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GRAPHS DRAWN WITH FEW CROSSINGS PER EDGE

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We show that if a graph of v vertices can be drawn in the plane so that every edge crosses at most $k>0$ others, then its number of edges cannot exceed $4.108\sqrt{k}v$. For $k \leq 4$, we establish a better bound, $(k+3)(v-2)$, which is tight for $k=1$ and 2. We apply these estimates to improve a result of Ajtai et al. and Leighton, providing a general lower bound for the crossing number of a graph in terms of its number of vertices and edges.

1. Introduction

Given a simple graph G, let $v(G)$ and $e(G)$ denote its number of vertices and edges, respectively. We say that G is *drawn* in the plane if its vertices are represented by distinct points of the plane and its edges are represented by Jordan arcs connecting the corresponding point pairs but not passing through any other vertex. Throughout this paper, we only consider *drawings* with the property that any two arcs have at most one point in common. This is either a common endpoint or a common interior point where the two arcs properly cross each other. We will not make any notational distinction between vertices of G and the corresponding points in the plane, or between edges of G and the corresponding Jordan arcs.

We address the following question. What is the maximum number of edges that a simple graph of v vertices can have if it can be drawn in the plane so that every edge crosses at most k others? For $k=0$, i.e. for *planar* graphs, the answer is $3v-6$. Our first theorem generalizes this result to $k \leq 4$. The case $k=1$ has been discovered independently by Bernd Gärtner, Torsten Thiele, and Günter Ziegler (personal communication).

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Theorem 1. *Let G be a simple graph drawn in the plane so that every edge is crossed by at most k others. If $0 \le k \le 4$, then we have*

$$e(G) \le (k+3)(v(G)-2).$$

For $k=0$, 1, 2, the above bound cannot be improved (see Remark 2.3 at the end of the next section.)

The *crossing number* $\mathrm{cr}(G)$ of a graph G is the minimum number of crossing pairs of edges, over all drawings of G in the plane.

Ajtai et al. [1] and, independently, Leighton [4] obtained a general lower bound for the crossing number of a graph, which found many applications in combinatorial geometry and in VLSI design (see [5], [6], [8]). Our next result, whose proof is based on Theorem 1, improves the bound of Ajtai et al. by roughly a factor of 2.

Theorem 2. *The crossing number of any simple graph G satisfies*

$$\mathrm{cr}(G) \ge \frac{1}{33.75}\frac{e^3(G)}{v^2(G)} - 0.9v(G) > 0.029\frac{e^3(G)}{v^2(G)} - 0.9v(G).$$

Theorem 3. *Let G be a simple graph drawn in the plane so that every edge is crossed by at most k others, for some $k \ge 1$. Then we have*

$$e(G) \le \sqrt{16.875k}\,v(G) \approx 4.108\sqrt{k}\,v(G).$$

Theorems 2 and 3 do not remain true if we replace the constants 0.029 and 4.108 by 0.06 and 1.92, respectively (see Remarks 3.2 and 3.3).

In the last section, we use the ideas of Székely [8] to deduce some consequences of Theorem 2.

2. Proof of Theorem 1

First we need a lemma for *multigraphs*, i.e., for graphs that may have multiple edges. In a *drawing* of a multigraph, any two non-disjoint edges either share only endpoints or have precisely one point in common, at which they properly cross.

Let M be a multigraph drawn in the plane so that every edge crosses at most k other edges. Let M' be a sub-multigraph of M with the largest number of edges such that in the drawing of M' (inherited from the drawing of M), no two edges cross each other. We say that M' is a *maximum plane sub-multigraph* of M, and its faces will be denoted by $\Phi_1, \Phi_2, \dots, \Phi_m$. Let $|\Phi_i|$ denote the number of edges of M' along the boundary of Φ_i, where every edge whose both sides belong to the interior of Φ_i is counted twice. It follows from the maximality of M' that every edge e of M which does not belong to M' (in short $e \in M - M'$) crosses at least one edge of M'. The closed portion between an endpoint of e and the nearest crossing of e with an edge of M' is called a *half-edge*. Thus, every edge of $M-M'$ contains two half-edges. Every half-edge lies in a face Φ and intersects at most $k-1$ other half-edges and an edge of Φ_i (not counting the incidences at the vertices of M). Let $h(\Phi_i)$ denote the number of half-edges in Φ_i.

Lemma 2.1. *Let $0 \le k \le 4$ and let M be a multigraph drawn in the plane so that every edge crosses at most k others. Let M' be a maximum plane sub-multigraph of M, and let Φ denote a face with $|\Phi| = s \ge 3$ sides in M', whose boundary is connected.*

Then the number of half-edges in Φ satisfies

$$h(\Phi) \le (s-2)(k+1) - 1.$$

Proof. We proceed by induction on s. First, let $s=3$ and denote the vertices of Φ by A, B, and C. Let a, b, and c denote the number of half-edges in Φ emanating from A, B, and C, respectively. We have to show that $a+b+c$, the total number of half-edges in Φ, is at most k. For $k=0$, there is nothing to prove. We check the cases $k=1,2,3,4$, separately.

• $\boldsymbol{k=1}$: If $a=b=c=0$, we are done. Assume without loss of generality that $a \ge 1$. But then $a=1$, because all half-edges in Φ emanating from A intersect the edge BC. Since any half-edge in Φ emanating from B or C would create another intersection on the half-edge starting from A, we obtain $b=c=0$. Hence, $a+b+c=1$.

• $\boldsymbol{k=2}$: Suppose without loss of generality that $a \ge 1$. Clearly, $a \le 2$. If $a=1$, the unique half-edge in Φ emanating from A intersects all half-edges coming from B and C. So $1+b+c=a+b+c \le 2$. If $a=2$, any half-edge from B would intersect both half-edges emanating from A and the edge AC, which is impossible. Hence, $b=0$. Similarly, $c=0$, and $a+b+c=2$.

• $\boldsymbol{k=3}$: Just like before, we can exclude all cases when $a+b+c>3$, except for the case $a=b=2$ and $c=0$. Now let e_1 and e_2 denote the edges containing the two half-edges in Φ emanating from A. Both of them intersect the two half-edges starting from B and the edge BC. So they cannot cross any other edge. Removing BC from M' and adding e_1 and e_2, we would obtain a larger plane sub-multigraph of M, contradicting the maximality of M'.

• $\boldsymbol{k=4}$: We can again exclude all cases when $a+b+c>4$, with the exception of the case $a=2$, $b=3$, $c=0$. As before, let e_1 and e_2 denote the edges containing the two half-edges in Φ emanating from A. Now both e_1 and e_2 are intersected by the three half-edges emanating from B and by the edge BC. Hence, there are no other edges crossing them, and the number of edges of M' can be increased by replacing BC with e_1 and e_2. Contradiction.

Now let $s>3$, and suppose that the lemma has already been proved for faces with fewer than s sides. Let $A_1, A_2, \ldots, A_s$ denote the sequence of vertices of Φ, listed in clockwise order. In this sequence, the same vertex may occur several times (as many times as it is visited during a full clockwise tour around the boundary of Φ). For simplicity, let $A_0 = A_s$ and $A_{s+1} = A_1$.

We call an open arc *empty* if it does not intersect any half-edge in Φ.

Case 1. Assume that there is a half-edge $e=A_iE$ in Φ, where E is an interior point of the side A_jA_{j+1}, and either

(i) the arc $A_jE\subseteq A_jA_{j+1}$ is empty and $i\neq j-1$, or

(ii) the arc $EA_{j+1}\subseteq A_jA_{j+1}$ is empty and $i\neq j+2$.

By symmetry, we can suppose that e satisfies (i).

Apply to M the following transformations.

(1) Delete all edges of M except the edges belonging to the boundary of Φ.

(2) Add all half-edges lying in Φ.

(3) Without introducing any crossing, replace each side f of Φ, which is encountered twice when we trace the boundary of Φ, by two parallel edges, f_1 and f_2, running very close to f. Every half-edge originally ending at f will now have an endpoint either on f_1 or on f_2. The resulting cell will be denoted by $\overline{\Phi}$. Clearly, we have $|\overline{\Phi}|=|\Phi|$.

(4) Extend every half-edge g ending at a side f of the boundary of $\overline{\Phi}$ to an edge, g', by adding a short arc and a new vertex beyond f, outside $\overline{\Phi}$. Add some pairwise disjoint short edges outside $\overline{\Phi}$, which cross only the extra arc, so that g' determines precisely k crossings.

The resulting multigraph drawing, $\overline{M}$, obviously satisfies the following properties.

(a) $\overline{\Phi}$ is a cell of some maximum plane sub-multigraph $\overline{M}'\subseteq\overline{M}$.

(b) The number of half-edges of $\overline{M}$ lying in $\overline{\Phi}$ is the same as the number of half-edges of M lying in Φ.

(c) Every edge of $\overline{M}$ crosses at most k other edges.

By (b), it is sufficient to bound $h(\overline{\Phi})$, the number of half-edges of $\overline{M}$ within $\overline{\Phi}$.

For simplicity, let e also denote the half-edge of $\overline{M}$, which corresponds to the half-edge $e=A_iE$ of M. Let $E_1,E_2,\dots,E_l$ denote the intersection points of e with the half-edges $e_1,e_2,\dots,e_l$ (resp.) in $\overline{M}$, emanating from A_j, such that E_1 is nearest to A_i.

Let $\overline{M}_e$ denote the multigraph drawing obtained from $\overline{M}$ by replacing the edge of $\overline{M}$ containing e with a new edge $e'=A_iA_j$ running very close to e between A_i and E_1 and very close to e_1 between E_1 and A_j. (e' must not cross e_1 and the boundary of $\overline{\Phi}$. If $l=0$, set $E_1=E$ and $e_1=A_jA_{j+1}$.)

Since A_jE was empty in M, every edge of $\overline{M}$, which crosses e', also crosses e. Thus, $\overline{M}_e$ also satisfies the condition that each of its edges crosses at most k others. Clearly, $\overline{M}'\cup e'$ is a maximum plane sub-multigraph of $\overline{M}_e$, in which e' divides $\overline{\Phi}$

into two faces $\overline{\Phi}'$ and $\overline{\Phi}''$ with s' and s'' sides, respectively, where $3 \le s', s'' < s$, $s'+s''=s+2$.

With the exception of e, every half-edge of $\overline{M}$, which lies in $\overline{\Phi}$, corresponds to a half-edge of $\overline{M}_e$ lying either in $\overline{\Phi}'$ or in $\overline{\Phi}''$. By the induction hypothesis,

$$h(\Phi) = h(\overline{\Phi})$$
$$= h(\overline{\Phi}')+h(\overline{\Phi}'')+1 \le (s'-2)(k+1)-1+(s''-2)(k+1)-1+1 = (s-2)(k+1)-1,$$

as required.

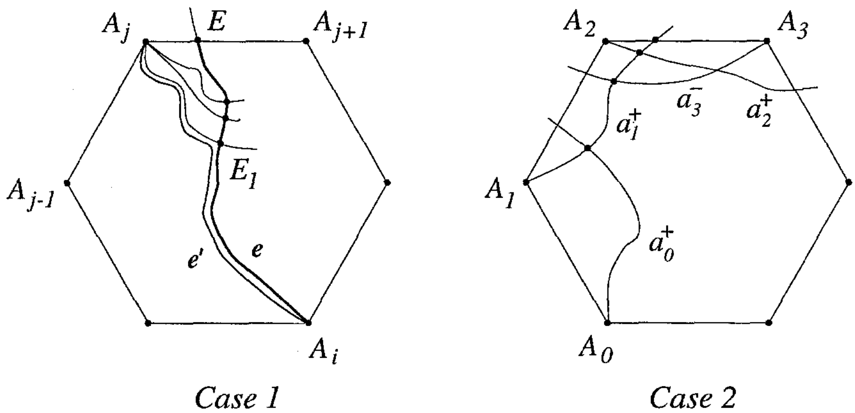

Figure 1

Case 2. Assume that there is no half-edge in Φ that satisfies the condition of Case 1.

Then, for any non-empty side A_iA_{i+1} of Φ, the half-edge a_{i-1}^+ (resp. a_{i+2}^-) whose intersection with A_iA_{i+1} is closest to A_i (resp. closest to A_{i+1}) starts at the vertex A_{i-1} (resp. A_{i+2}).

Since any side of Φ intersects at most k half-edges, if there are two empty sides of Φ, then $h(\Phi) \le (s-2)k \le (s-2)(k+1)-1$. So we can suppose that Φ has at most one empty side. Since $s>3$, there are three consecutive non-empty sides, say, A_1A_2, A_2A_3, and A_3A_4.

Then a_1^+ must intersect a_0^+, a_2^+, a_3^-, and the side A_2A_3. Similarly, a_4^- must intersect a_5^-, a_3^+, a_2^+, and the side A_2A_3. This is clearly impossible if $k=1,2$ or 3.

For $k = 4$, let e_1 and e_2 denote the edges of M containing a_1^+ and a_4^-, respectively. Both of these edges cross three half-edges and the side A_2A_3 of Φ, so neither of them can cross any further edges. Removing the edge A_2A_3 from M' and adding e_1 and e_2, we would obtain a plane sub-multigraph of M, whose number of edges is larger than the number of edges of M'. This contradicts the maximality of M', completing the proof of Lemma 2.1. ∎

For any face Φ with at least 3 sides, let $t(\Phi)$ denote the number of triangles in a triangulation of Φ.

Lemma 2.2. *Let Φ be any face of M' with $|\Phi| \geq 3$ sides. Then the number of half-edges of Φ satisfies*

$$h(\Phi) \leq t(\Phi)k + |\Phi| - 3.$$

Proof. If the boundary of Φ is connected, then $t(\Phi) = |\Phi| - 2$. Hence, by Lemma 2.1, $h(\Phi) \leq (|\Phi| - 2)(k+1) - 1 = t(\Phi)k + |\Phi| - 3$.

For any face Φ, the number of half-edges in Φ is at most $|\Phi|k$, because every side of Φ intersects at most k half-edges. If the boundary of Φ is not connected, then $t(\Phi) \geq |\Phi|$. Therefore, in this case, we have $h(\Phi) \leq |\Phi|k \leq t(\Phi)k + |\Phi| - 3$. ∎

Now we are ready to prove Theorem 1. Suppose that a simple graph G is drawn in the plane with at most k crossings on each edge. Let G' be a maximum plane subgraph of G. Denote the faces of G' by $\Phi_1, \Phi_2, \ldots \Phi_m$. To triangulate Φ_i, we need at least $|\Phi_i| - 3$ edges. Therefore,

$$e(G') \leq 3v - 6 - \sum_{i=1}^{m} (|\Phi_i| - 3).$$

Every edge of $G - G'$ gives rise to two half-edges. So, Lemma 2.2 yields that

$$e(G - G') \leq \frac{1}{2} \sum_{i=1}^{m} \left(t(\Phi_i)k + |\Phi_i| - 3\right).$$

Summing up the last two inequalities and noticing that the total number of triangles satisfies $\sum_i t(\Phi_i) = 2v(G) - 4$, we obtain

$$\begin{aligned} e(G) &\leq 3v(G) - 6 + \frac{1}{2} \sum_{i=1}^{m} \left(t(\Phi_i)k - (|\Phi_i| - 3)\right) \\ &\leq 3v(G) - 6 + (v(G) - 2)k = (k+3)(v(G) - 2), \end{aligned}$$

which completes the proof of Theorem 1. ∎

Remark 2.3. For $k=0$, the bound $e \leq 3v - 6$ is tight for any triangulation.

For $k = 1$, the bound $e \leq 4v - 8$ obtained, is also tight, provided that $v \geq 12$. First we show that for every $v \geq 12$ there is a planar graph with v vertices, all of whose faces are quadrilaterals, and no two faces share more than one edge. Indeed, Figure 2 illustrates that such graphs exist for $v = 8, 13, 14, 15$. Once we have an example G with v vertices, we can construct another one with $v + 4$ vertices, by replacing some face of G by the 8-point example in Figure 2. Notice that if we add both diagonals of each face (including the external face), then we obtain a graph with $4v - 8$ edges such that along each edge there is at most one crossing.

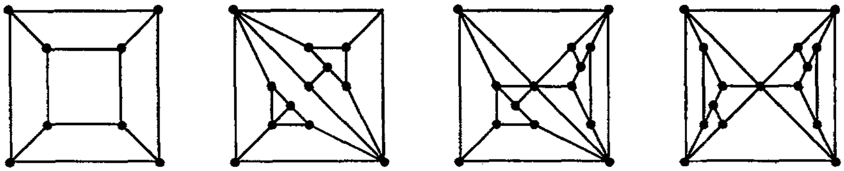

Figure 2

For $k = 2$, the bound $e \leq 5v - 10$ is sharp for all $v \geq 50$ such that $v \equiv 2 \pmod 3$. For simplicity, we only exhibit a construction for $v \equiv 5 \pmod{15}$. First we construct a planar graph whose faces are pentagons and two faces have at most one edge in common. For $v = 20$, such a graph is shown in Figure 3. The number of vertices of such an example G can be increased by 15, by replacing some face of G with the graph depicted in Figure 3. Notice that if we add all 5 diagonals of each face to G, then we obtain a graph with v vertices and $5v - 10$ edges, in which every edge crosses at most 2 others.

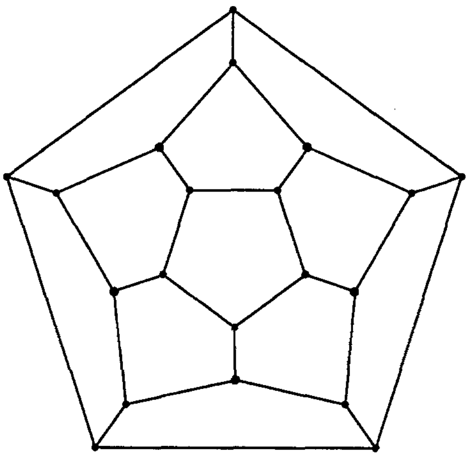

Figure 3

3. Proofs of Theorems 2 and 3

In this section, we slightly improve the best known general lower bound on the crossing number of a graph, due to Ajtai et al. [1] and Leighton [4]. Our proof is based on the following consequence of Theorem 1.

Corollary 3.1. *The crossing number of any simple graph G with at least 3 vertices satisfies*

$$cr(G) \geq 5e(G) - 25v(G) + 50.$$

Proof. If $e(G) \le 3v(G) - 6$, then the statement is void. Assume $e(G) > 3v(G) - 6$.

It follows from Theorem 1 that if $e(G) > (k+3)(v(G)-2)$, then G has an edge crossed by at least $k+1$ other edges ($k \le 4$). Deleting such an edge, we obtain by induction on $e(G)$ that the number of crossings is at least

$$\sum_{k=0}^{4} [e(G) - (k+3)(v(G)-2)] = 5e(G) - 25v(G) + 50.$$ ∎

Proof of Theorem 2. Let G be a simple graph drawn in the plane with $\mathrm{cr}(G)$ crossings, and suppose that $e(G) \ge 7.5v(G)$.

Construct a *random* subgraph $G' \subseteq G$ by selecting each vertex of G independently with probability $p = 7.5v(G)/e(G) \le 1$, and letting G' be the subgraph induced by the selected vertices. The expected number of vertices of G', $E[v(G')] = pv(G)$. Similarly, $E[e(G')] = p^2 e(G)$. The expected number of crossings in the drawing of G' inherited from G is $p^4 \mathrm{cr}(G)$, and the expected value of the crossing number of G' is even smaller.

By Corollary 3.1, $\mathrm{cr}(G') \ge 5e(G') - 25v(G')$ for every G'. Taking expectations,

$$p^4 \mathrm{cr}(G) \ge E[\mathrm{cr}(G')] \ge 5E[e(G')] - 25E[v(G')] = 5p^2 e(G) - 25pv(G).$$

This implies that

$$\mathrm{cr}(G) \ge \frac{1}{33.75} \frac{e^3(G)}{v^2(G)}, \tag{1}$$

whenever $e(G) \ge 7.5v(G)$. In fact, using Corollary 3.1 in the range $e(G) < 7.5v(G)$, it is easy to check that the slightly weaker inequality

$$\mathrm{cr}(G) \ge \frac{1}{33.75} \frac{e^3(G)}{v^2(G)} - 0.9v(G)$$

is valid for every simple graph G. ∎

Proof of Theorem 3. For $k \le 4$, the result is weaker than the bounds given in Theorem 1.

So let $k > 4$, and consider a drawing of G such that every edge crosses at most k others. Let C denote the number of crossings in this drawing. If $e(G) < 7.5v(G)$, then there is nothing to prove. If $e(G) \ge 7.5v(G)$, then using the stronger form (1) of Theorem 2, we obtain

$$\frac{1}{33.75} \frac{e^3(G)}{v^2(G)} \le \mathrm{cr}(G) \le C \le \frac{e(G)k}{2}.$$

Consequently,

$$e(G) \le \sqrt{16.875}\sqrt{k}v(G).$$ ∎

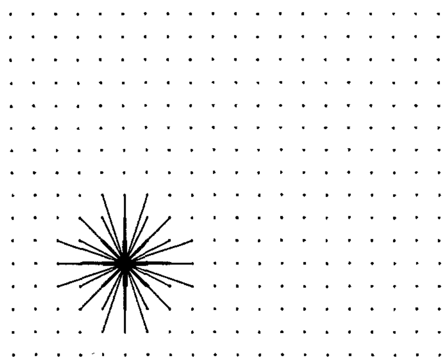

Figure 4

Remark 3.2. The bound given in Theorem 2 is asymptotically tight, apart from the values of the constants. The best construction we found is the following.

Let $v \ll e \ll v^2$. Let $V(G)$ be a set of v points arranged in a slightly perturbed unit square grid of size $\sqrt{v} \times \sqrt{v}$, so that the points are in general position. Let $d = \sqrt{2e/\pi v}$, so that $d^2\pi = 2e/v$.

Connect two points by a straight-line segment if and only if their distance is at most d. Then $v(G) = v$, $e(G) \approx vd^2\pi/2 = e$.

To count the number of crossings in G, let $S(a) = \{(x,y) | 1 \le x, y \le a\}$, and for any two segments $(u_1,u_2),(v_1,v_2)$, $(u_1,u_2) \otimes (v_1,v_2)$ means that the two segments cross each other. Then the number of crossings in G is

$$= \frac{1}{8}\left|\left\{(u_1,u_2,v_1,v_2) \in [V(G)]^4 \mid \|u_1-u_2\|, \|v_1-v_2\| \le d,\ (u_1,u_2)\otimes(v_1,v_2)\right\}\right|$$

$$\approx \frac{1}{8} \int\limits_{u_1 \in S(\sqrt{v})} \int\limits_{\substack{u_2 \in S(\sqrt{v}) \\ \|u_1-u_2\| \le d}} \int\limits_{v_1 \in S(\sqrt{v})} \int\limits_{\substack{v_2 \in S(\sqrt{v}) \\ \|v_1-v_2\| \le d \\ (u_1,u_2)\otimes(v_1,v_2)}} 1 \quad dv_2 dv_1 du_2 du_1 = \frac{2\pi}{27} vd^6 \left(1+o(1)\right).$$

Thus,

$$\operatorname{cr}(G) \le \frac{2\pi}{27} vd^6(1+o(1)) \approx \frac{16}{27\pi^2}\frac{e^3}{v^2} \approx .06\frac{e^3}{v^2}.$$

J. Spencer [7] showed that the limit

$$c = \lim \frac{v^2}{e^3} \min_{|V(G)|=v, |E(G)|=e} \operatorname{cr}(G)$$

exists, as $v \to \infty$ and $v \ll e \ll v^2$. By our results, $.06 \ge c \ge .029$.

Remark 3.3. The bound obtained in Theorem 3 is also asymptotically tight. Consider the same construction as in Remark 3.2, but now set $d = \sqrt[4]{3k/2}(1-o(1))$, as k tends to infinity. Just like above, it can be shown that no edge crosses more than k other edges. The number of edges

$$e(G) = v(G)\frac{d^2\pi}{2}(1-o(1)) = v(G)\sqrt{k}\frac{\sqrt{3}\pi}{\sqrt{8}}.$$

Thus, we have

$$1.92(1-o(1))\sqrt{k}v(G) < \max e(G) < 4.108\sqrt{k}v(G),$$

where the maximum is taken over all simple graphs with $v(G)$ vertices that have a drawing with at most k crossings per edge.

4. Three further applications

Using Székely's method (see [8]) and Theorem 2, we can improve the coefficient of the main term in the Szemerédi–Trotter theorem [9], [2].

Theorem 4.1. *Given m points and n lines in the Euclidean plane, the number of incidences between them is at most $2.57m^{2/3}n^{2/3}+m+n$.*

Proof. We can assume that every line and every point is involved in at least one incidence, and that $n \geq m$, by duality. Since the statement is clearly true for $m=1$ we have to check it only for $m \geq 2$. Define a graph G drawn in the plane such that the vertex set of G is the given set of m points, and join two points with an edge drawn as a straight line segment if the two points are consecutive along one of the lines. Let I denote the total number of incidences between the given m points and n lines. Then $v(G)=m$ and $e(G)=I-n$. Since every edge belongs to one of the n lines, $\mathrm{cr}(G) \leq \binom{n}{2}$. Applying Theorem 2 to G, we obtain that

$$\frac{1}{33.75}\frac{(I-n)^3}{m^2} - 0.9m \leq \mathrm{cr}(G) \leq \binom{n}{2}.$$

Using that $n \geq m \geq 2$, some calculation shows that

$$I - n \leq \sqrt[3]{16.875m^2n^2 + 30.375m^3} \leq \sqrt[3]{16.875}n^{2/3}m^{2/3} + m.$$

Therefore,

$$I \leq 2.57m^{2/3}n^{2/3} + m + n. \qquad \blacksquare$$

Remark 4.2. As Erdős pointed out fifty years ago, the order of magnitude of the bound in Theorem 4.1 cannot be improved. To see this, one can take n points arranged in a unit square grid of size $\sqrt{n}\times\sqrt{n}$ and consider the m most "populous" lines.

More precisely, for any fixed $1>\varepsilon>0$, take all lines which contain at least $\varepsilon\sqrt{n}$ of the points. Then, for the number of lines m we have

$$m \approx 4\sqrt{n}\sum_{r=1}^{1/\varepsilon}\sum_{\substack{s<r\\(r,s)=1}}(r+s-2rs\varepsilon) = 6\sqrt{n}\sum_{r=1}^{1/\varepsilon} r\phi(r) - 4\sqrt{n}\varepsilon\sum_{r=1}^{1/\varepsilon} r^2\phi(r) \approx \frac{6\sqrt{n}}{\pi^2\varepsilon^3}.$$

Here $\phi(n)$ denotes Euler's function and we used the formula $\sum_{r=1}^{N}\phi(r)\approx 3N^2/\pi^2$ (see e. g. [3]). By similar calculations, for the number of incidences I we get

$$I \approx 4n\sum_{r=1}^{1/\varepsilon}\sum_{\substack{s<r\\(r,s)=1}}(1-rs\varepsilon^2) = 4n\sum_{r=1}^{1/\varepsilon}\phi(r) - 2n\varepsilon^2\sum_{r=1}^{1/\varepsilon} r^2\phi(r) \approx \frac{3n}{\pi^2\varepsilon^2}.$$

Comparing the last two expressions, we obtain

$$I \approx cn^{2/3}m^{2/3} \qquad \text{with} \qquad c=\sqrt[3]{\frac{3}{4\pi^2}}\approx 0.42.$$

We can also generalize Theorem 2 for multigraphs with bounded edge-multiplicity, improving the constant in Székely's result [8].

Theorem 4.3. *Let G be a multigraph with maximum edge-multiplicity m. Then*

$$cr(G) \geq \frac{1}{33.75}\frac{e^3(G)}{mv^2(G)} - 0.9m^2v(G).$$

Proof. Define a random simple subgraph G' of G as follows. For each pair of vertices v_1, v_2 of G, let $e_1, e_2, \dots e_k$ be the edges connecting them. With probability $1-k/m$, G' will not contain any edge between v_1 and v_2. With probability k/m, G' contains precisely one such edge, and the probability that this edge is e_i is $1/m$ $(1\leq i\leq k)$.

Applying Theorem 2 to G' and taking expectations, the result follows. ∎

Remark 4.4. Let G be a graph drawn in the plane. A subset of the edges of G is said to form a *plane subgraph* of G, if it contains no two crossing edges. Let F_n denote the maximum number of plane subgraphs of a graph with n vertices.

Ajtai et al. [1] used their general lower bound for crossing numbers to obtain that $F_n\leq 10^{13n}$.

Using our results, we can improve this estimate. Let $F_n(m)$ stand for the maximum number of plane subgraphs that a graph of n vertices and m edges drawn in the plane can have. By (1), if $m\geq 7.5n$, then there is an edge e of G crossing at least $\frac{m^2}{16.875n^2}$ other edges. The number of plane subgraphs that do not contain e, is at most $F_n(m-1)$, while the number of those which do contain e cannot exceed $F_n(\lfloor m-\frac{m^2}{16.875n^2}\rfloor)$. Therefore,

$$F_n(m) \le F_n(m-1) + F_n\left(\left\lfloor m - \frac{m^2}{16.875n^2}\right\rfloor\right) \quad \text{if} \quad m \ge 7.5n.$$

Similarly, by Theorem 1 we have

$$\begin{aligned} F_n(m) &\le F_n(m-1) + F_n(m-6) && \text{if} \quad m \ge 7n, \\ F_n(m) &\le F_n(m-1) + F_n(m-5) && \text{if} \quad m \ge 6n, \\ F_n(m) &\le F_n(m-1) + F_n(m-4) && \text{if} \quad m \ge 5n, \\ F_n(m) &\le F_n(m-1) + F_n(m-3) && \text{if} \quad m \ge 4n, \\ F_n(m) &\le F_n(m-1) + F_n(m-2) && \text{if} \quad m \ge 3n. \end{aligned}$$

Clearly,

$$F_n(m) \le 2^{3n} \qquad \text{if} \quad m \le 3n.$$

Using these inequalities, we get that

$$F_n = F\left(\binom{n}{2}\right) \le 53000^n < 10^{4.73n}.$$

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