

In the Dnestr Notebook [5] L. A. Bokut' raised the following question: "Is an arbitrary recursively presented associative (Lie) algebra over a prime field embeddable in a finitely presented associative (Lie) algebra?" We will show that for associative rings and algebras this question has an affirmative answer.

Suppose that K is either a commutative associative finitely generated ring with unity or a finitely generated extension of a prime field. We will show that any associative K -algebra with a recursively enumerable set of defining relations can be embedded in a finitely presented associative K -algebra.

In this paper the expression " K -algebra" means "associative K -operator ring." The unity of an algebra, if there is one, is not fixed in the signature. We denote the set of all positive integers by N . For a K -algebra A and elements $a_1, \dots, a_n \in A$, we denote by $A[a_1, \dots, a_n]$ the subalgebra of A generated by a_1, \dots, a_n . For K -algebra A and B , the direct sum of these K -algebras is denoted by $A \times B$, and $A \subseteq B$ signifies that A is a subalgebra of B . For a set of symbols X , we denote by $L[X]$ the set of all nonempty associative words in the alphabet X , and by $K[X]$ the free K -algebra with set X of free generators. The elements of $K[X]$ are linear combinations of the form $\sum_{w \in L[X]} \alpha_w w$, where the $\alpha_w \in K$ and almost all are equal to zero.

LEMMA 1. Suppose A is an arbitrary K -algebra $a_1, \dots, a_n, b_1, \dots, b_n \in K$, and $\varphi: A \rightarrow A$ is a mapping such that $\varphi(a_i) = b_i$ ($i=1, \dots, n$), and the following conditions are satisfied:

- 1) φ is an endomorphism of A as a K -module;
- 2) the restriction of φ to the subalgebra $A[a_1, \dots, a_n]$ is a homomorphism into A .

Then in some K -algebra containing A as a subalgebra the following system of equations in the unknowns $x, y, z, \beta_1, \dots, \beta_n$ is solvable:

$$\begin{aligned} xa_i z y &= b_i, \\ a_i z &= z \beta_i, \\ x z \beta_i \beta_j &= b_i x z \beta_j \quad (i, j = 1, \dots, n). \end{aligned}$$

Proof. Suppose $A = K[X]/I$, where I is an ideal of the free algebra $K[X]$. It is convenient to assume that preimages of the elements a_i, b_i ($i=1, \dots, n$) under the canonical homomorphism $K[X] \rightarrow A$ are chosen to be distinct letters of X , which we also denote by a_i, b_i ($i=1, \dots, n$). Furthermore, we assume that for each word $w \in L[X]$ there is chosen in $L[X]$ a word w' such that the equality $\varphi(w+I) = w'+I$ holds in $K[X]/I$. We again denote

this w' by $\varphi(w)$. Finally, in view of condition 2) of the lemma we may assume that we always have

$$\varphi(a_{i_1} a_{i_2} \dots a_{i_k}) = b_{i_1} b_{i_2} \dots b_{i_k}.$$

Suppose $X_1 = X \cup \{x, y, z, \beta_1, \dots, \beta_n\}$ and I_1 is the ideal of the free K -algebra $K[X_1]$, generated by the set

$$\begin{aligned} & I \cup \{xwz\beta y - \varphi[w\beta(\bar{a})] / w \in L[X] \text{ or } w = \emptyset, \\ & \beta \in L[\beta_1, \dots, \beta_n] \text{ or } \beta = \emptyset, w\beta \neq \emptyset\} \cup \\ & \cup \{a_i z - z\beta_i / i = 1, \dots, n\} \cup \\ & \cup \{xz\beta_i \beta_j - b_i xz\beta_j / i, j = 1, \dots, n\}. \end{aligned}$$

Here and below, $\beta(\bar{a})$ denotes the result of replacing each other β in the word β_i by a_i ($i = 1, \dots, n$). To prove the lemma it suffices to show that $I_1 \cap K[X] = I$.

Suppose $W \in I_1 \cap K[X]$. We will prove that $W \in I$. The element W can be written as a linear combination, with coefficients in K , of elements of the form

$$\begin{aligned} & v_1, v_2, & w_1 (xwz\beta y - \varphi[w\beta(\bar{a})]) w_2, \\ & w_1 (a_i z - z\beta_i) w_2, & w_1 (xz\beta_i \beta_j - b_i xz\beta_j) w_2, \end{aligned} \quad (1)$$

where v_1, v_2, w_1, w_2 are certain words in $L[X_1]$ or are empty words. We may assume that v_1 does not end, and that v_2 does not begin, with a letter of X .

Let us imagine x and y to be left and right parentheses. In a word $w \in L[X_1]$ with properly arranged parentheses, these are naturally divided into pairs $\langle x, y \rangle$, a left parenthesis x and its corresponding right parenthesis y . By the depth of occurrence of some such pair $\langle x, y \rangle$ in w we mean the difference between the number of parentheses y and the number of the parentheses x to the left of the x in the considered pair $\langle x, y \rangle$ in w . If w has a pair of parentheses with depth of occurrence s but no pair with depth of occurrence $s+1$, then the number $s+1$ is called the rank of w . If w has no parentheses, its rank is zero.

It is easy to see that if in an expression for W we group together the monomials with properly arranged parentheses, then we again obtain a linear combination of the elements (1). Since W itself is an element of $K[X]$, the monomials with improperly arranged parentheses cancel. We may therefore assume that W is a linear combination of the elements (1), where parentheses are properly arranged in all monomials.

Let s be the largest number such that an expression for W contains words of rank s . If $s=0$, then $W \in I$. Suppose $s>0$. We will show that W has a representation of the same form in which all words have rank less than s . The proof of the lemma then follows by induction.

A word $w \in L[X_1]$ with properly arranged parentheses is called proper if its rank is either less than s , or is equal to s and for any pair $\langle x, y \rangle$ of a depth $s-1$ the part of

the word w from x to y for these x, y has the form

$$xuz\beta y,$$

where $u \in L[X]$, or $u = \emptyset$, $\beta \in L[\beta_1, \dots, \beta_n]$, or $\beta = \emptyset, u\beta \neq \emptyset$.

It is easy to see that for each element of the form (1), all words of which it is a linear combination are proper or improper simultaneously. If in the representation for W we collect the improper words and cancel them, we again obtain a representation for W (since $W \in K[X]$). In view of what was said, this representation is a linear combination of elements of the form (1).

In each word of rank W in this representation of S , for each pair $\langle x, y \rangle$ of depth $S-1$ we replace the subwords $xuz\beta y$ for these x, y by $\varphi[u\beta(\bar{a})]$. We again obtain a representation of W .

We will show that the new representation is, as before, a linear combination of elements of the form (1). Since the ranks of words in the new representation are less than S , this will prove the lemma.

Consider a summand $v_1 u v_2$, where $u \in I$. Suppose $u = \sum_i \alpha_i w_i$, where $\alpha_i \in K$. Clearly, the ranks of all of the words $v_1 w_i v_2$ are the same. If their common rank is equal to S , then under the replacement described above the words $w_i \in L[X]$ are affected only when $v_1 \neq v_1', x, v_2 \neq x\beta y v_2'$. In the remaining cases, we obviously obtain a sum of the same form. Then after the replacement we have

$$\sum_i \alpha_i v_1'' \varphi[w_i \beta(\bar{a})] v_2'' = v_1'' \left(\sum_i \alpha_i \varphi[w_i \beta(\bar{a})] \right) v_2''.$$

But $\sum_i \alpha_i w_i \beta(\bar{a}) \in I$, and since φ is an endomorphism of A as a K -module, it follows that $\sum_i \alpha_i \varphi[w_i \beta(\bar{a})] \in I$.

Consider a summand $w_1 (xwz\beta y - \varphi[w\beta(\bar{a})]) w_2$. If in the word $w_1 \varphi[w\beta(\bar{a})] w_2$ the subword $\varphi[w\beta(\bar{a})]$ occurred within a pair of depth $S-1$, then clearly in the word $w_1 xwz\beta y w_2$ the pair $\langle x, y \rangle$ would have depth S , which is impossible. Therefore, obviously, after the replacement the expression under consideration either vanishes or keeps the same form.

Consider a summand $w_1 (a_i x - x \beta_i) w_2$. It suffices to look at the case where $w_1 \neq w_1', x w_1', w_2 \neq \beta y w_2'$ and this pair $\langle x, y \rangle$ has depth $S-1$. Then after the replacement we obtain

$$w_1'' \varphi[w a_i \beta(\bar{a})] w_2'' - w_1'' \varphi[w(\beta_i \beta)(\bar{a})] w_2''.$$

This expression vanishes, since $(\beta_i \beta)(\bar{a}) \neq a_i \beta(\bar{a})$.

Finally, consider a summand $w_1 (x x \beta_i \beta_j - b_i x \beta_j) w_2$. Again, it suffices to look at the case where $w_2 \neq \beta y w_2'$ and the considered pair $\langle x, y \rangle$ has depth $S-1$. After the replacement we obtain

$$w_1' \varphi[(\beta_i \beta_j)(\bar{a})] w_2'' - w_1' b_i \varphi[(\beta_j \beta)(\bar{a})] w_2''.$$

This expression also vanishes, since by choice of $\varphi(w)$ we have

$$\varphi[(\beta_i \beta_j)(\bar{a})] = b_i b_j b_i \dots b_i \neq b_i \varphi[(\beta_j \beta)(\bar{a})],$$

if $\beta \equiv \beta_{i_1} \dots \beta_{i_k}$.

The lemma is proved.

Remark 1 (in some sense a converse to Lemma 1). Let $\text{Hom}_n(\bar{a}, \bar{b}, \bar{\theta})$, where $\bar{a} = (a_1, \dots, a_n)$, $\bar{b} = (b_1, \dots, b_n)$, $\bar{\theta} = (\beta_1, \dots, \beta_n, x, y, z)$, denote the system of equalities in the statement of Lemma 1. We will show that if $\text{Hom}_n(\bar{a}, \bar{b}, \bar{\theta})$ is satisfied in some K -algebra A for certain elements $\bar{a}, \bar{b}, \bar{\theta}$, then the mapping $\varphi: A[a_1, \dots, a_n] \rightarrow A$, defined by the rule $\varphi(a) = xay$ for $a \in A[a_1, \dots, a_n]$, is a K -algebra homomorphism for which $\varphi(a_i) = b_i$ ($i=1, \dots, n$). It suffices to observe that

$$\varphi(a_{i_1} a_{i_2} \dots a_{i_k}) = xa_{i_1} a_{i_2} \dots a_{i_k} y = xz\beta_{i_1} \beta_{i_2} \dots \beta_{i_k} y = b_{i_1} b_{i_2} \dots b_{i_k} = \varphi(a_{i_1}) \varphi(a_{i_2}) \dots \varphi(a_{i_k})$$

Remark 2. In the sequel, we will apply Lemma 1 in the following situation. Suppose A, B are K -algebras, A is the subalgebra of B , generated by elements $a_1, \dots, a_n \in B$, $\varphi: A \rightarrow B$ is a K -algebra homomorphism for which $\varphi(a_i) = b_i$ ($i=1, \dots, n$), and A as a K -module is a direct summand of B , i.e., the K -module B contains a K -submodule C such that $A + C = B$. Then there exists a K -algebra containing B as a subalgebra in which the system $\text{Hom}_n(\bar{a}, \bar{b}, \bar{\theta})$ is solvable.

The assertion that for associative rings $A \subseteq B$, where A is a direct summand of the additive group of B , an additive homomorphism $\varphi: A \rightarrow B$ is defined in an extension of B by a rule $\varphi(a) = xay$, is due to Taitlin [1].

LEMMA 2. Suppose $f(i, j)$ is a recursive function defined for all $i, j = 1, 2, \dots$ ($i+j$) such that $f(i, j) = f(j, i)$. Suppose that $Y \subseteq \mathbb{N}^2$ is some recursively enumerable set and A is a K -algebra with generators x, y, z and defining relations

$$\{xy^i z + xy^j z = xy^{f(i,j)} z / i+j; i, j = 1, 2, \dots\} \cup \{xy^i z = xy^j z / \langle i, j \rangle \in Y\}.$$

Then there exists a K -algebra B with the following properties:

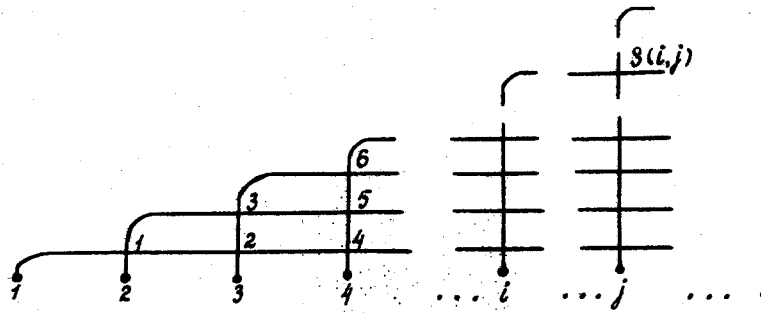
a) A is a subalgebra of B and, as a K -module, is a direct summand of the K -module B ;

b) B has a finite number of generators and a recursively enumerable set of defining relations, one of which has the form $\alpha + \beta = \gamma$ and the others are word equalities in the generator symbols.

Proof. For each $i=1, 2, \dots$ we define a function $n_i: \mathbb{N} \rightarrow \mathbb{N}$, and for $i+j; i, j=1, 2, \dots$, we define $s(i, j) \in \mathbb{N}$ as follows:

$$\begin{aligned} \text{if } 1 \leq i < j & \text{, then } n_j(i) = \frac{1}{2}(j-1)(j-2) + i; \\ \text{if } 1 \leq j \leq i & \text{, then } n_j(i) = n_{i+1}(j); \\ \text{if } 1 \leq i < j & \text{, then } s(i, j) = s(j, i) = n_j(i). \end{aligned}$$

The definition of these functions is illustrated by the following diagram.



Here from the natural numbers $1, 2, 3, \dots$ there emanate lines, each of which intersects each other line in exactly one point. The points of intersection are enumerated by the natural numbers. These numbers define the functions n_i and S . That is to say, the numbers on the line emanating from i are $n_i(1), n_i(2), \dots$. For example, $2, 3, 6, 9$ are the values $n_3(1), n_3(2), n_3(3), n_3(4)$ respectively. The number appearing at the intersection of the line emanating from i and the line emanating from j is $S(i, j)$. Thus, for any $n \in \mathbb{N}$ there exists a unique pair $1 < i < j$ such that $n = S(i, j) = S(j, i)$. Furthermore, $n = n_j(i) = n_i(j-1)$ and the functions n_i are one-to-one.

As generators of the desired K -algebra B we take the symbols $x, y, z, u, \alpha, \beta, \gamma$. As the set of defining relations we take the set of equalities

$$\{xy^i z = xy^j z / \langle i, j \rangle \in Y\} \cup \{xy^i z = xu^{n_i(j)} \varepsilon_{ij} z / i, j = 1, 2, \dots\} \cup \\ \cup \{xy^{f(i,j)} z = xu^{s(i,j)} \gamma z / i \neq j; i, j = 1, 2, \dots\} \cup \{\alpha + \beta = \gamma\}.$$

Here ε_{ij} is equal to α if $i+j$ is even, and is equal to β if $i+j$ is odd.

Note first that the defining relations of B imply those of A . Indeed, if $i < j$, then it follows from the relations of B that $xy^i z + xy^j z = xu^{n_i(j-i)} \varepsilon_{ij-1} z + xu^{n_j(i)} \varepsilon_{ij} z = xu^{s(i,j)} (\varepsilon_{ij} + \varepsilon_{j-1}) z = xu^{s(i,j)} \gamma z = xy^{f(i,j)} z$.

Now consider the ideal I_i of the free K -algebra $K[x, y, z, u, \alpha, \beta, \gamma]$ generated by the relations of B . Any element of this ideal can be written as a linear combination of elements of the form

$$\begin{aligned} &w_1 (xy^i z + xy^j z - xy^{f(i,j)} z) w_2, \quad w_1 (xy^i z - xy^j z) w_2 / \langle i, j \rangle \in Y / \\ &w_1 (xy^i z - xu^{n_i(j)} \varepsilon_{ij} z) w_2, \quad w_1 (xy^{f(i,j)} z - xu^{s(i,j)} \gamma z) w_2 \quad (i \neq j), \\ &w_1 (\alpha + \beta - \gamma) w_2, \end{aligned} \tag{1}$$

where w_1, w_2 are words in $x, y, z, u, \alpha, \beta, \gamma$, possibly empty.

By a reduction of a word $w \in L[x, y, z, u, \alpha, \beta, \gamma]$ we mean the simultaneous replacement in it:

of all subwords of the form $xu^n \alpha z$ by $xy^i z$, where i is such that $n = n_i(j)$ and $i+j$ is given.

of all subwords of the form $xu^r\beta z$ by $xy^i z$, where i is such that $r=n_i(j)$ and $i+j$ is odd;

of all subwords of the form $xu^r\gamma z$ by $xy^{f(i,j)} z$, where $i+j$ are such that $r=s(i,j)$.

A word w without subwords of the indicated form is called reduced. Let us now assume that some element $W \in K[x, y, z, u, \alpha, \beta, \gamma]$ is a linear combination of reduced words (in particular, belongs to $K[x, y, z]$) and at the same time belongs to I . Then W can be represented as a linear combination of elements of the forms (1). Reducing all words in such a representation, we obviously again obtain an expression for W in which there occur terms (1) of only the first, second, and perhaps the last forms. If $W \in K[x, y, z]$, then, cancelling the words in this representation in which the letters u, α, β, γ occur, we obtain that $W \in I$, the ideal generated in $K[x, y, z]$ by the defining relations of A . This proves that A is a subalgebra of B . Now let R be the submodule of the K -module B generated in B by the images under the canonical homomorphism of all reduced words in $x, y, z, u, \alpha, \beta, \gamma$, in which there must necessarily be occurrences of u, α, β, γ . Since any element of $K[x, y, z, u, \alpha, \beta, \gamma]$ modulo I is equal to a linear combination of reduced words, the K -module B is the sum of A and R . This sum is direct. Indeed, suppose $W_1 + W_2 \in I$, where $W_1 \in K[x, y, z]$, and W_2 is a linear combination of reduced words in $x, y, z, u, \alpha, \beta, \gamma$ in which u, α, β, γ occur. Then, by what has been proved, $W_1 + W_2$ can be represented as a linear combination of expressions of the first, second, and last forms in (1). Collecting the words in x, y, z and all the others separately, we obtain $W_1 \in I, W_2 \in I$.

The lemma is proved.

A small modification of an assertion proved by Mal'tsev [2] is the following

LEMMA 3. Suppose A is an arbitrary, at most countable K -algebra and $A = \{a_1, a_2, \dots\}$ is an arbitrary enumeration of its elements. Then there exist a K -algebra B containing A as a subalgebra and elements $a, b, c \in B$, such that $a_i = ab^i c$ ($i=1, 2, \dots$).

Proof. Consider $A_1 = A \times K[x]$. There exists an endomorphism $\varphi: A_1 \rightarrow A_1$ of A_1 as a K -module under which $\varphi(x^n) = a_n$ ($n=1, 2, \dots$). By Lemma 1, there exist a K -algebra $B \supseteq A_1$, and elements $u, v \in B$ such that $\varphi(a) = uav$ for all $a \in A_1$. Clearly, B is the desired algebra.

THEOREM. Suppose K is a commutative associative ring with unity or a field that is finitely generated over its prime subfield. Suppose A is an arbitrary associative K -algebra with a recursively enumerable set of defining relations. Then there exists a finitely generated associative K -algebra with a finite set of defining relations in which A is contained as a subalgebra.

Proof. Suppose $A = \{a_1, a_2, \dots\}$ is an enumeration of the elements of A in which each element of A is repeated at least twice. By Lemma 3, there exists a K -algebra B with three generators a, b, c such that $A \subseteq B$ and $a_i = ab^i c$ ($i=1, 2, \dots$). Since in the ring or field K the set of all true equalities in the generators of K is recursively enumerable, we may assume that B is also recursively presented, i.e., the set of all equalities in the elements a, b, c that are true in B is recursively enumerable. Moreover, we may assume that

the original algebra A has a finite set of generators, say a_1, \dots, a_k .

Let $Y = \{ \langle i, j \rangle \mid ab^i c = ab^j c \}$. Then $Y \subseteq N^2$ is a recursively enumerable set. Let $f(i, j)$ be a recursive function, defined for all $i, j \in N$ with $i \neq j$, such that $f(i, j) = f(j, i)$ and

$$ab^i c + ab^j c = ab^{f(i, j)} c \text{ in } B.$$

Consider the K -algebra S with generators x, y, z and defining relations

$$\{ xy^i z + xy^j z = xy^{f(i, j)} z \mid i \neq j; i, j = 1, 2, \dots \} \cup \{ xy^i z = xy^j z \mid \langle i, j \rangle \in Y \}.$$

By Lemma 2, there exists a K -algebra $S_1 \supseteq S$ in which the K -module S is a direct summand and which has an enumerable set of defining relations

$$\Sigma (x, y, z, u, \alpha, \beta, \gamma) \cup \{ \alpha + \beta = \gamma \}, \quad (1)$$

where Σ contains only word equalities in $x, y, z, u, \alpha, \beta, \gamma$.

There exists $\varphi: B \times S_1 \rightarrow B \times S_1$ such that $\varphi(x) = a, \varphi(y) = b, \varphi(z) = c$, φ is an endomorphism of $B \times S_1$ as a K -module, and the restriction of φ to S is a K -algebra homomorphism. By Lemma 1, there is K -algebra $S_2 \supseteq B \times S_1$ in which we can solve the system of equations

$$\text{Hom}_2 (x, y, z, a, b, c, \bar{\theta}_1). \quad (2)$$

Now consider the semigroup G with generators $x', y', z', u', \alpha', \beta', \gamma'$ and set of defining relations $\Sigma(x', y', z', u', \alpha', \beta', \gamma')$. By a theorem of Murskii [3], there exists a semigroup $G_1 \supseteq G$ with a finite set of generators and a finite set of defining relations

$$\Sigma_1 (x', y', z', u', \alpha', \beta', \gamma', \bar{\theta}_2). \quad (3)$$

Let $K[G_1]$ denote the semigroup K -algebra of the semigroup G_1 . Clearly, $K[G] \subseteq K[G_1]$ and $K[G]$, as a K -module, is a direct summand of $K[G_1]$. By Lemma 1, there is a K -algebra $S_3 \supseteq S_2 \times K[G_1]$ in which we can solve the system

$$\text{Hom}_3 (x', y', z', u', \alpha', \beta', \gamma', x, y, z, u, \alpha, \beta, \gamma, \bar{\theta}_3). \quad (4)$$

Suppose $g: N^2 \rightarrow N$ is a recursive function such that $ab^i c + ab^j c = ab^{g(i, j)} c$ in the K -algebra B . If H is the semigroup with generators ρ, q, r and defining relations $\rho q^i r + \rho q^j r = \rho q^{g(i, j)} r$ ($i, j \in N$), then, by the theorem of Murskii [3], there exists a semigroup $H_1 \supseteq H$ with a finite set of defining relations

$$\Sigma_2 (\rho, q, r, \bar{\theta}_3). \quad (5)$$

Again, by Lemma 1, there exists a K -algebra $S_4 \supseteq S_3 \times K[H_1]$ in which we can solve the system

$$\text{Hom}_4 (\rho, q, r, a, b, c, \bar{\theta}_4). \quad (6)$$

Finally, suppose $\alpha_1, \dots, \alpha_l \in K$ are generators of the ring or field K . Recall that a_1, \dots, a_k generate A . There exist $j_{ts} \in N$ such that

$$\alpha_t a b^{j_{ts}} c = a b^{j_{ts}} c \quad (t = 1, \dots, l; s = 1, \dots, k). \quad (7)$$

Now consider the K -algebra defined by the set of relations Σ^* , the union of all of the system (1)-(7). It is easy to see that this K -algebra on the elements $a_i = ab^i c$ ($i=1, 2, \dots$) satisfies all equalities of the K -algebra A and only these. This proves the theorem.

COROLLARY. If K is a commutative associative finitely generated ring with unity or a field that is finitely generated over its prime subfield, then there exists a 2-generator finitely presented associative K -algebra in which any associative K -algebra with a recursively enumerable set of defining relations can be isomorphically embedded. In particular, when $K = \mathbb{Z}$, there exists a 2-generator finitely presented associative ring A in which any associative ring with a recursively enumerable set of defining relations can be isomorphically embedded. Among the defining relations of A only one has the form $\omega_1 + \omega_2 = \omega_3$, and the others are word equalities in the generators.

The proof is analogous to that of Higman [4] for groups, owing to the fact that any countable K -algebra can be embedded in a 2-generator algebra (see [2]) and in the case $K = \mathbb{Z}$ the equalities (7) are superfluous.

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