

## Tight Lower Bounds on the Size of a Maximum Matching in a Regular Graph

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**Abstract.** In this paper we study tight lower bounds on the size of a maximum matching in a regular graph. For  $k \geq 3$ , let  $G$  be a connected  $k$ -regular graph of order  $n$  and let  $\alpha'(G)$  be the size of a maximum matching in  $G$ . We show that if  $k$  is even, then  $\alpha'(G) \geq \min \left\{ \left( \frac{k^2+4}{k^2+k+2} \right) \times \frac{n}{2}, \frac{n-1}{2} \right\}$ , while if  $k$  is odd, then  $\alpha'(G) \geq \frac{(k^3-k^2-2)n-2k+2}{2(k^3-3k)}$ . We show that both bounds are tight.

**Key words.** Lower bounds, matching number, regular graph.

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### 1. Introduction

Two edges in a graph  $G$  are *independent* if they are not adjacent in  $G$ . A set of pairwise independent edges of  $G$  is called a *matching* in  $G$ , while a matching of maximum cardinality is a *maximum matching*. The number of edges in a maximum matching of  $G$  is called the *matching number* of  $G$  which we denote by  $\alpha'(G)$ . In this paper we study tight lower bounds on the size of a maximum matching in a connected regular graph, that is, in a graph in which every vertex has the same degree. Matchings in graphs are extensively studied in the literature (see, for example, the survey articles by Plummer [5] and Pulleyblank [6]). For a graph  $G$  and a set  $S \subseteq V(G)$ , the subgraph induced by  $S$  is denoted by  $G[S]$ .

For  $k \geq 2$ , let  $G$  be a connected  $k$ -regular graph of order  $n$ . If  $k = 2$ , then  $G$  is a cycle  $C_n$  and  $\alpha'(G) \geq (n-1)/2$  with this bound achieved when  $G$  is an odd cycle. Hence in what follows, we assume that  $k \geq 3$ . When  $k = 3$ , Biedl et al. [2] proved that  $\alpha'(G) \geq (4n-1)/9$ . When  $k = 4$  and  $n \geq 6$ , Lichiardopol [4] has shown that  $\alpha'(G) \geq 7n/17$ , and if  $k \geq 5$ , then  $\alpha'(G) \geq ((3k-6)n)/(7k-13)$ . In this paper, we generalize the result of Biedl et al. [2] when  $k = 3$  to all values  $k \geq 3$ . Our results improve those of Lichiardopol [4].

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### 2. Preliminary Results

We need the following result of Berge [1] about the matching number of a graph, which is sometimes referred to as the Tutte-Berge formulation for the matching number.

**Theorem 1 (Berge [1]).** *For every graph  $G$ ,*

$$\alpha'(G) = \min_{X \subseteq V(G)} \frac{1}{2} (|V(G)| + |X| - \text{oc}(G - X)),$$

where  $\text{oc}(G - X)$  denotes the number of odd components of  $G - X$ .

We shall also need the following two observations.

**Observation 1.** *Every graph has an even number of vertices of odd degree.*

**Observation 2.** *Let  $k > 1$  and let  $G$  be a graph with  $\Delta(G) \leq k$ . If*

$$\sum_{x \in V(G)} (k - d(x)) < k,$$

then  $|V(G)| \geq k + 1$ .

*Proof.* Let  $G$  be defined as in the observation and let  $p = \sum_{x \in V(G)} (k - d(x))$  and  $n = |V(G)|$ . As the maximum degree in  $G$  is at most  $n - 1$  we note that  $p \geq n(k - (n - 1))$ , which implies that  $n(k + 1 - n) < k$ . As this is equivalent to  $0 < (n - 1)(n - k)$  we note that  $n > k$ , which implies the observation.  $\square$

### 3. Main Results

We shall show that:

**Theorem 2.** *For  $k \geq 2$  even, if  $G$  is a connected  $k$ -regular graph of order  $n$ , then*

$$\alpha'(G) \geq \min \left\{ \left( \frac{k^2 + 4}{k^2 + k + 2} \right) \times \frac{n}{2}, \frac{n - 1}{2} \right\},$$

and this bound is tight.

**Theorem 3.** *For  $k \geq 3$  odd, if  $G$  is a connected  $k$ -regular graph of order  $n$ , then*

$$\alpha'(G) \geq \frac{(k^3 - k^2 - 2)n - 2k + 2}{2(k^3 - 3k)},$$

and this bound is tight.

**4. Proof of Theorem 2**

First we present a proof of the lower bound of Theorem 2. If  $k = 2$ , then  $\alpha'(G) = \alpha'(C_n) \geq (n - 1)/2$ , and the theorem holds. Hence we may assume that  $k \geq 4$ . Let  $X$  be a set of vertices in  $G$  such that  $(n + |X| - \text{oc}(G - X))/2$  is minimum and let  $M$  be a maximum matching in  $G$ . By Theorem 1 we note that  $|M| = (n + |X| - \text{oc}(G - X))/2$ . If  $X = \emptyset$ , then  $|M| \geq (n - 1)/2$ , and we are done, so assume that  $X \neq \emptyset$ .

By Observation 1, we note that there is no odd component in  $G - X$  with exactly one edge into  $X$  in  $G$ , as it would have only one vertex of odd degree in  $G - X$ . Therefore every odd component has at least two edges into  $X$  in  $G$ . Let  $y_k$  denote the number of odd components in  $G - X$  that have at least  $k$  edges into  $X$  in  $G$  and let  $y_2$  denote the number of odd components in  $G - X$  that have less than  $k$  edge into  $X$  in  $G$ . If there are  $d(X, V - X)$  edges between  $X$  and  $V(G) - X$ , then we obtain the following, by the above,

$$k|X| \geq d(X, V - X) \geq 2y_2 + ky_k.$$

Hence,  $|X| \geq y_k + 2y_2/k \geq 2y_2/k$ . By Observation 2, we note that any odd component in  $G - X$  with less than  $k$  edges into  $X$  in  $G$  must contain at least  $k + 1$  vertices. This implies that  $n \geq |X| + y_2(k + 1)$ . We therefore obtain the following.

$$\begin{aligned} |M| &= \frac{1}{2}(n + |X| - \text{oc}(G - X)) \\ &= \left(\frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2}\right) + \left(\frac{k - 2}{k^2 + k + 2} \times \frac{n}{2}\right) + \left(\frac{|X| - (y_2 + y_k)}{2}\right) \\ &\geq \left(\frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2}\right) + \left(\frac{k - 2}{k^2 + k + 2} \times \frac{|X| + y_2(k + 1)}{2}\right) + \left(\frac{\frac{2y_2}{k} - y_2}{2}\right) \\ &\geq \left(\frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2}\right) + \left(\frac{k - 2}{k^2 + k + 2} \times \frac{\frac{2y_2}{k} + y_2(k + 1)}{2}\right) - y_2 \left(\frac{k - 2}{2k}\right) \\ &= \left(\frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2}\right) + \left(\frac{y_2(k - 2)}{2k(k^2 + k + 2)} \times (k^2 + k + 2 - (k^2 + k + 2))\right) \\ &= \frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2}. \end{aligned}$$

This establishes the lower bound of Theorem 2. The following proposition shows that the lower bound of Theorem 2 is tight.

**Proposition 4.** *For every integer  $p \geq 2$  and every even integer  $k \geq 4$ , there exists a connected  $k$ -regular graph  $G_p^k$  of order  $p(k^2 + k + 2)/2$  satisfying*

$$\alpha'(G_p^k) = \left(\frac{k^2 + 4}{k^2 + k + 2}\right) \times \frac{|V(G_p^k)|}{2}.$$

*Proof.* Let  $G_p^k$  be the connected  $k$ -regular graph of order  $p(k^2 + k + 2)/2$  defined as follows. Let  $X_p$  be any  $k$ -regular connected multigraph of order  $p$  (i.e., we allow multiple edges). Note that  $|E(X_p)| = kp/2$ . For every edge  $uv \in E(X_p)$ , add a complete graph,  $C_{uv}$ , on  $k + 1$  new vertices, delete the edge  $uv$  and any edge  $xy \in E(C_{uv})$ , and then add the edges  $ux$  and  $vy$ . Once we have done this for all edges in  $X_p$ , we note that the resulting graph,  $G_p^k$ , is a connected  $k$ -regular graph of order  $|V(G_p^k)| = |V(X_p)| + (k+1)kp/2 = p(1+k(k+1)/2) = p(k^2+k+2)/2$ . Furthermore by deleting the vertices  $V(X_p)$  from  $G_p^k$  we obtain  $kp/2$  odd components (namely all the  $C_{uv}$ 's). Therefore, by Theorem 1,

$$\begin{aligned} \alpha'(G_p^k) &\leq \frac{1}{2} \left( |V(G_p^k)| + |V(X_p)| - \text{oc}(G_p^k - V(X_p)) \right) \\ &= \frac{1}{2} \left( \frac{1}{2} p(k^2 + k + 2) + p - \frac{1}{2} kp \right) \\ &= \frac{p}{4} \left( (k^2 + k + 2) + 2 - k \right) \\ &= \left( \frac{k^2 + 4}{k^2 + k + 2} \right) \times \frac{p}{4} (k^2 + k + 2) \\ &= \left( \frac{k^2 + 4}{k^2 + k + 2} \right) \times \frac{|V(G_p^k)|}{2}. \end{aligned}$$

However the lower bound of Theorem 2 shows that  $\alpha'(G_p^k) \geq \left( \frac{k^2+4}{k^2+k+2} \right) \times \frac{|V(G_p^k)|}{2}$ , and so the desired result of the proposition follows.  $\square$

We remark that the  $(n - 1)/2$  bound in the statement of Theorem 2 is only included as it is necessary when  $n$  is very small or  $k = 2$ . As an immediate consequence of Theorem 2, we have the following slightly weaker result.

**Corollary 1.** *For  $k \geq 4$  even, if  $G$  is a connected  $k$ -regular graph of order  $n$ , then*

$$\alpha'(G) \geq \left\lfloor \left( \frac{k^2 + 4}{k^2 + k + 2} \right) \times \frac{n}{2} \right\rfloor,$$

*and this bound is tight.*

*Proof.* Observe that if  $n \geq k + 3 + \frac{8}{k - 2}$ , then

$$\frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2} \leq \frac{n - 1}{2},$$

while if  $n < k + 3 + \frac{8}{k - 2}$ , then  $\frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2} - \frac{n - 1}{2} \in \left( 0, \frac{1}{2} \right)$ , whence

$$\left\lfloor \frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2} \right\rfloor = \left\lfloor \frac{n - 1}{2} \right\rfloor.$$

Hence for all  $k \geq 4$  and  $n \geq k + 1$ , we have that

$$\left\lfloor \frac{k^2 + 4}{k^2 + k + 2} \times \frac{n}{2} \right\rfloor \leq \left\lfloor \frac{n - 1}{2} \right\rfloor.$$

The desired result now follows from Theorem 2. □

### 5. Proof of Theorem 3

In order to prove Theorem 3, we shall need the following definition.

**Definition 5.1.** Let  $k \geq 3$  be an odd integer and let  $H_{k+2}$  be the graph with vertex set  $V(H_{k+2}) = \{w_1, w_2, \dots, w_{k+2}\}$  and containing all possible edges except the following:

$$\left( \bigcup_{i=0}^{(k-3)/2} \{w_{2i+1}w_{2i+2}\} \right) \cup \{w_k w_{k+2}, w_{k+1}w_{k+2}\}.$$

Note that the degree of every vertex in  $H_{k+2}$  is  $k$ , except for  $w_{k+2}$  which has degree  $k - 1$ .

We are now in a position to prove Theorem 3.

**Proof of Theorem 3.** Let

$$\phi(k, n) = \frac{(k^3 - k^2 - 2)n - 2k + 2}{2(k^3 - 3k)}.$$

We wish to show that  $\alpha'(G) \geq \phi(k, n)$ . For a subset  $X \subset V(G)$ , we define

$$M_X = \frac{1}{2} (n + |X| - \text{oc}(G - X)).$$

We proceed further with the following claim.

*Claim.* If  $X$  is an independent set of vertices in  $G$  such that every component  $H$  in  $G - X$  has odd order and the number of edges between  $X$  and  $V(H)$  is either one or  $k$ , then  $M_X \geq \phi(k, n)$ .

*Proof.* Let  $y_k$  denote the number of components in  $G - X$  that have  $k$  edges into  $X$  in  $G$  and let  $V_k$  denote the set of all vertices that lie in such components. Let  $y_1$  denote the number of components in  $G - X$  that have one edge into  $X$  in  $G$  and let  $V_1$  denote the set of all vertices that lie in such components. Note that the following holds.

- (1)  $|X| = y_k + y_1/k$ . This is the case as if there are  $d(X, V - X)$  edges between  $X$  and  $V(G) - X$ , then we note that  $k|X| = d(X, V - X) = ky_k + y_1$ .
- (2)  $n \geq |X| + y_k + (k + 2)y_1$ . This follows from  $|V_k| \geq y_k$  and from Observations 1 and 2, which imply that  $|V_1| \geq (k + 2)y_1$ .

(3)  $y_k \geq \frac{y_1 - k}{k(k-2)}$ . Note that  $G - V_k$  has exactly  $|X|$  distinct components. By adding to  $G - V_k$  a component from  $G[V_k]$  and the resulting  $k$  edges from that component into  $X$ , we decrease the number of components by at most  $k - 1$ . As  $G$  is connected we must therefore have  $y_k(k - 1) \geq |X| - 1$ . Using (1), we get  $y_k(k - 1) \geq y_k + y_1/k - 1$ , and so solving for  $y_k$  implies (3).

By (1) we note that if  $y_1 = 0$ , then  $M_X = n/2$  which implies that the claim holds in this case. Hence we may assume that  $y_1 \geq 1$ . By (1), we can bound  $M_X$  as follows:

$$\begin{aligned} M_X &= \frac{1}{2} (n + |X| - oc(G - X)) \\ &\geq \frac{1}{2} \left( n + \frac{y_1}{k} + y_k - y_1 - y_k \right) \\ &= \frac{1}{2} \left( n + \frac{y_1(1 - k)}{k} \right). \end{aligned}$$

Now,

$$\begin{aligned} \frac{1}{2} \left( n + \frac{y_1(1 - k)}{k} \right) &\geq \frac{(k^3 - k^2 - 2)n - 2k + 2}{2(k^3 - 3k)} \\ \Downarrow \\ \frac{y_1(1 - k)}{k} + \frac{2(k - 1)}{k(k^2 - 3)} &\geq \left( \frac{(k^3 - k^2 - 2) - (k^3 - 3k)}{k(k^2 - 3)} \right) n \\ \Downarrow \\ \frac{y_1(1 - k)(k^2 - 3) + 2(k - 1)}{k(k^2 - 3)} &\geq \left( \frac{-k^2 + 3k - 2}{k(k^2 - 3)} \right) n \\ \Downarrow \\ -y_1(k - 1)(k^2 - 3) + 2(k - 1) &\geq -(k - 1)(k - 2)n \\ \Downarrow \\ n(k - 2) &\geq y_1(k^2 - 3) - 2 \\ \Downarrow \\ n &\geq \frac{y_1(k^2 - 3)}{k - 2} - \frac{2}{k - 2}. \end{aligned}$$

However, by (1), (2) and (3) we get the following bound on  $n$ .

$$\begin{aligned} n &\geq |X| + y_k + (k + 2)y_1 \\ &\geq \left( y_k + \frac{y_1}{k} \right) + y_k + (k + 2)y_1 \\ &\geq \frac{2(y_1 - k)}{k(k - 2)} + \frac{y_1}{k} + (k + 2)y_1 \\ &= y_1 \times \frac{2 + (k - 2) + k(k - 2)(k + 2)}{k(k - 2)} - \frac{2k}{k(k - 2)} \\ &= \frac{y_1(k^2 - 3)}{k - 2} - \frac{2}{k - 2}. \end{aligned}$$

Hence we have established that

$$M_X \geq \frac{1}{2} \left( n + \frac{y_1(1-k)}{k} \right) \geq \frac{(k^3 - k^2 - 2)n - 2k + 2}{2(k^3 - 3k)} = \phi(k, n),$$

as desired. □

We now return to the proof of Theorem 3. Let  $X$  be a set of vertices in  $G$  such that  $M_X$  is minimum. By Theorem 1,  $\alpha'(G) = M_X$ . We wish to show that  $M_X \geq \phi(k, n)$ . Suppose, to the contrary, that  $M_X < \phi(k, n)$ . Then,  $X$  does not satisfy the conditions of Claim 5. Hence we may assume that at least one of the following conditions holds.

*Condition (a).*  $X$  is not independent.

*Condition (b).* There are even components in  $G - X$ .

*Condition (c).* There are odd components in  $G - X$  which do not have one or  $k$  edges into  $X$ .

Now define a graph  $Z_X$  and integers  $z_X$ ,  $p_X$ , and  $r_X$  as follows.

- If  $G - X$  has odd components which do not have one or  $k$  edges into  $X$  or even components, then let  $Z_X$  be the graph consisting of all such components. Furthermore, let

$$z_X = |V(Z_X)|.$$

- For each  $x \in X$ , let  $p_x$  denote the degree of  $x$  in  $G[V(Z_X) \cup X]$ ; that is,  $p_x$  is the number of vertices in  $V(Z_X) \cup X$  adjacent to  $x$  in  $G$ . Let

$$p_X = \sum_{x \in X} p_x.$$

- Let  $r_X$  denote the number of components in  $G - E(G[X]) - V(Z_X)$  (i.e., we delete all edges in  $G[X]$  and all vertices in  $Z_X$ ).

For notational convenience, in what follows we denote  $Z_X$ ,  $z_X$ ,  $p_X$ , and  $r_X$  simply by  $Z$ ,  $z$ ,  $p$ , and  $r$ .

Let  $G_X$  be the graph obtained from  $G - V(Z)$  by deleting all edges in  $G[X]$ , if any. Hence,  $G_X$  can be obtained from  $G$  by the following three steps.

*Step (a).* Delete all edges in  $G[X]$ , if any.

*Step (b).* Delete all vertices in even components in  $G - X$ , if such components exist.

*Step (c).* Delete all vertices in odd components in  $G - X$  which do not have one or  $k$  edges into  $X$ , if such components exist.

Notice that in steps (b) and (c) we have deleted  $z$  vertices, while in step (c) we have deleted  $oc(Z)$  components. Furthermore we remark that the graph  $G_X$  consists of  $r$  components and

$$p = \sum_{x \in V(G_X)} (k - d_{G_X}(x)).$$

Now for every vertex  $x$  in  $G_X$  with  $d_{G_X}(x) < k$  (necessarily,  $x \in X$ ), add  $k - d_{G_X}(x)$  copies of the subgraph  $H_{k+2}$  to  $G_X$  and in each added copy of  $H_{k+2}$ , join the vertex named  $w_{k+2}$  in  $H_{k+2}$  to  $x$  (see Definition 5.1). The resulting graph, which we call  $G^*$ , is clearly  $k$ -regular and still consists of  $r$  components. Also note that  $X$  is a subset of  $V(G^*)$ , as we didn't delete any vertices from  $X$  in steps (a), (b) or (c). As  $X$  satisfies the conditions of Claim 5 in the graph  $G^*$ , we know that the following holds (note that we actually apply Claim 5 to each of the  $r$  components in  $G^*$ ).

$$\frac{1}{2} (|V(G^*)| + |X| - \text{oc}(G^* - X)) \geq \frac{(k^3 - k^2 - 2) |V(G^*)| - 2r(k - 1)}{2(k^3 - 3k)}.$$

Since  $|V(G^*)| = p(k + 2) + n - z$  and  $\text{oc}(G^* - X) = \text{oc}(G - X) - \text{oc}(Z) + p$ , this implies that

$$\begin{aligned} M_X + \frac{1}{2} (p(k + 2) - z - (p - \text{oc}(Z))) \\ \geq \phi(k, n) + \frac{(k^3 - k^2 - 2) (p(k + 2) - z) - 2(r - 1)(k - 1)}{2(k^3 - 3k)}. \end{aligned}$$

Thus since  $\phi(k, n) - M > 0$  (by assumption), we know that the following holds:

$$\begin{aligned} \frac{1}{2} (p(k + 2) - z - (p - \text{oc}(Z))) &> \frac{(k^3 - k^2 - 2) (p(k + 2) - z) - 2(r - 1)(k - 1)}{2(k^3 - 3k)} \\ \Leftrightarrow p \left( k + 1 - \frac{(k + 2)(k^3 - k^2 - 2)}{k^3 - 3k} \right) + \text{oc}(Z) &> z \left( 1 - \frac{k^3 - k^2 - 2}{k^3 - 3k} \right) - 2(r - 1) \left( \frac{k - 1}{k^3 - 3k} \right) \\ \Leftrightarrow p(4 - k - k^2) + \text{oc}(Z)(k^3 - 3k) &> z(k^2 - 3k + 2) - 2(r - 1)(k - 1). \end{aligned}$$

Note that  $p > 0$  since otherwise we were done by Claim 5 (it can also be seen as steps (a), (b) and (c) all decrease the degree of at least one vertex in  $X$  and  $p = \sum_{x \in V(G_X)} (k - d_{G_X}(x))$ ). As  $G$  is connected we note that every component in  $G_X$  has at least one vertex that has degree less than  $k$ . Therefore we must have  $p \geq r$ , which by the above implies the following.

$$\begin{aligned} p(4 - k - k^2) + \text{oc}(Z)(k^3 - 3k) &> z(k^2 - 3k + 2) - 2(p - 1)(k - 1) \\ \Leftrightarrow \text{oc}(Z)(k^3 - 3k) &> z(k^2 - 3k + 2) + p(k^2 - k - 2) + 2(k - 1). \end{aligned}$$

Note that if  $\text{oc}(Z) = z = p = 0$ , then the last inequality is false as  $0 > 2(k - 1)$  is false. We now consider what happens to this last inequality when each of the steps (a), (b) and (c), respectively, are performed.

- (a) When step (a) is performed, each edge in  $G[X]$  contributes two to  $p$  and zero to each of  $z$  and  $\text{oc}(Z)$ .
- (b) When step (b) is performed, each even component in  $G - X$  contributes at least one to  $p$ , at least two to  $z$ , and zero to  $\text{oc}(Z)$ .



(c) When step (c) is performed, we consider the following two possibilities.

Consider the case when we delete an odd component,  $C$ , of  $G - X$  with less than  $k$ , but with at least two, edges into  $X$ . Then,  $C$  contributes at least two to  $p$  and one to  $\text{oc}(Z)$ . By Observation 2, we note that  $|V(C)| \geq k + 1$ . As  $|V(C)|$  is odd we must have  $|V(C)| \geq k + 2$ , and so  $C$  contributes at least  $k + 2$  to  $z$ . Therefore we increase the left-hand side above with  $(k^3 - 3k)$  and we increase the right-hand side with at least  $(k + 2)(k^2 - 3k + 2) + 2(k^2 - k - 2) = (k^3 - 3k) + (k^2 - 3k)$ . As  $k \geq 3$ , this means we increase the right-hand side by at least the same as the left-hand side.

Now consider the case when we delete a component,  $C$ , of  $G - X$  with more than  $k$  edges into  $X$ . In this case  $C$  contributes at least  $k + 1$  to  $p$  and at least one to  $z$ , which implies that again we increase the right-hand side by at least as much as the left-hand side.

As  $\text{oc}(Z)(k^3 - 3k) \geq z(k^2 - 3k + 2) + p(k^2 - k - 2) + 2(k - 1)$  was false before we performed any of the steps (a), (b) or (c) (i.e., when  $\text{oc}(Z) = z = p = 0$ ), then by the above it must remain false as we never increase the left-hand side by more than the right-hand side. This is the desired contradiction. Hence,  $M_X \geq \phi(k, n)$  as claimed.  $\square$

The following proposition shows that the lower bound of Theorem 3 is tight.

**Proposition 5.** *For every integer  $p \geq 1$  and every odd integer  $k \geq 3$ , there exists a connected  $k$ -regular graph  $H_p^k$  of order  $p(k^3 - 3k) + k^2 + 2k + 1$  satisfying*

$$\alpha'(H_p^k) = \frac{(k^3 - k^2 - 2)|V(H_p^k)| - 2k + 2}{2(k^3 - 3k)}.$$

*Proof.* Let  $T_p^k$  be the bipartite graph with partite sets  $V_1 = \{w_1, w_2, \dots, w_p\}$  and  $V_2 = \{v_1, v_2, \dots, v_{p(k-1)+1}\}$ , and where  $N(w_i) = \{v_{(k-1)(i-1)+1}, v_{(k-1)(i-1)+2}, \dots, v_{(k-1)(i-1)+k}\}$  for all  $i = 1, 2, \dots, p$ . Note that the degree of all the  $w_i \in V_1$  is  $k$  and the degree of all vertices in  $V_2$  is one except for the vertices in  $\{w_{(k-1)+1}, w_{2(k-1)+1}, \dots, w_{(p-1)(k-1)+1}\}$  which have degree two. Now for every vertex  $x$  in  $V(T_p^k)$  with  $d_{T_p^k}(x) < k$ , add  $k - d_{T_p^k}(x)$  copies of the subgraph  $H_{k+2}$  to  $T_p^k$  and in each added copy of  $H_{k+2}$ , join the vertex named  $w_{k+2}$  in  $H_{k+2}$  to  $x$  (see Definition 5.1). Let  $H_p^k$  denote the resulting graph. Then,  $H_p^k$  is a connected  $k$ -regular graph of order

$$\begin{aligned} |V(H_p^k)| &= |V_1| + |V_2| + (|V_2|(k - 1) - (p - 1))(k + 2) \\ &= p + (p(k - 1) + 1) + ((p(k - 1) + 1)(k - 1) - (p - 1))(k + 2) \\ &= p(1 + (k - 1) + (k - 1)^2(k + 2) - (k + 2)) + 1 + (k - 1)(k + 2) + (k + 2) \\ &= p(k^3 - 3k) + k^2 + 2k + 1. \end{aligned}$$

Furthermore by deleting the vertices  $V_2$  from  $H_p^k$  we obtain  $p + (p(k - 1) + 1)(k - 1) + (p - 1)$  odd components. Therefore, by Theorem 1,

$$\begin{aligned}
 \alpha'(H_p^k) &\leq \frac{1}{2} \left( |V(H_p^k)| + |V_2| - \text{oc}(H_p^k - V_2) \right) \\
 &= \frac{1}{2} \left( (p(k^3 - 3k) + (k+1)^2) + (p(k-1) + 1) \right. \\
 &\quad \left. - (p + (p(k-1) + 1)(k-1) + (p-1)) \right) \\
 &= \frac{1}{2} \left( (p(k^3 - 3k + (k-1) - 1 - (k-1)^2 + 1) + (k+1)^2 - 1 - (k-1) + 1) \right) \\
 &= \frac{1}{2} \left( p(k^3 - k^2 - 2) + k^2 + k + 2 \right) \\
 &= \left( \frac{k^3 - k^2 - 2}{k^3 - 3k} \times \frac{p(k^3 - 3k) + k^2 + 2k + 1}{2} \right) \\
 &\quad + \frac{(k^3 - 3k)(k^2 + k + 2) - (k^2 + 2k + 1)(k^3 - k^2 - 2)}{2(k^3 - 3k)} \\
 &= \frac{(k^3 - k^2 - 2) |V(H_p^k)| - 2k + 2}{2(k^3 - 3k)}.
 \end{aligned}$$

However the lower bound of Theorem 3 shows that

$$\alpha'(H_p^k) \geq \frac{(k^3 - k^2 - 2) |V(H_p^k)| - 2k + 2}{2(k^3 - 3k)},$$

and so the desired result of the proposition follows. □

Using a shorter proof than that of Theorem 3 it is possible to prove the following weaker result. If  $k \geq 1$  is odd and  $G$  is a  $k$ -regular graph of order  $n$ , then  $\alpha'(G) \geq \frac{k^2+k+2}{k^2+2k+1} \times \frac{n}{2}$ . This can also be shown as a consequence of Theorem 3 as  $\lceil \frac{k^2+k+2}{k^2+2k+1} \times \frac{n}{2} \rceil \leq \lceil \frac{(k^3-k^2-2)n-2k+2}{2(k^3-3k)} \rceil$  holds for all  $n$  and  $k$  we consider. However this weaker result can be useful in some cases as for some values of  $n$  and  $k$  we actually have  $\frac{k^2+k+2}{k^2+2k+1} \times \frac{n}{2} > \frac{(k^3-k^2-2)n-2k+2}{2(k^3-3k)}$ . It is for example used in the paper [3].

### 6. Closing Remarks

Let  $G$  be a connected  $k$ -regular graph. As  $\alpha'(G) \leq \frac{n}{2}$  we note that our results imply that  $\frac{\alpha'(G)}{n} \rightarrow \frac{1}{2}$  when  $k \rightarrow \infty$ . In comparison, the bound,  $n(3k - 6)/(7k - 13)$ , in [4] is always less than  $3n/7$ . For small values of  $k$ , we summarize our results in the accompanying table.

G is a connected $k$ -regular graph						
$k$	3	4	5	6	7	8
$\alpha'(G) \geq$	$\frac{4n-1}{9}$	$\min \left\{ \frac{5n}{11}, \frac{n-1}{2} \right\}$	$\frac{49n-4}{110}$	$\min \left\{ \frac{5n}{11}, \frac{n-1}{2} \right\}$	$\frac{73n-3}{161}$	$\min \left\{ \frac{17n}{37}, \frac{n-1}{2} \right\}$

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