

Weight enumerators of a class of linear codes

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Abstract Recently, linear codes constructed from defining sets have been studied widely and they have many applications. For an odd prime p , let $q = p^m$ for a positive integer m and Tr_m the trace function from \mathbb{F}_q onto \mathbb{F}_p . In this paper, for a positive integer t , let $D \subset \mathbb{F}_q^t$ and $D = \{(x_1, x_2) \in (\mathbb{F}_q^*)^2 : \text{Tr}_m(x_1 + x_2) = 0\}$, we define a p -ary linear code \mathcal{C}_D by

$$\mathcal{C}_D = \left\{ \mathbf{c}(a_1, a_2) : (a_1, a_2) \in \mathbb{F}_q^2 \right\},$$

where

$$\mathbf{c}(a_1, a_2) = \left(\text{Tr}_m \left(a_1 x_1^2 + a_2 x_2^2 \right) \right)_{(x_1, x_2) \in D}.$$

We compute the weight enumerators of the punctured codes \mathcal{C}_D .

Keywords Linear codes · Weight distribution · Gauss sums

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1 Introduction

Let \mathbb{F}_p be the finite field with p elements, where p is an odd prime. An $[n, k, d]$ linear code \mathcal{C} over \mathbb{F}_p is a k -dimensional subspace of \mathbb{F}_p^n with minimum distance d . Let A_i denote the number of codewords with Hamming weight i the code \mathcal{C} of length n . The weight enumerator of \mathcal{C} is defined by $1 + A_1z + A_2z^2 + \cdots + A_nz^n$. The sequence $(1, A_1, A_2, \dots, A_n)$ is called the *weight distribution* of the code \mathcal{C} . The weight distribution of the linear code is an important subject in coding theory. However, it is difficult to compute the weight distribution of a linear code in general.

Recently, the weight enumerators of linear codes were studied in [1, 2, 4–6, 9–12, 15–18] with the help of exponential sums in some cases. Ahn, Ka and Li [1] defined a class of linear codes as follows. Let $D' = \{(x_1, x_2, \dots, x_t) \in \mathbb{F}_q^t \setminus \{(0, 0, \dots, 0)\} : \text{Tr}_m(x_1 + x_2 + \cdots + x_t) = 0\}$. A p -ary linear code $\mathcal{C}_{D'}$ is defined by

$$\mathcal{C}_{D'} = \left\{ \mathbf{c}(a_1, a_2, \dots, a_t) : (a_1, a_2, \dots, a_t) \in \mathbb{F}_q^t \right\},$$

where

$$\mathbf{c}(a_1, a_2, \dots, a_t) = \left(\text{Tr}_m \left(a_1x_1^2 + a_2x_2^2 + \cdots + a_tx_t^2 \right) \right)_{(x_1, x_2, \dots, x_t) \in D'}.$$

They determined the complete weight enumerators of $\mathcal{C}_{D'}$. Yang and Yao [17] generalized the results of Ahn, Ka and Li [1]. They defined $D_b = \{(x_1, x_2, \dots, x_t) \in \mathbb{F}_q^t : \text{Tr}_m(x_1 + x_2 + \cdots + x_t) = b\}$ for any $b \in \mathbb{F}_p^*$ and determined the complete weight enumerator of a class of p -ary linear codes given by

$$\mathcal{C}_{D_b} = \left\{ \mathbf{c}(a_1, a_2, \dots, a_t) : (a_1, a_2, \dots, a_t) \in \mathbb{F}_q^t \right\},$$

where

$$\mathbf{c}(a_1, a_2, \dots, a_t) = \left(\text{Tr}_m \left(a_1x_1^2 + a_2x_2^2 + \cdots + a_tx_t^2 \right) \right)_{(x_1, x_2, \dots, x_t) \in D_b}.$$

In this paper, we define

$$D = \left\{ (x_1, x_2) \in \left(\mathbb{F}_q^* \right)^2 : \text{Tr}_m(x_1 + x_2) = 0 \right\} \quad (1)$$

and a p -ary linear code \mathcal{C}_D by

$$\mathcal{C}_D = \left\{ \mathbf{c}(a_1, a_2) : (a_1, a_2) \in \mathbb{F}_q^2 \right\}, \quad (2)$$

where

$$\mathbf{c}(a_1, a_2) = \left(\text{Tr}_m \left(a_1x_1^2 + a_2x_2^2 \right) \right)_{(x_1, x_2) \in D}.$$

The purpose of this paper is to compute the weight enumerators of the punctured codes \mathcal{C}_D .

Minimal linear codes can be used to construct secret sharing schemes with interesting access structures [7,8]. The codes presented in this paper are minimal in the sense of Ding and Yuan [7,8]. We shall explain it at the end of this paper in detail.

2 Preliminaries

Let p be an odd prime and $q = p^m$ for a positive integer m . For any $a \in \mathbb{F}_q$, we can define an additive character of the finite field \mathbb{F}_q as follows:

$$\psi_a : \mathbb{F}_q \longrightarrow \mathbb{C}^*, \psi_a(x) = \zeta_p^{\text{Tr}_m(ax)},$$

where $\zeta_p = e^{\frac{2\pi\sqrt{-1}}{p}}$ is a p -th primitive root of unity and Tr_m denotes the trace function from \mathbb{F}_q onto \mathbb{F}_p . It is clear that $\psi_0(x) = 1$ for all $x \in \mathbb{F}_q$. Then ψ_0 is called the trivial additive character of \mathbb{F}_q . If $a = 1$, we call $\psi := \psi_1$ the canonical additive character of \mathbb{F}_q . It is easy to see that $\psi_a(x) = \psi(ax)$ for all $a, x \in \mathbb{F}_q$. The orthogonal property of additive characters is given by

$$\sum_{x \in \mathbb{F}_q} \psi_a(x) = \begin{cases} q, & \text{if } a = 0, \\ 0, & \text{if } a \in \mathbb{F}_q^*. \end{cases}$$

Let $\lambda : \mathbb{F}_q^* \rightarrow \mathbb{C}^*$ be a multiplicative character of \mathbb{F}_q^* . Now we define the Gauss sum over \mathbb{F}_q by

$$G(\lambda) = \sum_{x \in \mathbb{F}_q^*} \lambda(x)\psi(x).$$

Let $q - 1 = sN$ for two positive integers $s > 1, N > 1$ and α be a fixed primitive element of \mathbb{F}_q . Let $\langle \alpha^N \rangle$ denote the subgroup of \mathbb{F}_q^* generated by α^N . The *cyclotomic classes* of order N in \mathbb{F}_q are the cosets $C_i^{(N,q)} = \alpha^i \langle \alpha^N \rangle$ for $i = 0, 1, \dots, N - 1$. We know that $|C_i^{(N,q)}| = \frac{q-1}{N}$. The Gaussian periods of order N are defined by

$$\eta_i^{(N,q)} = \sum_{x \in C_i^{(N,q)}} \psi(x).$$

Suppose that η is the quadratic character of \mathbb{F}_q^* and η_p is the quadratic character of \mathbb{F}_p^* . For $z \in \mathbb{F}_p^*$, it is easily checked that

$$\eta(z) = \begin{cases} 1, & \text{if } m \text{ is even,} \\ \eta_p(z), & \text{if } m \text{ is odd.} \end{cases} \tag{3}$$

Lemma 1 [3,13] *Suppose that $q = p^m$ where p is an odd prime and $m \geq 1$. Then*

$$G(\eta) = (-1)^{m-1} \sqrt{(p^*)^m} = \begin{cases} (-1)^{m-1} \sqrt{q}, & \text{if } p \equiv 1 \pmod{4}, \\ (-1)^{m-1} (\sqrt{-1})^m \sqrt{q}, & \text{if } p \equiv 3 \pmod{4}, \end{cases}$$

where $p^* = \left(\frac{-1}{p}\right)p = (-1)^{\frac{p-1}{2}} p$.

Lemma 2 [13] *If q is odd and $f(x) = a_2x^2 + a_1x + a_0 \in \mathbb{F}_q[x]$ with $a_2 \neq 0$, then*

$$\sum_{x \in \mathbb{F}_q} \zeta_p^{\text{Tr}_m(f(x))} = \zeta_p^{\text{Tr}_m(a_0 - a_1^2(4a_2)^{-1})} \eta(a_2)G(\eta).$$

Lemma 3 [14] *When $N = 2$, the Gaussian periods are given by*

$$\eta_0^{(2,q)} = \begin{cases} \frac{-1+(-1)^{m-1}\sqrt{q}}{2}, & \text{if } p \equiv 1 \pmod{4}, \\ \frac{-1+(-1)^{m-1}(\sqrt{-1})^m\sqrt{q}}{2}, & \text{if } p \equiv 3 \pmod{4}, \end{cases}$$

and $\eta_1^{(2,q)} = -1 - \eta_0^{(2,q)}$.

3 Weight enumerators of the linear codes of \mathcal{C}_D

In this section, we present the weight distribution of the linear code \mathcal{C}_D defined by (1) and (2), where

$$D = \left\{ (x_1, x_2) \in \left(\mathbb{F}_q^*\right)^2 : \text{Tr}_m(x_1 + x_2) = 0 \right\}.$$

To get the length of \mathcal{C}_D , we need the following lemma.

Lemma 4 *Denote $n_c = |\{x_1, x_2, \in \mathbb{F}_q^* : \text{Tr}_m(x_1 + x_2) = c\}|$ for each $c \in \mathbb{F}_p$. Then*

$$n_c = \begin{cases} (p^m - 1)^2 + p - 1, & \text{if } c = 0, \\ \frac{(p^m - 1)^2 - 1}{p}, & \text{if } c \neq 0. \end{cases}$$

Proof By the orthogonal property of additive characters, we have

$$\begin{aligned} n_c &= \sum_{x_1, x_2 \in \mathbb{F}_q^*} \frac{1}{p} \sum_{y \in \mathbb{F}_p} \zeta_p^{y(\text{Tr}_m(x_1 + x_2) - c)} \\ &= \frac{(q - 1)^2}{p} + \frac{1}{p} \sum_{y \in \mathbb{F}_p^*} \zeta_p^{-yc} \sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(yx_1)} \sum_{x_2 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(yx_2)}. \end{aligned}$$

Thus, we get the desired results. □

By Lemma 4 it is easy to see that the length of \mathcal{C}_D is $n_0 = \frac{(p^m-1)^2+p-1}{p}$.

For a codeword $\mathbf{c}(a_1, a_2)$ of \mathcal{C}_D and $\rho \in \mathbb{F}_p^*$, let $N_0 := N(a_1, a_2)$ be the number of components $\text{Tr}_m(a_1x_1^2 + a_2x_2^2)$ of $\mathbf{c}(a_1, a_2)$ which are equal to 0. Then

$$\begin{aligned} N_0 &= \sum_{x_1, x_2 \in \mathbb{F}_q^*} \left(\frac{1}{p} \sum_{y \in \mathbb{F}_p} \zeta_p^{y \text{Tr}_m(x_1+x_2)} \right) \left(\frac{1}{p} \sum_{z \in \mathbb{F}_p} \zeta_p^{z \text{Tr}_m(a_1x_1^2+a_2x_2^2)} \right) \\ &= \frac{1}{p^2} \sum_{x_1, x_2 \in \mathbb{F}_q^*} \left(1 + \sum_{y \in \mathbb{F}_p^*} \zeta_p^{y \text{Tr}_m(x_1+x_2)} \right) \left(1 + \sum_{z \in \mathbb{F}_p^*} \zeta_p^{z \text{Tr}_m(a_1x_1^2+a_2x_2^2)} \right) \\ &= \frac{(p^m-1)^2}{p^2} + \frac{1}{p^2}(\Omega_1 + \Omega_2 + \Omega_3), \end{aligned} \tag{4}$$

where

$$\Omega_1 = \sum_{y \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(yx_1)} \sum_{x_2 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(yx_2)} = p - 1,$$

$$\Omega_2 = \sum_{z \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_1x_1^2)} \sum_{x_2 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_2x_2^2)},$$

and

$$\Omega_3 = \sum_{y, z \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_1x_1^2+yx_1)} \sum_{x_2 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_2x_2^2+yx_2)}.$$

We are going to determine the values of Ω_2 and Ω_3 in Lemmas 5 and 6. To simplify formulas, denote $G_i = G(\eta)\eta(a_i)$ for $i \in \{1, 2\}$.

Lemma 5 *If $a_1 = 0$ and $a_2 = 0$, then*

$$\Omega_2 = (p^m - 1)^2(p - 1).$$

(1) *If m is even, then*

$$\Omega_2 = \begin{cases} (p^m - 1)(p - 1)(G_1 - 1), & \text{if } a_1 \neq 0, a_2 = 0, \\ (p^m - 1)(p - 1)(G_2 - 1), & \text{if } a_1 = 0, a_2 \neq 0, \\ (p - 1)(G_1G_2 - G_1 - G_2 + 1), & \text{if } a_1 \neq 0, a_2 \neq 0. \end{cases}$$

(2) *If m is odd, then*

$$\Omega_2 = \begin{cases} -(p^m - 1)(p - 1), & \text{if } a_1 \neq 0, a_2 = 0 \text{ or if } a_1 = 0, a_2 \neq 0, \\ (p - 1)(G_1G_2 + 1), & \text{if } a_1 \neq 0, a_2 \neq 0. \end{cases}$$

Proof When $a_1 = 0$ and $a_2 = 0$, it is obvious that Ω_2 is equal to $(p^m - 1)^2(p - 1)$.

If $a_1 \neq 0$ and $a_2 = 0$, then by the orthogonal property of additive characters, we have

$$\begin{aligned} \Omega_2 &= \sum_{z \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_1x_1^2)} \sum_{x_2 \in \mathbb{F}_q^*} 1 \\ &= (q - 1) \sum_{z \in \mathbb{F}_p^*} \left(\sum_{x_1 \in \mathbb{F}_q} \zeta_p^{\text{Tr}_m(za_1x_1^2)} - 1 \right). \end{aligned}$$

By Lemma 2, we obtain

$$\begin{aligned} \Omega_2 &= (q - 1) \sum_{z \in \mathbb{F}_p^*} (G(\eta)\eta(za_1) - 1) \\ &= (p^m - 1)G_1 \sum_{z \in \mathbb{F}_p^*} \eta(z) - (p^m - 1)(p - 1). \end{aligned}$$

By (3), we get the results. Similarly, we compute the value of Ω_2 when $a_1 = 0$ and $a_2 \neq 0$.

If $a_1 \neq 0$ and $a_2 \neq 0$, then by Lemma 2, we get

$$\begin{aligned} \Omega_2 &= \sum_{z \in \mathbb{F}_p^*} \left(\sum_{x_1 \in \mathbb{F}_q} \zeta_p^{\text{Tr}_m(za_1x_1^2)} - 1 \right) \left(\sum_{x_2 \in \mathbb{F}_q} \zeta_p^{\text{Tr}_m(za_2x_2^2)} - 1 \right) \\ &= \sum_{z \in \mathbb{F}_p^*} (G(\eta)\eta(za_1) - 1)(G(\eta)\eta(za_2) - 1) \\ &= G_1G_2 \sum_{z \in \mathbb{F}_p^*} \eta(z^2) - G_1 \sum_{z \in \mathbb{F}_p^*} \eta(z) - G_2 \sum_{z \in \mathbb{F}_p^*} \eta(z) + (p - 1). \end{aligned}$$

By (3), we get the results. □

To simplify results, we denote $G(\eta)G(\eta_p)$ by G and $A_i = \eta(a_i)\eta_p(-\text{Tr}_m(a_i^{-1}))$ for $i \in \{1, 2\}$.

Lemma 6 *If $a_1 = 0$ and $a_2 = 0$, then*

$$\Omega_3 = (p - 1)^2.$$

Suppose that m is even.

(1) *If $a_1 \neq 0$ and $a_2 = 0$, then*

$$\Omega_3 = \begin{cases} -(p - 1)^2G_1 + (p - 1)^2, & \text{if } \text{Tr}_m(a_1^{-1}) = 0, \\ (p - 1)G_1 + (p - 1)^2, & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0. \end{cases}$$

(2) If $a_1 = 0$ and $a_2 \neq 0$, then

$$\Omega_3 = \begin{cases} -(p-1)^2 G_2 + (p-1)^2, & \text{if } \text{Tr}_m(a_2^{-1}) = 0, \\ (p-1)G_2 + (p-1)^2, & \text{if } \text{Tr}_m(a_2^{-1}) \neq 0. \end{cases}$$

(3) If $a_1 \neq 0$ and $a_2 \neq 0$, then

$$\Omega_3 = \begin{cases} (p-1)^2(G_1G_2 - G_1 - G_2 + 1), & \text{if } \text{Tr}_m(a_1^{-1}) = 0 \text{ and } \text{Tr}_m(a_2^{-1}) = 0, \\ (p-1)(-G_1G_2 + G_1 - (p-1)G_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_2^{-1}) = 0, \\ (p-1)(-G_1G_2 + G_2 - (p-1)G_1 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) = 0 \text{ and } \text{Tr}_m(a_2^{-1}) \neq 0, \\ (p-1)((p-1)G_1G_2 + G_1 + G_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0, \text{Tr}_m(a_2^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_1^{-1} + a_2^{-1}) = 0, \\ (p-1)(-G_1G_2 + G_1 + G_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0, \text{Tr}_m(a_2^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_1^{-1} + a_2^{-1}) \neq 0. \end{cases}$$

Suppose that m is odd.

(1) If $a_1 \neq 0$ and $a_2 = 0$, then

$$\Omega_3 = \begin{cases} (p-1)^2, & \text{if } \text{Tr}_m(a_1^{-1}) = 0, \\ -(p-1)GA_1 + (p-1)^2, & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0. \end{cases}$$

(2) If $a_1 = 0$ and $a_2 \neq 0$, then

$$\Omega_3 = \begin{cases} (p-1)^2, & \text{if } \text{Tr}_m(a_2^{-1}) = 0, \\ -(p-1)GA_2 + (p-1)^2, & \text{if } \text{Tr}_m(a_2^{-1}) \neq 0. \end{cases}$$

(3) If $a_1 \neq 0$ and $a_2 \neq 0$, then

$$\Omega_3 = \begin{cases} (p-1)((p-1)G_1G_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) = 0 \text{ and } \text{Tr}_m(a_2^{-1}) = 0, \\ (p-1)(-G_1G_2 - GA_1 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_2^{-1}) = 0, \\ (p-1)(-G_1G_2 - GA_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) = 0 \text{ and } \text{Tr}_m(a_2^{-1}) \neq 0, \\ (p-1)((p-1)G_1G_2 - GA_1 - GA_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0, \text{Tr}_m(a_2^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_1^{-1} + a_2^{-1}) = 0, \\ (p-1)(-G_1G_2 - GA_1 - GA_2 + (p-1)), & \text{if } \text{Tr}_m(a_1^{-1}) \neq 0, \text{Tr}_m(a_2^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_1^{-1} + a_2^{-1}) \neq 0. \end{cases}$$

Proof We only compute the value of Ω_3 for the case $a_1 \neq 0$ and $a_2 \neq 0$. One can compute the other cases similarly. By the orthogonal property of additive characters, we have

$$\begin{aligned} \Omega_3 &= \sum_{y, z \in \mathbb{F}_p^*} \sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_1x_1^2+yx_1)} \sum_{x_2 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_2x_2^2+yx_2)} \\ &= \sum_{y, z \in \mathbb{F}_p^*} \left(\sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_1x_1^2+yx_1)} - 1 \right) \left(\sum_{x_1 \in \mathbb{F}_q^*} \zeta_p^{\text{Tr}_m(za_1x_1^2+yx_1)} - 1 \right). \end{aligned}$$

By Lemma 2, we obtain

$$\begin{aligned} \Omega_3 &= \sum_{y, z \in \mathbb{F}_p^*} \left(\zeta_p^{\text{Tr}_m(-y^2((4a_1)^{-1}+(4a_2)^{-1}))} G(\eta)^2 \eta(a_1a_2) - \zeta_p^{\text{Tr}_m(-y^2(4a_1)^{-1})} G(\eta) \eta(za_1) \right. \\ &\quad \left. - \zeta_p^{\text{Tr}_m(-y^2(4a_2)^{-1})} G(\eta) \eta(za_2) + 1 \right) \\ &= G_1 G_2 \sum_{z \in \mathbb{F}_p^*} \sum_{y \in \mathbb{F}_p^*} \zeta_p^{-y^2(4z)^{-1} \text{Tr}_m(a_1^{-1}+a_2^{-1})} - G_1 \sum_{z \in \mathbb{F}_p^*} \eta(z) \sum_{y \in \mathbb{F}_p^*} \zeta_p^{-y^2(4z)^{-1} \text{Tr}_m(a_1^{-1})} \\ &\quad - G_2 \sum_{z \in \mathbb{F}_p^*} \eta(z) \sum_{y \in \mathbb{F}_p^*} \zeta_p^{-y^2(4z)^{-1} \text{Tr}_m(a_2^{-1})} + (p-1)^2. \end{aligned} \tag{5}$$

If one of $\text{Tr}_m(a_1^{-1})$, $\text{Tr}_m(a_2^{-1})$ and $\text{Tr}_m(a_1^{-1}+a_2^{-1})$ is zero, then it is easy to compute the term corresponding to it. We only consider the case of $\text{Tr}_m(a_1^{-1}) \neq 0$, $\text{Tr}_m(a_2^{-1}) \neq 0$ and $\text{Tr}_m(a_1^{-1} + a_2^{-1}) \neq 0$. Other cases can be computed similarly. From (5) we have

$$\begin{aligned} \Omega_3 &= G_1 G_2 \sum_{z \in \mathbb{F}_p^*} \left(\sum_{y \in \mathbb{F}_p} \zeta_p^{-y^2(4z)^{-1} \text{Tr}_m(a_1^{-1}+a_2^{-1})} - 1 \right) \\ &\quad - G_1 \sum_{z \in \mathbb{F}_p^*} \eta(z) \left(\sum_{y \in \mathbb{F}_p} \zeta_p^{-y^2(4z)^{-1} \text{Tr}_m(a_1^{-1})} - 1 \right) \\ &\quad - G_2 \sum_{z \in \mathbb{F}_p^*} \eta(z) \left(\sum_{y \in \mathbb{F}_p} \zeta_p^{-y^2(4z)^{-1} \text{Tr}_m(a_2^{-1})} - 1 \right) + (p-1)^2. \end{aligned}$$

By Lemma 2, we obtain

$$\begin{aligned} \Omega_3 &= G_1 G_2 \sum_{z \in \mathbb{F}_p^*} (\eta_p(-(4z)^{-1}) \eta_p(\text{Tr}_m(a_1^{-1} + a_2^{-1}) G(\eta_p) - 1) \\ &\quad - G_1 \sum_{z \in \mathbb{F}_p^*} \eta(z) (\eta_p(-(4z)^{-1}) \eta_p(\text{Tr}_m(a_1^{-1}) G(\eta_p) - 1) \\ &\quad - G_2 \sum_{z \in \mathbb{F}_p^*} \eta(z) (\eta_p(-(4z)^{-1}) \eta_p(\text{Tr}_m(a_2^{-1}) G(\eta_p) - 1) + (p-1)^2 \end{aligned}$$

$$\begin{aligned}
 &= G_1 G_2 \sum_{z \in \mathbb{F}_p^*} (\eta_p(z) \eta_p(-\text{Tr}_m(a_1^{-1} + a_2^{-1})G(\eta_p) - 1) \\
 &\quad - G_1 \sum_{z \in \mathbb{F}_p^*} \eta(z) (\eta_p(z) \eta_p(-\text{Tr}_m(a_1^{-1})G(\eta_p) - 1) \\
 &\quad - G_2 \sum_{z \in \mathbb{F}_p^*} \eta(z) (\eta_p(z) \eta_p(-\text{Tr}_m(a_2^{-1})G(\eta_p) - 1) + (p - 1)^2.
 \end{aligned}$$

By (3), we get the result. □

Then by Lemmas 5 and 6, we obtain the values of N_0 . To get the frequency of each composition, we need the following lemmas.

Lemma 7 [1, Lemma 3.4] *For any $c \in \mathbb{F}_p$, let*

$$m_c = |\{a \in \mathbb{F}_q^* : \text{Tr}_m(a^{-1}) = c\}|.$$

Then we have

$$m_c = \begin{cases} p^{m-1} - 1, & \text{if } c = 0, \\ p^{m-1}, & \text{if } c \neq 0. \end{cases}$$

Lemma 8 [1, Lemma 3.5] *For any $c \in \mathbb{F}_p$, let*

$$n_{i,c} = |\{a \in \mathbb{F}_q^* : \eta(a) = i \text{ and } \text{Tr}_m(a^{-1}) = c\}|, \quad i \in \{-1, 1\}.$$

(1) *If m is even, then*

$$n_{i,c} = \begin{cases} \frac{1}{2p}(q - p + i(p - 1)G(\eta)), & \text{if } c = 0, \\ \frac{1}{2p}(q - iG(\eta)), & \text{if } c \neq 0. \end{cases}$$

(2) *If m is odd, then*

$$n_{i,c} = \begin{cases} \frac{1}{2p}(q - p), & \text{if } c = 0, \\ \frac{1}{2p}(q + i\eta_p(-c)G), & \text{if } c \neq 0. \end{cases}$$

Proof If $c = 0$, then we get the result from [1, Lemma 3.5] with $t = 1$. If $c \neq 0$, then by the orthogonal property of additive characters, we have

$$\begin{aligned}
 n_{1,c} &= \sum_{a \in C_0^{(2,q)}} \frac{1}{p} \sum_{x \in \mathbb{F}_p} \zeta_p^{x(\text{Tr}_m(a^{-1})-c)} \\
 &= \sum_{a \in C_0^{(2,q)}} \frac{1}{p} \left(\sum_{x \in \mathbb{F}_p^*} \zeta_p^{x(\text{Tr}_m(a^{-1})-c)} + 1 \right)
 \end{aligned}$$

$$= \frac{1}{p} \left(\sum_{x \in \mathbb{F}_p^*} \zeta_p^{-cx} \sum_{a \in C_0^{(2,q)}} \zeta_p^{\text{Tr}_m(a^{-1}x)} + \frac{q-1}{2} \right). \tag{6}$$

Assume that m is even, then 2 divides $\frac{q-1}{p-1}$ and so $\mathbb{F}_p^* \subseteq C_0^{(2,q)}$. By (6) we obtain

$$n_{1,c} = \frac{1}{p} \left(\sum_{x \in \mathbb{F}_p^*} \zeta_p^{-cx} \eta_0^{(2,q)} + \frac{q-1}{2} \right).$$

Thus, we get the results. Also the case of $n_{-1,c}$ is proved similarly.

Now suppose that m is odd, then $\mathbb{F}_p^* = \{\mathbb{F}_p^* \cap C_0^{(2,q)}\} \cup \{\mathbb{F}_p^* \cap C_1^{(2,q)}\}$. i.e., $|\mathbb{F}_p^* \cap C_0^{(2,q)}| = |\mathbb{F}_p^* \cap C_1^{(2,q)}| = \frac{p-1}{2}$. By (6) we obtain

$$\begin{aligned} n_{1,c} &= \frac{1}{p} \left(\sum_{x \in \mathbb{F}_p^* \cap C_0^{(2,q)}} \zeta_p^{-cx} \sum_{a \in C_0^{(2,q)}} \zeta_p^{\text{Tr}_m(a^{-1}x)} + \sum_{x \in \mathbb{F}_p^* \cap C_1^{(2,q)}} \zeta_p^{-cx} \sum_{a \in C_0^{(2,q)}} \zeta_p^{\text{Tr}_m(a^{-1}x)} + \frac{q-1}{2} \right) \\ &= \frac{1}{p} \left(\sum_{x \in \mathbb{F}_p^* \cap C_0^{(2,q)}} \zeta_p^{-cx} \eta_0^{(2,p)} + \sum_{x \in \mathbb{F}_p^* \cap C_1^{(2,q)}} \zeta_p^{-cx} \eta_1^{(2,p)} + \frac{q-1}{2} \right). \end{aligned}$$

If $-c \in C_0^{(2,p)}$, then we have

$$n_{1,c} = \frac{1}{p} \left(\eta_0^{(2,p)} \eta_0^{(2,q)} + \eta_1^{(2,p)} \eta_1^{(2,q)} + \frac{q-1}{2} \right).$$

If $-c \in C_1^{(2,p)}$, then we have

$$n_{1,c} = \frac{1}{p} \left(\eta_1^{(2,p)} \eta_0^{(2,q)} + \eta_0^{(2,p)} \eta_1^{(2,q)} + \frac{q-1}{2} \right).$$

It is easily checked that $\eta_0^{(2,p)} \eta_0^{(2,q)} + \eta_1^{(2,p)} \eta_1^{(2,q)} = \frac{G+1}{2}$ and $\eta_1^{(2,p)} \eta_0^{(2,q)} + \eta_0^{(2,p)} \eta_1^{(2,q)} = \frac{-G+1}{2}$. Thus, we get the results. Also $n_{-1,c}$ is computed similarly. This completes the proof. \square

Lemma 9 [1, Lemma 3.7] *Suppose that m is odd, let*

$$n'_{i,j} = |\{a \in \mathbb{F}_q^* : \eta(a) = i \text{ and } \eta_p(-\text{Tr}_m(a^{-1})) = j\}|, \quad i, j \in \{-1, 1\}.$$

Then we have

$$n'_{i,j} = \frac{p-1}{4p} (q + ijG).$$

Table 1 The weight distribution of C_D for $m = 2$

Weight	Frequency
0	1
$p(p - 1)^2$	$(p + 1)(p - 1)$
$(p - 1)(p^2 \pm p - 1)$	$(p \pm 1)(p - 1)$
$(p - 1)(p^2 \pm p - 2)$	$\frac{(p^2 \mp 1)(p - 1)}{2}$
$(p - 1)^2(p + 1)$	$2p(p - 1)^2$
$(p - 1)(p^2 - 2)$	$p^2(p - 1)(p - 2)$

Table 2 The weight distribution of C_D for even $m \geq 4$

Weight	Frequency
0	1
$(p - 1) \left(p^{2(m-1)} - p^{m-2} \pm p^{\frac{m-4}{2}}(p^m - p) \right)$	$p^m \pm p^{\frac{m-2}{2}}(p - 1)$
$(p - 1) \left(p^{2(m-1)} - p^{m-2} \pm p^{\frac{3m-4}{2}} \right)$	$(p - 1) \left(p^{m-1} \pm p^{\frac{m-2}{2}} \right)$
$(p - 1) \left(p^{2(m-1)} - 2p^{m-2} - p^{m-1} \pm 2p^{\frac{m-2}{2}} \right)$	$\frac{\left(p^{m-1} - 1 \pm p^{\frac{m-2}{2}}(p - 1) \right)^2}{4}$
$(p - 1) \left(p^{2(m-1)} - 2p^{m-2} \pm p^{\frac{m-4}{2}} \right)$	$(p - 1) \left(p^{2(m-1)} - p^{m-1} \pm p^{\frac{3m-4}{2}}(p - 1) \right)$
$(p - 1)(p^{2(m-1)} - 2p^{m-2} + p^{m-1})$	$\frac{(p^{m-1} - 1)(p^m - 1)}{2}$
$(p - 1)(p^{2(m-1)} - 2p^{m-2} - p^{m-1})$	$\frac{p^{m-2}(p - 1)(p^m + 1)}{2}$
$(p - 1)(p^{2(m-1)} - 2p^{m-2})$	$p^{2(m-1)}(p - 1)(p - 2)$

Theorem 1 Let C_D be a linear code defined by (1) and (2) where $D = \{(x_1, x_2) \in (\mathbb{F}_q^*)^2 : \text{Tr}_m(x_1 + x_2) = 0\}$. Suppose that m is even. If $m = 2$, then the weight distribution of C_D is given by Table 1 and the code C_D has parameters $\left[\frac{(p^2 - 1)^2 + p - 1}{p}, 4, (p - 1)(p^2 - p - 2) \right]$. If $m \geq 4$, then the weight distribution of C_D is given by Table 2 and the code C_D has parameters $\left[\frac{(p^{m-1})^2 + p - 1}{p}, 2m, (p - 1)(p^{2(m-1)} - p^{m-2} - p^{\frac{3m-4}{2}}) \right]$.

Proof Recall that $N_0 = \frac{(q-1)^2}{p^2} + \frac{1}{p^2}(\Omega_1 + \Omega_2 + \Omega_3)$. We employ Lemmas 5 and 6 to compute N_0 .

Assume that $a_1 \neq 0$ and $a_2 = 0$. If $\text{Tr}_m(a_1^{-1}) = 0$, then we obtain

$$N_0 = \frac{(q - 1)^2}{p^2} + \frac{(p - 1)}{p^2}((q - p)G(\eta)\eta(a_1) + p - q + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}((q-p)(G(\eta) - 1) + 1), & \text{if } \eta(a_1) = 1, \\ \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}((q-p)(G(\eta) + 1) - 1), & \text{if } \eta(a_1) = -1. \end{cases}$$

Now the frequencies are $n_{1,0}$ and $n_{-1,0}$ in Lemma 8, respectively.

If $\text{Tr}_m(a_1^{-1}) \neq 0$, then we obtain

$$N_0 = \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(qG(\eta)\eta(a_1) + p - q + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(q(G(\eta) - 1) + p + 1), & \text{if } \eta(a_1) = 1, \\ \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(q(G(\eta) + 1) - p - 1), & \text{if } \eta(a_1) = -1. \end{cases}$$

Now the frequencies are $\sum_{c \in \mathbb{F}_p^*} n_{1,c}$ and $\sum_{c \in \mathbb{F}_p^*} n_{-1,c}$, respectively. If $a_1 = 0$ and $a_2 \neq 0$, then we also have the same weights and the same frequencies with the case of $a_1 \neq 0$ and $a_2 = 0$.

Now, assume that $a_1 \neq 0$ and $a_2 \neq 0$.

If $\text{Tr}_m(a_1^{-1}) = \text{Tr}_m(a_2^{-1}) = 0$, then we obtain

$$N_0 = \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2\eta(a_1a_2) - pG(\eta)\eta(a_1) - pG(\eta)\eta(a_2) + p + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2 - 2pG(\eta) + p + 1), & \text{if } \eta(a_1) = 1 \text{ and } \eta(a_2) = 1, \\ \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(pG(\eta)^2 - p - 1), & \text{if } \eta(a_1a_2) = -1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2 + 2pG(\eta) + p + 1), & \text{if } \eta(a_1) = -1 \text{ and } \eta(a_2) = -1. \end{cases}$$

Now the frequencies are $(n_{1,0})^2$, $2n_{1,0}n_{-1,0}$, $(n_{-1,0})^2$, respectively.

If $\text{Tr}_m(a_1^{-1}) \neq 0$ and $\text{Tr}_m(a_2^{-1}) = 0$, then we have

$$N_0 = \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(pG(\eta)\eta(a_2) - p - 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(pG(\eta) - p - 1), & \text{if } \eta(a_2) = 1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta) + p + 1), & \text{if } \eta(a_2) = -1. \end{cases}$$

Now the frequencies are $(\sum_{c \in \mathbb{F}_p^*} m_c) n_{1,0}$ and $(\sum_{c \in \mathbb{F}_p^*} m_c) n_{-1,0}$, in Lemmas 7 and 8, respectively.

If $\text{Tr}_m(a_1^{-1}) = 0$ and $\text{Tr}_m(a_2^{-1}) \neq 0$, then we have the same weights and the same frequencies with the case of $\text{Tr}_m(a_1^{-1}) \neq 0$ and $\text{Tr}_m(a_2^{-1}) = 0$.

If $\text{Tr}_m(a_1^{-1}) \neq 0$, $\text{Tr}_m(a_2^{-1}) \neq 0$ and $\text{Tr}_m(a_1^{-1} + a_2^{-1}) = 0$, then we have

$$N = \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2 \eta(a_1 a_2) + p + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2 + p + 1), & \text{if } \eta(a_1 a_2) = 1, \\ \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(pG(\eta)^2 - p - 1), & \text{if } \eta(a_1 a_2) = -1. \end{cases}$$

Now the frequencies are $\sum_{c \in \mathbb{F}_p^*} n_{1,c}^2 + n_{-1,c}^2$ and $2 \sum_{c \in \mathbb{F}_p^*} n_{1,c} n_{-1,c}$, respectively.

If $\text{Tr}_m(a_1^{-1}) \neq 0$, $\text{Tr}_m(a_2^{-1}) \neq 0$ and $\text{Tr}_m(a_1^{-1} + a_2^{-1}) \neq 0$, then we have

$$N_0 = \frac{(q-1)}{p^2} + \frac{(p-1)}{p^2}(p + 1).$$

And the frequency is

$$|\{a_1, a_2 \in \mathbb{F}_q^* : \text{Tr}_m(a_1^{-1}) \neq 0, \text{Tr}_m(a_2^{-1}) \neq 0, \text{ and } \text{Tr}_m(a_1^{-1} + a_2^{-1}) \neq 0\}|.$$

It is equal to $T - \sum_{c \in \mathbb{F}_p^*} (m_c m_{-c})$, where $T = |\{(a_1, a_2) \in (\mathbb{F}_q^*)^2 : \text{Tr}_m(a_1^{-1}) \neq 0 \text{ and } \text{Tr}_m(a_2^{-1}) \neq 0\}|$. By Lemmas 4, 7 we get

$$T = \left(\sum_{c \in \mathbb{F}_q^*} m_c \right)^2 = \frac{(p-1)^2 q^2}{p^2} \text{ and } \sum_{c \in \mathbb{F}_p^*} (m_c m_{-c}) = \frac{(p-1) q^2}{p^2}.$$

Thus we compute the frequency.

Since the Hamming weight of $\mathbf{c}(a_1, a_2)$ is equal to $W_H(\mathbf{c}(a_1, a_2)) = n_0 - N_0$, we immediately have the desired results. \square

Example 1 (1) Let $p = 3$ and $m = 2$. Then $q = 9$ and $n = 22$. By Theorem 1, the code C_D is a $[22, 4, 8]$ linear code. Its weight enumerator is

$$1 + 10x^8 + 4x^{10} + 18x^{14} + 8x^{12} + 24x^{16} + 8x^{20} + 8x^{22},$$

which is checked by Magma.

(2) Let $p = 5$ and $m = 2$. Then $q = 25$ and $n = 116$. By Theorem 1, the code C_D is a $[116, 4, 72]$ linear code. Its weight enumerator is

$$1 + 52x^{72} + 16x^{76} + 24x^{80} + 300x^{92} + 160x^{96} + 48x^{112} + 24x^{116},$$

which is checked by Magma.

(3) Let $p = 3$ and $m = 4$. Then $q = 81$ and $n = 2134$. By Theorem 1, the code C_D is a $[2134, 8, 1278]$ linear code. Its weight enumerator is

$$\begin{aligned} &1 + 48x^{1278} + 32x^{1284} + 100x^{1356} + 738x^{1368} \\ &+ 256x^{1380} + 1080x^{1416} + 1458x^{1422} \\ &+ 1728x^{1428} + 1040x^{1476} + 20x^{1596} + 60x^{1602}, \end{aligned}$$

which is checked by Magma.

Theorem 2 Let C_D be a linear code defined by (1) and (2) where $D = \{(x_1, x_2) \in (\mathbb{F}_q^*)^2 : \text{Tr}_m(x_1 + x_2) = 0\}$. Suppose that m is odd and $m \geq 3$. Then the weight distribution of C_D is given by Table 3 and the code C_D has parameters $\left[\frac{(p^m-1)^2+p-1}{p}, 2m, (p-1)(p^{2(m-1)} - 2p^{m-2} - p^{m-1} - 2p^{\frac{m-3}{2}}) \right]$.

Proof Recall that $N_0 = \frac{(q-1)^2}{p^2} + \frac{1}{p^2}(\Omega_1 + \Omega_2 + \Omega_3)$. We employ Lemmas 5 and 6 to compute N_0 .

Suppose that $a_1 \neq 0$ and $a_2 = 0$.

If $\text{Tr}_m(a_1^{-1}) = 0$, then we obtain

$$N_0 = \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(p-q+1).$$

Now the frequency is m_0 in Lemma 7.

If $\text{Tr}_m(a_1^{-1}) \neq 0$, then we obtain

$$N_0 = \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(-GA_1 + p - q + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(G - p + q - 1), & \text{if } A_1 = 1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(G + p - q + 1), & \text{if } A_1 = -1. \end{cases}$$

Table 3 The weight distribution of C_D for odd $m \geq 3$

Weight	Frequency
0	1
$(p - 1)(p^{2(m-1)} - p^{m-2})$	$2(p^{m-1} - 1)$
$(p - 1)\left(p^{2(m-1)} - p^{m-2} \pm p^{\frac{m-3}{2}}\right)$	$(p - 1)\left(p^{m-1} \pm p^{\frac{m-1}{2}}\right)$
$(p - 1)(p^{2(m-1)} - 2p^{m-2} + p^{m-1})$	$\frac{(p^{2m-1} - p^{m-1} - p^m + 1)}{2}$
$(p - 1)(p^{2(m-1)} - 2p^{m-2} - p^{m-1})$	$\frac{(p^{m-1} - 1)^2}{2}$
$(p - 1)\left(p^{2(m-1)} - 2p^{m-2} \pm p^{\frac{m-3}{2}}\right)$	$(p - 1)\left(p^{m-1} \pm p^{\frac{m-1}{2}}\right)(p^{m-1} - 1)$
$(p - 1)\left(p^{2(m-1)} - 2p^{m-2} - p^{m-1} \pm 2p^{\frac{m-3}{2}}\right)$	$\frac{p^{m-1}(p-1)\left(p^{\frac{m-1}{2}} \pm 1\right)^2}{4}$
$(p - 1)\left(p^{2(m-1)} - 2p^{m-2} \pm 2p^{\frac{m-3}{2}}\right)$	$\frac{p^{m-1}(p-1)(p-2)\left(p^{\frac{m-1}{2}} \pm 1\right)^2}{4}$
$(p - 1)(p^{2(m-1)} - 2p^{m-2})$	$\frac{p^{m-1}(p-1)(p-2)(p^{m-1}-1)}{2}$

Now the frequencies are $n'_{1,1} + n'_{-1,-1}$ and $n'_{1,-1} + n'_{-1,1}$ in Lemma 9, respectively. If $a_1 = 0$ and $a_2 \neq 0$, then we have the same values and the same frequencies with the case of $a_1 \neq 0$ and $a_2 = 0$.

Now, assume that $a_1 \neq 0$ and $a_2 \neq 0$.

If $\text{Tr}_m(a_1^{-1}) = \text{Tr}_m(a_2^{-1}) = 0$, then we obtain

$$N_0 = \frac{(q - 1)^2}{p^2} + \frac{(p - 1)}{p^2} (pG(\eta)^2 \eta(a_1 a_2) + p + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2} (pG(\eta)^2 + p + 1), & \text{if } \eta(a_1 a_2) = 1, \\ \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2} (pG(\eta)^2 - p - 1), & \text{if } \eta(a_1 a_2) = -1. \end{cases}$$

Now the frequencies are $n^2_{1,0} + n^2_{-1,0}$ and $2n_{1,0}n_{-1,0}$ in Lemma 8, respectively.

If $\text{Tr}_m(a_1^{-1}) \neq 0$ and $\text{Tr}_m(a_2^{-1}) = 0$, then we have

$$N_0 = \frac{(q - 1)^2}{p^2} + \frac{(p - 1)}{p^2} (-GA_1 + p + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(G - p - 1), & \text{if } A_1 = 1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(G + p + 1), & \text{if } A_1 = -1. \end{cases}$$

Now the frequencies are $m_0(n'_{1,1} + n'_{-1,-1})$ and $m_0(n'_{1,-1} + n'_{-1,1})$, respectively. If $\text{Tr}_m(a_1^{-1}) = 0$ and $\text{Tr}_m(a_2^{-1}) \neq 0$, then we have the same values and the same frequencies with the case of $\text{Tr}_m(a_1^{-1}) \neq 0$ and $\text{Tr}_m(a_2^{-1}) = 0$.

If $\text{Tr}_m(a_1^{-1}) \neq 0$, $\text{Tr}_m(a_2^{-1}) \neq 0$ and $\text{Tr}_m(a_1^{-1} + a_2^{-1}) = 0$, then we have

$$N_0 = \frac{(q - 1)^2}{p^2} + \frac{(p - 1)}{p^2}(pG(\eta)^2\eta(a_1a_2) - GA_1 - GA_2 + p + 1).$$

Since $\text{Tr}_m(a_1^{-1} + a_2^{-1}) = 0$, we have $A_1A_2 = \eta(a_1a_2)\eta_p(-1)$.

Assume that $p \equiv 1 \pmod{4}$. Then,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2 - 2G + p + 1), & \text{if } \eta(a_1a_2) = A_1 = A_2 = 1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(pG(\eta)^2 + 2G + p + 1), & \text{if } \eta(a_1a_2) = 1 \text{ and } A_1 = A_2 = -1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(-pG(\eta)^2 + p + 1), & \text{if } \eta(a_1a_2) = A_1A_2 = -1. \end{cases}$$

Now the frequencies are $\sum_{c \in C_0^{(2,p)}} n_{1,c}n_{1,-c} + \sum_{c \in C_1^{(2,p)}} n_{-1,c}n_{-1,-c}$, $\sum_{c \in C_1^{(2,p)}} n_{1,c}n_{1,-c} + \sum_{c \in C_0^{(2,p)}} n_{-1,c}n_{-1,-c}$, $\sum_{c \in C_0^{(2,p)}} n_{1,c}n_{1,-c} + \sum_{c \in C_1^{(2,p)}} n_{-1,c}n_{-1,-c}$ and $\sum_{c \in C_0^{(2,p)}} n_{-1,c}n_{1,-c} + \sum_{c \in C_1^{(2,p)}} n_{-1,c}n_{1,-c}$, in Lemma 8, respectively.

In the case of $p \equiv 3 \pmod{4}$, we compute similarly.

If $\text{Tr}_m(a_1^{-1}) \neq 0$, $\text{Tr}_m(a_2^{-1}) \neq 0$ and $\text{Tr}_m(a_1^{-1} + a_2^{-1}) \neq 0$, then we have

$$N_0 = \frac{(q - 1)^2}{p^2} + \frac{(p - 1)}{p^2}(-GA_1 - GA_2 + p + 1).$$

Thus,

$$N_0 = \begin{cases} \frac{(q-1)^2}{p^2} - \frac{(p-1)}{p^2}(2G - p - 1), & \text{if } A_1 = A_2 = 1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)}{p^2}(2G + p + 1), & \text{if } A_1 = A_2 = -1, \\ \frac{(q-1)^2}{p^2} + \frac{(p-1)(p+1)}{p^2}, & \text{if } A_1A_2 = -1. \end{cases}$$

We compute the frequency for the case of $A_1A_2 = -1$. We compute similarly for the other cases. From Lemma 8, the frequency is

$$\begin{aligned}
 & \sum_{c \in C_0^{(2,p)}} \frac{1}{2p} (\eta_p(c)G + q) \sum_{\substack{d \in C_0^{(2,p)} \\ d \neq -c}} \frac{1}{2p} (\eta_p(d)G + q) \\
 & + \sum_{c \in C_0^{(2,p)}} \frac{1}{2p} (\eta_p(c)G + q) \sum_{\substack{d \in C_1^{(2,p)} \\ d \neq -c}} \frac{1}{2p} (-\eta_p(d)G + q) \\
 & + \sum_{c \in C_1^{(2,p)}} \frac{1}{2p} (-\eta_p(c)G + q) \sum_{\substack{d \in C_0^{(2,p)} \\ d \neq -c}} \frac{1}{2p} (\eta_p(d)G + q) \\
 & + \sum_{c \in C_1^{(2,p)}} \frac{1}{2p} (-\eta_p(c)G + q) \sum_{\substack{d \in C_1^{(2,p)} \\ d \neq -c}} \frac{1}{2p} (-\eta_p(d)G + q), \\
 & = \frac{1}{4p^2} (G + q)^2 \frac{p-1}{4} (p-1 + p-1 + p-3 + p-3), \\
 & = \frac{(p-1)(p-2)}{4p^2} (G + q)^2.
 \end{aligned}$$

□

Example 2 (1) Let $p = 3$ and $m = 3$. Then $q = 27$ and $n = 226$. By Theorem 2, the code C_D is a $[226, 6, 128]$ linear code. Its weight enumerator is

$$\begin{aligned}
 & 1 + 18x^{128} + 32x^{132} + 72x^{136} + 18x^{146} + 96x^{148} + 72x^{150} + 192x^{152} \\
 & + 84x^{154} + 16x^{156} + 24x^{158} + 104x^{168},
 \end{aligned}$$

which is checked by Magma.

(2) Let $p = 5$ and $m = 3$. Then $q = 125$ and $n = 3076$. By Theorem 2, the code C_D is a $[3076, 6, 2352]$ linear code. Its weight enumerator is

$$\begin{aligned}
 & 1 + 400x^{2352} + 288x^{2360} + 900x^{2368} + 1200x^{2452} + 1920x^{2456} + 3600x^{2460} \\
 & + 2880x^{2464} + 2700x^{2468} + 80x^{2476} + 48x^{2480} + 120x^{2484} + 1488x^{2560},
 \end{aligned}$$

which is checked by Magma.

4 Concluding remarks

Let w_{min} and w_{max} be the minimum and maximum nonzero weight of linear code C_D , respectively. We recall that if

$$\frac{w_{min}}{w_{max}} > \frac{p-1}{p},$$

then all nonzero codewords of code C_D are minimal (see [8]).

By Theorem 1, we easily check

$$\frac{w_{min}}{w_{max}} = \frac{(p-1) \left(p^{2(m-1)} - p^{m-2} - p^{\frac{3m-4}{2}} \right)}{(p-1) \left(p^{2(m-1)} - p^{m-2} + p^{\frac{3m-4}{2}} \right)} > \frac{p-1}{p},$$

where even $m \geq 4$. Moreover, by Theorem 2 we easily check

$$\frac{w_{min}}{w_{max}} = \frac{(p-1) \left(p^{2(m-1)} - 2p^{m-2} - p^{m-1} - 2p^{\frac{m-3}{2}} \right)}{(p-1) \left(p^{2(m-1)} - 2p^{m-1} + p^{m-1} \right)} > \frac{p-1}{p},$$

where odd $m \geq 3$.

Hence, the linear codes in this paper satisfy $w_{min}/w_{max} > (p-1)/p$ for $m \geq 3$, and can be used to get secret sharing schemes with interesting access structures.

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