ORIGINAL ARTICLE

# **Optimal tolerance allocation using a multiobjective particle swarm optimizer**

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Abstract Particle swarm optimizers are routinely utilized in engineering design problems, but much work remains to take advantage of their full potential in the combined areas of sensitivity analysis and tolerance synthesis. In this paper, a novel Pareto-based multiobjective formulation is proposed to enhance the operations of a particle swarm optimizer and systematically distribute tolerances among various components of a mechanical assembly. The enhanced algorithm relies on nonlinear sensitivity analysis and the statistical root sum squares model to simultaneously optimize product performance criteria, the manufacturing cost, and the stack-up tolerance. It is shown that the proposed algorithm can accomplish its optimization task by successfully shifting nominal values of design parameters instead of the expensive tightening of component tolerances. Several numerical experiments for optimal design of a stepped bar assembly were conducted, which highlight the advantages of the proposed methodology.

**Keywords** Concurrent engineering · Multiobjective optimization · Particle swarm optimization · Sensitivity analysis · Tolerance allocation

# **1** Introduction

Optimum tolerance allocation is an effective robust design and concurrent engineering tool to reduce manufacturing

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costs while increasing the overall robustness of the product [1-3].

Evolutionary algorithms have proven successful in handling many real-world multiobjective concurrent engineering problems [4–7]. For example, an evolutionary radial basis function network was proposed to model a robust design optimizer, and the developed system was used to study a number of computationally expensive multiobjective optimizations problems [8].

Recently, approaches based on particle swarm optimization (PSO) have received a great deal of attention in engineering design because of their simplicity and fast convergence rates [9–11]. A multiobjective particle swarm optimizer was developed to evolve generations of hyperrectangular particles where dimensional tolerances were treated as intervals. This optimizer did not directly minimize manufacturing costs but rather attempted to widen design tolerances by maximizing hyper-rectangular design volumes in successive generations [12]. Several other improved PSO systems have been proposed and successfully used in challenging engineering design problems [13–17], but none of them included any sensitivity analysis study while performing parameter design.

The major goal of sensitivity analysis is to identify the sensitive dimensions of parameters in a product or process design. Sensitivity analysis is an important component in robust engineering design where the values of design parameters are prone to many sources of uncertainty and variation. It is, therefore, necessary to understand the sensitivity of the product or process responses to the perturbation in the model inputs [18]. An ant colony optimizer was developed to perform sensitivity analysis and tolerance allocation in concept design by optimizing the manufacturing cost and the product yield using an aggregated form of design objectives [19]. It has, however,

been shown that the traditional weighted aggregated method suffers from many shortcomings, the most important of which is its assumption of convexity of the Pareto region, which is generally invalid in many real-world engineering design problems. Furthermore, Pareto fronts are typically not uniform, that is, the obtained solutions are not uniformly distributed on the Pareto front given a set of evenly distributed weights [20]. In another approach based on stochastic integer programming, a quality loss function was developed to include cost tolerances, manufacturing costs, and design constraints [21]. The traditional approaches to robust tolerance design, however, generally do not scale up as the complexity of the problem domain increases [3].

A PSO-based approach was proposed for worst-case circuit design of LC high-pass filters [22] but with the main underlying assumption that all the individual worst-case tolerance limits occur at the same time. This shortcoming can potentially result in unnecessarily tight tolerances and high manufacturing costs [3]. A PSO system was reported for achieving the multiple objectives of minimum quality loss function and minimum manufacturing cost for the machining tolerance allocation of an over running clutch assembly [23]. This system utilized the weighted aggregated method, which as explained before, cannot successfully identify all the trade-off solutions along the Pareto front.

In this paper, a novel Pareto-based multiobjective formulation is proposed to enhance the operations of a particle swarm optimizer and optimally distribute tolerances among various components of a mechanical system. The enhanced algorithm relies on nonlinear sensitivity analysis and the statistical root sum squares (RSS) model to simultaneously optimize product performance criteria, the manufacturing cost, and the stack-up tolerance. It is shown that the proposed algorithm can accomplish its optimization task by efficiently exploring the Pareto front and successfully shifting nominal values of design parameters instead of the expensive tightening of component tolerances. Furthermore, the algorithm is able to operate in a mixed mode where some of the design tolerances must be kept fixed due to manufacturing restrictions while others are allowed to widen.

The remainder of the paper is organized as follows. Sections 2 and 3 provide brief descriptions of sensitivity analysis and multiobjective optimization, respectively. Section 4 introduces the particle swarm optimization methodology along with the implemented enhancements for more efficient coverage of the feasible design region. Section 5 is the application of the proposed methodology to multiobjective concurrent engineering design of a stepped bar assembly. And finally, Section 6 is the summary and conclusions.

## 2 Sensitivity analysis

The major goal of robust engineering design is to minimize functional sensitivity of a product or process to parametric changes caused by uncontrollable operational and manufacturing conditions. Sensitivity analysis is an important component of robust design in that it utilizes various optimization methods to identify the sensitive dimensions of parameters in a design, and more importantly, it aims to ensure that effects on product performance brought by changes in the design parameters are minimal.

There are various approaches to performing sensitivity analysis. For instance, mathematical sensitivity analysis, component manufacturing process output distribution sensitivity analysis, nonlinear sensitivity analysis, and statistical sensitivity analysis, to name a few. The following is a brief description of the use of nonlinear sensitivity analysis in concept design as it directly relates to the focus of this work. More thorough discussions of the wide topic of sensitivity analysis, however, can be found elsewhere [1, 2, 18, 19, 21].

#### 2.1 Nonlinear sensitivity analysis

The functional relationship between the independent design parameters  $x_1, x_2, ..., x_n$  and the dependent product or process response Y can be expressed as:

$$Y = f(x_1, x_2, \dots, x_n) \tag{1}$$

In many robust engineering design applications, however, the above relationship is not linear in nature. Therefore, small changes in the response may be expressed by a Taylor's series expansion as given by [2]:

$$\Delta Y = \sum_{i} = \frac{\delta f}{\delta x_{i}} \Delta x_{i} + \frac{1}{2} \sum_{i} \sum_{j} \frac{\delta^{2} f}{\delta x_{i} \delta x_{j}} \Delta x_{i} \Delta x_{j} + \cdots$$
(2)

The common practice for performing nonlinear sensitivity analysis is to evaluate the first term from the above Taylor series expansion and to manually calculate a sensitivity S by taking the partial derivative of the functional relationship between the dependent dimension and the independent component dimensions as  $S_i = \delta Y / \delta x_i$ .

Manufacturing process variations in the real-world situations manifest themselves as distributions around the nominal design variable values and may cause the deviation of model response around its nominal value. This model deviation can be mathematically expressed as:

$$\delta Y = f(S_1 \delta x_1, S_2 \delta x_2, \dots, S_n \delta x_n) \tag{3}$$

Here,  $\delta Y$  is the product response variation, and  $S_i$  and  $\delta x_i$  are the sensitivity and deviation for the *i*th design variable, respectively.

The variance of the product response function can then be expressed as a sum of the design variable variances as attenuated by their respective sensitivity coefficients:

$$\mathrm{d}Y^2 = \sum_{i=1}^n (S_i \times \mathrm{d}x_i)^2 \tag{4}$$

A careful examination of the above relation reveals two design opportunities to control the deviation of the product response Y. First, dY can be kept small by tightening manufacturing tolerances  $dx_i$ . And second,  $S_i$  can be reduced to make dY insensitive to independent design parameter variations. The first approach, which was traditionally used, is generally more expensive as it requires tighter manufacturing tolerances and protection from aging and the environment. The second way to control the product or process deviation is more practical as it can be accomplished without tightening manufacturing tolerances but by simply altering the value of  $S_i \times dx_i$ .

The above approach is very helpful in avoiding the cost associated with quality improvement based on buying down variance by tolerance tightening. In fact, any number of conceptual designs can be analyzed very early in the concept development stage using various optimization techniques to minimize critical dimension variance [18, 19, 21].

# 3 Multiobjective optimization

Multiobjective optimization (MO) is a methodology for finding optimal solutions to multivariate problems with multiple, often conflicting, objectives [4, 5]. The main goal of MO is to find the optimum input parameter vector, which results in some desired combination of maximization, minimization, or nominalization of the involved product or process responses. Mathematically, MO attempts to optimize the *p*-dimensional vector function *F* of objective responses  $F(X) = [f_1(X), \ldots, f_p(X)]^T$  where  $X = [x_1, \ldots, x_n]^T$  is an *n*-dimensional vector of design variables and *p* is the number design objectives for a product. The problem can be stated formally as follows (*l* is the inequality and *m* is the equality constraints).

$$\begin{array}{ll} \text{Minimize/Maximize}: & f_i(X) \text{ for } i=1,\ldots,p\\ \text{Subject to}: & g_j(X) \leq 0 \text{ for } j=1,\ldots,l\\ & h_k(X)=0 \text{ for } k=1,\ldots,m \end{array}$$

Since MO problems often have conflicting objectives, it is virtually impossible to find any one ideal solution. Instead, MO produces Pareto-optimal solutions. A design vector is Pareto optimal if there is no other design vector that optimizes one criterion without causing the simultaneous degradation of at least one other criterion [24]. Without loss of generality, it can be stated that the goal is to maximize p

objective responses considering the following key definitions [4, 5, 13]:

Definition 1. A vector  $u = (u_1, \dots, u_p)$  is said to be inferior to (dominated by) vector  $v = (v_1, \dots, v_p)$  if and only if  $\forall i \in \{1, \dots, p\},$  $v_i \ge u_i \land \exists i \in \{1, \dots, p\} | v_i > u_i.$ 

Definition 2. Vectors u and v are said to be *noninferior* to each other if neither u is inferior to v nor v is inferior to u.

The main challenge of MO is to develop efficient algorithms that can quickly converge to well-spread Pareto fronts in the feasible, nondominated design regions. This and other pertinent issues relating to the specifics of MO can be found elsewhere [14–17].

## 4 Particle swarm optimization

PSO is a stochastic global optimization approach, and its main strength is in its simplicity and fast convergence rates. The following is a brief introduction to PSO [9–11]:

A total of p particles are randomly distributed throughout the feasible design region, where  $X_i^t$  is the position of a particle *i* representing a design scenario at time *t*. The position of the particle can be updated in the following manner:

$$X_i^{t+1} = X_i^t + V_i^{t+1}$$
(5)

with a velocity  $V_i^{t+1}$  which is calculated as:

$$\mathbf{V}_{i}^{t+1} = \omega V_{i}^{t} + c_{1} r_{1} \left( P_{i}^{t} - X_{i}^{t} \right) + c_{2} r_{2} \left( P_{g}^{t} - X_{i}^{t} \right)$$
(6)

Here, the point  $P_i^t$  is the best local solution found up to now (time t) and represents the cognitive contribution to the search vector  $V_i^{t+1}$ . The point  $P_g^t$  is the best global solution found among all particles so far and forms the social contribution to the velocity vector. Random numbers  $r_1$  and  $r_2$  are uniformly distributed in the interval [0,1]. The cognitive and social scaling factors  $c_1$  and  $c_2$  are typically selected such that  $c_1 \times c_1$  and  $c_2 \times c_2$  have a mean of 1 so that the particles overshoot the attraction points  $P_i^t$  and  $P_g^t$  half the time, thereby allowing wider search fronts [25]. The variable  $\omega$  is the inertia weight and is typically chosen in the range of [0,1]. A larger inertia weight facilitates global exploration, and a smaller inertia tends to facilitate local exploration. Therefore,  $\omega$  is a critical factor for the convergence behavior of PSO and is used to promote global exploration of the search space [11].

The cognitive learning factor is computed by the term  $c_1r_1(P_i^t - X_i^t)$  in Eq. 6, and it is the short-term memory of a particle representing the particle's inclination to repeat past



Fig. 1 The nearest neighbor density estimator

behavior that has proven to be successful for that particular particle.

The social learning factor, on the other hand, is computed by the term  $c_2 r_2 \left(P_g^t - X_i^t\right)$  in Eq. 6, and it is the peer pressure of a particle representing the particle's inclination to imitate or emulate the behavior of other particles that are successful; it is the influence of a particle's neighbors.

During the course of the PSO algorithm, swarm particles are evaluated, and based on those evaluations, the new velocities and positions of the particles are updated. The basic equations (Eqs. 5 and 6) for PSO were utilized in this work, but a few key elements were modified in accordance with enhancements described in the following sections.

# 4.1 Crowding distance computation

The nearest neighbor density estimator quantifies how crowded the closest neighbors of a given particle are in



Fig. 2 The fly-back mechanism



Fig. 3 Stepped bar assembly

the objective space. As illustrated in Fig. 1, this measure is estimated by the area of the largest cuboid formed by using the two nearest neighbors of particle i as the vertices [11].

Before computing the crowding distance  $(d_c)$  for particles on the boundary of the feasible region, the entire nondominated population is first sorted based on the increasing values of their objective functions, one at a time. A  $d_c^j$  measure for a particular particle  $P_i^j$  is then defined as the average distance of its two nearest neighbors  $P_{i-1}^j$  and  $P_{i+1}^j$  along the dimension of a specific objective function *j*. The total  $d_c$  value for a particle is then computed as  $\sum_{j=1}^{M} d_c^j$ , where *M* is the total number of design objective functions. In the case of the two extreme solutions with the highest and lowest objective function values, the  $d_c$ measure is set to infinity so the two boundary points will always be selected [4, 26].

#### 4.2 Elitist mutation

Although one of the greatest advantages of PSO-based approaches are their simplicity—both conceptually and from the standpoint of implementation—these stochastic can also experience difficulty controlling population diversity while dealing with multiple-objective optimization problems [4].



Fig. 4 A cost-tolerance curve



Fig. 5 Discovered Pareto fronts. a Crowding distances. b No crowding distance

The loss of diversity along the Pareto front can potentially be avoided by utilizing the mutation operation.

The mutation operator implemented in this work is based on the elitist-mutation mechanism that was proposed to improve the performance of the particle swarm algorithm [16]. The main objective here is to best identify the true Pareto-optimal front by promoting diversity among the particles maintained in the repository. In its initial phase, the elitist-mutation operator replaces the infeasible solutions in the repository with the least crowded ones, and in its later phases, it will exploit the sparsely explored regions along the Pareto front. The main steps to mutation are outlined below:

- 1. Sort and index the particles in the repository in ascending order based on their value of a randomly selected objective.
- 2. Use the crowding distance metric to sort the current solutions in the repository in descending order and randomly select a solution in the top 10%.

3. Apply mutation to a predefined number of particles in the current population.

## 4.3 Maintaining feasibility

For the PSO algorithm to maintain feasibility in the population kept in the repository, an intuitive approach is to fly back a particle to its previous position when it is outside the feasible region [27]. This is called the *fly-back* mechanism, and it is a particularly useful strategy when solving engineering optimization problems with multiple geometric constraints. Each time a constraint is violated, the particle in flight is made to revisit its previous position, allowing effective exploration of the search space along the boundaries of the feasible region.

In this study, a fly-back mechanism to the feasible region was implemented as illustrated in Fig. 2.

Figure 2 illustrates how a particle  $X_i$  might have flown into the infeasible design region  $(V_i^k)$  in the *k*th iteration of



Fig. 6 Average functional behavior of particles. a Average dY profile of population. b Average cost profile of population





Fig. 7 Best functional behavior of particles. a Minimum dY profile of population. b Minimum cost profile of population

the algorithm. Assuming that the position of the best particle  $P_g$  does not change, a fly back to the particle's best position in the previous iteration ( $P_i$ ) will ensure that the direction of the new velocity ( $V_i^{k+1}$ ) will point to the feasible space boundary but it will be closer to the global best  $P_g$ . Experimental results have shown that this mechanism is particularly suitable for mechanical design problems where optimal solutions lie on or near the feasible Pareto front [27]. The main guarantee of employing the fly back is that starting with an initial population of feasible solutions, each population generated in the subsequent iterations of the algorithm is also feasible.

Furthermore, in order to examine feasibility in a given population, there has to be a mechanism to compare two solutions in the repository. A simple yet elegant method based on NSGA-II was implemented in this work as outlined below [24]:

- Definition 3. A solution  $S_i$  is said to dominate a solution  $S_j$  if:
  - 1.  $S_i$  is feasible and  $S_i$  is infeasible.
  - 2.  $S_i$  and  $S_j$  are both infeasible but  $S_i$  violates fewer design constraints.

3.  $S_i$  and  $S_j$  are both feasible but  $S_i$  dominates  $S_j$  (see Definition 1).

## 4.4 Performance measures

The performance of an optimizer can be judged along two categories: *efficiency*, which is a measure of computational effort to obtain optimal solutions (e.g., number of function evaluations, CPU time, etc.), and *effectiveness*, which measures the accuracy and convergence of the obtained solutions [4]. A crucial consideration regarding effectiveness of a multiobjective optimizer is that not only the optimizer should converge to Pareto-optimal solutions but it should also maintain diversity along the feasible region [4, 5, 16, 24].

Since no single performance measure can handle both requirements for making progress toward the Paretooptimal front and for evaluating the spread or diversity of the obtained solutions, in this study, it was decided to utilize three performance metrics, which together better evaluate the overall effectiveness of a PSO-based multiobjective optimizer. Each of the utilized effectiveness



Fig. 8 Constraint profiles of the population. a Average Y (target=166.67). b Minimum Y (target=166.67)



Fig. 9 a-f Design factor contributions to functional variance

metrics is described below. Please note that in the ensuing discussions, the set of discovered solutions is referred to as Q and  $P^*$  is the true Pareto-optimal set.

# 4.4.1 Generational distance

This performance metric measures the proximity of the obtained Pareto-optimal solutions Q to the true Pareto-optimal solutions  $P^*$  [4, 5, 16]:

$$GD = \frac{\left(\sum_{i=1}^{|\mathcal{Q}|} d_i^p\right)^{1/p}}{|\mathcal{Q}|} \tag{7}$$

 Table 1
 Two obtained solutions by the PSO (experiment 1)

The Euclidean distance-based metric  $d_i$  (p=2) is the phenotypic distance between each member i of the obtained set Q and the closest member in  $P^*$  to that member:

$$d_{i} = \min_{k=1}^{|P^{*}|} \sqrt{\sum \left(f_{m}^{(i)} - f_{m}^{*(k)}\right)^{2}}$$
(8)

Here,  $f_m^{*(k)}$  is the value of the *m*th objective function of the *k*th in *P*\*. Observe that a successful swarm should discover solutions that produce a small value for GD.

Variable	Value	Tolerance	Sensitivity	Deviation	Factor contribution
<i>x</i> <sub>1</sub>	22	0.2	-0.047	0.009	0.07%
<i>x</i> <sub>2</sub>	300	0.4	-0.007	0.003	0.007%
$y_1$	67	0.3	0.860	0.258	52.8%
<i>y</i> <sub>2</sub>	82	0.5	0.139	0.069	3.82%
h	42	0.1	1.001	0.100	7.93%
r	55	0.2	1.056	0.211	35.35%
				<i>Y</i> =166.67, d <i>Y</i> =0	.355, Cost=\$12.71
$x_1$	23	0.2	-0.006	0.001	0.001%
<i>x</i> <sub>2</sub>	204	0.4	-0.001	0.0006	0.0003%
$y_1$	109	0.3	0.813	0.243	50.06%
<i>y</i> <sub>2</sub>	110	0.5	0.186	0.093	7.32%
h	45	0.1	1.000	0.100	8.41%
r	57	0.2	1.008	0.201	34.19%
				<i>Y</i> =166.67, d <i>Y</i> =0	.344, <i>Cost</i> =\$12.71

Table 2 The best solution obtained by a traditional technique

Variable	Value	Tolerance	Sensitivity	Deviation	Factor contribution	
<i>x</i> <sub>1</sub>	70	0.2	-0.212	0.042	1.15%	
<i>x</i> <sub>2</sub>	338	0.4	-0.145	0.058	2.07%	
$y_1$	65	0.3	0.932	0.279	49.24%	
<i>y</i> <sub>2</sub>	128	0.5	0.108	0.054	1.82%	
h	35	0.1	1.025	0.102	6.29%	
r	65	0.2	1.251	0.250	39.43%	
				<i>Y</i> =166.67, d <i>Y</i> =0.3	<i>Y</i> =166.67, d <i>Y</i> =0.396, Cost=not available	

# 4.4.2 Spread

The spread metric ( $\Delta$ ) is used as an indicator of how well the solutions are distributed along the PSO's discovered Pareto front Q [4]:

$$\Delta = \frac{\sum_{m=1}^{M} d_m^e + \sum_{i=1}^{|Q|} |d_{i-}\overline{d}|}{\sum_{m=1}^{M} d_m^e + |Q|\overline{d}}$$
(9)

where  $d_i$  is the Euclidean distance between a solution *i* and its nearest member in *Q*, and  $\overline{d}$  defined as:

$$\overline{d} = \frac{1}{|Q|} \sum_{i=1}^{|Q|} d_i \tag{10}$$

is the mean value of these distances. The parameter  $d_m^e$  is the distance between the extreme solutions in Q and  $P^*$ relating to an objective function m where there are a total of M response functions in the problem. An algorithm finding a smaller value of  $\Delta$  (i.e., closer to zero) is better able to identify a diverse set of nondominated solutions.

# 4.4.3 Set coverage

This metric can be used as a measure of the relative spread of solutions between two sets of solutions A and

 Table 3 Two obtained solutions by the PSO (experiment 2)

*B*. The set coverage C(A,B) calculates the proportion of solutions in *B*, which are weakly dominated by solutions in *A* [4, 5]:

$$C(A,B) = \frac{|\{b \in B | \exists a \in A : a \le b\}|}{|B|}$$
(11)

If C(A,B)=1, all solutions in *B* are weakly dominated by *A*, whereas if C(A,B)=0, none of the solutions in *B* are weakly dominated by *A*. Thus,  $C(Q, P^*)$  should always be 0.

# 5 Design of a stepped bar assembly

The optimum design of a stepped bar assembly, as shown in Fig. 3, has been previously attempted using traditional calculus-based sensitivity analysis [2] and continuous ant colony optimization [19].

# 5.1 The design formulation

Consider the assembly shown Fig. 3. The height of the center of the cylinder, Y, is the only critical dimension in this problem, and it is dependent on the value of design

Variable	Value	Tolerance	Sensitivity	Deviation	Factor contribution
<i>x</i> <sub>1</sub>	30	0.09	-0.117	0.0106	0.82%
<i>x</i> <sub>2</sub>	243	0.21	-0.025	0.0054	0.21%
$y_1$	93	0.07	0.820	0.0575	24.23%
<i>Y</i> <sub>2</sub>	123	0.11	0.179	0.0197	2.85%
h	9	0.07	1.010	0.0720	36.70%
r	58	0.06	1.153	0.0690	35.17%
				<i>Y</i> =165.24, d <i>Y</i> =0	.12, Cost=\$20.00
<i>x</i> <sub>1</sub>	20	0.09	-0.021	0.002	0.029%
<i>x</i> <sub>2</sub>	206	0.21	-0.007	0.002	0.020%
$y_1$	88	0.07	0.736	0.052	21.71%
<i>y</i> <sub>2</sub>	93	0.11	0.264	0.029	6.91%
h	11	0.07	1.000	0.070	40.13%
r	66	0.06	1.029	0.062	31.18%
				Y = 166.56, dY = 0	11, Cost=\$20.00

Variable Value Tolerance Sensitivity Deviation Factor contribution 73 0.09 -0.1850.017 N/A  $x_1$ 336 0.21 -0.0040.009 N/A  $x_2$ 70 0.07 0.976 0.066 N/A  $y_1$ 120 0.002 N/A 0.11 0.023  $y_2$ 42 0.07 h 1.017 1.017 N/A 60 0.06 1.208 1.208 N/A r Y=171.35, dY=0.12

Table 4 The best solution obtained by an ant colony optimizer

variables  $x_1$ ,  $x_2$ ,  $y_1$ ,  $y_2$ , h, and r. The smaller the deviation of Y is, the higher the quality of the product [2].

$$Y = y_1 + (r+h)\sec\theta + (r-x_1)\tan\theta$$
(12)

Where,

$$\tan \theta = (y_2 - y_1)/(x_2 - x_1) \tag{13}$$

$$\sec \ \theta = \sqrt{(1 + \tan^2 \theta)} \tag{14}$$

The goal of the problem is to minimize the deviation of the critical dimension (dY) by changing the design variables while keeping the target value of the critical dimension *Y* at or as close as possible to 166.67.

Based on the discussions in Section 2, the variance of the critical dimension can be expressed as the sum of n component variances as attenuated by their respective sensitivity coefficients:

$$\mathrm{d}Y^2 = \sum_{i=1}^{n} \left( S_i \times \mathrm{d}x_i \right)^2 \tag{15}$$

The individual sensitivities can be calculated by taking the partial derivative of the functional relationship between

the critical assembly dimension Y and the noncritical component dimensions  $x_1$ ,  $x_2$ ,  $y_1$ ,  $y_2$ , h, and r. Hence,

$$S_{1} = \frac{\partial Y}{\partial x_{1}} = \frac{(r+h)(y_{2}-y_{1})^{2}}{\sec \theta (x_{2}-x_{1})^{3}} - \frac{y_{2}-y_{1}}{x_{2}-x_{1}} + \frac{(r-x_{1})(y_{2}-y_{1})}{(x_{2}-x_{1})^{2}}$$
(16)

$$S_{2} = \frac{\partial Y}{\partial x_{2}} = -\frac{(r+h)(y_{2}-y_{1})^{2}}{\sec \theta (x_{2}-x_{1})^{3}} - \frac{(r-x_{1})(y_{2}-y_{1})}{(x_{2}-x_{1})^{2}}$$
(17)

$$S_3 = \frac{\partial Y}{\partial y_1} = 1 - \frac{(r+h)(y_2 - y_1)}{\sec \theta (x_2 - x_1)^2} - \frac{(r-x_1)}{(x_2 - x_1)}$$
(18)

$$S_4 = \frac{\partial Y}{\partial y_2} = \frac{(r+h)(y_2 - y_1)}{\sec \theta (x_2 - x_1)^2} - \frac{(r-x_1)}{(x_2 - x_1)}$$
(19)

Table 5 An optimum solution obtained by the PSO (experiment 3)

Variable	Value	Tolerance	Sensitivity	Deviation	Factor contribution
<i>x</i> <sub>1</sub>	58	0.09	-0.269	0.0242	3.49%
<i>x</i> <sub>2</sub>	379	0.48	-0.001	0.0006	0.002%
<i>y</i> <sub>1</sub>	89	0.07	0.995	0.0696	28.78%
V2	176	0.46	0.004	0.0021	0.028%
h	38	0.07	1.036	0.0725	31.20%
r	40	0.06	1.307	0.0784	36.48%
				<i>Y</i> =164.70, d <i>Y</i> =0	.13, Cost=\$17.80



Fig. 10 Total tolerance vs. cost and variation. a Total tolerance vs. variance. b Total tolerance vs. cost

$$S_5 = \frac{\partial Y}{\partial h} = \sec\theta \tag{20}$$

$$S_6 = \frac{\partial Y}{\partial r} = \sec\theta + \tan\theta \tag{21}$$

#### 5.2 Least cost-tolerance allocation

The final production cost of any mechanical assembly, including the one shown in Fig. 3, is significantly impacted by the specified tolerances on the dimensions of the manufactured parts. Tight tolerances can result in excessive process costs, while loose tolerances can lead to waste and assembly problems [1]. Figure 4 shows a typical manufacturing cost-tolerance curve.

Generally, there are two models to estimate how individual tolerances in a mechanical assembly stack up during the manufacturing process: the worst-case model (WCM) and the statistical RSS model [3]. The WCM computes the product's final tolerance ( $T_{tot}$ ) by adding the individual tolerances ( $T_i$ ) as  $\sum T_i$ . The main assumption here is that all the individual worst-case tolerance limits occur at the same time, and that can result in unnecessarily tight tolerances and a higher manufacturing cost.

In the RSS model, the individual tolerances are assumed to follow a normal distribution, and the final tolerance is calculated as:

$$T_{\rm tot} = \left(\sum T_i^2\right)^{0.5} \tag{22}$$

Clearly, the RSS approach allows for looser tolerances than the WCM method and, therefore, results in lower manufacturing costs.

A substantial amount of research has been carried out regarding optimal tolerance allocation using cost-tolerance functions, and various functions have been proposed to describe the cost-tolerance relationship [1, 2]. Assuming that for a component *i* the constant coefficient  $A_i$  (\$) represents the fixed costs, such as tooling, setup, prior operations, etc., and that the  $B_i$  (\$) term represents the cost of producing the



Fig. 11 RSS tolerance profiles of the population. a Best RSS tolerances. b Average RSS tolerance



Fig. 12 Performance measures of a set coverage (SC), b generational distance (GD), and c spread (SP) for one run of the optimizer over 100 generations

component dimension to a specified tolerance  $T_i$ , the reciprocal power cost-tolerance function for the each component of the assembly can be calculated as [3]:

$$C_i = A_i + B_i / T_i^{K_i} \tag{23}$$

For some integer power,  $K_i$  is to be selected from standard data sets depending upon the required manufacturing process.

# 5.3 Multiobjective formulation of the problem

Previously, this problem was attempted by means of an ant colony optimizer, which aggregated the product variation dY and the minimum manufacturing cost *C* into a single term using the weighted aggregation method as follows [19]:

$$F(X) = \sum_{i=1}^{m} w_i f_i(X) \tag{24}$$

Where  $f_i$  is the *i*th objective function, *m* is the number of objectives to be optimized,  $w_i \ge 0$  with  $\sum_{i=1}^{m} w_i = 1$  is the weight assigned by designers to objective  $f_i$  representing its degree of importance, and *X* is the *n*-dimensional vector of

design parameters. It has, however, been shown that the traditional weighted aggregated method suffers from many shortcomings, the most important of which is its assumption of convexity of the Pareto region, which is generally an invalid assumption in many real-world engineering design problems [5]. Furthermore, Pareto fronts are typically not uniform, that is, the obtained solutions are not uniformly distributed on the Pareto front given a set of evenly distributed weights [20].

The particle swarm optimization method utilized in this study was enhanced to perform a multimodal Pareto search for the combination of tolerances that minimize the total cost function for the assembly within the feasible region while also minimizing the assembly response variance dY and maximizing the total RSS tolerance. The multiobjective problem can formally be stated as shown below.

Optimize  $F = [f_1, f_2, f_3]^T$ , where the goal of optimization is to:

Minimize 
$$f_1 = dY = \sqrt{\left(\sum_{i=1}^n (S_i \times dx_i)2\right)}$$
 (25)

Table 6 Summary of performance measures for 50 statistically independent trials

Performance statistics	Set coverage (SC)	Generational distance (GD)	Spread (SP)
Best	0.00979	0.00100	0.14327
Worst	0.05816	0.00671	0.34456
Mean	0.04556	0.00175	0.24281
Standard deviation	0.01463	0.00084	0.00857

Table 7 Comparison of optimization methods

Technique	Optimal design $(x_1, x_2, y_1, y_2, r, h)$	Y	dY	Cost (\$)
Traditional calculus-based	$(70\pm0.2, 338\pm0.4, 65\pm0.3, 128\pm0.5, 35\pm0.1, 65\pm0.2)$	166.67	0.39	_
PSO <sup>1</sup>	$(22\pm0.2, 300\pm0.4, 67\pm0.3, 82\pm0.5, 42\pm0.1, 55\pm0.2)$	166.67	0.35	12.71
	(23±0.2, 204±0.4, 109±0.3, 110±0.5, 45±0.1, 57±0.2)	166.67	0.34	12.71
Ant colony optimizer	$(73\pm0.09, 336\pm0.21, 70\pm0.07, 120\pm0.11, 42\pm0.07, 60\pm0.06)$	171.35	0.12	_
PSO <sup>2</sup>	$(30\pm0.09, 243\pm0.21, 93\pm0.07, 123\pm0.11, 9\pm0.07, 58\pm0.06)$	165.24	0.12	20.00
	$(20\pm0.09, 206\pm0.21, 88\pm0.07, 93\pm0.11, 11\pm0.07, 66\pm0.06)$	166.57	0.11	20.00
PSO <sup>3</sup>	$(58\pm0.09, 379\pm0.48, 89\pm0.07, 176\pm0.46, 38\pm0.07, 40\pm0.06)$	164.70	0.13	17.80
PSO <sup>4</sup>	$(53\pm0.42, 287\pm0.46, 53\pm0.29, 82\pm0.43, 26\pm0.23, 83\pm0.21)$	166.70	0.42	11.00
	$(51\pm0.36, 315\pm0.49, 42\pm0.13, 88\pm0.40, 64\pm0.10, 58\pm0.10)$	166.61	0.21	13.83

$$\text{Minimize} f_2 = C_M(T) = \sum_{i=1}^n \left( A_i \times B_i / T_i^{K_i} \right)$$
(26)

Minimize 
$$f_3 = T_{RSS} = \left(\sum_{i=1}^{n} T_i^2\right)^{0.5}$$
 (27)

Subject to:

*Y* is to be on or as close as possible to the target value  $Y_{\text{Target}} = 166.67$ 

 $x_2 > x_1$  and  $y_2 > y_1$ 

 $0.001 \le T_i \le 0.5$  for i=1,...,n, where *n* is the number of toleranced dimensions

Clearly, the three objectives  $f_1$ ,  $f_2$ , and  $f_3$  are conflicting in nature, and their optimization will require a thorough examination of the trade-off solutions along the Pareto front.

#### 5.4 The experiments and results

In the experiments conducted, the enhanced PSO used 30 particles over 100 generations, maintained a repository of 500 particles (maximum size), and performed mutation at the rate of 0.5. The global best particle was selected from the top 10% sorted repository and replaced one of the nondominated solutions in the bottom 10% of the repository. To account for statistical fluctuations, the reported results were averaged over three statistically independent runs of the algorithm (the overall computation was very fast, and each experiment consisting of 100 generations typically took anywhere from 1 to 2 s).

Figure 5 depicts two Pareto fronts of the functional variance (dY) vs. the manufacturing cost  $(C_M)$  discovered by the optimizer after 100 generations with and without the crowding distance computations, respectively. In terms of solution diversity, it is clearly demonstrated in Fig. 5a that the utilization of the crowding distance computation allows a more evenly distributed exploration of the Pareto front.

A true indication of an optimizer's accuracy is the average behavior of it solutions as the swarm of particles explores the boundaries of the Pareto front across the feasible region. As shown in Fig. 6, as generations of solutions evolve, the average behavior (the statistical mean over three independent runs) of each of the population stays focused, gradually converging toward the trade-off optimum values both for product variance d*Y* and the cost function  $C_{\rm M}$ .

The profile of the best discovered (minimum-valued) solutions for the two objective functions are depicted in Fig. 7. A cursory examination of the range of Cost and dY functions shown in Fig. 5 illustrates the efficiency of convergence of the optimizer.

The optimization constraint required that the height of the center of the cylinder (Y) remains as close as possible to the target value of 166.67. The method of handling constraint violations via the fly-back mechanism to maintain feasibility was discussed in Section 4.3. Figure 8 shows the statistical means of three independent runs for the average and minimum product response Y of an entire swarm of particles.

The main goal of tolerance allocation is to re-distribute the "tolerance budget" within an assembly, systematically tightening tolerances on less expensive processes and loosening tolerances on costly processes, for a net reduction in cost [1]. This goal can very economically be accomplished by designers by simply moving the nominal values of the independent design parameters to less sensitive design regions. Figure 9 depicts the effects of the six design parameters ( $x_1$ ,  $x_2$ ,  $y_1$ ,  $y_2$ , h, r) on the overall product variation calculated in one run of the algorithm.

A careful examination of factor contributions shown in Fig. 9 reveals that design variables  $x_1$ ,  $y_1$ , and r have the largest leverage on the overall product variation. Therefore, by changing the dimensions of the independent variables and *not* the deviations on their manufacturing (tolerances), it will be possible to improve the design quality with little or no additional cost. This point is best demonstrated by the

two solutions obtained from the first conducted experiment where the algorithm was allowed to modify the design nominal values but not the tolerances (see Table 1).

In the first experiment, the tolerances for the six design parameters were fixed at 0.2, 0.4, 0.3, 0.5, 0.1, and 0.2, respectively, so that the performance of the PSO algorithm can be directly compared to that of a traditional calculusbased technique as shown in Table 2 [2]. Clearly, the PSO's solutions are more desirable as they better represent the optimal portions of the trade-off design region. Further, the solutions reported in Table 1 are but two of dozens of optimal trade-off solutions available in the repository.

To further judge the true mettle of the PSO, its performance was compared to that of an ant colony optimizer that used the method of weighted aggregation to collapse its objective functions into a single global one [19]. In this second experiment, the tolerances were fixed at 0.09, 0.21, 0.07, 0.11, 0.07, and 0.06, respectively, so that the two approaches can be compared on a more equal footing. Table 3 demonstrates two solutions obtained by the PSO algorithm ( $Y_{\text{Target}}=166.67$  in both experiments).

Table 4 shows the best nominal solutions obtained by the ant colony optimizer [19]. Note that factor contributions (percent) were not reported, and manufacturing cost comparisons could not be made simply due to the differences between the two approaches' method of cost computation (choice of *A* and *B* values in Eq. 26). Cost computations notwithstanding, the solutions obtained by the PSO have better optimized both the objective (minimization of d*Y*) and the constraint ( $Y=Y_{Target}$ ).

In another attempt to further reduce the incurred cost associated with design scenarios reported in Table 3 while maintaining on or near-target performance levels, a third experiment was conducted as follows. Design factor contributions to overall product variation (see Fig. 9) indicate that variables  $x_1$ ,  $y_1$ , and r have the most leverage on variance. It was, therefore, decided to start the optimization task with the same set of tolerances used in the second experiment (0.09, 0.21, 0.07, 0.11, 0.07, 0.06) but this time have the optimizer widen the tolerances for the remaining parameters ( $x_2$ ,  $y_2$ , and h) to which the assembly appeared less sensitive. Table 5 shows one of many optimal solutions that were obtained in this third experiment.

The trade-off solution shown in Table 5 results in an 11% reduction in the manufacturing cost compared to the solutions discovered in experiment 2. However, this reduction will cause an increase of 18% in the overall product variation d*Y*. The designers can carefully examine the trade-off solutions and determine if they are acceptable for their specific needs.

The optimization tasks performed in the three experiments reported so far only relied on the minimization of  $cost (C_M)$  and variance (dY) as the tolerances where kept at

fixed levels for comparative purposes. In the last experiment reported in this section, the total RSS tolerance ( $T_{RSS}$ ) was allowed to be maximized (see Section 5.3) simultaneously as product variation and manufacturing cost were each minimized. Figure 10 depicts the Pareto fronts obtained in this manner.

As expected and shown in Fig. 10a and b, allocating wider tolerances to the components of the stepped bar assembly significantly decreased the manufacturing cost while adding to the overall product variation. The advantage of this type of analysis is that the designers can readily re-distribute tolerances by tightening tolerances on less costly processes while widening tolerances on more expensive ones. Figure 11 depicts how the algorithm to allocated wider tolerances as each generation of particles was undergone the various evolutionary operators. It must be noted that the widest-possible  $T_{RSS}=1.24$ , where all individual tolerances reach their highest levels.

To further assess the performance of the swarm optimizer, the performance measures of generational distance (GD), set coverage (SC), and spread (SP) were calculated as described in Section 4.4. To calculate the SC and GD metrics, a set of approximately 16,000 (a full-factorial experiment with five design levels per independent variable) uniformly spaced true Pareto-optimal solutions ( $P^*$ ) was calculated a priori and used in all the conducted experiments. Figure 12 depicts the performance of one run of the algorithm over 100 generations.

As noted earlier, it is desired for the optimizer to identify solution regions in which all the performance measures are as close as possible to zero. Figure 12 demonstrates that the algorithm is able to converge very quickly and identify optimal solutions starting around the generation number 20.

To better examine the quality of the nondominated solutions obtained by the optimizer, a total of 50 random experiments were conducted, and statistical information about the collected performance measures for the final (100th) generation of each independent trial is depicted in Table 6.

It can be seen from Table 6 that the optimizer is able to converge to the true Pareto-optimal set with good distribution of nondominated solutions. The low values of GD and SC metrics indicate that the optimizer was able to explore and exploit solution regions (Q) very closed to the true Pareto-optimal front  $P^*$ .

Finally, Table 7 tabulates the results of four model experiments conducted in this study ( $PSO^1$ ,  $PSO^2$ ,  $PSO^3$ , and  $PSO^4$ ) and compares them to the solutions obtained through other optimization methods (see Tables 1, 2, 3, 4, and 5).

The solutions obtained in the last experiment (PSO<sup>4</sup>) were able to maintain the performance very close to  $Y_{\text{Target}}$  while minimizing the cost and allocating widest-possible

tolerances. Designers can in turn decide whether or not the operational costs due to the projected product deviations are worth considering. The important point, therefore, is that the multiobjective particle swarm optimizer *does* provide the designers with the ability to have access to myriad Pareto-optimal solutions.

# **6** Conclusions

Experience has shown that the benefits of reducing operational costs of a product due to its functional variations far outweigh the benefits of selecting more expensive and efficient components. Designers rely on tolerance allocation, which is an effective design tool to reduce the overall cost of production while ensuring the product performance stays on target.

Various types of evolutionary algorithms have recently been applied to the task of tolerance allocation in engineering design problems. Given cost tolerances for individual components of an assembly, the optimizer can systematically search the design space for specific combinations of tolerances that minimize the cost while satisfying the overall design objectives and constraints. In this paper, an enhanced particle swam optimizer was successfully utilized to perform tolerance allocation in multiple-objective design of a stepped bar mechanical assembly. The optimizer was designed to search the Pareto-optimal solution space in one of two possible modes. First, it was allowed to optimize design objectives and constraints by modifying both the tolerances and the nominal values of the independent design variables. And second, in order to avoid expensive manufacturing costs associated with tighter tolerances, the optimizer was only allowed to modify the nominal component dimensions with the assumption that the component tolerances were fixed.

Comparative studies revealed that the swarm optimizer outperformed a traditional calculus-based method and ant colony optimizer in identifying optimal solutions. The optimizer's use of an effective fly-back mechanism along the boundaries of the Pareto front and the computation of crowding distances ensured that the repository of the nondominated particles represented the widest-possible coverage of the Pareto front. The advantages of the proposed methodology in allocating widest-possible tolerances for reducing manufacturing costs while minimizing functional variations and maintaining on or near-target performance levels were demonstrated in several numerical experiments.

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