Distribution of 0 and 1 in the highest level of primitive sequences over $Z/(2^e)$ (II)

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Abstract The distribution of 0 and 1 is studied in the highest level a_{r-1} of primitive sequences over $Z/(2^r)$. It is proved that the proportion of 0 (or 1) in one period of a_{r-1} is between 40% and 60% for $e \ge 8$.

Keywords: linear recurring sequence, primitive sequence, highest levelZ sequence, distribution of 0 and 1.

LET $f(x) = x^n + c_{n-1}x^{n-1} + \dots + c_0$ be a monic polynomial over $Z/(2^e)$. The sequence $a = (a_0, a_1, \dots)$ satisfying the linear recursion

$$a_{i+n} = -(c_0 a_i + c_1 a_{i+1} + \dots + c_{n-1} a_{i+n-1}), \quad i = 0, 1, 2, \dots$$
 (1)

is called a linear recurring sequence over $Z/(2^{\epsilon})$ generated by f(x). $G(f(x))_{\epsilon}$ denotes the set of all sequences over $Z/(2^{\epsilon})$ generated by f(x) and $G'(f(x))_{\epsilon} = \{a \in G(f(x))_{\epsilon} | a \neq 0 \mod 2\}$. Recursion (1) is equivalent to $f(x)a = 0 = (0, 0, \cdots)$, where x is the shift operator, that is, $xa = (a_1, a_2, a_3, \cdots)$. Sequence a over $Z/(2^{\epsilon})$ has a unique binary decomposition $a = a_0 + a_1 2 + \cdots + a_{\epsilon-1} 2^{\epsilon-1}$, where $a_i = (a_{i0}, a_{i1}, \cdots)$ is a binary sequence with $a_{ij} = 0$ or 1, and a_i is called the *i*th level sequence of $a_{\epsilon-1}$ the highest level of a.

For a monic polynomial f(x) over $Z/(2^{\epsilon})$ with degree n, if f(0) (i.e. c_0) is invertible, then the period of $f(x) \operatorname{per}(f(x))_e \leq 2^{\epsilon^{-1}}(2^n - 1)$ by ref. [1]. If $\operatorname{per}(f(x)) = 2^{\epsilon^{-1}}(2^n - 1)$, f(x) is called a primitive polynomial over $Z/(2^{\epsilon})$ with degree n and sequences in $G'(f(x))_e$ are called primitive sequences generated by f(x). Ref. [2] provides a coefficient criterion for primitiveness of polynomial. Ref. [3] has shown the following entropy-preservation theorem with significance of cryptography: let f(x) be a primitive polynomial over $Z/(2^{\epsilon})$, $a, b \in G(f(x))_e$. If $a_{\epsilon^{-1}} = b_{\epsilon^{-1}}$, then a = b.

Let f(x) be a primitive polynomial over $Z/(2^{\epsilon})$ with degree n, $a \in G'(f(x))_{\epsilon}$. Then the period $per(a_k)$ of the kth level of a is 2^kT by ref. [4], where $T = 2^n - 1$. By ref. [1] or [4], $x^{2^{\ell-1}T} - 1 = 2^d h_d(x) \mod f(x)$ over $Z/(2^{\epsilon})$ for $1 \le d \le e^{-1}$, where $h_d(x)$ is a polynomial over $Z/(2^{\epsilon})$ with its degree less than n and $h_d(x) \ne 0 \mod 2$. Ref. [5] provided the following result.

Proposition A. Let f(x) be a primitive polynomial over $Z/(2^{\epsilon})$ with degree $n, T = 2^{n} - 1, a \in G'(f(x))_{\epsilon}, d = \lfloor e/2 \rfloor, s = h(x) a \mod 2^{d}$. Then the proportion λ of 0 (or 1) in $a_{\epsilon-1}$ satisfies

$$\frac{1}{2}-\frac{N(s, 0)}{2^d T}\leqslant\lambda\leqslant\frac{1}{2}+\frac{N(s, 0)}{2^d T},$$

where N(s, 0) denotes the number of 0 in one period of $s, h(x) = \begin{cases} h_d(x), & \text{if } e = 2d, \\ h_{d+1}(x), & \text{otherwise } e = 2d+1. \end{cases}$

By applying the result, ref. [5] has proved that when e is sufficiently large, most of a_{e-1} are of good distribution of 0 and 1. For example, let e = 32 and f(x) be primitive over $Z/(2^e)$. Then the pro-

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portion of sequences in G'(f(x)), with 0.498 046 875 $\leq \lambda \leq 0.501$ 953 125 is 99.6% at least.

In the above results, we cannot estimate the distribution of a little of $a_{e^{-1}}$ because it is difficult to compute the number of 0 in one period of $s = h(x)a \mod 2^d$. In this note, we give an upper bound of the number of 0 in one period of primitive sequences over $Z/(2^4)$ and the proportion of 0 in one period of $a_{e^{-1}}$ is between 40% and 60% for $e \ge 8$.

First we give the generalization of Proposition A.

Theorem 1. Let f(x) be a primitive polynomial over $Z/(2^{\epsilon})$ with degree $n, T = 2^{n} - 1, a \in G(f(x))_{\epsilon}, 1 \leq d \leq e/2, s = h_{\epsilon-d}(x)a \mod 2^{d}$. Then the proportion λ of 0 (or 1) in $a_{\epsilon-1}$ satisfies

$$\frac{1}{2}-\frac{N(s, 0)}{2^d T}\leqslant\lambda\leqslant\frac{1}{2}+\frac{N(s, 0)}{2^d T},$$

where N(s, 0) denotes the number of 0 in one period of s.

Lemma 1. Let f(x) be a primitive polynomial over F_2 with degree n. $a = (a_0, a_1, \cdots)$ and $b = (b_0, b_1, \cdots)$ are two distinct primitive sequences generated by f(x). $S_{a, b}(u, v) = \{i \mid a_i = u, b_i = v, 0 \le i \le 2^n - 2\}$, $M_{a, b}(u, v) = |S_{a, b}(u, v)|$. Then $M_{a, b}(0, 0) = 2^{n-2} - 1$, $M_{a, b}(0, 1) = M_{a, b}(1, 0) = M_{a, b}(1, 1) = 2^{n-2}$.

Lemma 2. Let f(x) be a primitive polynomial over F_2 with degree $n, n \ge 3$. a, b and c are pairwise relatively distinct sequences generated by f(x) and $a + b + c \ne 0$. For $u, v, w \in F_2$, set $S_{a, b, c}(u, v, w) = |i|a_i = u, b_i = v, c_i = w, 0 \le i \le 2^n - 2|, M_{a, b, c}(u, v, w) = |S_{a, b, c}(u, v, w)|$. Then

$$M_{a, b, c}(u, v, w) = \begin{cases} 2^{n-3} - 1, & \text{if } u = v = w = 0, \\ 2^{n-3}, & \text{otherwise.} \end{cases}$$

By applying Lemma 1, Lemma 2 and the relation of leval sequences, we get the following result.

Lemma 3. Let f(x) be a primitive polynomial over $Z/(2^4)$ with degree $n, n \ge 3$. a is a primitive sequence generated by f(x) over $Z/(2^4)$. If $h_2(x) \ne 1 \mod 2$, then the number of 0 in one period of a satisfies

$$N(a, 0) \leq 15 \cdot 2^{n-3} - 8$$

By Theorem 1 and Lemma 3, we have the following theorem.

Theorem 2. Let f(x) be a primitive polynomial over $Z/(2^{\epsilon})$ with degree $n, n \ge 3$, $e \ge 8$. a is a primitive sequence generated by f(x). $h_2(x)$ satisfies the conditions of Lemma 3. Then the proportion of 0 (or 1) in a_{r-1} satisfies

$$\frac{1}{2} - \frac{15 \cdot 2^{n-6} - 1}{2(2^n - 1)} \leq \lambda \leq \frac{1}{2} + \frac{15 \cdot 2^{n-6} - 1}{2(2^n - 1)}.$$

Corollary 1. The conditions are as those in Theorem 2. Then the proportion of 0 in a_{e-1} satisfies 38.281 25% $< \lambda < 61.718$ 75%.

To estimate λ more accurately, we shall improve Lemma 3. First we give the following Lemma 4 and Lemma 5.

Lemma 4. Let f(x) be a primitive polynomial over F_2 with degree $n, n \ge 4$, a, b, c and d are primitive sequences generated by f(x). The symbol $M_{a, b, c}(u, v, w, x)$ is defined as in Lemma 2. If a, b, c, d are pairwise relatively distinct sequences and any of a + b + c, a + b + d, a + c + d, b + c + d and a + b + c + d is not 0, then

$$M_{a, b, c, d}(u, v, w, x) = \begin{cases} 2^{n-4} - 1, & \text{if } u = v = w = 0, \\ 2^{n-4}, & \text{otherwise.} \end{cases}$$

Lemma 5^[6]. Let *a* be a primitive sequence of degree *n* over $Z/(2^{e})$. Then the number of 0 in one period of *a* satisfies $N(a, 0) \leq 2^{n} + 2^{n/2} - 2$.

Applying Lemma 4, Lemma 5 and the relation among level sequences, we get the following lemma. Lemma 6. Under the condition of Lemma 3, let $n \ge 4$, $h_1(x)h_2(x) \ne 1 \mod(2, f(x))$, $(1+h_1(x))h_2(x) \ne 1 \mod(2, f(x))$. Then the number of 0 in one period of a satisfies

$$N(a, 0) \leq \min\{13 \cdot 2^{n-3} - 8, 12 \cdot 2^{n-3} + 2^{n/2} - 2\}.$$

By Theorem 1 and Lemma 6, we have the following theorem.

Theorem 3. Let f(x) be a primitive polynomial of degree *n* over $Z/(2^e)$, $n \ge 4$, $e \ge 8$. *a* is a primitive sequence over $Z/(2^e)$ generated by f(x). If $h_1(x)$ and $h_2(x)$ satisfy the conditions of Lemma 6, then the proportion of 0 (or 1) in $a_{e^{-1}}$ satisfies

$$\frac{1}{2} - \frac{6 \cdot 2^{n-4} + 2^{\frac{n}{2}-1} - 1}{2^2(2^n - 1)} \leq \lambda \leq \frac{1}{2} + \frac{6 \cdot 2^{n-4} + 2^{\frac{n}{2}-1} - 1}{2^2(2^n - 1)},$$
$$\frac{1}{2} - \frac{13 \cdot 2^{n-6} - 1}{2(2^n - 1)} \leq \lambda \leq \frac{1}{2} + \frac{13 \cdot 2^{n-6} - 1}{2(2^n - 1)}.$$

Corollary 2. The condition is as that in Lemma 3. Then the proportion λ of 0 (or 1) in a_{r-1} satisfies

() 40.157 9% $< \lambda < 59.842$ 1% if $n \leq 7$;

(||) 40.250 5% $< \lambda < 59.749$ 5% if $n \ge 8$;

(|||) 40.624 9% < λ < 59.375 1% if $n \ge 30$.

Reference [5] proved that most of a_{e-1} have good ditributions. If e is sufficiently large, most of a_{e-1} have good ditributions. Corollary 2 insures other a_{e-1} against bad distributions.

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