such trees. Regarding the 1-factorization  $GK_{2n}$ , however, they proved a much stronger result, stated as follows.

**Theorem 1** [\[3](#page-12-0), Thm. 3] *If*  $2n - 1$  *is a prime number of the form* 8*m* + 7 *for some integer m, then there exists a full set of n disjoint spanning trees for*  $K_{2n}$  *that are orthogonal to GK*2*n.*

In this paper, we develop a technique for finding spanning tree decompositions that are orthogonal to rotational 1-factorizations of *K*2*n*. Applying our technique, we prove the following.

<span id="page-1-0"></span>**Theorem 2** *For any integer n* > 2*, there exists a full set of n disjoint spanning trees for*  $K_{2n}$  *that are orthogonal to*  $GK_{2n}$ *.* 

Because our methods exploit the rotational nature of  $GK_{2n}$ , we believe that there are similar applications to other rotational families of 1-factorizations. In Sect. [7,](#page-10-0) we prove Theorem [2](#page-1-0) and offer directions for further application of the results.

#### **2 Terminology for Edges in Rotational Subgraphs of** *K***2***<sup>n</sup>*

We will be considering many different subgraphs of the complete graph  $K_{2n}$ . We draw such graphs in **rotational form**, meaning vertices 1 through  $2n - 1$  are spaced evenly around a circle in clockwise manner, and vertex 2*n* is placed at the center. An example of a subgraph drawn in rotational form appears in Example [1](#page-1-1) above.

**Definition 1** Suppose  $K_{2n}$  is drawn in rotational form. For any edge  $\{a, b\}$ , if  $2n \notin$  ${a, b}$ , we define its **length** by<br>length ${a, b}$ 

length
$$
\{a, b\}
$$
 = min $\{|a - b|, 2n - 1 - |a - b|\}$ ,

and if  $2n \in \{a, b\}$ , we say the edge has length 0.

**Definition 2** Suppose  $K_{2n}$  is drawn in rotational form. For any edge  $\{a, b\}$  of nonzero length, we define its **center** (denoted center{*a*, *b*}) to be the unique vertex  $x \neq 2n$ such that

$$
length{a, x} = length{x, b}.
$$

For any edge {*a*, 2*n*} of length 0, we define the center to be *a*.

*Example 2* A subgraph *G* of  $K_8$  is drawn in rotational form. The length and center are given for each edge.



Note that, in rotational form, any edge of  $K_{2n}$  is uniquely determined by its center and length.

#### **3 Starter Graphs and Rotational Families**

To obtain a rotational decomposition of  $K_{2n}$ , we begin with starter graphs, which are graphs drawn in rotational form that contain one edge of each length.

**Definition 3** Fix any integer  $n > 0$  and any *n*-tuple of integers  $(c_0, \ldots, c_{n-1})$  satisfying  $0 < c_l < 2n$  for each  $l$   $(0 \le l < n)$ . We define the **starter graph**, denoted  $SG(c_0, \ldots, c_{n-1})$ , to be the subgraph of  $K_{2n}$  with a single edge of length *l* and center *c<sub>l</sub>*, for each  $l$  ( $0 \le l < n$ ).



*Example 3* The following are examples of starter graphs.

<span id="page-3-0"></span>**Definition 4** Suppose *G* is a subgraph of  $K_{2n}$  with edge set *E*. Let  $\rho$  denote the permutation of the vertex set that cyclically permutes  $(1, 2, 3, \ldots, 2n - 1)$  and fixes 2*n*. We define  $\rho(G)$  to be the subgraph of  $K_{2n}$  with edge set

$$
\rho(E) = \{ \{ \rho(a), \rho(b) \} \mid \{a, b\} \in E \}.
$$

Equivalently, vertices *a*, *b* are adjacent in  $\rho(G)$  if and only if the vertices  $\rho^{-1}(a)$ ,  $\rho^{-1}(b)$  are adjacent in *G*. Observe that, for any edge {*a*, *b*} and integer *i*, the edge  $\{\rho^i(a), \rho^i(b)\}\$  has the same length as  $\{a, b\}$  and has center  $c + i$ , where  $c$  is the center of  $\{a, b\}$ .

Using the above, we can now introduce the notion of a rotational family of subgraphs.

**Definition 5** Fix any integers *n*,  $d > 0$  and let *G* be any subgraph of  $K_{2n}$ . We define the **rotational family**  $\mathcal{F}^d_G$  **generated by**  $G$  to be the set

$$
\mathcal{F}_G^d := \{G, \rho(G), \ldots, \rho^{d-1}(G)\}
$$

<span id="page-3-1"></span>where  $\rho$  is the permutation given in Definition [4.](#page-3-0)

**Proposition 1** *Fix any integer n* > 0 *and any n-tuple of integers*  $(c_0, \ldots, c_{n-1})$ *satisfying*  $0 < c_l < 2n$  *for each l*  $(0 \le l < n)$ *. Let*  $G = SG(c_0, \ldots, c_{n-1})$ *. Then the following hold.*

- (i) *The rotational family*  $\mathcal{F}_G^{2n-1}$  *forms a decomposition of*  $K_{2n}$ *.*
- (ii) If G is a 1-factor, then  $\rho^i(G)$  is a 1-factor for every integer i, and the rotational  $f$ *family*  $\mathcal{F}_G^{2n-1}$  *forms a rotational 1-factorization of*  $K_{2n}$ *.*
- *Proof* (i). By way of contradiction, fix any  $l$  ( $0 \le l < n$ ) and suppose that for some *i*, *j*  $(0 \le i \le j \le 2n - 1)$ , the graphs  $\rho^{i}(G)$  and  $\rho^{j}(G)$  in  $\mathcal{F}_{G}^{2n-1}$  share an edge of length *l*. Then  $c_l + i \equiv c_l + j \pmod{2n - 1}$ , contradicting our choice of *i*, *j*. So the graphs in  $\mathcal{F}_G^{2n-1}$  are pairwise edge-disjoint and partition the edges of  $K_{2n}$ .
- (ii). By Definition [4,](#page-3-0) the map  $\rho$  is an isomorphism between *G* and  $\rho(G)$ . So every graph in  $\mathcal{F}_G^{2n-1}$  is isomorphic to *G*, and the result follows.

#### **4 Three Families of Rotational 1-Factorizations**

<span id="page-4-0"></span>We will illustrate our results using the following three families of rotational 1 factorizations, the first two of which are fairly common in the literature (see  $[5,8]$  $[5,8]$ ).

**Definition 6** Fix any integer  $n > 0$ . The 1-factorization  $G K_{2n}$  of  $K_{2n}$  is the rotational family  $\mathcal{F}_G^{2n-1}$  generated by  $G = SG(c_0, \ldots, c_{n-1})$  where

$$
c_l = 1 \text{ for all } l \ (0 \le l < n).
$$



**Definition 7** (Adapted from [\[8](#page-12-6)]) Fix any integer  $n > 0$ . The 1-factorization  $W K_{2n}$ of  $K_{2n}$  is the rotational family  $\mathcal{F}_G^{2n-1}$  generated by  $G = SG(c_0, \ldots, c_{n-1})$  where  $c_i$ is given by the following table, depending on the form of *n* and *l*. (When appropriate, entries are to be read modulo  $2n - 1$ .)





**Definition 8** Fix any integer  $n > 0$ . The 1-factorization  $HK_{2n}$  of  $K_{2n}$  is the rotational family  $\mathcal{F}_G^{2n-1}$  generated by the graph  $G = SG(c_0, \ldots, c_{n-1})$  where  $c_l$  is given by the following table, depending on the form of *n* and *l*:





## **5 Opposing Pair Graphs and Their Rotations**

To construct spanning trees that are orthogonal to a given rotational 1-factorization, we use rotational families of graphs that are built using opposing pairs of edges. We begin this section by defining these terms.

**Definition 9** Suppose  $K_{2n}$  is drawn in rotational form and let  $e_1$ ,  $e_2$  be any pair of edges with centers  $c_1$ ,  $c_2$ . We define the **distance** between them to be

$$
dist(e_1, e_2) = length{c_1, c_2},
$$

and edges at distance *n* − 1 are said to be **opposing**. Observe that, for any edge pair  ${e_1, e_2}$  and integer *i*, we have dist( $\rho^i(e_1), \rho^i(e_2)$ ) = dist( $e_1, e_2$ ).

**Definition 10** Suppose  $K_{2n}$  is drawn in rotational form and let  $e_1$ ,  $e_2$  be any pair of edges with centers  $c_1$ ,  $c_2$ . If  $c_1 \neq c_2$ , we define the **direction** of the pair  $e_1$ ,  $e_2$  to be the vertex

$$
\operatorname{dir}(e_1, e_2) = \operatorname{center}\{c_1, c_2\}.
$$

For the case  $c_1 = c_2$ , we set dir( $e_1, e_2$ ) =  $c_1$ . Observe that, for any edge pair  $\{e_1, e_2\}$ and integer *i*, we have  $\text{dir}(\rho^{i}(e_1), \rho^{i}(e_2)) = \text{dir}(e_1, e_2) + i$ .

*Example 4* A subgraph *G* of  $K_8$  is drawn in rotational form below:



With this terminology in place, we can now define the graphs of interest.

**Definition 11** Fix any integers  $n > t \geq 0$  and any *n*-tuple of integers  $(d_0, \ldots, d_{n-1})$ satisfying  $0 < d_i < 2n$  for each  $i \ (0 \leq i < n)$ . We define the **opposing pair graph**, denoted  $OPG_t(d_0, \ldots, d_{n-1})$ , to be the subgraph of  $K_{2n}$  with a single edge of length *t* and center  $d_t$ , and, for each  $i \neq t$ , an opposing pair of edges of length *i* and direction *di* . We refer to *t* as the **exceptional length** of the OPG.

<span id="page-6-0"></span>*Example 5* The following are examples of opposing pair graphs:



<span id="page-6-1"></span>Let us next observe that any opposing pair graph can be expressed as a union of a pair of starter graphs.

**Proposition 2** *Fix any integers n* >  $t \ge 0$  *and any n-tuple of integers*  $(d_0, \ldots, d_{n-1})$ *satisfying*  $0 < d_i < 2n$  *for each i*  $(0 \le i < n)$ *. Let*  $G = OPG_t(d_0, \ldots, d_{n-1})$ *. Then* 

$$
G = SG(a_0,\ldots,a_{n-1}) \cup SG(b_0,\ldots,b_{n-1}),
$$

*where*  $a_i, b_i \equiv d_i \pm |n/2|$  (*mod*  $2n - 1$ *), respectively, for each i*  $\neq$  *t, and where*  $a_t = b_t = d_t$ .

*Proof* Let  $X^-$ ,  $X^+$  be the two starter graphs in the union above. For each  $i \neq t$ , the graph  $X^-$  has a single edge of length *i* and center  $a_i$ , and  $X^+$  has a single edge of length *i* and center  $b_i$ . These two edges form an opposing pair of length *i* and direction *d<sub>i</sub>*, as desired. When *i* = *t*, both graphs  $X^-$  and  $X^+$  share a single edge of length *t* and center *d*. So  $X^- \cup X^+$  gives the desired opposing pair graph *G* and center  $d_t$ . So  $X^- \cup X^+$  gives the desired opposing pair graph G.

Notice that any opposing pair graph in  $K_{2n}$  has exactly  $2n - 1$  edges, which is the same as the number of edges required for a spanning tree of  $K_{2n}$ . Indeed, the graph *OPG*<sub>4</sub>(1, 2, 9, 3, 6), which is the left-most graph depicted in Example [5,](#page-6-0) is a spanning tree for  $K_{10}$ . In general, however, an opposing pair graph need not be acyclic. By a standard result about spanning trees [\[7](#page-12-7), p. 68], if *S* is *any* set of  $2n - 1$  edges in  $K_{2n}$ , then *S* will be a spanning tree for  $K_{2n}$  iff *S* is acyclic and iff *S* forms a connected graph on the vertex set.

If we find an opposing pair graph that forms a spanning tree for  $K_{2n}$ , then it is useful to note that each of its rotations will be spanning trees as well. But unlike the starter graphs in Proposition [1,](#page-3-1) the rotations of an opposing pair graph do not form a decomposition of  $K_{2n}$ . If we stop rotating in time, however, we can come fairly close, as the following result indicates.

<span id="page-7-0"></span>**Proposition 3** *Fix any integers n* >  $t \ge 0$  *and any n-tuple of integers*  $(d_0, \ldots, d_{n-1})$ *satisfying*  $0 < d_i < 2n$  *for each i*  $(0 \le i < n)$ *. Let*  $G = OPG_t(d_0, \ldots, d_{n-1})$ *. Then the following hold.*

- (i) *The graphs in the rotational family*  $\mathcal{F}_G^{n-1}$  *are pairwise edge-disjoint.*
- (ii) *The set of edges in*  $K_{2n}$  *that are not contained in any graph of*  $\mathcal{F}_G^{n-1}$  *form a subgraph with exactly* 2*n* − 1 *edges.*
- (iii) If G is a spanning tree for  $K_{2n}$ , then every rotation of G is a spanning tree for *K*<sub>2*n*</sub>. In particular, every graph in the rotational family  $\mathcal{F}_G^{n-1}$  is a spanning tree *for*  $K_{2n}$ *.*
- *Proof* (i). By way of contradiction, suppose that for some  $0 \le i \le j \le n 1$ , the graphs  $\rho^{i}(G)$  and  $\rho^{j}(G)$  share an edge of length  $x \neq t$ . Then by Proposition [2,](#page-6-1)

$$
d_x \pm \lfloor n/2 \rfloor + i \equiv d_x \pm \lfloor n/2 \rfloor + j \pmod{2n-1}.
$$

But then  $j - i \equiv 0, \pm (n - 1)$ , which is impossible for these values of *i*, *j*. Similarly, if  $\rho^{i}(G)$  and  $\rho^{j}(G)$  share an edge of length *t*, then  $d_{t} + i \equiv d_{t} + j$  contradicting our choice of *i*, *j*.

(ii). Each of the *n* − 1 graphs in  $\mathcal{F}_G^{n-1}$  has  $2n - 1$  edges, so by (i), their union has  $(n-1)(2n-1)$  edges. Since  $K_{2n}$  has  $n(2n-1)$  edges, the result follows.

(iii). By Definition [4,](#page-3-0) the map  $\rho$  is an isomorphism between *G* and  $\rho(G)$ . So every graph in  $\mathcal{F}_G^{n-1}$  is isomorphic to *G*, and the result follows.

The graph with  $2n - 1$  edges, mentioned in part (ii) of the above proposition, deserves special attention. We make the following definition. -

**Definition 12** Fix any integers  $n > t \ge 0$  and any *n*-tuple of integers  $(d_0, \ldots, d_{n-1})$ satisfying  $0 < d_i < 2n$  for each  $i$  ( $0 \le i < n$ ). Let  $G = OPG_t(d_0, ..., d_{n-1})$ . We define the graph *G* by

$$
\widetilde{G}=K_{2n}\backslash\left(\bigcup\nolimits_{H\in\mathcal{F}_G^{n-1}}H\right).
$$

In other words,  $G$  is the complementary graph in  $K_{2n}$  of the union of the rotational family  $\mathcal{F}_G^{n-1}$ .

<span id="page-8-0"></span>*Example 6* The complementary graphs *G* for the graphs of Example [5:](#page-6-0)



Note that when  $G = OPG_t(d_0, \ldots, d_{n-1})$ , the graph *G* has exactly *n* edges of length *t* and a single edge of each length  $i \neq t$  (see Example [6\)](#page-8-0). Occasionally it happens that both *G* and *G* are spanning trees for *K*2*n*. When this occurs, *G* is of great use in constructing an orthogonal spanning tree decomposition.

#### **6 Orthogonality and Opposing Pair Graphs**

<span id="page-8-1"></span>**Theorem 3** Fix any integers  $n > t \geq 0$  and fix any two n-tuples of integers  $(c_0, ..., c_{n-1})$  *and*  $(d_0, ..., d_{n-1})$  *satisfying* 0 < *c<sub>i</sub>*, *d<sub>i</sub>* < 2*n for each i* (0 ≤ *i* < *n*). *Let*  $S = SG(c_0, \ldots, c_{n-1})$  *and*  $G = OPG_t(d_0, \ldots, d_{n-1})$ *. Suppose S* is a 1-factor. *Then the following hold.*

- (i) *If* G is orthogonal to the 1-factorization  $\mathcal{F}_S^{2n-1}$ , then every rotation of G is *orthogonal to*  $\mathcal{F}_{S}^{2n-1}$ *. In particular, every graph in the rotational family*  $\mathcal{F}_{G}^{n-1}$  *is orthogonal to*  $\mathcal{F}_S^{2n-1}$ . *s*  $\mathcal{F}_{S}^{2n-1}$ , then every rotation of G is graph in the rotational family  $\mathcal{F}_{G}^{n-1}$  is  $S^{2n-1}$ , then the graph  $\widetilde{G}$  is also orthog-
- (ii) *If G is orthogonal to the 1-factorization*  $\mathcal{F}_{S}^{2n-1}$ , then the graph *onal to*  $\mathcal{F}_{S}^{2n-1}$ .
- *Proof* (i) If  $\rho(G)$  is not orthogonal to  $\mathcal{F}_{S}^{2n-1}$ , there exists an integer *j* such that  $\rho^{j}(S) \in \mathcal{F}_{S}^{2n-1}$  and where  $\rho^{j}(S)$  shares more than one edge with  $\rho(G)$ . But then by Definition [4,](#page-3-0) *G* shares more than one edge with  $\rho^{j-1}(S) \in \mathcal{F}_S^{2n-1}$ , so that *G* is not orthogonal to  $\mathcal{F}_S^{2n-1}$ . By contrapositive,  $\rho(G)$  is orthogonal whenever *G* is, so by repeated application of  $\rho$ , the result follows.
- (ii) By (i), each of the  $n-1$  rotations of *G* in  $\mathcal{F}_G^{n-1}$  shares one edge with each of the 2*n* − 1 rotations of *S* in the 1-factorization  $\mathcal{F}_{S}^{2n-1}$ . By Proposition [3\(](#page-7-0)i), these rotations of *G* are edge-disjoint. This leaves exactly one edge from each 1-factor in  $\mathcal{F}_{S}^{2n-1}$  as the set of edges of  $\widetilde{G}$ , as desired. □

Our next goal is to characterize which opposing pair graphs  $OPG_t(d_0, \ldots, d_{n-1})$ are orthogonal to the 1-factorization generated by a given starter  $SG(c_0, \ldots, c_{n-1})$ . To obtain the characterization, we begin with the following lemma, which concerns a partition of a set of integers.

<span id="page-9-1"></span>**Lemma 1** *Fix any integer n*  $\geq 2$  *and subsets A, B*  $\subseteq$  {1, 2, ..., 2*n* − 1} *such that*  $|A| = |B| = n - 1$  *and where* 

$$
B \equiv \{a + n \mid a \in A\} \pmod{2n - 1}.
$$

*Then the following are equivalent.*

(i)  $A \cup B = \{1, 2, \ldots, 2n - 2\}.$ 

(ii)  $A = \{n, n+1, \ldots, 2n-2\}$  *and*  $B = \{1, 2, \ldots, n-1\}$ .

*Proof* Since (ii) clearly implies (i), it remains to consider the converse. Suppose (i) holds. Then  $0 \notin A$  and it suffices to show that, for any  $i$  ( $0 \le i \le n - 2$ ), if  $i \notin A$ , then  $i + 1 \notin A$ . To this end, suppose  $i \notin A$ . Then  $i + n \notin B$ , so  $i + n \in A$ , forcing  $i + 2n \in B$ . But  $i + 2n \equiv i + 1$ , so  $i + 1 \notin A$ . *i* + 2*n* ∈ *B*. But *i* + 2*n* ≡ *i* + 1, so *i* + 1 ∉ *A*.

With the above lemma in place, we now characterize the opposing pair graphs that are orthogonal to a given rotational 1-factorization. Without loss of generality, we may restrict our attention to opposing pair graphs that satisfy  $d_t = c_t$ , in view of Theorem [3\(](#page-8-1)i) above.

<span id="page-9-2"></span>**Theorem 4** *Fix any integers*  $n > t \geq 0$  *and fix any two n-tuples of integers*  $(c_0, ..., c_{n-1})$  *and*  $(d_0, ..., d_{n-1})$  *satisfying* 0 < *c<sub>i</sub>*, *d<sub>i</sub>* < 2*n for each i* (0 ≤ *i* < *n*)*. Let*  $S = SG(c_0, \ldots, c_{n-1})$  *and*  $G = OPG_t(d_0, \ldots, d_{n-1})$ *. Suppose S* is a 1-factor *and suppose*  $d_t = c_t$ *. Then the graph G is orthogonal to the 1-factorization*  $\mathcal{F}_S^{2n-1}$  *if and only if*

$$
\{d_i-c_i+(-1)^n\lfloor n/2\rfloor\}_{i\neq i}\equiv \{1,2,\ldots,n-1\}\ \ (\mathrm{mod}\ 2n-1).
$$

*Proof* First note that an edge of length *x* and center *y* is in  $\rho^{i}(S)$  if and only if

<span id="page-9-0"></span>
$$
y - c_x \equiv i \pmod{2n - 1}.
$$
 (1)

The graph *G* has a single edge of length *t* with center  $d_t = c_t$ , and so it shares this edge with the starting 1-factor  $S = \rho^0(S)$ . By Proposition [2,](#page-6-1) for each length  $l \neq t$ , *G* has a pair of edges of length *l* and centers  $d_l \pm |n/2|$ . Now *G* is orthogonal to  $\mathcal{F}_{S}^{2n-1}$  if and only if it shares exactly 1 edge with each of the other rotations  $\rho^{i}(S)$  $(1 \le i \le 2n - 2)$  in  $\mathcal{F}_{S}^{2n-1}$ . By [\(1\)](#page-9-0), this occurs if and only if

$$
A \cup B = \{1, 2, \ldots, 2n - 2\}
$$

where *A*, *B* are subsets of  $\{1, 2, ..., 2n - 1\}$  satisfying  $A = \{d_i - c_i - |n/2|\}_{i \neq t}$  and  $B \equiv \{d_i - c_i + \lfloor n/2 \rfloor\}_{i \neq t} \pmod{2n-1}.$ 

When *n* is even,  $B \equiv \{a + n \mid a \in A\}$  (mod  $2n - 1$ ). So by Lemma [1,](#page-9-1) *G* is orthogonal to  $\mathcal{F}_{S}^{2n-1}$  if and only if  $B = \{1, 2, ..., n-1\}$  as desired. When *n* is odd,  $A \equiv \{b + n \mid b \in B\}$  (mod 2*n* − 1). So by swapping the roles of *A* and *B* in Lemma [1,](#page-9-1) *G* is orthogonal to  $\mathcal{F}_{S}^{2n-1}$  if and only if *A* = {1, 2, ..., *n* − 1} as desired.

Since there are exactly  $(n - 1)!$  ways to pair up the elements of the sets appearing on the two sides of the congruence in Theorem [4,](#page-9-2) we have the following corollary.

**Corollary 1** *Fix any integers*  $n > t \ge 0$  *and fix any two n-tuples of integers*  $(c_0, \ldots, c_{n-1})$  *and*  $(d_0, \ldots, d_{n-1})$  *satisfying*  $0 < c_i$ ,  $d_i < 2n$  *for each i*  $(0 \le i < n)$ *. Let*  $S = SG(c_0, \ldots, c_{n-1})$  *and*  $G = OPG_t(d_0, \ldots, d_{n-1})$ *. Suppose* S is a 1-factor. *Then there are exactly* (*n* − 1)! *different opposing pair graphs orthogonal to the 1 factorization*  $\mathcal{F}_S^{2n-1}$  *that have a single edge of length t and center*  $d_t = c_t$ *.* 

# <span id="page-10-0"></span>**7 Application to Rotational 1-Factorizations** -

<span id="page-10-1"></span>**Definition 13** Fix any integer  $n > 2$ . We define  $DGK_{2n}$  to be the spanning tree decomposition of  $K_{2n}$  given by  $\mathcal{F}_G^{n-1} \cup {\widetilde{G}}$ , where  $\mathcal{F}_G^{n-1}$  is the rotational family generated by the graph  $G = OPG_{n-1}(d_0, \ldots, d_{n-1})$ , and where  $d_l$  is given by the following table, depending on the form of *n* and *l*. (When appropriate, entries are to be read modulo  $2n - 1$ .)





 $OPG_{10}(11, 12, 10, 13, 9, 14, 8, 15, 7, 16, 1)$ 

*Proof of Theorem* [2](#page-1-0) Using the criteria given in Theorems [4](#page-9-2) and [3,](#page-8-1) it is easily verified that the spanning tree decomposition  $DGK_{2n}$  of Definition [13](#page-10-1) is orthogonal to the 1factorization  $GK_{2n}$  given in Definition [6.](#page-4-0) We also observe that each graph in  $D GK_{2n}$ has exactly  $2n - 1$  edges. So to prove that these graphs are trees, it suffices to check that they are connected which is straightforward using Definition 13 that they are connected, which is straightforward using Definition [13.](#page-10-1)

We suspect that similar families of spanning tree decompositions can be found that are orthogonal to other families of 1-factorizations, as the following result suggests.

**Proposition 4** *For every integer n* (2 < *n* < 11)*, there exist orthogonal spanning tree decompositions of K*2*<sup>n</sup> for the 1-factorizations W K*2*<sup>n</sup> and H K*2*n.* -

*Proof* For each value of *n*, the table below gives an opposing pair graph *G* that is a spanning tree  $K_{2n}$  and that is orthogonal to  $WK_{2n}$  (respectively  $HK_{2n}$ ). For each of these, the complementary graph *G* is also a tree. Accordingly, there is a spanning tree decomposition of  $K_{2n}$  given by  $\mathcal{F}_G^{n-1} \cup \{ \widetilde{G} \}$ , where  $\mathcal{F}_G^{n-1}$  is the rotational family generated by the graph *G*, and where each tree in the decomposition is orthogonal to  $W K_{2n}$  (respectively  $HK_{2n}$ ), as desired.



### **References**

- <span id="page-12-5"></span>1. Brualdi, R.A., Hollingsworth, S.: Multicolored trees in complete graphs. J. Combin. Theory Ser. B **68**, 310–313 (1996)
- <span id="page-12-3"></span>2. Cameron, P.J.: Parallelisms of Complete Designs. London Math. Soc., Lecture Note Series, vol. 23. Cambridge Univ. Press, Cambridge (1976)
- <span id="page-12-0"></span>3. Krussel, J., Marshall, S., Verrall, H.: Spanning trees orthogonal to 1-factorizations of *K*2*n*. ARS Comb. **57**, 77–82 (2000)
- <span id="page-12-4"></span>4. Mazzuoccolo, G., Rinaldi, G.: *k*-pyramidal one-factorizations. Graphs Combin. **23**, 315–326 (2007)
- <span id="page-12-1"></span>5. Mendelsohn, E., Rosa, A.: One-factorizations of the complete graph - a survey. J. Graph Theory **9**, 43–65 (1985)
- <span id="page-12-2"></span>6. Wallis, W.D.: One-Factorizations, volume 390 of Mathematics and Its Applications. Springer (1997)
- <span id="page-12-7"></span>7. West, D.B.: Introduction to Graph Theory, 2nd edn. Prentice-Hall, Upper Saddle River (2001)
- <span id="page-12-6"></span>8. Wolff, K.E.: Fast-blockplaene. Mitt. math. Sem. Giessen **102**, 72–73 (1973)