

The Undecidability of the First-order Theory of Diagonalizable Algebras

Abstract. The undecidability of the first-order theory of diagonalizable algebras is shown here.

Introduction

A *diagonalizable algebra* is a pair $\langle \mathcal{A}, \tau \rangle$, where \mathcal{A} is a Boolean algebra, and τ is a unary operation on \mathcal{A} satisfying the following properties:

$$\tau 1 = 1, \quad \tau(x \cdot y) = \tau x \cdot \tau y, \quad \tau(\tau x \rightarrow x) = \tau x.^1$$

The diagonalizable algebras constitute an equational class, which we call *DA*. An important aspect of the diagonalizable algebras is the following:

Let $\varphi \equiv \{p_n\}_{n \in \omega}$ be a sequence of sentences of Peano Arithmetic *PA*; let us define a mapping φ^* from the set of the polynomials of *DA* into the set of *PA* sentences in the following inductive manner:

$$\begin{aligned} \varphi^* 0 &= \ulcorner 0 = 1 \urcorner; & \varphi^* 1 &= \ulcorner 0 = 0 \urcorner; & \varphi^* x_i &= p_i; & \varphi^*(f+g) &= \varphi^* f \vee \varphi^* g; \\ \varphi^*(vf) &= \neg \varphi^* f; & \varphi^* \tau f &= \text{Theor } \overline{\varphi^* f}. \end{aligned}$$

Then, if $f = 1$ is valid in *DA*, for every φ as above, $\varphi^* f$ is a theorem of *PA*. (See R. MAGARI, [3]).

In [6], R. SOLOVAY shows that, conversely, if, for every φ , $\varphi^* f$ is a theorem of *PA*, then $f = 1$ is valid in *DA*.

In [5], F. MONTAGNA extends MAGARI's result to the first-order theory \mathcal{T} of the diagonalizable algebras. More precisely, the author defines, for every φ as above, a mapping φ° from the set of the formulas of \mathcal{T} into the set of the sentences of *PA*, and shows that, if A is a theorem of \mathcal{T} , then, for every φ , $\varphi^\circ A$ is a theorem of *PA*.

It is known (see C. BERNARDI, [1]) that *DA* (that, is, the set of the equations which are valid in all diagonalizable algebras) is decidable, therefore, the set of the polynomials f such that, for every φ , $\varphi^* f$ is a theorem of *DA*, is decidable.

In this paper, we show that, on the contrary, the first-order theory \mathcal{T} of diagonalizable algebras is undecidable. The decision problem remains

¹ In the following, $+$, \cdot , ν denote the operations of join, meet, and complementation; $x \rightarrow y$ is an abbreviation for $\nu x + y$.

open for the set of the formulas A of \mathcal{T} such that, for every φ , $\varphi^\circ A$ is a theorem of PA .

By [7], in order to show the claim it suffices to find a model $\mathcal{A} = \langle A, \tau \rangle$ of DA , and formulas $F(x)$, $\psi(x, y)$, $\vartheta_0(x, y, z)$, $\vartheta_1(x, y, z)$ of the signature of \mathcal{A} , so that, putting $\mathcal{F} = \{x \in A : \vDash_{\mathcal{A}} F(x)\}$, $\approx = \{\langle x, y \rangle \in A^2 : \vDash_{\mathcal{A}} \psi(x, y)\}$, $\dagger = \{\langle x, y, z \rangle \in A^3 : \vDash_{\mathcal{A}} \vartheta_0(x, y, z)\}$ and $\times = \{\langle x, y, z \rangle \in A^3 : \vDash_{\mathcal{A}} \vartheta_1(x, y, z)\}$, the following conditions hold:

1) \approx is a congruence relation in the structure $\mathfrak{F} = \langle \mathcal{F}, \dagger, \times \rangle$.

2) The quotient \mathfrak{F}/\approx is isomorphic to the standard model $\mathfrak{N} = \langle N, +, \cdot \rangle$ of the natural numbers.

Now, let us define \mathcal{A} as follows:

Let X be the set $\omega \cup \{\langle \{i, j\}, n \rangle : i, j, n \in \omega, i \neq j, i, j < n\}$. Let R be the binary relation on X defined by: xRy iff $\exists i, j, n : i \neq j; i, j < n, x = \langle \{i, j\}, n \rangle$ and $y = i$ or $y = j$ or $y > n$. Let us define the mapping τ from $\mathcal{P}(X)$ into $\mathcal{P}(X)$ by $\tau A = \{x \in X : \forall z \in X, \text{ if } xRz, \text{ then } z \in A\}$.

Now, let \mathcal{A} be the algebra $\langle \mathcal{P}(X), \cup, \cap, \mathcal{C}, X, \emptyset, \tau \rangle$. Since R is a transitive and reverse well-founded relation on X , \mathcal{A} is a diagonalizable algebra (see R. MAGARI, [4]). It is easily seen that $\tau \emptyset = \omega$, $\tau^2 \emptyset = X$. Now, we define

$$Ix = "x \cdot \tau x"$$

$$Op x = "Ix = x"$$

$$Cn x = "\forall u \forall v [Op u \wedge Op v \wedge u + v \geq x \wedge u \cdot v \cdot x = 0 \rightarrow u \cdot x = 0 \vee v \cdot x = 0]"$$

$$At x = "\neg x = 0 \wedge \forall u (u < x \rightarrow u = 0)"$$

$$Cmp(x, y) = "Cn x \wedge x \leq y \wedge \forall u [Cn u \wedge u \leq y \rightarrow \neg x < u]."$$

It is easily seen that the set $\{B \subseteq X : \vDash_{\mathcal{A}} Op B\}$ constitutes a topology on X . We denote this topology by T . Of course, for every $B, C \subseteq X$, $\vDash_{\mathcal{A}} Cn B$ iff B is connected, $\vDash_{\mathcal{A}} Cmp(B, C)$ iff B is a connected component of C ; finally $B \subseteq X$ is open iff for each $x, y \in X$, if $x \in B$, and xRy , then $y \in B$.

DEFINITION 1. Let B be a non-empty subset of X . We define T_{1B} = $\{C \subseteq B : \exists D \subseteq X [Op D \wedge C = X \cap D]\}$.

Of course, T_{1B} is a topology on B ; moreover, it is easily seen that for every $C \subseteq B$, C is open in T_{1B} iff, for every $x, y \in B$, if $x \in C$ and xRy , then $y \in C$.

DEFINITION 2. Let Y be a subset of X , and let a, b be elements of Y ; we define $a R_Y b$ iff there is a finite sequence $x_0 \dots x_n$ of elements of Y such that $x_0 \equiv a$, $x_n \equiv b$, and, for every i , either $x_i R x_{i+1}$ or $x_{i+1} R x_i$.

R_Y is an equivalence relation on Y .

LEMMA 1. Let Y be a subset of X . Then, Y is connected iff for every $a, b \in Y$ $a R_Y b$.

PROOF. Suppose that Y is connected; let x be an arbitrary element of Y , and let \bar{Y} be the set $\{y \in Y: xR_y y\}$. Then, \bar{Y} is open in the topology T_{1Y} . Moreover, $Y - \bar{Y}$ is also open in T_{1Y} ; in fact, if $a \in Y - \bar{Y}$, $b \in Y$ and aRb , then $b \in \bar{Y}$. Otherwise, we would have $xR_Y b$, and then $xR_Y a$, $a \in \bar{Y}$, a contradiction. Since $(Y - \bar{Y}) \cup \bar{Y} = Y$ and \bar{Y} is non-empty, we can conclude that $Y - \bar{Y}$ is empty and that $Y = \bar{Y}$. Therefore, for every $y \in \bar{Y}$, $xR_Y y$.

On the other hand, suppose that, for every $x, y \in Y$ $xR_Y y$. Let A, B be non-empty subsets of Y such that $A \cup B = Y$; suppose that both A and B are open in T_{1Y} . Let $a \in A$, $b \in B$; since $a R_{\bar{Y}} b$, there is a sequence $x_0 \dots x_n$ of elements of Y such that $x_0 = a$, $x_n = b$, and, for every i , either $x_i R x_{i+1}$ or $x_{i+1} R x_i$. Let x_{i_0+1} be the first element of the sequence, which belongs to B ; then $x_{i_0} \in A$; since B is closed under R , if $x_{i_0+1} R x_{i_0}$ $x_{i_0} \in A \cap B$; since A is also closed under R , if $x_{i_0} R x_{i_0+1}$, $x_{i_0+1} \in A \cap B$; in both cases, $A \cap B$ is non-empty. Then, Y is connected. Q.E.D.

By Lemma 1, we can easily deduce that, if $Y \subseteq X$, the connected components of Y are exactly the equivalence classes determined by the relation R_Y .

DEFINITION 3. We define $\sigma x = \tau \nu x$; $Fx = x \leq \tau 0 \wedge \neg \sigma x = \sigma 1$;

$$\mathcal{F} = \{x \in A: \vdash_{\mathcal{A}} Fx\}.$$

LEMMA 2. Let Y be a subset of X ; then $\mathcal{F}Y$ is true iff Y is a finite subset of ω .

PROOF. We have $\sigma Y = \{x \in X: \exists y \in Y \ xRy\}$, $\sigma X = \mathcal{C}\tau 0 = \{\langle\{i, j\}, n\rangle: i, j, n \in \omega, i \neq j, i, j < n\}$. Suppose that Y is a finite subset of ω ; then $Y \subseteq \omega = \tau 0$; moreover, $\sigma Y \neq \sigma 1$; in fact, let $n = \max Y$, and let $z = \langle\{n+1, n+2\}, n+3\rangle$; by the definition of R , $zR x$ iff $x \geq n+1$; therefore, for every $x \in Y$, zRy ; so $z \in \sigma 1 - \sigma Y$, and $\mathcal{F}Y$ is true.

On the other hand, if $\mathcal{F}Y$ is true, then $Y \subseteq \tau 0 = \omega$; moreover Y is a finite set; indeed suppose, by contradiction, that Y is infinite; let $\langle\{i, j\}, n\rangle$ be an arbitrary element of $\sigma 1$. Since Y is infinite, there is an $m \in Y$, such that $n < m$; then $\langle\{i, j\}, n\rangle R m$ therefore $\langle\{i, j\}, n\rangle \in \sigma Y$. So, we would conclude $\sigma Y = \sigma 1$, which contradicts our hypothesis. Q.E.D.

DEFINITION 4. We define $\psi(u, v)$ to be the formula $Fu \wedge Fv \wedge [u = v \vee \exists z \forall w (Cmp(w, z) \rightarrow At(u \cdot w) \wedge At(v \cdot w))]$; moreover, we put $\approx = \{\langle x, y \rangle \in \mathcal{A}^2: \vdash_{\mathcal{A}} \psi(x, y)\}$.

LEMMA 3. For every $x, y \in \mathcal{A}$, we have $x \approx y$ iff x, y are finite subsets of ω , and $\bar{x} = \bar{y}$.

PROOF. If $A \approx B$, then both A and B are finite subsets of ω : moreover, either $A = B$, or there is a $Z \subseteq X$ such that, for every connected compo-

ment W of Z , both $W \cap (A - B)$ and $W \cap (B - A)$ have exactly one element; if $A = B$, then $\overline{A} = \overline{B}$; otherwise, let us consider an arbitrary $a \in A$; there is a unique connected component W_a of Z such that $W_a \cap (A - B) = \{a\}$; furthermore $W_a \cap (B - A)$ has exactly one element, say b_a ; now let us define for every $a \in A - B$, $fa = b_a$; it is easily seen that f is a one-to-one function from $A - B$ onto $B - A$; then $\overline{A - B} = \overline{B - A}$, therefore $\overline{A} = \overline{B}$.

Conversely, let A and B be finite subsets of ω such that $\overline{A} = \overline{B}$; if $A = B$, then the claim is obvious; otherwise, let $A - B = \{a_1 \dots a_n\}$, $B - A = \{b_1 \dots b_n\}$, and let $h = \max\{a_1 \dots a_n, b_1 \dots b_n\}$. Let Z be the set $\{a_1 \dots a_n, b_1 \dots b_n, (\{a_1, b_1\}, h+1) \dots (\{a_n, b_n\}, h+1)\}$. By Lemma 1, it is easily seen that the connected component of Z are $W_1 = \{a_1, b_1, (\{a_1, b_1\}, h+1)\}$, ..., $W_n = \{a_n, b_n, (\{a_n, b_n\}, h+1)\}$, therefore for every W_i , $W_i \cap (A - B)$ and $W_i \cap (B - A)$ have exactly one element. Therefore $A \approx B$. Q.E.D.

DEFINITION 5. We define $\vartheta_0(x, y, z) = "Fu \wedge Fy \wedge Fz \wedge \exists u (Fu \wedge \psi(u, y) \wedge u \cdot x = 0 \wedge \psi(z, x+u))"$ and $\dagger = \{\langle x, y, z \rangle \in \mathcal{S}^3 : \models_{\mathcal{S}} \vartheta_0(x, y, z)\}$.

It is easily seen that, if x, y, z are finite subsets of ω , then we have $\dagger(x, y, z)$ iff $\overline{x} + \overline{y} = \overline{z}$.

DEFINITION 6. We define $\vartheta_1(u, v, w) = "Fu \wedge Fv \wedge Fw \wedge \exists x \exists y [\psi(x, u) \wedge \psi(y, \tau 0, w) \wedge y \geq x \wedge \forall z (Cmp(z, y) \rightarrow At(z \cdot x) \wedge \psi(\tau 0 \cdot z, v))]"$ and $\times = \{\langle x, y, z \rangle \in \mathcal{S}^3 : \models_{\mathcal{S}} \vartheta_1(x, y, z)\}$.

LEMMA 4. For every $A, B, C \subseteq X$, we have $\times(A, B, C)$ iff A, B, C are finite subsets of ω , and $\overline{A} \cdot \overline{B} = \overline{C}$.

PROOF. If $\times(A, B, C)$, then, by Lemma 2, A, B, C are finite subsets of ω ; moreover, there are $D, E \subseteq X$ such that $\overline{D} = \overline{A}$, $E \supseteq D$ and for every connected component F of E , $\overline{F \cap \tau 0} = \overline{B}$ and $F \cap D$ has exactly one element; therefore, E has exactly $\overline{D} = \overline{A}$ components; now, let $\overline{A} = n$, and let $F_1 \dots F_n$ be the components of E ; then $\overline{C} = \overline{E \cap \tau 0} = \sum_{i=1}^n \overline{F_i \cap \tau 0} = \sum_{i=1}^n \overline{B} = \overline{A} \cdot \overline{B}$.

On the other hand, suppose that A, B, C are finite subsets of ω , and $\overline{A} \cdot \overline{B} = \overline{C}$. Let $\overline{A} = n, \overline{B} = m$. Let $D = \{0, m, 2m, \dots, (n-1)m\}$, $\overline{F}_i = \{im, im+1, \dots, (i+1)m-1\}$ ($i = 0 \dots n-1$), $\overline{F}_i^1 = \{(\{h, k\}, nm) : h, k \in \overline{F}_i\}$, $F_i = \overline{F}_i \cup F_i^1$, $E = \bigcup_{i=0}^{n-1} F_i$. By Lemma 1, it is easily seen that the connected components of E are $F_0 \dots F_{n-1}$; moreover, $\overline{E \cap \tau 0} = mn = \overline{C}$, $E \supseteq D$, and for every i , $F_i \cap D = \{i\}$, $\overline{F_i \cap \tau 0} = m = \overline{B}$, therefore we can conclude $\times(A, B, C)$. Q.E.D.

Now, we are able to show the following theorem:

THEOREM 1. \mathcal{T} is undecidable.

PROOF. From Lemmas 1–4, it follows that, putting $\mathfrak{F} = \langle \mathcal{F}, +, \times \rangle$, \mathcal{F}/\approx is isomorphic to the standard model $N = \{N, +, \cdot\}$ of natural numbers, therefore the claim follows. Q.E.D.

REMARK. Let \mathcal{T}^n be $\mathcal{T} + \{\tau^n 0 = 1\}$. Since for $n \geq 2$ \mathcal{A} is also a model of \mathcal{T}^n , the above argument shows that, if $n \geq 2$, \mathcal{T}^n is undecidable.

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Received February 16, 1979.