# **THE SPECTRUM PROBLEM III: UNIVERSAL THEORIES\***

BY

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## *Dedicated to Professor Abraham Robinson*

#### ABSTRACT

We solve the classification problem and essentially the spectrum problem for universal theories (see [6] for discussion of the meaning of this). We first solve it for  $T$  such that if  $M_1$ ,  $M_2$  elementarily extend  $M_0$  and are independent over it, then over  $M_0 \cup M_1$  there is a prime model. This generalizes [2]. This was subsequently used and generalized for countable first order theories. (This will appear in [5].) But note that there the theory is countable and in the case of structure the model is prime over a non-forking tree of models; here the model is generated by the union (and the  $T$  not necessarily countable). The universality is used in

THEOREM. If *T* is stable, and complete then either  $(A)$  for every  $M_1 \leq M$  $(l = 0, 1, 2)$  models of T, if  $M_0 \subseteq M_1M_2$ ,  $\{M_1, M_2\}$  is independent over  $M_0$  (i.e.  $tp(M_1, M_2)$  is finitely satisfiable in  $M_0$ ), then the submodel of M which  $M_1 \cup M_2$ *generates is an elementary submodel of M, or* (B) *there is an unstable theory extending the universal part of T (we can replace universal by*  $\Sigma_2$  *and slightly more).* 

CONCLUSION. For any universal T: *Either* (a) for every model M of T there is a tree I with  $\leq \omega$  levels and submodels  $N_n$  ( $\eta \in I$ ) of power  $\leq 2^{|T|}$  (by [5], just  $\leq |T|$ ) such that (i) M is generated by  $\bigcup_{\eta \in I} N_{\eta}$ , (ii)  $\eta \leq v \Rightarrow N_{\eta} \subseteq N_{\nu}$ , (iii) if v is an immediate successor of  $\eta$  then tp(N<sub>u</sub>,  $\bigcup \{N_\rho : \rho \in I, \nu \not\leq \rho\}$ ) is finitely satisfiable in  $N_{\eta}$  (note that asking this just for quantifier-free formulas is enough). *Or* (b) for every cardinal  $\lambda > |T|$ , there are 2<sup> $\lambda$ </sup> non-isomorphic models for power A.

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In Sections 0-3, Th ( $\mathfrak{C}$ ) is assumed to be stable.  $\prec$  is elementary submodel.

### **§0. Canonization**

**This section can be avoided, if you avoid 1.5A (but not 1.5), 1.8 and 2.3(1) which are not used later.** 

**Here we quote some facts from the new edition of the author's book [5] using**  the definition from there (see Definition V 4.5). So  $\mathfrak C$  is a quite saturated model **of a complete first-order stable T.** 

**0.1.** CLAIM.  $({\mathbb{C}}^{eq})$  Suppose r is regular and  $\text{stp}\left(\bar{a}, A\right)$  is not orthogonal to r. *Then:* 

**(1)** *There is*  $E \in FE^m$  (acl A)  $(m = l(\bar{a}))$  *such that stp*  $(\bar{c}/E, \bar{b})$  *is cl<sup>2</sup>(r)-simple but not orthogonal to r.* 

(2) *Moreover* stp  $(\bar{c}/E, A)$  *contains a formula*  $\vartheta(y, \bar{b})$  *that is cl<sup>2</sup>(r)-simple.* 

(3) *Also, for every*  $\bar{b}'$  *realizing tp*( $\bar{b}$ , $\varnothing$ ),  $\vartheta$ ( $y, \bar{b}'$ ) *is cl<sup>3</sup>(r)-simple.* 

0.1A. FACT. There is  $\varphi(x, \bar{c})$  which is cl<sup>3</sup>(r)-regular but has completions not orthogonal to r.

Let  $\varphi(x, \bar{c})$  be any formula which has completions not orthogonal to r, but has no extension of smaller rank  $(R<sup>1</sup>(-, L, \infty))$  with this property.

Clearly such  $\varphi(x,\bar{c})$  exists. Now any such  $\varphi(x,\bar{c})$  is cl<sup>2</sup>(r)-regular -- this follows by

 $\oplus$  for every complete stationary p,  $\varphi(x,\bar{c}) \in p$ : if  $R(p,L,\infty)$  $R(\varphi(x, \bar{c}), L, \infty)$  then p is orthogonal to r; hence if p forks over  $\bar{c}$  then p is orthogonal to r.

0.1B. FACT. Every complete stationary  $\mathcal P$ -regular p not orthogonal to  $r \in \mathcal P$ is regular.

0.1C. FACT. There is  $\varphi(x, \bar{c})$  as in Fact 0.1A, such that for every  $\bar{c}'$  realizing tp  $(\bar{c}, \emptyset)$ ,  $\varphi(x, \bar{c}')$  is cl<sup>2</sup>(r)-regular.

0.2. CLAIM.  $(\mathbb{C}^{eq})$  *If tp*  $(\bar{a}, A)$  *is not orthogonal to some trivial regular type r* then *for some*  $b \in act(A \cup \overline{a})-act(A)$ *,* tp(*b, A) is cl<sup>3</sup>({r})-simple and*  $w_r(b, A) = 1$ . If tp ( $\bar{a}, A$ ) is regular we can replace simple by regular, and if T is *superstable then some*  $\varphi \in$  tp  $(b, A)$  *is cl<sup>3</sup>({r})-simple [regular].* 

# **§I. On Strong Elementary Submodels**

HYPOTHESIS.  $T$  is superstable.

1.1. DEFINITION. (1) We say that  $M\subset sN$  if  $M\subset N$  and for every  $\bar{a}\in M$ ,  $\bar{b} \in N$  there is  $\bar{b}' \in M$  realizing tp ( $\bar{b}, \bar{a}$ ). We define  $M \subseteq A, B \subseteq A$  similarly.

(2) We say that  $M \subseteq_{\alpha} N$  if  $M \subseteq N$ , and for every  $\bar{a} \in M$ ,  $\bar{b} \in N$  there is  $\bar{b}' \in M$ realizing stp  $(\bar{b}, \bar{a})$ . We define  $M \subseteq_a A$ ,  $B \subseteq_a A$  similarly.

1.2. CLAIM. (1) If  $A \subseteq B \subseteq C$ ,  $A \subseteq S$  *C* then  $A \subseteq S$ *B*.

(2) Let  $A \subseteq B_i$  for  $l < n$ ,  $\bar{a} \in A$ ,  $\bar{b}_l \in B_l$ ,  $\neq \varphi[\bar{b}_0, \ldots, \bar{b}_{n-1}, \bar{a}]$  and  $\{B_l : l < n\}$  is *independent over A. Then there are*  $\bar{b}'_0, \ldots, \bar{b}'_{n-1} \in A$ , such that:  $\forall \phi[\bar{b}'_0,\ldots,\bar{b}'_{n-1},\bar{a}]$  and  $\bar{b}'_i$  realizes tp( $\bar{b}_i,\bar{a}$ ) (and in fact it realizes tp  $(\bar{b}_i, \bar{a} \cup \bigcup_{m \leq i} \bar{b}'_m)$  and  $\{\bar{b}'_0, \ldots, \bar{b}'_{n-1}\}$  is independent over some  $\bar{a}'$ ,  $\bar{a} \subseteq \bar{a}' \subseteq A$ ).

(3) If  $N \subseteq N_1 \subseteq M$ ,  $N \subseteq A$ , M is  $\mathbf{F}_{n_0}^t$ -atomic over  $N_1 \cup A$ ,  $\{A, N_1\}$  is indepen*dent over N. Then*  $N_1 \subseteq_s M$  (in fact  $N_1 \subseteq_s M$  holds, M a set suffices).

(4) If  $A \subseteq B \subseteq C_{\tau}$ - $A \subseteq a$  C then  $A \subseteq a$  B.

(5) If  $A \subseteq_{a} B_{i}$   $(i < \alpha)$  and  $\{B_{i} : i < \alpha\}$  is independent over A, then  $A \subseteq a \bigcup_{i \leq a} B_i$ .

(6) If  $A \subseteq B$  then  $A \subseteq B$ .

PROOF. (1), (4), (5), (6) are easy.

(2) We can find a finite  $\bar{a}'$ ,  $\bar{a} \subseteq \bar{a}' \subseteq A$ , such that tp  $(\bar{b}_0 \land \cdots \land \bar{b}_{n-1}, A)$  does not fork over  $\bar{a}'$ . Now we define by induction on  $i < \omega$ ,  $\langle \bar{b}_i : l < n \rangle$ , such that  $\bar{b}_i \in A$ . For a given i we define  $\bar{b}_i^i$  by induction on  $l : \bar{b}_i^i$  realizes  $\text{tp }({\vec{b}}_i,{\vec{a}}'\cup\bigcup_{j. We can easily prove by induc$ tion on  $k < \omega$  that  ${\{\overline{b}_i : ni + l < k\}}$  is independent over  $\overline{a}'$ . As in the proof of [1; Ch. II, 2.17, p. 38, Ch. III, 2.13, pp. 98] we finish.

(3) Let  $\bar{a} \in N_1$ ,  $\bar{b} \in M$  and we should find  $\bar{b'} \in N_1$  realizing stp  $(\bar{b}, \bar{a})$ . As M is  $\mathbf{F}_{\kappa_0}^t$ -atomic over  $N_1 \cup A$ , there are  $\bar{b}_1 \in N_1$ ,  $\bar{b}_2 \in A$  such that  $\models \psi[\bar{b}, \bar{b}_1, \bar{b}_2]$  and  $\psi(\bar{x}, \bar{b_1} \bar{b_2})$  + tp ( $\bar{b}$ , N<sub>1</sub>  $\cup$  A); w.l.o.g.  $\bar{a} \subseteq b_1$ . For some  $\bar{b_0} \in N$ , tp ( $\bar{b_0}$ , N) does not fork over  $\bar{b}_0$  for  $l = 1,2$  and remember  $\{\bar{b}_1, \bar{b}_2\}$  is independent over N.

Now choose  $\bar{b}_2' \in N$  which realizes stp( $\bar{b}_2, \bar{b}_0$ ) (possible as  $N \subseteq_a A$ ). As  $\text{tp}(\bar{b}_1, A)$  does not fork over  $\bar{b}_0$ ,  $\bar{b}_2'$  realizes stp $(\bar{b}_2, \bar{b}_0 \cup \bar{b}_1)$ . So  $\bar{b}_0 \wedge \bar{b}_1 \wedge \bar{b}_2$ ,  $\bar{b}_0 \hat{b}_1 \hat{b}_2'$  realizes the same type, hence  $\hat{z} = (\exists \bar{x})\psi(\bar{x}, \bar{b}_1, \bar{b}_2')$ , and letting  $\bar{b}'$  realize  $\psi(\bar{x}, \bar{b}_1, \bar{b}_2)$ ,  $\bar{b}_0 \hat{b}_1 \hat{b}_2 \hat{b}_2 \hat{b}_3 \hat{b}_0 \hat{b}_1 \hat{b}_2 \hat{b}_1 \hat{b}_2 \hat{b}_1$  realizes the same type. But we can choose  $\bar{b}' \in N_1$ . So there is  $\bar{b}' \in N_1$  realizing tp( $\bar{b}, \bar{a}$ ) (as  $\bar{a} \subseteq \bar{b}_1$ ). In fact  $\bar{b}'$ realizes stp ( $\bar{b}$ ,  $\bar{a}$ ). For every  $E \in FE(\bar{b_1})$ , E is a formula over  $N_1$  (as it is almost over  $\bar{b}_1$ ) hence  $tp(\bar{b} \wedge \bar{b}_2, N_1) \vdash E(\bar{x}, \bar{y}; \bar{b}, \bar{b}_2)$ . But for every  $\Theta(\bar{x}, \bar{y}) \in$ tp  $(\bar{b} \cdot \bar{b}_2, N_1)$ , tp  $(\bar{b}, N_1 \cup \bar{b}_2)$   $\vdash \Theta(\bar{x}, \bar{b}_2)$ , hence  $\psi(\bar{x}, \bar{b}_1, \bar{b}_2)$   $\vdash \Theta(\bar{x}, \bar{b}_2)$ . We can conclude that  $\psi(\bar{x}, \bar{b}_1, \bar{b}_2') \vdash E(\bar{x}, \bar{b}_2'; \bar{b}, \bar{b}_2)$  [as  $\Theta(\bar{x}, \bar{y}) \stackrel{\text{def}}{=} E(\bar{x}, \bar{y}; \bar{b}, \bar{b}_2)$  is almost over  $\bar{b}_1$ and  $\bar{b}_2'$  realizes stp  $(\bar{b}_2, \bar{b}_1 \wedge \bar{b}_0)$ , clearly  $\psi(\bar{x}, \bar{b}_1, \bar{b}_2') \vdash \Theta(\bar{x}, \bar{b}_2')$  hence  $\psi(\bar{x}, \bar{b}_1, \bar{b}_2') \vdash$  $E(\bar{x}, \bar{b}_2; \bar{b}, \bar{b}_2)$ ], hence  $\bar{E}(b', \bar{b}_2; \bar{b}, \bar{b}_2)$ . As this holds for every  $E \in FE(\bar{b}_1)$ clearly  $\bar{b}' \hat{b}'$  realizes stp $(\bar{b} \hat{b} \hat{b}$ ,  $\bar{b}_2, \bar{b}_1)$ , hence  $\bar{b}'$  realizes stp $(\bar{b}, \bar{a})$ .

1.3. LEMMA. *Suppose*  $N \subseteq A$ ,  $N \subseteq aM$ ,  $\{M, A\}$  is independent over N:

(1) If  $tp(\bar{a}, A)$  is orthogonal to N or is  $\mathbf{F}_{\kappa_0}^c$ -isolated (for  $\mathbf{F}_{\kappa_0}^c$ -isolation see *Definition IV 2.7), then*  $tp(\bar{a}, A)$  +  $tp(\bar{a}, A \cup M)$ .

(2) If for each  $i < \alpha$ , tp ( $\bar{a}_i$ , A  $\cup \bigcup_{i < i} \bar{a}_i$ ) is orthogonal to N or  $\mathbf{F}_{\kappa_0}^c$ -isolated (for *F~o-isolation see Definition IV,* 2.7) *then* 

$$
\mathrm{tp}_*(\{\bar{a}_i : i < \alpha\}, A) \vdash \mathrm{tp}_*(\{\bar{a}_i : \leq \alpha\}, M \cup A).
$$

(3) In (1), (2) we can replace M by any B (if  $N \subseteq a$ B).

PROOF. (1) Trivial (for the second case note  $dcl(M\cup A)\cap acl(A)=$  $dcl(A)$ ).

(3) The same proof.

1.4. CLAIM. (1) (in  $\mathbb{S}^{eq}$ ) *Let*  $x \in \{t, a\}$ *. If p is an m-type over A, r is a regular type,* stp  $(\bar{a}, A)$  *is not orthogonal to r, r extends p, p is an*  $\mathbf{F}_{\mu}^{x}$ -type,  $A = |N|$ , N is  $\mathbf{F}_{\mu}^{\mathbf{x}}$ -compact, then  $\text{stp}_{\mu}(B,A)$  is not orthogonal to r where  $B =$ Cb (tp  $(\bar{a}, A \cup p(\mathbb{C}^{eq}))$ ).

(2) *Suppose*  $\bar{b} \in M$ ,  $\bar{a} \notin M$ ,  $\text{tp } (\bar{a}, M)$  is not orthogonal to some type q to which  $\psi(\bar{x}, \bar{b})$  belongs. Then for every model N including  $M \cup \bar{a}$ ,  $\psi(N, \bar{b}) \neq \psi(M, \bar{b})$ .

REMARK. The most interesting cases of (1) are  $\mathbf{F}_{\mu}^{\mathbf{x}} = \mathbf{F}_{\mathbf{x}_0}^{\mathbf{x}}$ , p finite (so N is just a model) and  $\kappa = \kappa_r(T)$ ,  $\text{Dom } p \mid \leq \kappa$  (even for any stable T).

PROOF. (1) We shall prove later:

1.4A. FACT. For every  $\overline{d}$ 

tp  $(\bar{a}, \bar{d} \cup A \cup p(\mathbb{C}^{eq}))$  does not fork over  $A \cup B$ whenever tp  $(\bar{a}, A \cup \bar{d})$  does not fork over A.

Let  $\lambda > |T|+|A|+||M||$  be regular, let  $N_0$  be an  $\mathbf{F}_\lambda^2$ -saturated model,  $A \subseteq N_0$ ; and w.l.o.g. r is a type over  $N_0$  and tp<sub>\*</sub>( $N_0$ ,  $A \cup \overline{a}$ ) does not fork over A.

By 1.4A, tp  $(\bar{a}, N_0 \cup p(\mathbb{S}^{eq}))$  does not fork over  $N \cup B$  hence over  $N_0 \cup B$ . Let  $p \subseteq r_0 \in S^m(N_0)$ ,  $r_0$  regular not orthogonal to r ( $r_0$  exists, of course). Let  $N_1$  be  $\mathbf{F}_{\lambda}^a$ -prime over  $N_0 \cup \bar{a}$ , hence it is  $\mathbf{F}_{\lambda}^a$ -atomic over  $N_0 \cup \bar{a}$  hence over  $N_0 \cup \bar{a} \cup B$ (note that  $B \subseteq N_1$ , as  $B \subseteq \text{acl}(N \cup \bar{a})$ ). As tp<sub>\*</sub> ( $\bar{a}$ ,  $N_0 \cup p(\mathbb{S}^{eq})$ ) does not fork over  $N_0 \cup B$  clearly tp<sub>\*</sub> ( $\bar{a}$ ,  $N_0 \cup p(N_1)$ ) does not fork over  $N_0 \cup B$  hence (by III, 0.1)  $tp_*(p(N_1), N_0 \cup B \cup \bar{a})$  does not fork over  $N_0 \cup B$ . By IV, 4.3 it is easy to see that  $p(N_1)$  is a  $\mathbf{F}_\lambda^a$ -atomic over  $N_0 \cup B$  ( $B \subseteq N_1$  as noted above). So if 1.4(1)'s conclusion fails, as for every  $\bar{b} \in B$ , as tp  $(\bar{b}, N_0)$  is a stationarization of stp  $(\bar{b}, A)$ , clearly it is orthogonal to r.

Note tp  $(\bar{b}, N_0)$  is orthogonal to  $r_0$ .

So stp<sub>\*</sub>(B, N<sub>0</sub>) is orthogonal to  $r_0$  and  $p(N_1)$  is  $\mathbf{F}_\lambda^a$ -atomic over  $N_0 \cup B$ .

Hence if  $N_2$  is  $\mathbf{F}_4^a$ -prime over  $N_0 \cup B$ ,  $r_0$  is not realized in  $N_2$  and for every  $\bar{c} \in p(N_1)$ , tp ( $\bar{c}$ ,  $N_0 \cup B$ ) is  $\mathbf{F}_n^a$ -isolated hence is realized in  $N_2$ . We can conclude that (as p is over  $N_0$ )  $N_1$  does not realize  $r_0$ . But as tp ( $\bar{a}$ ,  $N_0$ ) is not orthogonal to  $r_0$ ,  $r_0$  is realized in  $N_1$ , and  $p \subseteq r_0$ , contradiction.

PROOF OF FACT 1.4A. Why does it hold? As we have assumed tp  $(\bar{a}, A \cup \bar{d})$ does not fork over A, also tp( $\overline{d}$ , A  $\cup$   $\overline{a}$ ) does not fork over A, hence [as

 $B \subseteq \text{acl}(A \cup \overline{a})$ ] tp( $\overline{d}, A \cup \overline{a} \cup B$ ) does not fork over A, hence [as A is a model] is finitely satisfiable in A. Suppose the conclusion of 1.4A fails, i.e., for some  $\bar{e} \in A \cup \bar{d} \cup p(\mathbb{C}^{eq})$ , and  $\varphi$ ,  $\models \varphi[\bar{a}, \bar{e}]$ , but  $\varphi(x, \bar{e})$  forks over  $A \cup B$ . W.l.o.g. for some finite  $\Delta$ ,  $R(\varphi(\bar{x}, \bar{e}), \Delta, \aleph_0) \leq R(\text{tp}(\bar{a}, A), \Delta, \aleph_0)$ .

We know that for some large enough finite  $\kappa$ ,  $R[\varphi(\bar{x}, \bar{e}), \Delta, \kappa] =$  $R[\varphi(\bar{x},\bar{e}),\Delta,\mathbf{N}_0], \quad R[\text{tp}(\bar{a},A\cup B),\Delta,\kappa] = R[\text{tp}(\bar{a},A\cup B),\Delta,\mathbf{N}_0]. \quad \text{W.l.o.g.}$ for every  $\bar{e}'$ ,  $R[\varphi(\bar{x}, \bar{e}'), \Delta, \kappa] < R[\varphi(\text{tp}(\bar{a}, A \cup B), \Delta, \kappa]$ . W.l.o.g.  $\bar{e}$  =  $\bar{d}^{\wedge} \bar{a}'^{\wedge} \bar{e}_1^{\wedge} \cdots \bar{e}_n^{\wedge}$  where  $\bar{a}' \subset A$ ,  $\bar{e}_i \in p(\mathbb{C}^{eq})$ ; so  $\models \varphi(\bar{a}, \bar{d}, \bar{a}', \bar{e}_1, \ldots, \bar{e}_m^{\wedge})$  and p w.l.o.g. is a singleton or closed under conjunction. So  $\bar{d}$  realizes the type

$$
q = \left\{ (\exists \bar{z}_1, \ldots, \bar{z}_m) \bigg[ \varphi(\bar{a}, \bar{x}, \bar{a}', \bar{a}_1, \ldots, \bar{z}_m) \wedge \bigwedge_{l=1}^m \Theta(\bar{z}_l) \bigg] : \Theta(\bar{z}) \in p \right\}
$$

As q does not fork over A, q an  $\mathbf{F}_{\mu}^{\prime}$ -type, some  $\bar{d}' \in A$  realizes q. So for some  $\bar{e}'_1 \in p(\mathbb{C}^{eq}) \models \varphi[\bar{a}, \bar{d}', \bar{a}', \bar{e}'_1, \ldots]$ . But this is a contradiction as

$$
R[\operatorname{tp}(\bar{a}, A \cup p(\mathbb{C}^{\text{eq}}), \Delta, \mathbf{N}_0] = R[\operatorname{tp}(\bar{a}, A \cup B), \Delta, \mathbf{N}_0]
$$

 $= R[tp(\bar{a}, A \cup B), \Delta, \kappa] > R[\varphi(x, \bar{d}', \bar{a}', \bar{e}', \ldots), \Delta, \kappa].$ 

PROOF OF 1.4(2). By (1) with  $(p = {\varphi(\bar{x}, \bar{b})}, A = M)$ . We know that Cb(tp( $\bar{a}$ ,  $M \cup \varphi(\mathfrak{C}, \bar{b})$ ) is not contained in M, but it is contained in acl $(M \cup \varphi(N,\bar{b}))$ . Hence  $\varphi(N,\bar{b}) \not\subseteq M$ .

1.5. LEMMA. *Assume N*  $\subseteq_a M'$ , N  $\subseteq M' \subseteq M$  and  $m < \omega$ . Suppose  $\bar{c} \in M$ ,  $\bar{c} \notin M'$ ,  $l(\bar{c})=m$  and  $R[tp(\bar{c},M'),L,\infty]$  is minimal (under this condition). If  $tp~(\bar{c}, M')$  is not orthogonal to N then there is  $\bar{c}' \in M$ ,  $l(\bar{c}') = l(\bar{c})$ ,  $\bar{c}' \notin M'$ ,  $\text{tp } (\bar{c}', M')$  *does not fork over M and R*  $[\text{tp } (\bar{c}', M'), L, \infty] = R [\text{tp } (\bar{c}, M), L, \infty]$ .

PROOF. Let  $\psi(\bar{x}, \bar{b}) \in \text{tp}(\bar{c}, M'), \quad R^m[\psi(\bar{x}, \bar{b}), L, \infty] = R^m[\text{tp}(\bar{c}, M'), L, \infty]$ hence tp ( $\bar{c}$ , M') does not fork over  $\bar{b}$ . We work in  $\mathbb{C}^{eq}$ . We can choose  $\bar{a} \in N$ , such that tp  $(\bar{b} \wedge \bar{c}, N)$  does not fork over  $\bar{a}$  and stp  $(\bar{c}, \bar{b})$  is not orthogonal to  $\bar{a}$ (as tp  $({\bar c}, M')$ , stp  $({\bar c}, {\bar b})$  are parallel, stp  $({\bar c}, {\bar b})$  is not orthogonal to some  $q \in$  $\bigcup_{n} S^{n}(N)$  and w.l.o.g. q does not fork over  $\bar{a}$ ). Now as  $N \subseteq a M'$  there is a sequence  $\bar{b}' \in N$  realizing stp  $(\bar{b}, \bar{a})$  and for some  $\bar{c}', \bar{b}' \hat{c}'$  realizes stp  $(\bar{b} \hat{c}, \bar{a})$ . By V, 3.5 stp $(\bar{c}, \bar{b} \cup \bar{a})$ , stp $(\bar{c}', \bar{b}' \cup \bar{a})$  are not orthogonal (note that stp( $\bar{c}, \bar{b} \cup \bar{a}$ ), stp( $\bar{c}, \bar{b}$ ) and stp( $\bar{c}, M'$ ) are parallel). So tp( $\bar{c}, M'$ ) is not orthogonal to some type to which  $\psi(\bar{x}, \bar{b}')$  belongs, hence by 1.4(2) there is  $\bar{c}'' \in M$ ,  $\bar{c}'' \not\in M'$ ,  $\models \psi[\bar{c}'', \bar{b}']$ . So

 $R [\text{tp } (\bar{c}^{\prime\prime}, M^{\prime}), L, \infty] \leq R [\psi(\bar{x}, \bar{b}^{\prime}), L, \infty] = R [\psi(\bar{x}, \bar{b}), L, \infty] = R [\text{tp } (\bar{c}, M^{\prime}), L, \infty]$ .

By the hypothesis that  $R$ [tp( $\bar{c}$ , M'),  $L$ ,  $\infty$ ] is minimal equality holds, hence  $\text{tp}(\bar{c}'', M')$  does not fork over N. So  $\bar{c}''$  is as required.

1.5A. LEMMA. Suppose  $N \subset M$  and  $N \subset M' \subset M$  and  $m < \omega$ . Suppose  $\bar{c} \in M$ ,  $\bar{c} \notin M$ ,  $l(\bar{c}) = m$ , tp  $(\bar{c}, M')$  is not orthogonal to N and R<sup>m</sup> [tp ( $\bar{c}$ , M'), L,  $\infty$ ] *is minimal (under the previous constraints).* 

*Then there is*  $\bar{c}' \in M$ ,  $\bar{c}' \notin M'$ ,  $I(\bar{c}) = I(\bar{c}')$ , tp  $(\bar{c}', M')$  does not fork over N and

$$
R^m[\operatorname{tp}(\bar{c}',M'),L,\infty]=R^m[\operatorname{tp}(\bar{c},M'),L,\infty].
$$

REMARK. We can wave this lemma if in the decomposition theorems we omit  $2.3(1)$ .

PROOF. Let  $\varphi(\bar{x}, \bar{b}) \in \text{tp}(\bar{c}, M'), R'''[\varphi(\bar{x}, \bar{b}), L, \infty] = R'''[\text{tp}(\bar{c}, M'), L, \infty]$ . We work in  $\mathbb{C}^{eq}$ . Let  $r^*$  be a regular type not orthogonal to N and not orthogonal to tp ( $\bar{c}$ , M'), with  $R^1[r^*, L, \infty]$  minimal. By 0.1(1) (maybe replacing  $\bar{b}$  by  $\bar{b'} \subseteq \text{acl}\bar{b}$ ) there is a formula  $E = E(\bar{x}, \bar{y}, \bar{b})$  where  $\bar{b} \in M'$ , tp ( $\bar{c}$ , M') does not fork over  $\bar{b}$ , such that stp  $(\bar{c}/E, \bar{b})$  is  $\text{cl}^3(r^*)$ -simple not orthogonal to r. Moreover (see 0.1(2), (3)):

 $\oplus$  stp  $(\overline{c}/E, \overline{b})$  contains a formula  $\Theta(y, \overline{b})$  which is cl<sup>3</sup>( $r^*$ )-simple. Moreover, tp  $(\bar{b}, \emptyset) =$  tp  $(\bar{b}', \emptyset)$  implies  $\Theta(x, \bar{b}')$  is  $cl^3(r^*)$ -simple.

Let  $\varphi_0(y,\vec{b}) = (\exists \vec{x})[\varphi(\vec{x},\vec{b}) \wedge \vec{x}/E = y \wedge \Theta(y,\vec{b})]$ . We shall prove

(\*) there are  $\bar{b}' \in N$  realizing tp( $\bar{b}, \emptyset$ ), and  $c' \in M-M'$ ,  $\models \varphi_0[c', \bar{b}'],$ tp  $(c', M')$  not orthogonal to N (in fact, to  $r^*$ ).

If (\*) holds, there is  $\bar{c}'' \in M$  such that  $\varphi[\bar{c}'', \bar{b}'] \wedge \bar{c}''/E = c' \wedge \Theta(c', \bar{b})$ . Now tp( $\bar{c}''$ , M') is not orthogonal to N as  $c' \in \text{acl}(N \cup \{\bar{c}''\})$ , tp( $c', M'$ ) is not orthogonal to *N*. Now  $R^m$ [tp( $\bar{c}^n$ ,  $M'$ ),  $L$ ,  $\infty$ ]  $\leq R^m$ [ $\varphi(\bar{x}, \bar{b}')$ ,  $L$ ,  $\infty$ ] =  $R^m[\varphi(\bar{x}, \bar{b}), L, \infty]$ . Equality holds by the hypothesis " $R^m[\text{tp}(\bar{c}, M'), L, \infty]$  is minimal". ( $\bar{c}'' \not\in M'$  follows from: tp( $\bar{c}''$ , M') is not orthogonal to N.) Hence tp  $(\bar{c}'', M')$  does not fork over  $\bar{b}' \subset N$ , so we get our conclusion.

Now we shall prove  $(*)$ . For notational simplicity let E be the equality, and  $\varphi(\bar{x}, \bar{b}) \vdash \Theta(\bar{x}, \bar{b})$ . Choose  $\bar{a} \in N$  such that tp  $(\bar{b}, N)$  does not fork over  $\bar{a}$ . We have assumed that  $r^*$  (hence tp  $(\bar{c}, M')$ ) is not orthogonal to N, hence it is not orthogonal to some  $r \in S^*(N)$ . Also, let  $\bar{d}_0$  realize r, tp ( $\bar{d}_0$ , M) does not fork over N. W.l.o.g.  $r$  does not fork over  $\bar{a}$ .

Now there are  $\bar{b}_n$ ,  $\bar{c}_n$  ( $n < \omega$ ) such that  $\bar{b}_n \hat{c}_n$  realizes stp ( $\bar{b} \hat{c}_n$ , N),  $\{\bar{b}_n \hat{c}_n : 0 <$  $n < \omega$  is independent over  $(M \cup \overline{d}_0, N)$ ,  $\overline{b}_0^{\wedge} \overline{c}_0 = \overline{b}^{\wedge} \overline{c}$  (so  $\{\overline{b}_n^{\wedge} \overline{c}_n : n < \omega\}$  is independent over N). Note that  $\Theta(x,\bar{b}_n)$  is cl<sup>3</sup>(r)-simple.

For each n let  $\{\bar{c}_{n,i}: i < \omega\}$  be a family of sequences realizing stp  $(\bar{c}_n, \bar{b}_n \cup N)$ , independent over  $N \cup \overline{b}_n$ . Let  $\{\overline{d}^i : i < \omega\}$  be a family of sequences realizing tp  $(\bar{d}_0, N)$  independent over  $(N \cup \bigcup_{n} \bar{b}_n \cup \bigcup_{n,i} \bar{c}_{n,i}, N)$ . By V 2.7 for some k, l  $\text{tp} (\bar{d}^{0 \wedge \cdots \wedge \bar{d}^t}, N \cup \bar{c}_n), \text{tp} (\bar{c}_{n,0}^{\wedge \cdots \wedge \bar{c}_{n,k}}, N \cup \bar{c}_n)$  are not weakly orthogonal. Now w.l.o.g.  $l = 0$ .

Now by the proof of V 4.11 there is  $d_1, d_1 \in \text{acl}(N \cup \overline{d}_0) - N$ , tp $(d_1, N)$  not orthogonal to  $r^*$  and some  $n < \omega$ , and there is a formula  $\psi(x, \bar{a})$  such that

(a)  $\vdash \psi[d_1, \bar{a}];$ 

(b)  $d_i \in \text{acl}(\bar{a} \cup \bigcup_{l \leq n} \bar{b}_l \cup \bigcup_{l \leq n, i \leq i(l)} \bar{c}_{l,i}'$  where  $\bar{c}_{l,i}'$  realizes stp  $(\bar{c}_l, \bar{b}_l), i(l) \leq \omega$ (for  $l < n$ ).

(c)  $\psi(\mathfrak{C}^{eq}, \bar{a}) \subseteq \text{acl}(\bar{a} \cup \bigcup_{l \leq n} \bar{b}_l \cup \bigcup_{l \leq n} \varphi(\mathfrak{C}^{eq}, \bar{b}_l)).$ 

By 1.2(2) there are  $\bar{b}'_1 \in N$  ( $l < n$ ) realizing tp ( $\bar{b}_i$ ,  $\bar{a}$ ) such that

(d)  $\psi(\mathbb{C}^{eq}, \bar{a}) \subset \text{acl}(\bar{a} \cup \bigcup_{l \leq n} \bar{b}'_l \cup \bigcup_{l \leq n} \varphi(\mathbb{C}^{eq}, \bar{b}'_l)).$ 

This is not exactly a first order property, but if it holds then some first order formula witnesses it, by compactness. Note that  $\Theta(\bar{x}, \bar{b}_1)$  is  $cl^3(r^*)$ -simple (as  $\text{tr}(b'_{i},\emptyset) = \text{tr}(b,\emptyset)$ . Remember  $\text{tr}(d_{i},N)$  is not orthogonal to  $r^{*}, \psi(x, \bar{a}) \in$ tp  $(d_1, N)$ . Easily tp  $(\bar{c}, M' \cup \psi(\mathbb{C}^{eq}))$  does not fork over  $\bar{b} \cup N \cup \psi(\mathbb{C}^{eq})$  hence  $\mathrm{Cb}(tp(\bar{c}, M' \cup \psi(\mathbb{C}^{eq}))) = \mathrm{Cb}(tp(\bar{c}, \bar{b} \cup N \cup \psi(\mathbb{C}^{eq})))$  and the type of this set over  $M' \cup \psi(\mathbb{C}^{eq})$  does not fork over  $\bar{b} \cup N \cup \psi(\mathbb{C}^{eq})$ .

Hence by 1.4(1) in Cb(tp( $\bar{c}$ ,  $\bar{b} \cup N \cup \psi(\mathbb{C}^{eq})$ )) there is an element  $d_2 \notin \text{acl}(N \cup \overline{b})$  such that tp  $(d_2, N \cup \overline{b})$  is not orthogonal to  $r^*$ . As tp  $(\overline{c}, M')$ does not fork over  $\bar{b}$ , clearly tp(d<sub>2</sub>, M') does not fork over  $\bar{b} \cup N$ . So as  $d_2 \notin \text{acl}(\bar{b} \cup N)$  also  $d_2 \notin M'$  remembering that

 $\text{Cb}$  (tp  $(\bar{c}, \bar{b} \cup N \cup \psi(\mathbb{C}^{eq}, \bar{a})) \subseteq \text{acl}(\bar{b} \cup N \cup \psi(\mathbb{C}^{eq}, \bar{a})).$ 

Clearly  $d_2 \in M$ .

Now there are  $\bar{c}_{i,i} \in M$ ,  $d_2 \in \text{acl}(\bar{b} \cup N \cup \{\bar{c}_{i,i} : l, i\})$ ,  $\models \varphi[\bar{c}_{i,i}\bar{b}_i']$  (this is by (d)). Now by  $\bigoplus$ , if tp ( $\overline{c}'_{i,j}$ , M') is orthogonal to  $r^*$ , then tp ( $\overline{c}'_{i,j}$ ,  $M' \bigcup \{c_{k,j} : k < l < i$  or  $k = l, j < i$ ) is also orthogonal to  $r^*$ . If this holds for every *l*, *i* then  $tp_{\star}(\{\bar{c}_{i,i} : l, i\}, M')$  is orthogonal to  $r^*$ . Hence tp  $(d_2, M')$  is orthogonal to  $r^*$ , contradiction. Now if tp ( $\bar{c}'_{~\iota}$ , M') is not orthogonal to  $r^*$ , (\*) holds, and we are finished.

1.6. LEMMA. *Suppose N*  $\subseteq$ <sub>s</sub> M and N  $\subseteq$  M'  $\subseteq$  M, and m <  $\omega$ . Suppose  $\bar{c} \in M$ ,  $\bar{c} \notin M'$ ,  $l(\bar{c}) = m$ ,  $tp (\bar{c}, M')$  *is not orthogonal to N and R* [tp ( $\bar{c}$ , M'), L,  $\infty$ ] *is minimal.* 

*If tp*  $(\bar{c}, M')$  does not fork over N then it is regular.

REMARK. We can omit 1.6 if in  $\S$ 2 we waive the regularity, i.e., omit 2.2(b) and 2.4(3).

PROOF. Let  $\bar{b} \in N$ , tp ( $\bar{c}$ , M') does not fork over  $\bar{b}$ . We can choose  $\psi$  such that  $\forall \psi[\bar{c},\bar{b}], R[\psi(\bar{x},\bar{b}), L, \infty] = R[\text{tp}(\bar{c},M'), L, \infty]$ . Let  $\bar{c}_n$  realize stp  $(\bar{c},\bar{b})$  for  $n <$  $\omega$ ,  $\{\bar{c}_n : \langle \omega \rangle\}$  independent over  $(M, \bar{b})$ . If stp  $(\bar{c}, \bar{b})$  is not regular then there are m and  $\bar{c}^*_i$  ( $j < m$ ) and n such that  $\bar{c}^*_i$  realizes stp ( $\bar{c}, \bar{b}$ ), tp ( $\bar{c}^*_i$ ,  $\bar{b} \cup {\bar{c}_i : i \leq n}$ ) forks over  $\bar{b}$  and tp  $(\bar{c}, \bar{b} \cup {\bar{c}_i : i \leq n} \cup {\bar{c}_i^* : j < m})$  forks over  $\bar{b}$ . For each  $j < m$ for some  $n_i \leq n$ ,  $tp(\bar{c}_{n_i}, \bar{b} \cup \{\bar{c}_i : i \leq n_i\} \cup \{\bar{c}_i^*\})$  fork over  $\bar{b}$  and remember  $tp({\bar c}, {\bar b} \cup {\bar c}_i : l \leq n \} \cup {\bar c^*} : j < m$ }) forks over  ${\bar b}$ . By III 1.2(2), III 2.6(2), II 3.7 for some finite  $\Delta$ ,  $k$ 

$$
R\left[\operatorname{tp}(\bar{c}_{n_j}, \bar{b} \cup \{\bar{c}_i : i < n_j\} \cup \{\bar{c}^*\}\right], \Delta, k\right] < R\left[\operatorname{tp}(\bar{c}, \bar{b}), \Delta, k\right]
$$

and

$$
R[\text{tp }(\bar{c},\bar{b}),\Delta,k] = R[\text{tp }(\bar{c},\bar{b}),\Delta,\mathbf{N}_0] = R[\text{stp }(\bar{c},\bar{b}),\Delta,\mathbf{N}_0]
$$

and

$$
R[\operatorname{tp}(\bar{c}, \bar{b} \cup \{\bar{c}_i : i \leq n\} \cup \{\bar{c}_j^* : j < m\}), \Delta, k] < R[\operatorname{tp}(\bar{c}, \bar{b}), \Delta, k].
$$

These properties for fixed  $\bar{c}$ ,  $\bar{b}$  are expressed by first-order formulas, i.e. there are formulas which  $\bar{b}$ ,  $\bar{c}$ ,  $\bar{c}$ ,  $\bar{c}$ ,  $\bar{c}$  satisfy and imply this (see II 2.19). So by 1.2(2) we can define  $\bar{c}_0', \ldots, \bar{c}_n' \in N$  and then define  $\bar{c}_i^{**} \in M$   $(j < m)$  such that

- (i)  $\bar{c}_0', \ldots, \bar{c}_n'$  realizes tp  $(\bar{c}, \bar{b})$ .
- (ii)  $\{\bar{c}_0',\ldots,\bar{c}_n'\}$  is independent over  $\bar{b}$ .
- (iii)  $\psi(\bar{c}^{**},\bar{b})$  for  $(j < m)$ .
- (iv)  $R^m$  [tp( $\bar{c}_n$ ,  $\bar{b} \cup {\bar{c}}'_i$ :  $l \leq n_i$ }  $\cup {\bar{c}}^{**}$ :  $j < m$ },  $\Delta, k$ ]  $< R^m$  [stp( $\bar{c}, \bar{b}$ ),  $\Delta, k$ ].
- (v)  $R^m$  [tp  $(\bar{c}, \bar{b} \cup {\bar{c}}'; 1 \leq n) \cup {\bar{c}}^*$ ;  $j < m$ ,  $\Delta, k$ ]  $< R^m$  [stp  $(\bar{c}, \bar{b}), \Delta, k$ ].

By (v), tp( $\bar{c}, \bar{b} \cup \{\bar{c}'_i: l \leq n\} \cup \{\bar{c}^{**}_i: j < m\} \subseteq \text{tp}(\bar{c}, M' \cup \{\bar{c}^{**}_i: j < m\})$  forks over  $\bar{b}$ . But tp( $\bar{c}$ , M') does not fork over  $\bar{b}$ , hence  $\bar{c}_{i(0)}^{**} \notin M'$  for some  $j(0)$ . As tp ( $\bar{c}$ , M'), tp ( $\bar{c}^{**}_{i(0)}, M'$ ) are not orthogonal, and the first does not fork over N, the second is not orthogonal to N. For notational simplicity assume  $n = n_{\text{ion}}$ .

By (iv) (and (i)) tp  $(\bar{c}'_n, \bar{b} \cup {\bar{c}'_i : l < n} \cup {\bar{c}^*_{i(0)} \atop n \geq 0})$  forks over  $\bar{b}$ .

(Note that as  $\bar{c}'_n$ ,  $\bar{c}$  realizes the same type over  $\bar{b}$ ,

$$
R^{m}[\operatorname{ktp}(\bar{c}',\bar{b}),\Delta,k]=R^{m}[\operatorname{stp}(\bar{c}',\bar{b}),\Delta,k]=R^{m}[\operatorname{tp}(\bar{c},\bar{b}),\Delta,k]).
$$

Hence by (ii)  $tp(\bar{c}'_n, \bar{b} \cup {\bar{c}'_i : l < n} \cup {\bar{c}''_{i(0)}})$  forks over  $\bar{b} \cup {\bar{c}'_i : l < n}$ . Hence tp  $(\bar{c}_{i(0)}^{**}, \bar{b} \cup \{c'_i: l \leq n\})$  forks over  $\bar{b} \cup \{\bar{c}'_i: l < n\}$  and hence over  $\bar{b}$ . So

$$
R\left[\text{tp}\left(\bar{c}_{j(0)}^{\ast\ast},M'\right),L,\infty\right]\leq R\left[\text{tp}\left(\bar{c}_{j(0)}^{\ast\ast},\bar{b}\cup\{\bar{c}'_i: l\leq n\}\right),L,\infty\right]
$$
\leq R^{\prime\prime}\left[\psi(\bar{x},\bar{b}),L,\infty\right].
$$
$$

This contradicts the minimality of the rank of  $\varphi(x,\bar{b})$ . Hence stp  $(\bar{c},\bar{b})$  is regular, and we are finished.

1.7. CLAIM. (1) For any  $A \subseteq M$  there is  $N \subseteq aM$ , such that  $A \subseteq N$ ,  $||N|| \leq$  $A + \lambda(T)$ .

(2) For any  $A \subset M$  there is  $N \subset M$  such that  $a \subset N$ ,  $||N|| \leq |A| + |D(T)|$ .

PROOF. Trivial.

REMARK. We can replace  $M$  by  $B$ . The next lemma will not be used in the sequal.

1.8. LEMMA.  $(\mathbb{C}^{eq})$  *Suppose*  $N \subseteq A \subseteq M$ ,  $N \subseteq_a M$ ,  $\bar{a} \in M$ ,  $\bar{a} \notin A$ ,  $r =$ stp  $(\bar{a}, A)$  is regular and trivial, but not orthogonal to N. Then there is a  $\bar{a}' \in M$ *such that stp*  $(\bar{a}', A)$  *is regular but not orthogonal to stp*  $(\bar{a}, A)$  *and does not fork over N.* 

PROOF. W.l.o.g.  $A = \text{acl}(A)$ . By 0.2 there is  $b \in \text{acl}(A \cup \bar{a}) - \text{acl}(A)$  (hence  $b \in M$ ) such that: tp(b, A) is regular but not orthogonal to stp( $\bar{a}$ , A) (and hence it is trivial, too) and some  $\varphi(x, \bar{c}_0) \in \text{tp}(b, A)$  is cl<sup>3</sup>(r)-regular. W.l.o.g. tp  $(\bar{b}, A)$  does not fork over  $\bar{c}_0$ . Choose  $\bar{d} \in N$  such that tp  $(\bar{c}_0^{\wedge} \langle b \rangle, N)$  does not fork over  $\overline{d}$ . Now as  $N \subseteq aM$  we can choose  $\overline{c}'_0$ , and  $b' \in N$  such that  $\text{stp}(\bar{c}_0'^{\wedge}(b'),\text{dbar})=\text{stp}(\bar{c}_0'^{\wedge}(b),\bar{d})$ . By V. 3.4,  $\text{stp}(b,\bar{c}_0)$ ,  $\text{stp}(b',\bar{c}_0')$  are not orthogonal. By [2] 5.I1 (or more elaborately [5] X 7.1) there is *b"* realizing stp  $(b', \bar{c}_0')$  such that tp  $(b, \bar{c}_0 \cup \bar{c}_0' \cup \bar{d} \cup \{b''\})$  forks over  $\bar{c}_0 \cup \bar{c}_0' \cup \bar{d}$ . Hence easily there is  $b^* \in M$  satisfying  $\varphi(x, \bar{c}_0)$  such that tp  $(b, \bar{c}_0, \bar{c}_0 \cup \bar{d} \cup \{b^*\})$  forks over  $\bar{c}_0 \cup \bar{c}_0' \cup \bar{d}$ .

Hence tp  $(b, A \cup \{b^*\})$  forks over A, and thus tp  $(b^*, A)$  is not orthogonal to tp  $(b, A)$  (and tp  $(\bar{a}, A)$ ). As  $\varphi(x, \bar{c}_0)$  is cl<sup>3</sup>(r)-regular,  $\varphi(x, \bar{c}_0)$  is cl<sup>3</sup>(r)-regular also. So as tp  $(b^*, A)$  is not orthogonal to r, it is regular, also it does not fork over  $\bar{c}_0$ hence over N.

#### **§2. A prime atomic model over stable amalgamation is enough**

HYPOTHESIS. T is superstable, and *if*  $\{M_1, M_2\}$  is independent over M,  $M \subset M_t$  ( $l = 1,2$ ) *then* there is a model N,  $\mathbf{F'_{\kappa_0}}$ -prime and  $\mathbf{F'_{\kappa_0}}$ -atomic over  $M_1 \cup M_2$ .

2.1. LEMMA. *Suppose*  $\{N_n : \eta \in I\}$  is a non-forking tree (see III or [2] 3.2). *Then there is a model N, which is*  $\mathbf{F}_{\kappa_0}^t$ -prime and  $\mathbf{F}_{\kappa_0}^t$ -atomic over  $\bigcup_{n\in I}N_n$ .

PROOF. Let  $I = \{ \eta_{\alpha} : \alpha < |I| \}$  be such that for every  $\alpha$  and  $k < l(\eta_{\alpha})$ ,

 $\eta_{\alpha}$  |  $k \in {\eta_{\beta} : \beta < \alpha}$ . We define by induction, on  $\alpha > 0$ ,  $M_{\alpha}$  such that

$$
(*) \qquad \begin{cases} (1) & M_{\alpha} \text{ is } F'_{n_0}\text{-prime and } F'_{n_0}\text{-atomic over } \bigcup_{\beta < \alpha} N_{n_{\beta}} \\ & \text{and even over } M_{\gamma} \cup \bigcup_{\beta < \alpha} N_{n_{\beta}} \text{ for each } \gamma < \alpha. \\ (2) & M_{i} \ (i \leq \alpha) \text{ is increasing continuous.} \end{cases}
$$

For  $\alpha = 1$ , let  $M_{\alpha} = N_{m}$ ; for  $\alpha$  a limit take the union, and for  $\alpha = \beta + 1$  use the hypothesis  $(M, M_1, M_2, N)$  there correspond to  $N_{\eta_a}$   $((l(\eta_a)-1), N_{\eta_a}, M_\beta, M_\alpha)$ here). Why does this work? Note that  $(\bigcup_{\gamma<\alpha}N_{\eta_{\gamma}}, \bigcup_{\gamma<|I|}N_{\eta_{\gamma}})$  satisfies the Tarski-Vaught condition (see 3.2A below). Hence if  $M_a$  is  $\mathbf{F}_{\mathbf{m}_0}^i$ -prime and  $\mathbf{F}_{\mathbf{x}_0}^t$ -atomic over  $\bigcup_{\beta \leq \alpha} N_{\eta_{\beta}}$ , then necessarily  $\{M_{\alpha}, \bigcup_{\gamma \leq |I|} N_{\eta_{\gamma}}\}$  is independent over  $\bigcup_{\beta<\alpha}N_{\eta_{\beta}}$ , and if  $\bigcup_{\alpha<|I|}N_{\eta_{\gamma}}\subseteq N$ , F an embedding of  $M_{\alpha}$  into  $N$ , F  $\bigcup_{\beta<\alpha}N_{\eta_{\beta}}=$ the identity, then  $F \cup G$  is an elementary embedding, where G is the identity map of  $\bigcup_{\gamma \leq |I|} N_{\eta_{\gamma}}$ .

2.2. THE ATOMIC DECOMPOSITION LEMMA (in  $\mathbb{C}^{eq}$ ). Suppose T is superstable with the dop. Then for any pair of models  $N_1 \subseteq_a M$ , there are elements  $a_i \in M$  $(i < \alpha)$  and models  $M_i$  such that:

- (a)  $N_1 \subseteq_a M_i \subseteq_a M;$
- (b) tp  $(a_i, N_1)$  is regular;

(c)  $|M_i| = N_1 \cup \{a_i\} \cup \{b_{i,a} : a < a_i\}$ , for every  $\alpha$ ,  $b_{i,a} \notin A_{i,a}$  and one of the *following occurs (letting*  $A_{i,\alpha} = N_1 \cup \{a_i\} \cup \{b_{i,\beta} : \beta < \alpha\}$ *)*:

(c1) tp  $(b_{i,\alpha}, A_{i,\alpha})$  is  $\mathbf{F}_{\kappa_0}^c$ -isolated,

(c2) tp( $b_{i,\alpha}$ ,  $A_{i,\alpha}$ ) is *orthogonal to*  $N_1$ ;

(d) for no  $b \in M - M_i$  is tp  $(b, M_i)$  orthogonal to  $N_1$ ;

(e) *M* is  $\mathbf{F}_{\mathbf{m}_0}^t$ -prime and  $\mathbf{F}_{\mathbf{m}_0}^t$ -atomic over  $\bigcup_{i<\alpha}M_i$  (and  $\mathbf{F}_{\mathbf{m}_0}^t$ -minimal);

(f)  ${M_i : i < \alpha}$  *is independent over*  $N_1$ .

**PROOF.** Let  $I = \{a_i : i \leq \alpha^*\}$  be a maximal subset of M independent over  $N_1$ , of elements realizing regular types of  $N_1$  and for each  $i < \alpha^*$  define  $b_{i,\alpha}$ ,  $M_i$ ,  $\alpha_i$  as required in (c),  $b_{i,\alpha} \in M - A_{i,\alpha}$  and  $\alpha_i$  is maximal. So (b), (c) hold trivially.

Why is  $|N_1| \cup \{a_i\} \cup \{b_{i,\alpha} : \alpha < \alpha_i\}$  the universe of a submodel (elementary, of course)? See IV 2.21. Now (d) follows from the choice of  $\alpha_i$ .

Clearly (f) follows by 1.3(2) and (a) by 1.2(3) provided that (e) holds. Apply 2.1. Let M' be  $\mathbf{F}_{\kappa_0}^i$ -prime  $\mathbf{F}_{\kappa_0}^i$ -atomic model over  $\bigcup_{i \leq \alpha} M_i$ . So w.l.o.g. (3) M'  $\subseteq M$ .

The only missing point is  $M' = M$ .

If not, there is  $c \in M - M'$ ,  $R[\text{tp}(c, M'), L, \infty]$  is minimal. Then by 1.5: tp  $(c, M')$  is orthogonal to  $N_1$ , or tp  $(c, M')$  does not fork over  $N_1$ , hence by 1.6

tp  $(c, N_1)$  is regular. The latter case contradicts the maximality of I. In the former case, we can find  $N_1^*$ ,  $M_i^*$  such that:  $N_1 \subseteq N_1^*$ ,  $M_i \subseteq M_i^*$  are  $\mathbf{F}_{\mathbf{w}_0}^a$ -saturated, tp ( $N^*_1$ ,  $\bigcup_{i \leq \alpha} M_i$ ) does not fork over  $N_1$ , tp ( $M^*_i$ ,  $\bigcup_{i \neq i} M^*_i \cup N^*_1$ ) does not fork over  $N^*_1 \cup M_i$ . By 3.3 (next section) the pair  $(\bigcup_{i \leq \alpha} M_i, \bigcup_{i \leq \alpha} M_i^*)$  satisfies the Tarski-Vaught condition.

Let tp (c, M') not fork over some  $\bar{a} \in M'$ , then tp ( $\bar{a} \cup_{i \leq \alpha} M_i$ ) is  $\mathbf{F}_{\kappa_0}^i$ -isolated, hence tp ( $\bar{a}$ ,  $\bigcup_{i \leq a} M^*_i$ ) is  $F'_{\text{No}}$ -isolated. By [2] §2 (as T does not have the dop) for some *i*, tp (c, M') (equivalently, stp (c,  $\bar{a}$ )) is not orthogonal to  $M^*$ , hence (as  ${M^*, \bar{a}}$  is independent over  $M_i$ , the type is not orthogonal to  $M_i$ . For notational simplicity let  $i = 0$ . As  $M_0 \subseteq_a M'$  (by 1.2(3), as mentioned above), we can apply 1.5 and find  $c' \in M - M'$ , tp  $(c', M')$  does not fork over  $M_0$ . If tp  $(c', M_0)$  is not orthogonal to  $N_1$ , we can get a contradiction to the maximality of I. If tp (c',  $M_0$ ) is orthogonal to  $N_1$  we get a contradiction to the choice of  $M_0$ .

2.2A. ASSERTION. We can add the demand: for each i, tp( $b_{i,\alpha}, A_{i,\alpha}$ ) is orthogonal to every trivial regular type not orthogonal to  $N_1$ .

**PROOF.** The only problem is when  $A_{i,a_i}$  is not the universe of an  $M_i$ . Let  $\alpha = \alpha_i$ . If not there is a formula  $\varphi(x, \bar{b}), \bar{b} \in A_{i,\alpha}, \models (\exists x)\varphi(x,\bar{b})$  but for no  $c \in A_{i,\alpha} \models \varphi(c,\bar{b})$ . Choose such  $\varphi(x,\bar{b})$  with minimal  $R[\varphi(x,\bar{b}), L, \infty]$ , hence every q,  $\varphi(x, \bar{b}) \in q \in S^1(A_{i\alpha})$  is  $\mathbf{F}_{n_0}^c$ -isolated. Let  $c \in M$ ,  $\models \varphi[c, \bar{b}]$ , so tp  $(c, A_{i\alpha})$ is  $\mathbf{F}_{\mathbf{r}_0}^c$ -isolated. So necessarily some trivial regular r is not orthogonal to stp(c,  $A_{i\alpha}$ ) and not orthogonal to N<sub>1</sub>. We can find  $\overline{d} \in N_1$  such that tp  $(\bar{b} \wedge \langle c \rangle, N_1)$  does not fork over  $\bar{d}$ , and r is not orthogonal to  $\bar{d}$ . As  $N_1 \subseteq_a M$ there are  $\vec{b}', c' \in N_1$  such that  $\vec{b}' (c')$  realizes stp  $(\vec{b} (c), \vec{d})$ . It is easy to see that r is not orthogonal to stp  $(c, \bar{b}')$ . So by [2] 5.11 or [5] X 7.1 there is c'' realizing stp  $(c, \bar{d} \cup \bar{b} \cup \bar{b}')$  such that tp  $(c'', \bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c'\})$  fork over  $\bar{d} \cup \bar{b} \cup \bar{b}'$ . Hence tp  $(c', \bar{d} \cup \bar{b} \cup \bar{b'} \cup \{c''\})$  fork over  $\bar{d} \cup \bar{b} \cup \bar{b'}$ , so some  $c^* \in M$ ,  $\models \varphi(c^*, \bar{b})$  and  $tp(c',\overline{d}\cup\overline{b}\cup b'\cup\{c^*\})$  fork over  $\overline{d}\cup\overline{b}\cup\overline{b'}$  hence  $tp(c^*,\overline{d}\cup\overline{b}\cup\overline{b'}\cup\{c'\})$ fork over  $\bar{d} \cup \bar{b} \cup \bar{b}'$ . As  $\bar{d} \cup \bar{b} \cup \bar{b}' \cup \{c'\} \subseteq A_{i,a}$ , we get a contradiction to the choice of  $\varphi(x,\bar{b})$ .

REMARK. We have essentially used (and proved):

2.2B. FACT. Suppose  $A \subseteq B \subseteq M$ , and tp  $(\bar{a}, A)$  has an extension over B which forks over A. Then for every  $\varphi(x, \bar{b}) \in \text{tp}(\bar{a}, A)$  there is  $\bar{a}' \in |M|$  such that  $\models \varphi[\bar{a}', \bar{b}]$  and tp  $(\bar{a}', B)$  forks over A.

2.3. CLAIM. In 2.2, 2.2A:

(1) *If*  $N_{2,i} \subseteq M_i$ ,  $N_1 \cup \{a_i\} \subseteq N_{2,i}$  *then*  $N_{2,i} <_{N_1} M_i$ .

(2) If  $p \in S^m(M_i)$  is orthogonal to  $N_1$  and  $Dp(tp(a_i,N_i))$  is  $\langle \infty,$  then  $Dp(p) < Dp$  (tp  $(a_i, N_i)$ ).

(3) If  $\bar{a} \in M_i$ , Dp (tp  $(a_i, N_1)$ )  $\lt \infty$ ,  $N_{2,i}$  as in (1) then

$$
Dp(tp(\bar{a}, N_{2,i})) < Dp(tp(a_i, N_1)).
$$

PROOF. (1) By 1.5, 1.6 and I's maximality.

(2) Let  $N_1^*$  be  $\mathbf{F}_{\mathbf{x}_0}^*$ -saturated,  $N_1 \subseteq N_1^*$ ,  $\{N_1^*, M_i\}$  independent over  $N_1$ . Now for every  $\alpha < \alpha_i$  (see 2.2) tp ( $b_{i,\alpha}, A_{i,\alpha}$ ) is orthogonal to every regular  $p \in S^m(N_1^*)$ with depth  $\geq 1$ . [If p is orthogonal to N<sub>1</sub> then it is orthogonal to any type over  $N_1^*$ . Suppose p is not orthogonal to  $N_1$ , then by 2.2A, if tp ( $b_{i,\alpha}, A_{i,\alpha}$ ) is not orthogonal to p then p is not trivial. We finish remembering that by [2] 5.10, as T does not have the dop, any regular type of depth $>0$  is trivial.] So clearly  $tp_*(M_i,N_1\cup\{a_i\})$ , is orthogonal to every regular complete type of  $N_1^*$  of depth > 0. We can find  $\mathbf{F}_{\mathbf{R}_{\alpha}}^*$ -saturated  $M^*, N^* \cup M_i \subseteq M^*$  with this property. We can apply  $[2]$  3.2; so the conclusion of 2.3(2), (3) follows.

(3) See the proof of (2).

2.4. THE DECOMPOSITION LEMMA. *Suppose T is superstable without the dop. Then for any model M there is a tree*  $I(\subseteq^{\omega} \|M\|)$  *and*  $N_n$   $(\eta \in I)$ ,  $a_n$   $(\eta \in I^-)$ *such that:* 

(1)  $N_n \subseteq_a M$  (hence  $N_n \leq M$ );

(2)  $N_n \subseteq_{\alpha} \cup \{N_{\nu}; \eta \leq \nu\};$ 

(3)  $p_{n^{\hat{}}(i)} =$  tp  $(a_{n^{\hat{}}(i)}, N_n)$  is regular;

(4)  $tp_*(\bigcup \{N_v : \eta \land \langle i \rangle \leq \nu, \nu \in I\} \cup \{N_v : \text{not } \eta \land \langle i \rangle \leq \nu, \nu \in I\})$  does not fork *over*  $N_n$ ;

(5)  $\{a_{\eta^{\wedge}(i)}: \eta^{\wedge}(i) \in I\}$  *is a maximal subset of M independent of*  $N_{\eta}$ ;

(6)  $tp_*(\bigcup \{N_\nu : \eta^\wedge \langle i \rangle \leq \nu \in I\}, N_{\eta^\wedge \langle i \rangle})$  is orthogonal to  $N_\eta$ ;

(7) if  $Dp(p_n) < \infty$ ,  $\eta$   $\land$  (i)  $\in I$  then  $Dp(p_n) < Dp(p_n)$ .

PROOF. Just combine the proofs of [2] 3.2 and 2.2 (and 2.3(2), (3)).

Now it is no problem to compute the number of non-isomorphic models, as in [2], [3] (using the same depth function).

# **§3. Universal Theories**

- 3.1. DEFINITION. (1) We call cl a closure operation if:
	- (i) for every A,  $A \subseteq \text{cl } A \subseteq \text{acl } A$ , and for every function symbol F (of  $\mathfrak{S}$ ) and  $\bar{a} \in A$ ,  $F^{\mathfrak{G}}(\bar{a}) \in \text{cl } A$ ;
	- (ii) cl (cl (A)) = cl A, and  $A \subseteq B$  implies cl  $A \subseteq c$ l B;

(iii) the property " $a \in cl A$ " is preserved by an automorphism of  $\mathfrak{C}$ .

- (2) We call a closure operation cl local if in addition
	- (iv) for every  $b \in cl A$ , there are a formula  $\varphi(x, \bar{y})$  and a sequence  $\bar{a} \in A$ such that  $\models \varphi[b, \bar{a}]$  and:  $\varphi(b_1, \bar{a_1})$  implies  $b_1 \in \text{cl}(\bar{a_1})$ .
- (3) For a set of formulas  $\Phi$  (of the form  $\varphi(\bar{x}, \bar{y})$ ) let acl. be defined by:

 $\text{acl}_\Phi^1(A) = \bigcup {\{\overline{b} : \text{for some } \overline{a} \in A, \text{ and } \varphi(\overline{x}, \overline{y}) \in \Phi, \models \varphi[\overline{b}, \overline{a}],\}$ and  $\varphi(\bar{x}, \bar{a})$  is an algebraic formula}  $\cup A$ ,

$$
\operatorname{acl}_{\Phi}^{0}(A)=A,
$$

 $\operatorname{acl}_{\Phi}^{n+1}(A) = \operatorname{acl}_{\Phi}^{1}(\operatorname{acl}_{\Phi}^{n}(A)),$ 

$$
\operatorname{acl}_{\Phi}(A) = \bigcup_{n} \operatorname{acl}_{\Phi}^{n}(A).
$$

(4) We call cl a  $\Phi$ -closure operation if it is an operation and  $A \subseteq \text{cl } A \subseteq \text{acl}_\Phi A$ .

3.1A. CLAIM. (1) *Every*  $\text{acl}_\Phi$  *is a local closure operation and is*  $\text{acl}_\Psi^1$  *for some rF.* 

(2) *Every local closure operation is a*  $\Phi$ -closure operation for some  $\Phi$ , and in Definition  $3.1(2)$  there is a  $\varphi$  satisfying in addition: for some n,  $\mathbf{I} = (\mathbf{\nabla} \vec{v})(\mathbf{J}^{\leq n}\vec{x})\varphi(\vec{x}, \vec{v})$  *and*  $\varphi(\vec{x}, \vec{a})$  + tp ( $\vec{b}, \vec{a}$ ).

PROOF. Easy.

3.2. DEFINITION. We call  $\langle M_s : s \in I \rangle$  a stable system *if*  $M_s \subseteq \mathcal{K}$ , I a family of finite subsets of  $\bigcup_{s \in I} s$  closed under subsets,

 $s < t \Rightarrow M_s \subseteq M_t$ and for every  $s \in I$ ,  $tp_*(M_s, \bigcup_{s \in I} M_t)$  does not fork over

$$
A_s \stackrel{\text{def}}{=} \bigcup \{M_t : t \subseteq s, t \neq s\}.
$$

We implicitly assume  $Th(\mathfrak{C})$  is stable.

3.3. CLAIM. (1) If  $I = \{s_\alpha : \alpha < \alpha_0\}$ ,  $[s_\alpha \subseteq s_\beta \Rightarrow \alpha \leq \beta\}$ ;  $M_s < C$ , and  $tp_*(M_{s_0}, \bigcup_{j<\alpha} M_{s_j})$  does not fork over  $A_{s_0}$ , then  $\langle M_s : s \in I \rangle$  is a stable system of *models* 

(2) *If*  $\langle M_s : s \in I \rangle$  *is a stable system,*  $J \subseteq I$  *and*  $s \in I \land s \subseteq \cup J \Rightarrow s \in J$ , then  $(U_{s\in J}M_s, U_{s\in I}M_s)$  *satisfies the Tarski-Vaught condition (i.e. if*  $\bar{a} \in U_{s\in J}M_s$ ,  $\bar{b} \in \bigcup_{s \in I} M_s$ ,  $\mathfrak{E} \models \varphi[\bar{a},\bar{b}]$  then for some  $\bar{b}' \in \bigcup_{s \in J} M_s$ ,  $\mathfrak{E} \models \varphi[\bar{a},\bar{b}']$ ).

(3) *If*  $\langle M_s : s \in I \rangle$  *is a stable system,*  $N_s \prec M_s$ , tp  $(N_s, \bigcup_{i \in s, t \neq s} M_i)$  does not fork

*over*  $\bigcup_{i \in s, i \neq s} N_{i}$ ,  $[s < t \Rightarrow N_{s} \subseteq N_{i}]$  then  $\langle N_{s} : s \in I \rangle$  is a stable system and  $(U_{s\in I}N_{s}, U_{s\in I}M_{s})$  satisfies the Tarski-Vaught condition.

PROOF. Essentially like [4] 3.5. Since we do not want to assume that the reader is familiar with [4], we prove the claim completely:

3.3A. FACT. If  $\langle M_s : s \in I \rangle$  is a stable system  $J_0 \subset I$ ,  $J \subset I$ ,  $J_0$  is closed under subsets then  $tp_*(\bigcup_{i\in J}M_i,\bigcup_{s\in J_0}M_s)$  does not fork over  $\bigcup \{M_s : s \in J_0 \text{ and }$  $(\exists t \in J) s \subseteq t$ .

REMARK. If *J* is closed under subsets, the last set is  $\bigcup \{M_s : s \in J_0 \cap J\}$ .

PROOF. W.l.o.g. *J* is closed under subsets, and let  $J_1 = J \cap J_0$ . We can find a list  $\{s_{\alpha}: \alpha < \alpha^*\}$  of *I* (and  $\alpha_1 \leq \alpha_2 \leq \alpha_3 \leq \alpha$ ) such that  $[s_{\alpha} \subseteq s_{\beta} \Rightarrow \alpha \leq \beta],$  $J_0 = \{s_\alpha : \alpha < \alpha_2\}, J = \{s_\alpha : \alpha < \alpha_1, \text{ or } \alpha_2 \leq \alpha < \alpha_3\}.$  Clearly for  $\alpha < \alpha_3, \alpha \geq \alpha_2,$  $tp_*(M_{s_\alpha}, \bigcup_{\beta<\alpha}M_{s_\alpha})$  is included in  $tp_*(M_{s_\alpha}, \bigcup\{M_t:s_\alpha\not\subseteq t, t\in I\})$  hence does not fork over  $A_{s_\alpha}$ , but  $A_{s_\alpha} \subseteq \bigcup \{M_{s_\beta} : \beta < \alpha_1 \text{ or } \alpha_2 \leq \beta < \alpha\}$ . So  $tp_*(M_{s_\alpha}, \bigcup_{\beta < \alpha} M_{s_\beta})$ does not fork over  $\bigcup \{M_{ss} : \beta \leq \alpha_1, \text{ or } \alpha_2 \leq \beta \leq \alpha\}$ . By IV 3.2(1) we can conclude that  $tp_*(\cup \{M_{s_0} : \alpha_2 \leq \alpha < \alpha_3 \text{ or } \alpha < \alpha_1\}, \bigcup_{\beta < \alpha_2} M_{s_\beta}$ ) does not fork over  $\bigcup_{\beta \leq \alpha_1} M_{sg}$ , but this is as required.

3.3B. FACT. If  $S = \langle M_s : s \in I \rangle$  is a stable system  $\bar{a}_i \in M_{s(i)}$   $(l \leq n)$ ,  $t \subseteq \cup I$ and  $\vdash \varphi[\bar{a}_0,\ldots,\bar{a}_{n-1}]$  then we can find  $\bar{a}'_1 \in M_{s(l) \cap t}$  such that  $\models \varphi[\bar{a}'_0,\ldots,\bar{a}'_{n-1}]$ and  $s(l) \subseteq t \Rightarrow a'_i = a_l$ .

PROOF. W.l.o.g.  $s \subseteq s(l) \Rightarrow s \in \{s(m): m < l\}$ . We prove it by induction on n. For  $n = 0$  there is nothing to prove, and for  $n = 1$  note  $M_{s(t)\cap t}$  is an elementary submodel  $M_{s(t)}$ . So suppose we have proved for n and we shall prove for  $n + 1$ , i.e. for given  $\bar{a}_i \in M_{s(t)}$   $(l < n + 1)$ ,  $t \subseteq \cup I$  and  $\varphi$ . W.I.o.g. the  $s(l)$   $(l \leq n)$  are distinct and  $s(n) \not\subseteq s(l)$  for  $l < n$ . We concentrate on the case  $s(n) \not\subseteq t$ . As tp ( $\bar{a}_{s(n)}$ ,  $\bigcup_{l \leq n} M_{s(l)}$ ) does not fork over  $A_{s(n)}^s$  clearly  $\varphi(\bar{a}_0, \ldots, \bar{a}_{n-1}, \bar{x})$  does not fork over  $A_{s(n)}^s$  hence is realized in every model which includes  $A_{s(n)}^s$ . So for some type  $p = p(\bar{x}_i)_{i \leq \alpha}$  over  $A^s_{(n)}$  (infinitely many variables)  $p({\bar{x}}_0,..., {\bar{x}}_i, \ldots)$   $\vdash$  V<sub>i $\in \alpha$ </sub>  $\varphi({\bar{a}}_0, \ldots, {\bar{a}}_{n-1}, {\bar{x}}_i)$ . So for some  ${\bar{b}} \subseteq A^S_{s(n)}$  and  $\psi$  =  $\psi(\bar{x}_0,\bar{x}_1,\ldots,\bar{x}_k, b)$ 

$$
\vdash (\exists \bar{x}_0, x_1, \ldots, x_k) \psi(\bar{x}_0, \ldots, \bar{x}_k, \bar{b}),
$$
  

$$
\psi(\bar{x}_0, \ldots, \bar{x}_k, \bar{b}) \vdash \bigvee_{i \leq k} \varphi(\bar{a}_0, \ldots, \bar{a}_{n-1}, \bar{x}_i).
$$

As  $\bar{b} \subset A_{s(n)}^{s}$  ( $\forall s \subset s(n)$ ) [ $s \in \{s(l): l < n\}$ ], w.l.o.g.  $\bar{b} = \bar{b_0}^{\wedge} b_1^{\wedge} \cdots \hat{b}_{n-1}, \bar{b_l} \subseteq$  $M_{s(1)}$ ,  $[s(1) \not\subseteq s(n) \Rightarrow \bar{b}_i$  empty]. Now apply the induction hypothesis to  $\bar{a}_i \bar{b}_i \in$   $M_{s(l)}$  (for  $l < n$ ) and the formula

$$
(\exists \bar{x}_0,\ldots,\bar{x}_k)\psi(\bar{x}_0,\ldots,\bar{x}_k,\bar{b}_0,\ldots,\bar{b}_{n-1})\wedge(\forall \bar{x}_0,\ldots,\bar{x}_k)
$$

$$
\left[\psi(\bar{x}_0,\ldots,\bar{x}_k,\bar{b}_0,\ldots,\bar{b}_{n-1})\rightarrow \bigvee_{i\leq k}\varphi(\bar{a}_0,\ldots,\bar{a}_{n-1},\bar{x}_i)\right].
$$

So there are  $\bar{a}'_i \bar{b}'_i \in M_{s(i)\cap i}$   $(l < n)$  satisfying the above formula (and as in 3.3B). Now clearly  $\bar{b}'_0 \wedge \bar{b}'_1 \wedge \cdots \wedge \bar{b}'_{n-1} \subseteq A^s_{s(n)\cap i}$ , hence there are  $\bar{c}_0, \ldots, \bar{c}_k \in M_{s(n)\cap i}$ such that  $\vdash \psi[\bar{c}_0,\ldots,\bar{c}_k,\bar{b}'_0,\ldots,\bar{b}'_{n-1}]$ . So for some  $i \leq k \vdash \varphi[\bar{a}'_0,\ldots,\bar{a}'_{n-1},\bar{c}_i]$ . So  $\bar{a}'_0, \ldots, \bar{a}'_{n-1}, \bar{a}'_{n-1}, \bar{a}''_n = \bar{c}_i$  are as required.

PROOF OF 3.3. (1) An exercise in non-forking.

(2) Follows from Fact 3.3B.

(3) First we prove that  $\langle N_s : s \in I \rangle$  is a stable system. For every s,  $tp_*(M_s, \cup \{M_t : M_t \in I, s \not\subseteq t\})$  does not fork over  $\cup \{M_t : t \subseteq s, t \neq s\}$ , hence (as  $N_s \subseteq M_s$ ) also tp $*(N_s, \cup \{M_t : t \in I, s \nsubseteq t\})$  does not fork over  $\cup \{M_t : t \subseteq s, t \neq s\}.$ But  $tp_*(N_s, \cup \{M_t : t \subseteq s, t \neq s\})$  does not fork over  $\cup \{N_t : t \subseteq s, t \neq s\}$ . So by III 0.1,  $tp_*(N_s, \cup \{M_t : t \in I, s \nsubseteq t\})$  does not fork over  $\cup \{N_t : t \subseteq s, t \neq s\}$ . As  $N_t \subseteq$ M<sub>t</sub>, by monotonicity of non-forking we get the stability of the system ( $s \subseteq t \Rightarrow$  $N_s \subseteq N_t$  was assumed, and we know *I* is as required).

The Tarski-Vaught condition follows by Fact 3.3B and the following fact. Let  $j \notin \bigcup I, J = I \cup \{s \cup \{j\} : s \in I\}$ , and  $N_{s \cup \{j\}} = M_s$ .

3.3C. FACT.  $\langle N_s; s \in J \rangle$  is a stable system  $(J, N_s)$  as above).

PROOF. Let  $s_{\alpha}$  ( $\alpha < \alpha_0$ ) be as in 3.3(1), and define  $t_{\alpha}$  ( $\alpha < 2\alpha_0$ ) by:  $t_{2\alpha} =$  $s_{\alpha}, t_{2\alpha+1} = s_{\alpha} \cup \{j\}.$  Clearly  $J = \{t_{\alpha} : \alpha < 2\alpha_0\}$  and  $t_{\alpha} \subseteq t_{\beta} \Rightarrow \alpha \leq \beta.$  Now use 3.3(1): For  $\alpha$  even (=2 $\beta$ ) remember we have proved tp<sub>\*</sub>(N<sub>sp</sub>,  $\cup$ {M<sub>s</sub>; s  $\in$ I,  $s_{\beta} \not\subseteq s$ }) does not fork over  $\cup \{N_s : s \subseteq s_{\beta}, s \neq s_{\beta}\}\$  and this is what we need. For  $\alpha$  odd (= 2 $\beta$  + 1) remember tp<sub>\*</sub>( $M_{s_{\beta}}, \bigcup_{\gamma \leq \beta} M_{s_{\gamma}}$ ) does not fork over  $\cup \{M_s : s \subseteq$  $s_{\beta}$ ,  $s \neq s_{\beta}$ ). As  $N_{s_{\beta}} \subseteq M_{s_{\beta}}$ , by III 0.1 this gives  $tp_*(N_{t_{\alpha}}, \bigcup_{\gamma < \alpha} N_{t_{\gamma}}) = tp_*(M_{s_{\beta}},$  $\bigcup_{\gamma < \beta} M_{s_{\gamma}} \cup N_{s_{\beta}}$  does not fork over  $\bigcup \{M_s : s \subseteq s_{\beta}, s \neq s_{\beta} \} \cup N_{s_{\beta}} = \bigcup \{N_s : s \subseteq s_{\beta}\}$  $t_{\alpha}$ ,  $s \neq t_{\alpha}$ , and this is what we need.

3.3D. FACT. Suppose  $\langle M_s : s \in I \rangle$  is a stable system, and each  $M_s$  is  $\kappa$ compact,  $a_{i,j} \in M_{s(i)}$  for  $i < \alpha$ ,  $j < j_\alpha$ ,  $\sum_{\alpha} j_\alpha < \kappa$ , and p is a set of  $\leq \kappa$  formulas, in the variables  $x_{i,j}$   $(i < \alpha, j < j_\alpha)$  satisfied by the assignment  $x_{i,j} \mapsto a_{i,j}$ . If  $t \subseteq \cup I$ (not necessarily finite) then we can find  $a'_{i,j} \in M_{s(i) \cap t}[s(i) \subseteq t \implies a'_{i,j} = a_{i,j}]$  such that p is satisfied by the assignment  $x_{i,j} \mapsto a'_{i,j}$ .

PROOF. Use Fact 3.3B, and the observation: a set of formulas (not necessarily

over M) of power  $\lt k$  finitely satisfiable in a  $\kappa$ -compact model, is satisfiable in it.

3.3E. FACT. Suppose  $\langle M_s : s \in I \rangle$  is a stable system,  $t \subseteq \cup I$ , each  $M_s$  is  $\mathbf{F}^{\alpha}_{\kappa}$ -saturated and  $a_{i,j} \in M_{s(i)}$  for  $i < \alpha, j < j_i$ ,  $\sum_{i < \alpha} j_i < \kappa$ , then we can find  $a'_{i,j} \in M_{s(i) \cap b}$ , such that

 $\sup_*(\langle a'_{i,j}:i<\alpha, j\leq j_\alpha\rangle, \{a_{i,j}:s(i)\subseteq t\}\equiv \sup_*(\langle a_{i,j}:i<\alpha, j\leq j_\alpha\rangle, \{a_{i,j}:s(i)\subseteq t\}\rangle).$ 

PROOF. Left to the reader.

3.3F. FACT. If  $S = \langle M_s : S \in I \rangle$  is a stable system, then for any  $t \in I$ ,  $tp_*(M_{\alpha} \cup \{M_s : t \not\subset s\})$  is definable over  $A_{\alpha}^s$ .

This follows from 3.3(2).

3.4. THE MAIN THEOREM. *Suppose T is stable and cl is a*  $\Sigma_1$ -closure operation; *then at least one of the following holds:* 

(A) *If*  $M_0 < M_1$ ,  $M_2$ ,  $\{M_1, M_2\}$  independent over  $M_0$ , then  $cl$   $(M_1 \cup M_2) < \mathfrak{C}$ .

(B) *There is a set A = cl A such that the theory of*  $\mathfrak{C} \restriction A$  *is unstable*  $(\mathfrak{C} \restriction A$  *is the model*  $\Im$  restricted to the set A, which by Definition 3.1(1) is closed under *functions); moreover, the theory of*  $\mathfrak{C} \restriction A$  has the independence property (see [1] II §4). *In fact, we can have A* = cl( $\bigcup_{i\leq i_0}M_i$ ),  $i_0\neq 0$ .

REMARKS. (1) If cl is as above, by  $3.5(2)$  we can assume that in a counterexample to (A),  $M_1$ ,  $M_2$  are isomorphic over  $M_0$ , hence get rid of the predicate P in the proof.

(2) Really we do not need the order  $\leq$  of L, but then we have to work a little more. It is also quite reasonable that we can replace stable by "without the independence property," and then in (A) say " $tp_*(M_2, M_1)$  is finitely satisfiable in  $M_0$ ," but this was not checked.

(3) In conclusion 3.6(1) we shall show that when T is universal then  $(B)$ implies that some completion of  $T$  is unstable.

PROOF. Suppose  $M_0$ ,  $M_1$ ,  $M_2$  form a counterexample to (A), and we shall prove (B). Let  $\lambda = |T| + ||M_1|| + ||M_2||$  and choose a model  $L = (|L|, <, P, R)$ , P a one place predicate,  $\leq a$  (linear) order, R a symmetric and reflexive two-place relation,  $L \models (\forall x, y)(P(x) \equiv P(y) \rightarrow xRy)$ , which is a  $\lambda$ -homogeneous and  $\lambda$ universal (i.e. any isomorphism from one submodel of  $L$  onto another, both of power  $\lt \lambda$ , can be extended to an automorphism of L, and any model L' of power  $\leq \lambda$  satisfying the other conditions can be embedded into L; L may have the power  $> \lambda$ ; see e.g. [1] I 1.8).

We shall now define for every  $s \in I = \{t : t \text{ a finite subset of } L\}$  a model  $M_s < \mathcal{C}$  and isomorphism  $F_h$  for every  $h \in PI$  (see below), Dom  $h = s$ , such that:

- (a)  $\langle M_s : s \in I \rangle$  is a stable system.
- (b) Let

 $PI = \{h : \text{for some } s, t \in I, h \text{ is an isomorphism form } L \upharpoonright s \text{ on } L \upharpoonright t\}.$ 

Now for any  $h \in PI$  there will be an isomorphism  $F_h$  from  $M_s$  onto  $M_i$  (where  $s =$ Dom  $h$ ,  $t =$ Range(h)) such that:

 $(b, \alpha)$  if  $f \subseteq h$  then  $F_f \subseteq F_h$ ,

 $(b, \beta)$  if  $h_2 = h_0 h_1$ , then  $F_{h_2} = F_{h_0} F_{h_1}$ ,

 $(b, \gamma)$  if h is the identity on  $s =$  Dom h, then  $F_h$  is the identity mapping on  $M_s$ .

(c)  $M_{\emptyset} = M_0$ ,  $M_{\{u\}}$  is isomorphic to  $M_1$  over  $M_0$  if  $u \in P$  and  $M_{\{u\}}$  is isomorphic to  $M_2$  over  $M_0$  if  $u \notin P$  (note that  $\emptyset \in I$  and we deal exactly with the  $u \in L$ ).

We denote  $F_h^* = F_h \upharpoonright \bigcup \{M_s : s \subseteq \text{Dom } h, |s| \leq n\}$ . Now the definition is as follows: we define by induction on n, M<sub>s</sub> ( $s \in I$ ,  $|s| = n$ ) and  $F_{h}^{n}$  ( $h \in PI$ ) such that (a), (b), (c) hold in the relevant cases (restricting to  $F_h^*$  when appropriate). In the end we shall let  $F_h = \bigcup_n F_h^n$ .

For  $n = 0$  use (c), for  $n = 1$  use (c) and the facts on forking (see III 0.1). For  $n > 1$  use 3.3: let  $\{t_i : i < i(*)\}$  be a list of all  $t \in I$ ,  $|t| = n$ .

Now we define by induction on  $i < i(*)$  the model  $M_{i}$ . If there is no  $j < i$  such that  $L \upharpoonright t_i \cong L \upharpoonright t_i$ , choose  $M_{t_i}$  as any  $M < \mathfrak{S}$ ,

$$
A_{i_i} \stackrel{\text{def}}{=} \bigcup_{\substack{s \subseteq i_i \\ s \neq i_i}} M_s \subseteq M, \qquad ||M|| \leq |T| + |A_{i_i}| \leq \lambda,
$$

and such that  $tp_*(M_{i_0}, \cup \{M_s : s \in I, |s| < n\} \cup \bigcup_{j \leq i} M_{i_j}$  does not fork over  $A_{i_i}$ (which is possible by the extension property of non-forking, see III 0.1). If there is  $j < i$  such that  $L \upharpoonright t_i \cong L \upharpoonright t_j$ , choose minimal  $j = j(i)$ , and let  $h_i^n$  be the isomorphism from L  $\uparrow t_i$  onto L  $\uparrow t_i$  (it is unique as  $\leq$  is a linear order of L and  $t_i$ is finite). Now there is an elementary mapping  $H_i^*$  extending  $F_{h_i}^{n-1}$  (see above) and whose domain is  $M_{v_0}$ , and  $tp_*(\text{Range}(H_i^*) , \cup \{M_i : s \in I, |s| < n\} \cup$  $\bigcup_{\alpha < i} M_{l_\alpha}$ ) does not fork over Range  $(F_{l_i}^{n-1})$ ; but note Range  $(F_{l_i}^{n-1}) = A_{l_i}$ . Now let  $M_{\mu}$  = Range  $(H_{\iota}^{n})$ .

So we have defined all the  $M_k$  but still have to define  $F_k^n$  for  $h \in Pl$ . Let  $\alpha(0)$  <  $\cdots$  <  $\alpha(k-1)$  be a list of  $\{i : t_i \subseteq \text{Dom } h\}$ , and let  $t_{\beta(i)}$  be the range of h  $\mid t_{\alpha(l)}\mid$ . Now we can define  $F_h^n$  as  $F_h^{n-1} \cup \bigcup_{l \leq k} H_{\beta(l)}^n(H_{\alpha(l)}^r)^{-1}$ . It is easy to check that this is a well-defined one-to-one function with the suitable range and

domain.  $F_h^*$  is an elementary mapping by 3.3F and if g extends h, both in PI then  $F_{\kappa}^{n}$  extend  $F_{h}^{n}$ . So we have defined  $M_{s}$ ,  $F_{h}$  ( $s \in I, h \in PI$ ) as required.

Now let

$$
J = \{s \in I : \text{for every } u \neq v \in s, L \models uRv\}
$$

and for any  $S \subseteq L$ , let  $B_s = cl \cup \{M_s : s \subseteq S, s \in J\}$ . By the hypothesis (that  $M_0$ ,  $M_1$ ,  $M_2$  exemplify the failure of 3.4(A)), if  $u \in P$ ,  $v \notin P$ ,  $u, v \in L$  but not  $uRv$ then  $\mathfrak{C} \restriction B_{\{\mu,\nu\}}$  is not an elementary submodel of  $\mathfrak{C}$ . By 3.3(2), the pair  $(B_{\{\mu,\nu\}}, B_{\nu})$ satisfies the Tarski-Vaught condition (inside  $\mathfrak{C}$ ). Hence  $\mathfrak{C} \restriction B_L$  is not an elementary submodel of  $\mathfrak{C}$ . So there are  $\bar{a} \in B_L$ , and a first order formula  $\varphi$ , such that  $\mathbb{E}~|~\varphi[\bar{a}], B_L$   $\mathbb{E}~|~\neg~\varphi[\bar{a}],$  so for some  $n < \omega$  and finite  $t \subseteq L$ ,  $\varphi$  is a  $\Sigma_n$ -formula or  $\Pi_n$ -formula and  $\bar{a} \subseteq B_n$ . Among all possible  $\bar{a}$ ,  $\varphi$ , n, t choose an example with minimal n, and for the fixed n, a minimal  $|t|$ , and for the minimal n and |t|, maximal  $|\{s : s \subseteq t, s \in J\}|$ . It is easy to prove that  $\varphi$  cannot be quantifier-free, nor a  $\Pi_n$ -formula, so  $n \ge 1$ ,  $\varphi(\bar{x}) = (\exists \bar{y}) \psi(\bar{y}, \bar{x})$ ,  $\psi$  a  $\Pi_{n-1}$ formula. It is also very easy to see that necessarily for some  $u, v \in t$ ,  $\neg uRv$ . Hence  $B_i = c \mid (B_{i-i\mu} \cup B_{i-i\mu})$ . So w.l.o.g. there are  $\overline{b} \in B_{i-i\mu}$ ,  $\overline{c} \in B_{i-i\mu}$  so that  $\overline{a}$ is algebraic over  $\bar{b} \hat{ } \tilde{c}$ , in fact for some  $\Sigma_1$ -formula  $\Theta$ ,  $\models \Theta[\bar{a}, \bar{b}, \bar{c}]$  and  $\Theta(\bar{x}, \bar{b}, \bar{c})$  is algebraic. Let k be the number of  $\bar{a}'$  satisfying  $\Theta(\bar{x}, \bar{b}, \bar{c}) \wedge \varphi(\bar{x})$  (in  $\mathfrak{C}$ ) and let  $\varphi^*(\bar{y}, \bar{z})$  be defined as

$$
(\exists^{\geq k}\bar{x})(\Theta(\bar{x},\bar{y},\bar{z})\wedge\varphi(\bar{x})).
$$

Now for every set  $Y \subseteq \lambda \times \lambda$  we can define in L elements  $u_i, v_i$   $(i < \lambda)$  such that:

(a) In L, u, u<sub>i</sub> realize the same quantifier-free type over  $t-{u, v}$ ; and similarly v,  $v_i$  realize the same quantifier-free type over  $t-\{u, v\}$ .

 $(\beta)$   $u_iRv_i$  holds iff  $\langle i, j \rangle \in Y$ .

(y) If  $u < v$  then for every i,  $j < \lambda$ ,  $u_i < v_j$  (in L) and if  $v < u$  then for every i,  $j < \lambda$ ,  $v_i < u_i$  (in L).

Let  $g_i$  be the function with domain  $t-\{v\}$ ,  $g_i(u)=u_i$ ,  $g_i(u')=u'$  for  $u' \in t - \{u, v\}$ . Let  $h_i$  be the function with domain  $t - \{u\}$ ,  $h_i(v) = v_i$ ,  $h_i(v') = v'$ , for  $v' \in t - \{u, v\}$ . It is easy to check that  $g_i, h_i \in PI$  and  $g_i \cup h_j \in PI$  iff  $\neg u_i R v_j$ (iff  $\langle i, j \rangle \notin Y$ ). Let  $\overline{b_i} = F_{g_i}(\overline{b})$ ,  $\overline{c_j} = F_{h_i}(\overline{c})$ .

FACT A. tp  $(\bar{b} \wedge \bar{c}, \emptyset) =$  tp  $(\bar{b}_i \wedge \bar{c}_i, \emptyset)$  (in  $\mathfrak{C}$ ).

This is because  $F_{\rm g} \cup F_{\rm h}$  is an elementary mapping (by 3.3F), and  $F_{\mathbf{s}_i} \cup F_{\mathbf{h}_i}(\bar{b} \wedge \bar{c}) = \bar{b}_i \wedge \bar{c}_j.$ 

FACT B.  $(\mathfrak{C} \restriction B_L) \models \varphi^*[\bar{b}_i, \bar{c}_i]$  if  $u_i R v_i$ .

We have chosen k such that  $\mathbb{G} \models (\exists^{k} \bar{x})[\Theta(\bar{x}, \bar{b}, \bar{c}) \times \varphi(\bar{x})]$ , i.e.  $\mathbb{G} \models \varphi^*[\bar{b}, \bar{c}]$ , hence by Fact A,  $\mathfrak{C} \models \varphi^*[\bar{b}_i, \bar{c}_i].$ 

Now if  $\mathfrak{C} \upharpoonright B_L \models \neg \varphi^*[\bar{b}_i, \bar{c}_j],$  then  $\bar{b}_i \upharpoonright \bar{c}_i, \varphi^*, n, t^* = t \cup \{u_i, v_j\} - \{u, v\}$ contradict the choice of  $\bar{a}$ ,  $\varphi$ , n, t, i.e.  $\varphi^*$  is a  $\Sigma_n$ -formula as  $\Theta$  is a  $\Sigma_1$ -formula,  $\varphi$  a  $\Sigma_n$ -formula and  $n > 0$ , also  $|t^*| = |t|$ , however the maximality of  $|\{s : s \subseteq t, s \in J\}|$  is contradicted.

**FACT C.** 
$$
(\mathcal{C} \restriction B_L) \models \neg \varphi^*[\bar{b}_i, \bar{c}_j]
$$
 if  $\neg u_i R v_j$ .

If  $\bar{a}' \in B_L$ ,  $\mathfrak{C} \restriction B_L \models \Theta[\bar{a}', \bar{b}, \bar{c}] \land \varphi[\bar{a}']$  then  $\mathfrak{C} \models \Theta[\bar{a}', \bar{b}, \bar{c}]$ ,  $\mathfrak{C} \models \varphi[\bar{a}']$  (by the minimality of n, as  $\varphi$  is a  $\Sigma_n$ -formula; for  $\Theta$  — trivially). Hence the set of  $\bar{a}' \in B_L$ for which  $(\mathbb{C} \restriction B_L) \models \Theta[\bar{a}', \bar{b}, \bar{c}] \land \varphi[\bar{a}']$  is a subset of the set of  $\bar{a}' \in B_L$  for which  $\mathfrak{E} \models \Theta[\bar{a}', \bar{b}, \bar{c}] \wedge \varphi[\bar{a}']$  which is a proper subset of the set of  $\bar{a}' \in \mathfrak{C}$  for which  $\mathfrak{E} \models \Theta[\bar{a}', \bar{b}, \bar{c}] \land \varphi[\bar{a}']$  (as witnessed by  $\bar{a}$ ). So we have just proved  $({\mathcal{C}} \restriction B_L) \models \neg \varphi^*[\bar{b}, \bar{c}]$ . So it is enough to find an automorphism of  ${\mathcal{C}} \restriction B_L$  taking  $\bar{b} \hat{c}$  to  $b_i \hat{c}_i$ . Now, as we have noted before,  $g_i \cup h_j \in PI$ , and  $F_{g_i \cup h_j}(\bar{b} \hat{c}) =$  $b_i^{\wedge} c_j$ . By choice of L,  $g_i \cup h_j$  can be extended to an automorphism f of L. Now  ${F_{ft}}: t \in PI$  is a directed family of elementary mapping, hence its union, F, is an elementary mapping. We shall prove that  $F^* = F[B_L]$  is as required.

(i)  $F^*(\bar{b} \wedge \bar{c}) = \bar{b}_i \wedge c_i$  : as  $F_{g_i \cup h_i}$  belong to the family.

(ii) F maps  $\bigcup_{s \in J} M_s$  onto itself by the properties of the  $F_f$ 's, and as  $cl(\bigcup_{s\in I}M_s)\subseteq Dom F$ , clearly by the properties of cl,  $F^*$  has to map cl( $\bigcup_{s\in J}M_s$ ) onto itself, so it is an automorphism of  $\mathfrak{C} \restriction B_L$ .

So we have proved that

$$
\mathfrak{C} \restriction B_{\mathfrak{L}} \vDash \varphi^*[\bar{b}_i, \bar{c}_i] \quad \text{iff } \langle i, j \rangle \in Y
$$

where  $Y \subset \lambda \times \lambda$  was arbitrary. So for some Y, we get that  $\mathfrak{C} \restriction B_L$  has a theory with the independence properties.

We can get some more relevant facts.

3.5. CLAIM. (1) In 3.4 we can replace "cl is a  $\Sigma_1$ -closure operation" by "cl is *a ~-closure operation," provided that:* 

(\*) if  $A = cl(U_{i \leq \alpha} M_i)$ ,  $\bar{a} \in A$ ,  $\varphi(\bar{x}) \in \Phi$  and  $\mathbb{C} {\upharpoonright} A \models \varphi[\bar{a}]$  then  $\mathfrak{C} \models \varphi[\bar{a}].$ 

(2) *If* cl *is local, then* "(A, *B) satisfies the Tarski- Vaught condition" implies*  "(cl A, cl *B) satisfies the Tarski- Vaught condition."* 

(3) If 3.4(A) *holds then for every such*  $M_0 < M_1, M_2, c1(M_1 \cup M_2)$ acl  $(M_1 \cup M_2)$ .

PROOF. (1) We replace in the proof  $\Sigma_n$ ,  $\Pi_n$  by  $\Sigma'_n$ ,  $\Pi'_n$  defined below by induction on *n*, and then the proof of 3.4 proves the assertion.

 $\Sigma'_n$ : if  $n = 0$ ,  $\Sigma'_n$  is the set of quantifier-free formulas; if  $n = 1$ ,  $\Sigma_n'$  is the set  $\Phi$  of formulas satisfying (\*) of 3.5(1); if  $n > 1$ ,  $\Sigma_n'$  is the closure of  $\Pi_{n-1}$  by conjunctions, disjunctions and existential quantifiers.

 $\Pi'_n$  is the set of negations of  $\Sigma'_n$  formulas.

(2) Trivial.

(3) Trivial.

3.6. CONCLUSION. (1) If T is a stable universal theory (i.e., every completion of T is stable), then for every model  $\mathfrak C$  of T, 3.4(A) holds.

(2) We can replace "universal" by  $\Sigma_2$  (i.e. every sentence of T is a  $\Sigma_2$ , or even  $\Sigma_2^{\text{al}}$  (defined below).

3.7. DEFINITION. (1) Let  $\Pi_1^{\text{at}}$  be the closure of the family of quantifier-free formulas by: universal quantification, conjunctions, disjunctions, and  $(\exists^{[1,k]}x)$ ... which mean: there are at least one but no more than  $k \bar{x}$ 's satisfying ...

(2) Let  $\Sigma_2^{\text{al}} = \{(\exists \bar{y})\varphi(\bar{y}, \bar{x}) : \varphi(\bar{y}, \bar{x}) \in \Pi_1^{\text{al}}\}, \Sigma_1^{\text{al}} = \{\neg \varphi : \varphi \in \Pi_1^{\text{al}}\}.$ 

**PROOF.** So Th ( $(\mathcal{E} \restriction B_L)$  is unstable, so what? Here comes the main use of the assumption that the theory T is universal (this is the only use of that fact). Remember the definition of  $B_L$ : it was obtained by applying cl on a subset of  $\mathfrak{C}$ . By part (1) (i) of Definition 3.1 and part (1) (ii) the set  $B<sub>L</sub>$  is closed under the functions of the model  $\mathfrak G$  hence it is a submodel of  $\mathfrak G$ . But since T has  $\Pi_1$ axiomatization, really  $\mathfrak{C} \restriction B_L \models T$ . So we found a model of T which is unstable (even has the independence property), namely a completion of  $T = Th(\mathfrak{C} \restriction B_L)$ ) which is not stable.

(2) By adding constants, translating the  $\Sigma_2$  to  $\Pi_1$  axioms, i.e., when applying 3.5, the set  $B_L$  such that Th $(\mathfrak{C} \mid B_L)$  has the independence property, satisfies:  $B_L = cI B_L$  and  $B_L$  extend some elementary submodel of  $\mathfrak{C}$ . We can conclude that  $\mathfrak{C} \restriction B_L$  is a model of T.

The proof for  $\Sigma_2^{\text{al}}$  is similar.

# §4. Examples

### 4.1. *Unstable*

Let the language have one two-place function, the theory: empty. This  $T$  is a variety, not stable and  $I(N_{\alpha}, T) = 2^{N_{\alpha}}$ .

4.2. *Stable not Superstable* 

Let the language have  $\omega$  one-place function  $F_{n}$ , and T be

$$
\{(\forall x)[F_n(F_m(x))=F_m(F_n(x))=F_n(x)]:n\leq m<\omega\}.
$$

Clearly T is a variety stable not superstable,  $I(N_a, T) = 2^{k_a}$ . A "natural" model is  $({}^{\omega\geq}\lambda, F_0, \ldots, F_n, \ldots), F_n(\eta)=\eta \restriction n.$ 

4.3. *Superstabte with the dop* 

Let us define a model M: let G be an abelian group of order 2,  $|M| = G \times \omega$ , with the functions:



Let T be the set of sentences  $(\forall x_1 \dots)(\tau = \sigma)$ ,  $(\forall x_1 \dots)(\tau \neq \sigma)$  for terms  $\tau$ ,  $\sigma$ , which M satisfies.

T is universal, superstable with the dop.

#### 4.4. *SuperstabIe without the dop, but Deep*

If L contains just one-place function  $F$ ,  $T$  is empty, then  $T$  is a variety and as required.

4.5. DEFINITION. Call T suitable if  $I(\mathbf{N}_{\alpha}, T) \geq |\alpha| + \mathbf{N}_0$ ,  $I(\mathbf{N}_{\alpha}, T) \leq I(\mathbf{N}_{\beta}, T)$ for  $\alpha \leq \beta$ .

4.6. CLAIM. (1) For universal  $T_i$  ( $i \leq \alpha$ ) there is a universal  $T_i |T| =$  $\sum_{i \leq \alpha} |T_i|$  (the power of  $T_i$  is  $|L(T_i)| + \aleph_0$ ) such that: if each  $T_i$  is suitable then T is *too and* 

$$
I(\lambda, T) = \sum_{i \leq \alpha} I(\lambda, T_i).
$$

(2) *Similarly with*  $I(\lambda, T) = \prod_{i < \alpha} I(\lambda, T_i)$ .

REMARK. (1) The suitability hypothesis is just to simplify the computation, and anyhow we here encounter only such T's.

(2) E.g. in (2) we should of course write  $Min\{2^{\lambda}, \Pi_{i \leq \alpha} I(\lambda, T_i)\}\)$ , but we shall ignore this.

**PROOF.** (1) W.l.o.g.  $L(T_i)$  ( $i < \alpha$ ) are pairwise disjoint. Let  $c_i$  ( $i < \alpha$ ) be new individual constants,  $P$  a monadic predicate and let  $T$  consist of the following sentences:

 $F(x_1, x_2, ...) \neq x_1 \rightarrow \bigwedge_{l=1} P(x_l)$ 

$$
c_i \neq c_{\alpha} \rightarrow c_j = c_{\alpha} \qquad \text{for } i \neq j < \alpha.
$$
  
\n
$$
- P(c_i)
$$
  
\n
$$
c_i = c_{\alpha} \rightarrow (\forall \bar{x})R(\bar{x}) \qquad R \in L(T_i) \text{ a predicate.}
$$
  
\n
$$
c_i = c_{\alpha} \rightarrow (\forall x_1, \dots) [F(x_1, \dots) = x_1] \qquad F \in L(T_i) \text{ a function symbol.}
$$
  
\n
$$
c_i \neq c_{\alpha} \rightarrow \psi^P \qquad \psi \in T_i.
$$
  
\n
$$
R(x_1, \dots, x_n) \rightarrow \bigwedge_{i=1}^n P(x_i) \qquad R \in L(T_i) \text{ a predicate.}
$$

(2) Let  $\{P_i : i \leq \alpha\}$  be new monadic predicates. T will say: the  $P_i$ 's are pairwise disjoint,  $P_i$  is a model of  $T_i$ , and the predicates and function symbols of  $T_i$  are trivial when applied to elements not all of which are in  $P_i$ .

 $F \in L(T_i)$  a function symbol.

4.7. CLAIM. *For a universal*  $T_0$  there is a universal  $T_1$ ,  $|T_1| = |T_0|$ , such that if  $T_0$  is suitable,  $T_1$  is suitable too; and

$$
I(\lambda, T_1) = 2^{I(\lambda T_0)}.
$$

**PROOF.** W.l.o.g.  $T_0$  has no individual constants (replace *c* by  $F_c(x)$ , adding  $(\forall x, y)$   $[F_c(x) = F_c(y)]$  to  $T_0$ ).  $T_1$  "says": E is an equivalence relation, each equivalence class is a model of  $T_0$ , and

$$
P(x_1, \ldots, x_n) \to \bigwedge_{i=1}^n x_i E x_1,
$$
  

$$
F(x_1, \ldots, x_n) \neq x_1 \to \bigwedge_{i=1}^n x_i E x_1
$$

for every predicate P and function symbol F of  $T_0$ .

4.8. CLAIM. *Claims* 4.6 *and* 4.7 *hold for "quasi-varieties" and [or "totally transcendental universal" and for "total transcendental quasi-varieties" instead of "'universal" provided that for the quasi-variety cases the language has no function symbols.* 

PROOF. The problem is taking care of the "quasi-variety".

*Case I. 4.6(1)* We assume w.l.o.g.  $L(T_i)$  are pairwise disjoint and with no individual constant. Let  $P_i$  ( $i < \alpha$ ) be new monadic predicates. T will say: (a)  $P_i(x) \wedge P_j(y) \rightarrow \varphi$  for  $i \neq j$ ,  $\varphi$  an atomic formula,

(b)  $P_i(x) \to P_i(y)$  for  $i < \alpha$ , (c)  $R(x_1, \ldots, x_n) \to P_i(y)$  for R a predicate of  $L(T_i)$ , (d)  $\psi \in T_i$  for some  $i < \alpha$ .

Of course  $L(T) = \bigcup_{i \leq \alpha} L(T_i) \cup \{P_i : i \leq \alpha\}$ . Let M be a model of T. If  $P_{i}^{M} = \emptyset$  for  $i < \alpha$  by (c), (d) all the relations of M are empty so M is determined up to isomorphism by its cardinality.

If  $P_{(0)}^M \neq \emptyset$ ,  $P_{(0)}^M \neq \emptyset$  and  $i(0) \neq i(1)$  the model M is again determined up to isomorphism by its cardinality: by (a).

Lastly suppose  $P_{i(0)}^M \neq \emptyset$ ,  $P_i^M = \emptyset$  for  $j \neq i(0)$ . By (b)  $P_{i(0)}^M = |M|$ , by (c)  $R^M = \emptyset$  for  $R \in L(T_i)$ ,  $j \neq i(0)$ , and by (d)  $M \upharpoonright L(T_{i(0)})$  is a model of  $T_{i(0)}$ . As in the cases with inclusion we can conclude:

$$
I(\lambda, T) = 2 + \sum_{i \leq \alpha} I(\lambda T_i).
$$

*Case II.* 4.6(2) Let T say:

- (a)  $P_i(x) \times P_i(x) \rightarrow \varphi$ ,  $i \leq j \leq \alpha$ ,  $\varphi$  any atomic formula,
- (b)  $R(x_1, \ldots, x_n) \rightarrow P_i(x_i)$ , R an *n*-place predicate is  $L(T_i)$ ,  $l \in \{1, \ldots, n\}$ ,
- (c)  $n \bigwedge_{m} P_i(x_m) \wedge \bigwedge_{i=1}^{m} \psi_i \rightarrow \psi, \bigwedge_{i=1}^{m} \psi_i \rightarrow \psi$  an axiom of  $T_i$ .

*Case III.* 4.7 Similar.

REMARK. We can add to the definition of *suitable* 

(iii)  $I(\lambda, T \cup \{c = c\}) = I(\lambda, T)$  for c a new individual constant.

Then also for quasi-varieties we can allow the languages to have individual constants and get similar results.

4.9. CLAIM. (1) There is a variety  $T_{\kappa}$  ( $\kappa \ge \aleph_0$  a cardinal)  $|T_{\kappa}| = \kappa$ ,  $I(\mathbf{N}_{\alpha}, T) = (\vert \alpha \vert + \mathbf{N}_{0})^{2^{k}}$ .

(2) *There is a variety T, I*( $\mathbf{N}_{\alpha}$ , *T) is 1 for*  $\alpha > 0$  *and*  $\mathbf{N}_0$  *for*  $\alpha = 0$ ,  $|T| = \mathbf{N}_0$ .

(3) *There is a variety T,*  $I(\mathbf{N}_{\alpha}, T) = 1, |T| = \mathbf{N}_{0}$ .

(4) *For any n*  $\lt \omega$ , there is a variety *T*,  $I(\mathbf{N}_{\alpha}, T) = \mathbf{I}_{n}(|\alpha| + \mathbf{N}_{0})$ .

PROOF. (1) Let the language have  $\kappa$  one-place functions  $F_i$  ( $i \leq \kappa$ ), and

$$
T = \{(\forall x)F_i(F_i(x)) = x : i < \kappa\}.
$$

(2) The axioms of T are those of a vector space over the rationals (i.e., for each rational number there is a one-place function symbol for multiplication by it and of course we have the addition function).

 $(3)$  T is empty.

(4) When  $n > 0$  let T be  $(\forall x)F^{n+1}(x) = c$ . For  $n = 0$  let  $T = {(\forall x)F(F(x)) = \forall x \in F}$ x). A model of *T*,  $M = \langle |M|, F^M \rangle$  is determined up to isomorphism by the following two cardinals:  $|\{a \in M : F(a) = a\}|$  and  $|\{a \in M : F(a) \neq a\}|$ .

4.10. CLAIM. (1) *There is a countable variety T, I*( $\mathbf{N}_{\alpha}$ , *T*) = **2**<sub>*m*</sub>(( $|\alpha| + \mathbf{N}_0$ )<sup>2<sup>*r*</sup><sup>0</sup>).</sup> (2) *There is a countable variety T*,  $I(\mathbf{N}_{\alpha}, T) = \mathbf{I}_{m}((|\alpha| + \mathbf{N}_{0})^{\mathbf{N}_{0}})$ .

PROOF. (1) For  $m = 0$ , let T consist of

$$
(\forall x)F_n(F_n(x))=x \ \ (\text{for } n<\omega).
$$

For  $m > 0$  let T consist of  $(G, F_n)$  are unary functions):

$$
G^{m+2}(x) = G^{m+1}(x),
$$
  
\n
$$
F_n(F_n(x)) = x,
$$
  
\n
$$
F_n(G(x)) = G(x),
$$
  
\n
$$
G(F_n(x)) = G(x) \text{ for } n < \omega.
$$

(2) First let  $m = 0$ , and  $T_0$  will consist of  $(\forall x)[F(F^{-1}(x)) = x = F^{-1}(F(x))]$ (i.e., essentially one unary one-to-one function). So a model  $M$  is characterized by the cardinals

 $\lambda_n(M) = |\{x \in M : n \geq 1 \text{ is minimal such that } F^n(x) = x\}/E|$ 

where

$$
xEy \stackrel{\text{def}}{=} (\exists n) \Big( x = F^n(y) \vee y = F^n(x) \Big).
$$

**(E** is an equivalence relation by our axioms, F" is the nth power of **F.)** 

For  $m > 0$ ,  $T = T_m$  will consist of  $(F, G, G^{-1})$  are unary functions)

$$
F^{m+1}(x) = F^{m+2}(x),
$$
  
\n
$$
G(F(x)) = F(G(x)),
$$
  
\n
$$
F(G(x)) = F(x),
$$
  
\n
$$
G(G^{-1}(x)) = x = G^{-1}(G(x)).
$$

Note that if  $M \vDash T$ ,  $A_i$  = Range  $(F<sup>i</sup>)$ , then G is the identity on  $A_i$ : for  $x \in A_i$ , there is y,  $x = F(y)$  so  $x = F(y) = F(G(y)) = G(F(y)) = G(x)$ .

4.11. CONJECTURE. For every variety T, either  $I(\lambda, T) = 2^{\lambda}$  for  $\lambda > |T|$  or  $I(\mathbf{N}_{\alpha}, T) < \mathbf{I}_{\omega}(|\alpha| + |T|)$  for every  $\alpha$ .

## **§5. Counterexample**

5.1. EXAMPLE. For each  $n < \omega$  there is a universal theory  $T_n$ , which has an unstable completion, but if  $\varphi(\bar{x}, \bar{y})$  is an unstable formula in such a completion, then  $\varphi$  is not  $\Sigma_n$ .

We first prove:

5.2. CLAIM. We can define by induction on n a theory  $T_n$  such that:

(A)  $T_n$  is universal, countable  $\aleph_0$ -stable, and in fact every completion of it has a *finite depth* (see [2, §4]), *in fact depth n, and I(* $\mathbf{N}_{\alpha}$ *, T<sub>n</sub>)* =  $\mathbf{I}_{n}$  ( $\alpha$  | +  $\mathbf{N}_{0}$ ).

(B)  $T_n$  is complete for  $\bigcup_{m \leq n} (\Sigma_m \cup \Pi_m)$  sentences.

(C)  $T_n$  is not complete for  $\Sigma_n$ -sentences, but only sentences from  $\Sigma_n \cup \Pi_m$  are *needed for completion, i.e. if*  $T$  *is a completion of*  $T_n$  *then* 

$$
T_n \cup \{\varphi : \varphi \in T, \varphi \in \Sigma_n \text{ or } \Pi_n \text{ sentence}\}\
$$

*is complete.* 

(D) *For every completion T of T., there is a complete universal countable*   $\aleph_0$ -stable theory  $S(T)$  of depth n (maybe in a larger language),  $T \subseteq S(T)$ ,  $S(T)$  as *in* (A), *such that*  $I(\mathbf{N}_{\alpha}, S(T)) \leq \mathbf{I}_{n}(|\alpha| + \mathbf{N}_{0}).$ 

(E)  $T_n^* = T_n \cup \{ \Theta : \Theta \text{ a } \Sigma_n \text{-sentence consistent with } T_n \}$  is a consistent theory.

(F)  $L_n$ , the language of  $T_n$ , as well as the languages of  $S(T)$  from  $(D)$ , consists *of predicates and one-place functions only.* 

(G) The language of  $S(T)$  (for T a completion of  $T_n$ ) is  $L_n^+$  (the language does *not depend on the theory T, only on n), and for every sentence*  $\psi \in L(T_n)$  there is a *sentence*  $\Theta_{\psi}^{n} \in L_{n}$  *such that for every completion T of*  $T_{n}$  [ $\psi \in T \leftrightarrow \Theta_{\psi}^{n} \in S(T)$ ].

**PROOF.** For  $n = 1$ . We let  $L_0$  contain the equality sign only, and  $T_0 =$  $\{(\forall x)x = x\}.$ 

Clearly  $T_n^*$  is (equivalent to)  $\{(\exists x_1 \cdots x_k) \land_{i \leq j} x_i \neq x_j : k < \omega\}$  (hence  $T_n^*$  is consistent, i.e. (E) holds). Also any other completion of  $T_0$  is the theory of a finite model, i.e. it is

$$
T_0^l = \left\{ \{ (\exists x_1 \cdots x_l) \} \bigwedge_{i < j} x_i \neq x_j \right\} \cup \left\{ (\forall x_1 \cdots x_{l+1}) \bigvee_{i < j} x_i = x_j \right\} \qquad \text{for some } l.
$$

It is now easy to prove (D): e.g., for  $T = T_0^k$ 

$$
S(T) = \{ (\forall x)(\forall y) F_n(x) = F_n(y) : n < k \}
$$
  

$$
\cup \left\{ (\forall x_1, \dots, x_{l+1}), \bigvee_{i < j} x_i = x_j \right\}
$$
  

$$
\cup \{ (\forall x) F_n(x) \neq F_m(x) : n < m < k \}.
$$

So  $L_n^+ = L_n \cup \{F_n : n < \omega\}$ . (We use one place functions instead of individual constants just for the convenience of the induction, formally  $-$  so that  $(F)$ holds.) The other parts are very easy, too.

*For n* + 1. So  $T_n$ ,  $L_n$ ,  $T_n^*$ ,  $S(T)$ ,  $L_n^+$  are defined and satisfy (A)-(F).

Let  $S^{\dagger}$  ( $l < \omega$ ) be each a copy of some theories  $S(T)$  such that every sentence in  $L_n$  consistent with  $T_n$  appears in infinitely many  $S^l$ , their languages are disjoint outside  $L_n$ , and  $F_n^m$ ,  $E_n \in S^{\perp}(l, m \lt \omega)$ . Now  $L_{n+1}$  consists of  $L_n \cup \{E_n\} \cup$  ${F_{\iota}^n : l < \omega} \cup \bigcup_{l<\omega} L(S^l)$ ,  $E_n$  a two place relation,  $F_n^m$  a one place function symbol (so it is still countable).  $T_{n+1}$  will consist of sentences saying the following:

(i) E is an equivalence relation, and for any function symbol F of  $L_n$ (necessarily one place by  $(F)$ ), and the choice of  $F<sub>i</sub><sup>n</sup>$ 

$$
(\forall x)[xE_nF(x)].
$$

(ii) Each  $E_n$ -equivalence class is a model of  $T_n$ , i.e. if  $\forall x_1 \cdots x_{k-1} \psi \in T_n$  ( $\psi$ quantifier-free) then

$$
(\forall x_0\cdots x_{k-1})\left(\bigwedge_{l
$$

(iii)  $F_n^i$  is really a hidden individual constant, i.e.

$$
\forall x \forall y (F'_n(x) = F'_n(y))
$$

and those "constants" are not  $E_n$ -equivalent, i.e.

$$
(\forall x)[\neg F'_{n}(x)E_{n}F'_{n}(x)] \qquad \text{for } l < k < \omega.
$$

(iv) The  $E_n$ -equivalence class of  $F_n^l(x)$  is a model of  $S^l$ . It is easy to check that  $T_{n+1}$  is as required.

(v) Every function symbol of  $S<sup>t</sup>$  not in  $L<sub>n</sub>$  is the identity outside the equivalence class which is a model of  $S<sup>i</sup>$  (and similarly for other non-logical symbols).

REMARK. In 5.2(A) we could add that  $T_n$  has only countably many completions.

PROOF OF 5.1. Let  $P$  be a place predicate,  $F$  a two-place function,  $E$  a two-place predicate, all not in  $T_n$  from 5.2. Now T will say:

(a) F is a pairing function from P into  $\neg P$ :

$$
(\forall xy)\bigg[\neg P(x) \vee \neg P(y) \rightarrow F(x, y) = x\bigg],
$$

$$
(\forall xy) \left[ P(x) \land P(y) \rightarrow \neg P(F(x, y)) \right],
$$
  

$$
(\forall x_1 x_2 y_1 y_2) \left[ P(x_1) \land P(y_1) \land P(x_2) \land P(y_2) \rightarrow \right]
$$
  

$$
\left[ F(x_1, y_1) = F(x_2, y_2) \equiv \left( x_1 = x_2 \land y_1 = y_2 \right) \right] \right]
$$

(b) E is an equivalence relation on  $\neg P$ , and each equivalence class contains at most one  $F(x, y)$ .

(c) Each equivalence class is a model of  $T_n$  and the unary functions of  $L(T_n)$ are the identity on  $\rightarrow P$ .

(d) If  $\neg P(F(x, y))$ , then  $F(x, y)$  is  $F_{n-1}^0(F(x, y))$ , i.e. it is one of individual constants of  $T<sub>n</sub>$  which we have hidden for technical reasons as one place functions.

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