



A note on the generalized Hamming weights of Reed–Muller codes

Peter Beelen¹

Received: 3 July 2018 / Revised: 8 August 2018 / Accepted: 8 August 2018 / Published online: 18 August 2018
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Abstract

In this note, we give a very simple description of the generalized Hamming weights of Reed–Muller codes. For this purpose, we generalize the well-known Macaulay representation of a nonnegative integer and state some of its basic properties.

Keywords Reed–Muller code · Macaulay decomposition · Generalized Hamming weight

Mathematics Subject Classification 11H71 · 94B27

1 Preliminaries

Let \mathbb{F}_q be the finite field with q elements and denote by $\mathbb{A}^m := \mathbb{A}^m(\mathbb{F}_q)$ the m -dimensional affine space defined over \mathbb{F}_q . This space consists of q^m points (a_1, \dots, a_m) with $a_1, \dots, a_m \in \mathbb{F}_q$. Let $T(m) := \mathbb{F}_q[x_1, \dots, x_m]$ denote the ring of polynomials in m variables and coefficients in \mathbb{F}_q . Further let $T_{\leq d}(m)$ be the set of polynomials in $T(m)$ of total degree at most d . A monomial $X_1^{\alpha_1} \dots X_m^{\alpha_m}$ is called reduced if $(\alpha_1, \dots, \alpha_m) \in \{0, 1, \dots, q-1\}^m$. Similarly a polynomial $f \in T(m)$ is called reduced if it is an \mathbb{F}_q -linear combination of reduced monomials. We denote the set of reduced polynomials by $T^{\text{red}}(m)$ and define $T_{\leq d}^{\text{red}}(m) := T_{\leq d}(m) \cap T^{\text{red}}(m)$.

One reason for considering reduced polynomials comes from coding theory. Indeed, Reed–Muller codes are obtained by evaluating certain polynomials in the points of \mathbb{A}^m , but the evaluation map

$$\text{Ev} : T(m) \rightarrow \mathbb{F}_q^m, \quad \text{defined by } \text{Ev}(f) = (f(P))_{P \in \mathbb{A}^m}$$

✉ Peter Beelen
pabe@dtu.dk

¹ Department of Applied Mathematics and Computer Science, Technical University of Denmark, 2800 Kongens Lyngby, Denmark

is not injective. However, its restriction to $T^{\text{red}}(m)$ is. In fact, the kernel of Ev consists precisely of the ideal $I \subset T(m)$ generated by the polynomials $x_i^q - x_i$ ($1 \leq i \leq m$). Working with reduced polynomials is simply a convenient way to take this into account, since for two reduced polynomials $f_1, f_2 \in T(m)$ the equality $f_1 + I = f_2 + I$ holds if and only if $f_1 = f_2$.

The Reed–Muller code $\text{RM}_q(d, m)$ is the set of vectors from $\mathbb{F}_q^{q^m}$ obtained by evaluating polynomials of total degree up to d in the q^m points of \mathbb{A}^m , that is to say:

$$\text{RM}_q(d, m) := \{(f(P))_{P \in \mathbb{A}^m} : f \in T_{\leq d}(m)\}.$$

By the above, we also have $\text{RM}_q(d, m) := \{(f(P))_{P \in \mathbb{A}^m} : f \in T_{\leq d}^{\text{red}}(m)\}$ and moreover, we have

$$\dim \text{RM}_q(d, m) = \dim T_{\leq d}^{\text{red}}(m). \tag{1}$$

Reed–Muller codes $\text{RM}_q(d, m)$ have been studied extensively for their elegant algebraic properties. Their generalized Hamming weights $d_r(\text{RM}_q(d, m))$ have been determined in [4] by Heijnen and Pellikaan. For a general linear code $C \subseteq \mathbb{F}_q^n$ these are defined as follows:

$$d_r(C) := \min_{D \subseteq C: \dim D=r} |\text{supp}(D)|,$$

where the minimum is taken over all r -dimensional \mathbb{F}_q -linear subspaces D of C and where $\text{supp}(D)$ denotes the support of D , that is to say

$$\text{supp}(D) := \{i : \exists (c_1, \dots, c_n) \in D, c_i \neq 0\}.$$

In case of Reed–Muller codes, there is a direct relation between generalized Hamming weights and the number of common solutions to systems of polynomial equations. Indeed, if $D \subset \text{RM}_q(d, m)$ is spanned by $(f_i(P))_{P \in \mathbb{A}^m}$ for $f_1, \dots, f_r \in T_{\leq d}^{\text{red}}(m)$, then $|\text{supp}(D)| = q^m - |Z(f_1, \dots, f_r)|$ where $Z(f_1, \dots, f_r) := \{P \in \mathbb{A}^m : f_1(P) = \dots = f_r(P) = 0\}$ denotes the set of common zeros of f_1, \dots, f_r in the m -dimensional affine space \mathbb{A}^m over \mathbb{F}_q . Therefore, if we define

$$\bar{e}_r^{\mathbb{A}}(d, m) := \max \left\{ |Z(f_1, \dots, f_r)| : f_1, \dots, f_r \in T_{\leq d}^{\text{red}}(m) \text{ linearly independent} \right\}, \tag{2}$$

then $d_r(\text{RM}_q(d, m)) = q^m - \bar{e}_r^{\mathbb{A}}(d, m)$. Note that $T^{\text{red}}(m)$ is a vector space over \mathbb{F}_q of dimension q^m and that a reduced polynomial has total degree at most $m(q - 1)$. Therefore $T^{\text{red}}(m) = T_{\leq m(q-1)}^{\text{red}}(m)$. This implies in particular that $\text{RM}_q(d, m) = \mathbb{F}_q^{q^m}$ for $d \geq m(q - 1)$. Therefore, we will always assume that $d \leq m(q - 1)$.

The result of Heijnen–Pellikaan in [4] on the value of $d_r(\text{RM}_q(d, m))$ can now be restated as follows, see for example [2].

$$\bar{e}_r^{\mathbb{A}}(d, m) = \sum_{i=1}^m \mu_i q^{m-i}, \tag{3}$$

where (μ_1, \dots, μ_m) is the r -th m -tuple in descending lexicographic order among all m -tuples $(\beta_1, \dots, \beta_m) \in \{0, 1, \dots, q - 1\}^m$ satisfying $\beta_1 + \dots + \beta_m \leq d$.

Following the notation in [4], we denote with $\rho_q(d, m)$ the dimension of $\text{RM}_q(d, m)$. Equation (1) implies that $\rho_q(d, m) = \dim(T_{\leq d}^{\text{red}}(m))$. In particular, we have

$$\rho_q(d, m) = \dim(T_{\leq d}(m)) = \binom{m+d}{d}, \quad \text{if } d \leq q - 1, \tag{4}$$

since $T_{\leq d}(m) = T_{\leq d}^{\text{red}}(m)$ if $d < q$. Here as well as later on we use the convention that $\binom{a}{b} = 0$ if $a < b$. In particular we have $\rho_q(d, m) = 0$ if $d < 0$. As shown in [1, §5.4], for the general case $d \leq m(q - 1)$, we have

$$\rho_q(d, m) = \dim(T_{\leq d}^{\text{red}}(m)) = \sum_{i=0}^d \sum_{j=0}^m (-1)^j \binom{m}{j} \binom{m-1+i-qj}{m-1}. \tag{5}$$

In this note, we will present an easy-to-obtain expression for $\bar{e}_r^{\Delta}(d, m)$ involving a certain representation of the number $\rho_q(d, m) - r$ that we introduce in the next section.

2 The d -th Macaulay representation with respect to q

Let d be a positive integer. The d -th Macaulay (or d -binomial) representation, of a nonnegative integer N is a way to write N as sum as certain binomial coefficients. To be precise

$$N = \sum_{i=1}^d \binom{s_i}{i},$$

where the s_i integers satisfying $s_d > s_{d-1} > \dots > s_1 \geq 0$. The usual convention that $\binom{a}{b} = 0$ if $a < b$, is used. For example, the d -th Macaulay representation of 0 is given by $0 = \sum_{i=1}^d \binom{i-1}{i}$. Given d and N the integers s_i exist and are unique. The Macaulay representation is among other things used for the study of Hilbert functions of graded modules, see for example [3]. It is well known (see for example [3]) that if N and M are two nonnegative integers with Macaulay representations given by (k_d, \dots, k_1) and (ℓ_d, \dots, ℓ_1) then $N \leq M$ if and only if $(k_d, \dots, k_1) \preceq (\ell_d, \dots, \ell_1)$, where \preceq denotes the lexicographic order.

For our purposes it is more convenient to define $m_i := s_i - i$. We then obtain

$$N = \sum_{i=1}^d \binom{m_i + i}{i}, \tag{6}$$

where m_i are integers satisfying $m_d \geq m_{d-1} \geq \dots \geq m_1 \geq -1$. The reason for this is that for $d \leq q - 1$ we have $\rho_q(d, m) = \binom{m+d}{d}$. Therefore, we can interpret Eq. (6) as a statement concerning dimensions of the Reed–Muller codes $\text{RM}_q(i, m_i)$. For a suitable choice of N , it turns out that the m_i completely determine the value of $\bar{e}_r^{\Delta}(d, m)$ if $d \leq q - 1$. For $d \geq q$, even though the dimension $\rho_q(d, m)$ is not longer given by

$\binom{m+d}{d}$, there exists a variant of the usual d -th Macaulay representation that turns out to be equally meaningful for Reed–Muller codes. Before stating this representation, we give a lemma.

Lemma 2.1 *Let $m \geq 1$ be an integer. We have*

$$\rho_q(d, m) = \sum_{i=0}^{\min\{d, q-1\}} \rho_q(d - i, m - 1).$$

Proof Any polynomial $f \in T(m)$ can be seen as a polynomial in the variable X_m with coefficients in $T(m - 1)$. This implies that $T(m) = \sum_{i \geq 0} X_m^i T(m)$, where the sum is a direct sum. Similarly we can write

$$T_{\leq d}^{\text{red}}(m) = \sum_{i=0}^{\min\{d, q-1\}} X_m^i T_{\leq d-i}^{\text{red}}(m - 1).$$

The result now follows. □

A consequence of this lemma is the following.

Corollary 2.2 *Let $d = a(q - 1) + b$ for integers a and b satisfying $a \geq 0$ and $1 \leq b \leq q - 1$. Further suppose that $m \geq a$. Then*

$$\rho_q(d, m) - 1 = \sum_{j=0}^{a-1} \sum_{\ell=0}^{q-2} \rho_q(d - j(q - 1) - \ell, m - j - 1) + \sum_{i=1}^b \rho_q(i, m - a - 1).$$

Proof This follows using Lemma 2.1 repeatedly. First applying the lemma to each sum within the double summation on the right-hand side, we see that

$$\begin{aligned} & \sum_{j=0}^{a-1} \sum_{\ell=0}^{q-2} \rho_q(d - j(q - 1) - \ell, m - j - 1) \\ &= \sum_{j=0}^{a-1} (\rho_q(d - j(q - 1), m - j) - \rho_q(d - (j + 1)(q - 1), m - j - 1)) \\ &= \rho_q(d, m) - \rho_q(d - a(q - 1), m - a) = \rho_q(d, m) - \rho_q(b, m - a). \end{aligned}$$

Using the same lemma to rewrite the single summation on the right-hand side in Eq. (9) we see that if $m > a$

$$\sum_{i=1}^b \rho_q(i, m - a - 1) = \rho_q(b, m - a) - \rho_q(0, m - a - 1) = \rho_q(b, m - a) - 1,$$

while if $m = a$, the single summation equals 0 and the double summation simplifies to $\rho_q(d, m) - 1$. In either case, we obtain the desired result □

We can now show the following.

Theorem 2.3 *Let $N \geq 0$ and $d \geq 1$ be integers and q a prime power. Then there exist uniquely determined integers m_1, \dots, m_d satisfying*

1. $N = \sum_{i=1}^d \rho_q(i, m_i)$,
2. $-1 \leq m_1 \leq \dots \leq m_d$,
3. for all i satisfying $1 \leq i \leq d - q + 1$, either $m_{i+q-1} > m_i$ or $m_{i+q-1} = m_i = -1$.

Proof We start by showing uniqueness. Suppose that

$$N = \sum_{i=1}^d \rho_q(i, m_i) = \sum_{i=1}^d \rho_q(i, n_i) \tag{7}$$

and the integers n_1, \dots, n_d and m_1, \dots, m_d satisfy the conditions from the theorem. First of all, if $m_d = -1$ or $n_d = -1$ then $N = 0$. Either assumption implies that $(m_d, \dots, m_1) = (-1, \dots, -1) = (n_d, \dots, n_1)$. Indeed $n_i \geq 0$ or $m_i \geq 0$ for some i directly implies that $N > 0$. Therefore we from now on assume that $m_d \geq 0$ and $n_d \geq 0$. To arrive at a contradiction, we may assume without loss of generality that $n_d \leq m_d - 1$.

Define e to be the smallest integer such that $n_e \geq 0$. Equation (7) can then be rewritten as

$$N = \sum_{i=1}^d \rho_q(i, m_i) = \sum_{i=e}^d \rho_q(i, n_i) \tag{8}$$

Condition 3 from the theorem implies that $n_{i-q+1} < n_i$ for all i satisfying $e \leq i \leq d$. Now write $d - e + 1 = a(q - 1) + b$ for integers a and b satisfying $a \geq 0$ and $1 \leq b \leq q - 1$. With this notation, we obtain that for any $0 \leq j \leq a - 1$ and $0 \leq \ell \leq q - 2$ we have that

$$n_{d-j(q-1)-\ell} \leq n_d - j \leq m_d - j - 1.$$

In particular choosing $j = a - 1$ and $\ell = 0$, this implies that $m_d \geq a + n_{q-1+b} \geq a + 1 + n_b \geq a$. Using these observations, we obtain from Eq. (7) that

$$\begin{aligned} \rho_q(d, m_d) \leq N &= \sum_{i=e}^d \rho_q(i, n_i) \\ &\leq \sum_{j=0}^{a-1} \sum_{\ell=0}^{q-2} \rho_q(d - j(q - 1) - \ell, m_d - j - 1) \\ &\quad + \sum_{i=1}^b \rho_q(e + i - 1, m_d - a - 1). \end{aligned} \tag{9}$$

Applying the same technique as in the proof of Corollary 2.2, we derive that

$$\sum_{j=0}^{a-1} \sum_{\ell=0}^{q-2} \rho_q(d - j(q - 1) - \ell, m_d - j - 1) = \rho_q(d, m_d) - \rho_q(b + e - 1, m_d - a)$$

and Eq. (9) can be simplified to

$$\rho_q(d, m_d) \leq \rho_q(d, m_d) - \rho_q(b + e - 1, m_d - a) + \sum_{i=1}^b \rho_q(e + i - 1, m_d - a - 1). \tag{10}$$

For $m_d = a$ the right-hand side equals $\rho_q(d, m_d) - 1$, leading to a contradiction. If $m_d > q$, Eq. (10) implies

$$\begin{aligned} \rho_q(b + e - 1, m_d - a) &\leq \sum_{i=1}^b \rho_q(e + i - 1, m_d - a - 1) \\ &= \sum_{j=0}^{b-1} \rho_q(e + b - 1 - j, m_d - a - 1) \\ &< \sum_{j=0}^{\min\{e+b-1, q-1\}} \rho_q(e + b - 1 - j, m_d - a - 1) \\ &= \rho_q(b + e - 1, m_d - a), \end{aligned}$$

where in the last equality we used Lemma 2.1. Again we arrive at a contradiction. This completes the proof of uniqueness of the d -th Macaulay representation with respect to q .

Now we show existence. Let d, N and q be given. We will proceed with induction on d . For $d = 1$, note that $\rho_q(1, m) = m + 1$ for any $m \geq -1$. Therefore, for a given $N \geq 0$, we can write $N = \rho_q(1, N - 1)$.

Now assume the theorem for $d - 1$. There exists $m_d \geq -1$ such that

$$\rho_q(d, m_d) \leq N < \rho_q(d, m_d + 1). \tag{11}$$

Applying the induction hypothesis on $N - \rho_q(d, m_d)$, we can find m_{d-1}, \dots, m_1 satisfying the conditions of the theorem for $d - 1$. In particular we have that

1. $N - \rho_q(d, m_d) = \sum_{i=1}^{d-1} \rho_q(i, m_i)$,
2. $-1 \leq m_1 \leq \dots \leq m_{d-1}$,
3. $m_{i+(q-1)} > m_i$ for all $1 \leq i \leq d - q$.

Clearly this implies that $N = \sum_{i=1}^d \rho_q(i, m_i)$, but it is not clear a priori that m_1, \dots, m_d satisfy conditions 2 and 3 as well. Conditions 2 and 3 would follow once we show that $m_d \geq m_{d-1}$ and either $m_d > m_{d-q+1}$ or $m_d = m_{d-q+1} = -1$. First of all, if $m_d = -1$, then $N = 0$ and $(m_d, \dots, m_1) = (-1, \dots, -1)$. Hence there

is nothing to prove in that case. Assume $m_d \geq 0$. From Eq. (11) and Lemma 2.1 we see that

$$N - \rho_q(d, m_d) < \rho_q(d, m_d + 1) - \rho_q(d, m_d) = \sum_{i=1}^{\min\{d, q-1\}} \rho_q(d - i, m_d). \tag{12}$$

First suppose that $d \leq q - 1$. First of all, Condition 3 is empty in that setting. Further, Eq. (12) implies

$$N - \rho_q(d, m_d) < \sum_{i=1}^d \rho_q(d - i, m_d) = \sum_{i=1}^{d-1} \rho_q(d - i, m_d) + 1$$

and hence

$$N - \rho_q(d, m_d) \leq \sum_{i=1}^{d-1} \rho_q(d - i, m_d) = \sum_{j=0}^{d-2} \rho_q(d - 1 - j, m_d) < \rho_q(d - 1, m_d + 1).$$

This shows that $m_{d-1} \leq m_d$ as desired.

Now suppose that $d \geq q$. In this situation Eq. (12) implies

$$N - \rho_q(d, m_d) < \sum_{i=1}^{q-1} \rho_q(d - i, m_d) = \sum_{j=0}^{q-2} \rho_q(d - 1 - j, m_d) < \rho_q(d - 1, m_d + 1).$$

Hence $m_{d-1} \leq m_d$ as before. Finally assume that $m_d \leq m_{d-q+1}$. Then by the previous and Condition 2, we have $m_d = m_{d-1} = \dots = m_{d-q+1}$. Hence $N \geq \sum_{i=0}^{q-1} \rho_q(d - i, m_d) = \rho_q(d, m_d + 1)$ which is in contradiction with Eq. (11). This concludes the induction step and hence the proof of existence. \square

We call the representation of N in the above theorem the d -th Macaulay representation of N with respect to q . One retrieves the usual d -th Macaulay representation letting q tend to infinity. We refer to (m_d, \dots, m_1) as the coefficient tuple of this representation. A direct corollary of the above is the following.

Corollary 2.4 *The coefficient tuple (m_d, \dots, m_1) of the d -th Macaulay representation with respect to q of a nonnegative integer N can be computed using the following greedy algorithm: The coefficient m_{d-i} can be computed recursively (starting with $i = 0$) as the unique integer $m_{d-i} \geq -1$ such that*

$$\rho_q(d - i, m_{d-i}) \leq N - \sum_{j=d-i+1}^d \rho_q(j, m_j) < \rho_q(d - i, m_{d-i} + 1).$$

Proof From the existence-part of the proof of Theorem 2.3 it follows directly that the given greedy algorithm finds the desired coefficients. \square

A further corollary is the following. As before \leq denotes the lexicographic order.

Corollary 2.5 *Suppose the N and M are two nonnegative integers whose respective coefficient tuples are (n_d, \dots, n_1) and (m_d, \dots, m_1) . Then*

$$N \leq M \text{ if and only if } (n_d, \dots, n_1) \leq (m_d, \dots, m_1).$$

Proof Assume $(n_d, \dots, n_1) \leq (m_d, \dots, m_1)$. It is enough to show the corollary in case $n_d < m_d$. We know from the previous corollary that n_d and m_d may be determined using the given greedy algorithm. In particular this implies that $n_d < m_d$ implies

$$N < \rho_q(d, n_d + 1) \leq \rho_q(d, m_d) \leq M.$$

Assume that $N \leq M$. We use induction on d . The induction basis is trivial: If $d = 1$, then $m_1 = M - 1$ and $n_1 = N - 1$. For the induction step, note that $N \leq M < \rho_q(d, m_d + 1)$ implies by the greedy algorithm that $n_d \leq m_d$. If $n_d < m_d$, we are done. If $n_d = m_d$, we replace N with $N - \rho_q(d, m_d)$ and M with $M - \rho_q(d, m_d)$ and use the induction hypothesis to conclude that $(n_d, \dots, n_1) \leq (m_d, \dots, m_1)$. \square

3 A simple expression for $\bar{e}_r^{\Delta}(d, m)$

We are now ready to state and prove the relation between the Macaulay representation with respect to q and $\bar{e}_r^{\Delta}(d, m)$.

Theorem 3.1 *For $1 \leq r \leq \rho_q(d, m)$, let the d -th Macaulay representation of $\rho_q(d, m) - r$ with respect to q be given by*

$$\rho_q(d, m) - r = \sum_{i=1}^d \rho_q(i, m_i).$$

Denoting the floor function as $\lfloor \cdot \rfloor$, we have

$$\bar{e}_r^{\Delta}(d, m) = \sum_{i=1}^d \lfloor q^{m_i} \rfloor.$$

Proof We know from Eq. (3) that we need to show that

$$\sum_{i=1}^d \lfloor q^{m_i} \rfloor = \sum_{i=1}^m \mu_i q^{m-i},$$

with (μ_1, \dots, μ_m) is the r -th element in descending lexicographic order among all m -tuples $(\beta_1, \dots, \beta_m)$ in $\{0, 1, \dots, q - 1\}^m$ satisfying $\beta_1 + \dots + \beta_m \leq d$. First of all note that since $r \geq 1$, we have $\rho_q(d, m) - r < \rho_q(d, m)$. In particular this implies that $m_d \leq m - 1$. Therefore the coefficients of the d -tuple (m_d, \dots, m_1) are

in $\{-1, 0, \dots, m - 1\}$. Now for $1 \leq i \leq m + 1$ define $\mu_i := |\{j : m_j = m - i\}|$. Since the d -tuple (m_d, \dots, m_1) is nonincreasing by Condition 2 from Theorem 2.3, we can reconstruct it uniquely from the $(m + 1)$ -tuple $(\mu_1, \mu_2, \dots, \mu_{m+1})$. Moreover, Condition 3 from Theorem 2.3, implies that $(\mu_1, \dots, \mu_m) \in \{0, 1, \dots, q - 1\}^m$, but note that μ_{m+1} could be strictly larger than $q - 1$. Further by construction we have $\mu_1 + \dots + \mu_m + \mu_{m+1} = d$, implying that $\mu_1 + \dots + \mu_m \leq d$. Note that μ_{m+1} is determined uniquely by (μ_1, \dots, μ_m) , since $\mu_0 = d - \mu_1 - \dots - \mu_m$. Therefore the correspondence between the d -tuples (m_d, \dots, m_1) of coefficients of the d -th Macaulay representations with respect to q of integers $0 \leq N < \rho_q(d, m)$ and the m -tuples $(\mu_1, \dots, \mu_m) \in \{0, 1, \dots, q - 1\}^m$ satisfying $\mu_1 + \dots + \mu_m \leq d$, is a bijection. Moreover by construction we have

$$\sum_{i=1}^d [q^{m_i}] = \sum_{j=1}^{m+1} \mu_j [q^{m-j}] = \sum_{j=1}^m \mu_j q^{m-j}.$$

What remains to be shown is that the constructed m -tuple coming from the integer $\rho_q(d, m) - r$ is in fact the r -th in descending lexicographic order. First of all, by Corollary 2.2 we see that for $r = 1$ and $d = aq + b$ that the m -tuple associated to $\rho_q(d, m) - 1$ equals $(q - 1, \dots, q - 1, b, 0, \dots, 0)$, which under the lexicographic order is the maximal m -tuple among all m -tuples $(\beta_1, \dots, \beta_m) \in \{0, 1, \dots, q - 1\}^m$ satisfying $\beta_1 + \dots + \beta_m \leq d$. Next we show that the conversion between d -tuples (m_d, \dots, m_1) to m -tuples (μ_1, \dots, μ_m) preserves the lexicographic order. Suppose therefore that $1 \leq r \leq s \leq \rho_q(d, m)$. We write $N := \rho_q(d, m) - s$ and $M := \rho_q(d, m) - r$. and denote their Macaulay coefficient tuples with (n_d, \dots, n_1) and (m_d, \dots, m_1) . Since $N \leq M$, Corollary 2.5 implies that $(n_d, \dots, n_1) \leq (m_d, \dots, m_1)$. Also, since these d -tuples are nonincreasing, this implies that their associated m -tuples (v_1, \dots, v_m) and (μ_1, \dots, μ_m) satisfy $(v_1, \dots, v_m) \leq (\mu_1, \dots, \mu_m)$. Indeed assuming without loss of generality that $v_1 < \mu_1$ we see that $m_i = n_i = m - 1$ for $d - v_1 \leq i \leq d$ but $n_i < m_i = m - 1$ for $i = v_1 + 1$. Now the desired result follows immediately. \square

Combining this theorem with the greedy algorithm in Corollary 2.4, it is very simple to compute values of $\bar{e}_r^{\mathbb{A}}(d, m)$ or equivalently of $d_r(\text{RM}_q(d, m))$. We illustrate this in the two following examples. The parameters in these example also occur in examples from [4].

Example 3.2 Let $q = 4, r = 8, d = m = 3$. Since $d \leq q - 1$, we may work with the usual Macaulay representation when applying Theorem 3.1. We have $\rho_q(d, m) = \binom{6}{3} = 20$ and hence

$$\rho_q(d, m) - r = 12 = \binom{5}{3} + \binom{2}{2} + \binom{1}{1} = \rho_4(3, 2) + \rho_4(2, 0) + \rho_4(1, 0)$$

is the 3-rd Macaulay representation of 12. Theorem 3.1 implies that $\bar{e}_8^{\mathbb{A}}(3, 3) = 4^2 + 4^0 + 4^0 = 18$ and hence $d_8(\text{RM}_4(3, 3)) = 64 - 18 = 46$ in accordance with Example 6.10 in [4].

Example 3.3 Let $q = 2$, $r = 10$, $d = 3$ and $m = 5$. We have $\rho_2(3, 5) = 26$ by Eq. (5) and hence applying the greedy algorithm from Corollary 2.4, we compute that

$$\rho_q(d, m) - r = 16 = 15 + 1 + 0 = \rho_2(3, 4) + \rho_2(2, 0) + \rho_2(1, -1)$$

is the 3rd Macaulay representation of 16 with respect to 2. Theorem 3.1 implies that $\bar{e}_{10}^{\mathbb{A}}(3, 3) = 2^4 + 2^0 = 17$ and hence $d_8(\text{RM}_2(3, 5)) = 32 - 17 = 15$ in accordance with Example 6.12 in [4].

Remark 3.4 Theorem 3.1 is somewhat similar in spirit as Theorem 6.8 from [4] in the sense that in both theorems a certain representation in terms of dimensions of Reed–Muller codes is used to give an expression for $d_r(\text{RM}_q(d, m))$. Where we studied decompositions of $\rho_q(d, m) - r$, in [4] the focus was on r itself. This suggests there may exist a duality between the two approaches, but the similarities seem to stop there. The representation in [4] is not the Macaulay representation with respect to q that we have used here. For us it is for example very important that each degree i between 1 and d occurs once in Theorem 2.3 (implying that the greedy algorithm terminates after at most d iterations), while this is not the case in Theorem 6.8 [4]. It could be interesting future work to determine if a deeper lying relationship between the two approaches exists.

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