Binding Domains: Anaphoric and Pronominal Pronouns in Categorial Grammar

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Abstract. In this paper we present a treatment for anaphoric pronouns and reflexives in a Type Logical Grammar. To this end, we introduce structural modalities into the left pronominal rule of the categorial calculus with limited contraction **LLC** [8]. Following a proposal due to Hepple [6], we also sketch an analysis for the long-distance anaphora *seg* from Icelandic.

Keywords: 3sg Reflexives \cdot 3sg Pronouns \cdot Binding Theory Type Logical Grammar \cdot Calculus **LLC**

1 Introduction

From a generative perspective, the licensing of pronominal expressions such as *he*, *him*, *himself* is determined by the so-called Principles A and B of the Binding Theory [3]. Principle A stipulates that an anaphor (reflexives and reciprocals) must be bound in its governing category (roughly, it must have a c-commanding local antecedent). Principle B stipulates that a pronoun must be free (i.e. not bound) within its governing category; notwithstanding, a pronoun can be bound from outside this syntactic domain. Thus a pronoun, unlike an anaphor, also admits a free reading. Principles A and B jointly imply a strict complementary distribution between pronouns and reflexives in *some* syntactic domains, as exemplified below:

- (1) John₁ admires himself₁/*him₁.
- (2) John₁'s father₂ loves $\lim_{1/*2}/\lim_{1/*2}$
- (3) The father₁ of John₂ loves $him_{2/*1}/himself_{1/*2}$.
- (4) John₁ believes $himself_1/he_1/him_1$ to love Mary.
- (5) John₁ says $he_{1/2}/*himself_1$ loves Mary.

The Binding Theory has been successively revisited to overcome some counterexamples. Complementary distribution is disconfirmed, on the one hand, in adjunct clauses, like in (6) below. On the other hand, languages like Icelandic, Dutch, German and Norwegian each contain an anaphoric form—sig, zich, sich, seg, respectively—that does not meet Principle A, as the former can be bound by a long-distance antecedent (cf. [22]).¹

- (6) John₁ glanced behind $him_1/himself_1$.
- (7) Jón₁ segir að María₂ telji að Haraldur₃ vilji að Billi₄
 John says that Maria believe.SBJ that Harold want.SBJ that Bill
 heimsæki sig_{1/2/3/4}.
 visit.SBJ SE-ANAPHOR
 'John says that Maria believes that Harold wants that Bill visit him/her.'

Several categorial—combinatory and type-logical—calculi have been proposed to deal with reflexives and anaphoric pronouns. Some of them treat multiple-binding into the lexicon (cf. [16,23]); others use a syntactic approach (cf. [6–8]). Working on a Type-Logical Grammar, Jaeger [8] develops the Lambek calculus with Limited Contraction (**LLC**) to syntactically process personal pronouns in a uniform way; he does not discriminate syntactically (nor semantically) among reflexives, anaphoric pronouns and pronominals. In other words, he does not take Principles A and B of the Binding Theory into account.

Our goal is to give a more accurate treatment of personal pronouns, taking as a starting point the pronominal connective | from Jaeger and the intuition behind its logical rules. Firstly, we modify the right rule of **LLC** to distinguish on the one hand the *free* (or pronominal) *use* from the *bound* (or anaphoric) *use* of a personal pronoun. Secondly, by using the (lexical) structural modality $\langle \rangle$ of [18] (and, analogously, for our []), we identify different syntactic domains for binding: we impose structural conditions into the pronominal left rule of **LLC** by using the corresponding (syntactic) structural modality [] (and also { }). Thus, on the other hand, we also distinguish anaphoric pronouns from reflexive anaphors. As a consequence, although we deal with reflexives, anaphoric pronouns and pronominals, our proposal is not intended to be a uniform approach. For reasons of space, we restrict ourselves to cases where the binder is a nominal phrase and the bindee (a pronoun or a reflexive) carries 3SG features.² Since our proposal is inspired by Jaeger's calculus, we also do not deal with cases in which the bindee precedes its binder.³

The structure of the paper is as follows. In Sect. 2, we present Jaeger's calculus **LLC** (in a sequent format) and we briefly discuss some questions related to the problem of overgeneration. In Sect. 3 we change the right pronominal rule of **LLC** to distinguish between a reflexive and a pronominal type-constructor. In Sect. 4, firstly we present our treatment for subject-oriented anaphors in several syntactic

¹ Although these languages contain this kind of simple (also weak) reflexive form, their syntactic behavior is not the same in all of them (cf. for example, [5]).

² Hence, we restrict ourselves to what some theories call anaphoric coreference, not binding (cf. [2,21]). Though it is generally accepted that reflexives and reciprocals behave in the same way with respect to binding conditions, their semantic value diverges. For this reason, we also do not deal with reciprocal anaphors.

³ However, a version of Jaeger's rules that also allows cases of cataphora is presented in [17].

domains and secondly, we deal with object-oriented anaphors in double-object constructions and prepositional complements. Finally, we sketch an analysis for long-distance anaphors from Icelandic. Section 5 concludes the paper. In the Appendix we sketch the principal cut for our new pronominal rules.

2 LLC Calculus

LLC is a conservative extension of the Lambek **L** calculus (without empty antecedents) [9]. Like **L**, **LLC** is free of structural rules. Jaeger's calculus treats resource multiplication syntactically. **LLC** extends the sequent calculus **L** by adding the anaphoric type-constructor |. The rules of the latter encode a restricted version of the structural rule of Contraction, thus allowing for multiplebinding (see Fig. 1). Despite incorporating this structural rule, **LLC**, as well as Lambek system, enjoys Cut elimination, decidability and the subformula property. Indeed, as the reader can check, all the formulas that occur in the premises of the two new rules for the anaphoric type-constructor are subformulas of the formulas that occur in their conclusion.

$$\frac{Y \Rightarrow M:A}{X,Y,Z,z:B|A,W \Rightarrow N[M/x][(zM)/y]:C} \mid \mathcal{L}$$

$$\frac{X, x_1 : B_1, Y_1, \dots, x_n : B_n, Y_n \Rightarrow N : C}{X, y_1 : B_1 | A, Y_1, \dots, y_n : B_n | A, Y_n \Rightarrow \lambda z . N[(y_1 z)/x_1] \dots [(y_n z)/x_n] : C | A} | \mathbf{R}$$

Fig. 1. Left and right rules for |

Note that when A is a basic type, the left premise of |L| is an instance of the identity axiom; thus the rule can be simplified, as shown in Fig. 2.⁴

$$\frac{X, x: A, Z, y: B, W \Rightarrow M: C}{X, x: A, Z, z: B|A, W \Rightarrow M[zx/y]: C} |L$$

Fig. 2. Simplified left rule for

Anaphoric expressions are then assigned a type B|A: it works as a type B in the presence of an antecedent of type A. The |L rule expresses the fact that for an anaphoric expression to be bound it needs an *antecedent* in the same

⁴ Jaeger is not only concerned with anaphoric pronouns but also with other anaphoric phenomena, such as ellipsis of VP.

premise, that is, in some *local* syntactic domain. Besides imposing an antecedent condition, this rule incorporates a restricted version of the structural rule of (long-distance) Contraction, in that the antecedent A for the anaphoric type B|A occurs in both premises of this rule.

Since personal pronouns take their reference from a nominal antecedent, they are assigned the syntactic anaphoric category n|n.⁵ In semantic terms, a pronoun denotes the identity function $\lambda x.x$ over individuals; the reference of a pronoun is identical with the reference of its antecedent.⁶

Since a pronominal type n|n can be constructed by using the |R and |L rules, and since anaphoric and pronominal pronouns are assigned the same syntactic type (and the same semantic category), the system can accurately recognize the free and the bound readings for a pronominal. Thus, for example, the system recognizes the double reading for he in (5), and so may assign it the saturated type s or the unsaturated (or functional) type s|n. The latter expresses the fact that a free pronoun occurs in the clause. In addition, the system can also derive the co-occurrence of bound pronouns and reflexives in syntactic domains in which complementary distribution fails, as exemplified in (6) above (see Fig. 3). It can also recognize the ungrammaticality of (8b) below, since the antecedent condition on |L is not fulfilled. Nevertheless, **LLC** also allows for the ungrammatical anaphoric readings in the following examples.⁷

- (8) a. John₁ saw himself₁/*him₁.
 - b. * Himself₁ saw John₁.
- (9) a. John talked to $Mary_1$ about $herself_1$.
 - b. * John talked about $Mary_1$ to $herself_1$.
- (10) John₁ saw *himself₁'s/his₁ mother.
- (11) a. John₁ believes himself₁ to kiss Mary.
 - b. * John₁ believes himself₁ kisses Mary.

$$\frac{n, (n \setminus s)/pp, pp/n, n \Rightarrow s}{n, (n \setminus s)/pp, pp/n, n | n \Rightarrow s} | \mathbf{L} \qquad \qquad \frac{n, (n \setminus s)/pp, pp/n, n \Rightarrow s}{n, (n \setminus s)/pp, pp/n, n | n \Rightarrow s | n} | \mathbf{R}$$

Fig. 3. Schematic derivation for $John_1$ glanced behind himsel f_1/him_1

Since we are looking for a more accurate treatment for the distribution of pronominal and anaphoric pronouns, we shall begin by distinguishing between an

 $^{^5}$ As usual, we use n for proper names, s for sentences, cn for common nouns and pp for prepositional phrases.

⁶ In this respect, Jaeger follows [6,7].

⁷ Everaert [4] uses these sentences to evaluate the scope and limits of several generative models for binding.

anaphoric connective for reflexives and a (possibly) non-anaphoric connective for personal pronouns like *he* and *him*. Later on, we shall draw a distinction between reflexives and bound pronouns.

3 Bound and Free Pronouns: Splitting the Pronominal Connective

For non-reflexive pronouns, we adopt Jaeger's left rule and the following left and right rules, which split the |R| rule of **LLC**. It is important to emphasize that these two new rules, like those of **LLC**, satisfy the subformula property: all the formulas that occurs in the premises of ||L| and ||R| are subformulas of the formulas that occurs in their conclusion. Given that the proof of Cut elimination for these rules requires using a limited version of the Expansion rule (see Appendix), we call our modified version of Jaeger's system **LLBE**: Lambek calculus with Bracketed Expansion (Fig. 4).

$$\frac{X, x: \llbracket A \rfloor, Z, y: B, W \Rightarrow M: C}{X, Z, z: B \|A, W \Rightarrow M[zx/y]: C} \parallel \mathbb{L} \qquad \qquad \frac{Y \Rightarrow N: C}{x: \llbracket A \rfloor, Y \Rightarrow \lambda x. N: C \|A} \parallel \mathbb{R}$$

Fig. 4. Left and right rules for ||

As can be noted, we split (an extremely simplified version of) the |R rule of **LLC** to obtain a second left rule.⁸ Hence, a pronominal type-constructor will have two left rules: |L and ||L.⁹ By breaking the |R rule of **LLC** we can more clearly show that free and bound are labels that result from the procedures by which we use a pronoun: we apply the rule of use |L to get a bound (or anaphoric) use of a pronoun, while in applying ||L we use a pronoun freely.¹⁰

The $||\mathbf{R}|$ rule compiles a restricted form of the structural rule of Expansion, as it introduces a formula that is a sub-formula of the pronominal type C||A. Given that we do not assume logical rules for the brackets $|| \parallel ||$, they can only be introduced (deleted) through the use of the $||\mathbf{R}| (||\mathbf{L}|)$ rule. Consequently, like in Jaeger's proposal, the rule of proof for a free pronoun goes hand in hand with its free use in **LLBE**. However, unlike the $||\mathbf{R}|$ rule of **LLC**, the $||\mathbf{R}|$ rule of **LLBE**

⁸ Strictly speaking, we split the $|\mathbf{R}|$ rule of **LLC** for the case where n = 1 and X is the empty sequence ϵ . As we shall show in the Appendix, the proof of principal Cut for the new rules requires using bracketed versions of the structural rules of Permutation and Expansion. In order to avoid a proof of a pronominal type C||A| for any type C, the antecedent type A of the rule $||\mathbf{R}|$ has to be left-peripheral.

⁹ From Sect. 4, the |L rule (for non-reflexive pronouns) will be renamed $||L_a$, and ||L will have to be read as $||L_p$. Though we shall retain |L for reflexives only, we will rename it $|L_a$ for the sake of uniformity.

¹⁰ As we shall see later, the formula B in the ||L rule will have a bracketed structure [B] in most cases.

does not simultaneously construct a pronominal type to the left and to the right sides of a sequent.

In addition, since **LLBE** contains two left pronominal rules (i.e. |L and ||L), we are able to characterize two anaphoric type-constructors: a reflexive type-constructor, which uses only the |L rule, and a pronominal type, which uses |L, ||R and also ||L.¹¹ By assigning different syntactic types for reflexives and pronouns—n|n and n||n, respectively—, and given that the |R rule of **LLC** for the case n = 1 may be derived by using ||L and ||R of **LLBE**, the latter system, like the former one, adequately recognizes grammatical sentences like those in (12–16a), whilst blocking the ungrammatical sentence in (16b) below (Fig. 5).

- (12) John₁ said $he_{1/2}$ runs.
- (13) John₁ said Mary likes $him_{1/2}$.
- (14) John₁ likes him_2 .
- (15) He_1 likes himself_1.
- (16) a. John₁ likes himself₁.
 - b. * John₁ likes himself₂.

$$\frac{ \begin{array}{c} \vdots \\ n, (n \backslash s)/n, n \Rightarrow s \\ \hline \hline [\boxed{ n \rfloor, n, (n \backslash s)/n, n \Rightarrow s \| n \\ n, (n \backslash s)/n, n \| n \Rightarrow s \| n \end{array} \| \mathbf{R}$$

Fig. 5. Derivation for $John_1$ likes him_2

Nevertheless, since |L is adopted for the pronominal type n||n| and also for the anaphoric type n|n, **LLBE** is not yet capable of separating bound (object) pronouns from reflexives.¹²

¹¹ A Type-Logical sequent calculus generally contains one left and one right rules for each type-constructor. Since in our proposal the reflexive type-constructor uses only a left rule, our approach is non-standard.

¹² In order to distinguish subject and object pronouns, we could assign the lifted type $(s||n)/(n\backslash s)$ to the former (cf. [17]). Although at first glance it would seem that a lifted type— $(s/n)\backslash(s||n)$ —is also adequate to categorize an object pronoun like *him*, it is not clear how we could deal with Exceptional Case Marked (ECM) constructions, in which the semantic argument of the embedded infinitive clause surfaces with accusative case. Indeed, if *him* were assigned $(s/n)\backslash(s||n)$ because of its surface form, it would combine with a verb phrase to the left, like a real object complement does. But if this were the case, the subject slot of the embedded complement clause would not be saturated and then, the sentential argument of the ECM verb would become unsaturated as well.

4 Reflexives and Bound Pronouns: Imposing Structure Through Bracket Modalities

4.1 Subject-Oriented Reflexives and Bound Pronouns

There are several syntactic domains where a reflexive can occur: in nominal object complements, prepositional object complements, adjunct clauses, and even in an embedded position within nominal phrases (the so-called NP anaphora). In some of them, complementary distribution is fully verified: when the reflexive and its antecedent are co-arguments of the same function, a bound pronoun is ruled out.

Propositional Complements. Complementary distribution is also verified in the opposite direction in (some) clauses selected by propositional verbs like *say* and *believe*. As is widely known, reflexives are ruled out and bound pronouns are licensed in (finite) propositional complements of both these verbs, as exemplified below:

- (17) John₁ said/believes *himself₁/he_{1/2} walks.
- (18) John₁ said/believes Mary hates $*himself_1/him_{1/2}$.

Anaphors within propositional complements have already been adequately analyzed in Categorial Grammar (cf. [6,15,23], a. o.). The correct binding relation in these complements is ensured by using a normal (or semantic) S4 modality \Box [13]. In these categorial proposals, reflexives and pronouns are assigned different pronominal types. The following modalized lexical entries capture the above-mentioned facts:

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\begin{array}{l} \mathbf{him/he} : \Box(\Box n \| n) \\ \mathbf{himself} : \Box(n | n) \\ \mathbf{say/believe} : \Box((n \setminus s) / \Box s) \\ \mathbf{walk} : \Box(n \setminus s) \\ \mathbf{hate} : \Box((n \setminus s) / n) \end{array}
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Nevertheless, authors have sometimes glossed over the fact that a reflexive may occur within a propositional complement if it occupies an embedded position:

(19) Max₁ said (that) the queen invited both Lucie and $himself_1/him_1$ for tea.

Similarly, it has not always been noted that *believe* licenses the occurrence of a reflexive in the subject position when the complement verb is in a non-finite form.¹³ This last fact is a specific case of a more general situation: in complements of Exceptional Case Marked (ECM) verbs, such as *believe* or *expect*, a reflexive is allowed, while a bound pronoun is ruled out in the subject (and also the

¹³ In passing, we point out that, unlike English, literary Spanish and Italian allow a nominative free or bound pronoun in non-finite complements of propositional verbs [11].

object) argument slot. Thus, in ECM constructions, the claimed complementary distribution is verified, as in other verb complements. Nevertheless, complementary distribution in ECM constructions is, in some sense, unexpected, since the reflexive in the subject position of non-finite complements is not a co-argument of the binder.

- (20) John₁ believes $himself_1/*him_1/*he$ to kiss Mary.
- (21) Lucie expects John₁ to like himself₁/*him₁.
- (22) Lucie₁ expects $herself_1/*her_1/*she$ to kiss John.

From this evidence, it seems important to differentiate the lexical entry for a propositional *believe* from the ECM *believe*, despite the fact that both verbs select a sentential (finite or non-finite) complement. Following [12, 14, 18], we shall use the structural modality $\langle \rangle$ to mark the (syntactic) argument positions of a verb; though we shall use it to mark not only the subject position, but also the object complement position. In order to distinguish the propositional verb *believe/say* from non-propositional verbs, we shall set a different bracket modality [] aside for the former.

A sample of a bracketed lexicon is given below:

walk : $\langle n \rangle \setminus s$ like/hate/kiss : $(\langle n \rangle \setminus s) / \langle n \rangle$ say/believe : $(\langle n \rangle \setminus s) / \lceil s \rfloor$ john/mary : nhim : $n \parallel n$ himself : $n \mid n$

The right rules for brackets $\langle \rangle$ and [] are given below. The rules for Lambek's slashes are applied to structured sequences [X] and $\{X\}$ of types.¹⁴ The structures [] and $\{ \}$ are then spread over the sequents when functional types A/B and $B \setminus A$ are built out of [B] and $\{B\}$, respectively. Hence, while the structural modality $\langle \rangle$ is a lexical mark, the insertion of the modalities [] and $\{ \}$, and thus the delimitation of syntactic domains, is a consequence of syntactic operations (Figs. 6 and 7).

$$\frac{X \Rightarrow A}{[X] \Rightarrow \langle A \rangle} \langle \rangle \mathbb{R} \quad \frac{X \Rightarrow A}{Y[B/A, X] \Rightarrow C} [/] \mathbb{L} \quad \frac{X \Rightarrow A}{Y[X, A \setminus B] \Rightarrow C} [\backslash] \mathbb{L}$$

Fig. 6. Rules for brackets $\langle \ \rangle$ and structured [] sequents

¹⁴ Generally, $\Delta[\Gamma]$ indicates a configuration Δ containing a distinguished configuration Γ of types. In our rules, X[Z] would indicate a sequence X with a distinguished structured sequence [Z] of types, and analogously for $\{Z\}$.

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$$\frac{X \Rightarrow A}{\{X\} \Rightarrow \lceil A \rfloor} \left[\ \rfloor \operatorname{R} \quad \frac{X \Rightarrow A}{Y \{B/A, X\} \Rightarrow C} \left\{ / \right\} \operatorname{L} \quad \frac{X \Rightarrow A}{Y \{X, A \setminus B\} \Rightarrow C} \left\{ \setminus \right\} \operatorname{L}$$

Fig. 7. Rules for brackets [] and structured $\{ \}$ sequents

$$\frac{[A], Z_1, [Z_2, B], W \Rightarrow C}{[A], Z_1, [Z_2, B|A], W \Rightarrow C} \quad [|] \mathcal{L}_a \qquad \frac{[X_1, A, X_2], Z_1, [Z_2, B], W \Rightarrow C}{[X_1, A, X_2], Z_1, [Z_2, B||A], W \Rightarrow C} \quad [||] \mathcal{L}_a$$

Fig. 8. Rules for (subject-oriented) reflexives and bound pronouns within a [] domain

$$[X_1, A, X_2], Z_1, \{Z_2, [Z_3, B], W\} \Rightarrow C X_1, A, X_2], Z_1, \{Z_2, [Z_3, B || A], W\} \Rightarrow C$$
 {|} L_a

$$\frac{[X_1, A, X_2], Z_1, \{Z_2, [Z_3, B], W\} \Rightarrow C}{[X_1, A, X_2], Z_1, \{Z_2, [Z_3, B ||A|, W\} \Rightarrow C} \{\|\} L_a$$

Fig. 9. Rules for bound pronouns and reflexives within a { } domain

We propose, consequently, bracketed versions for the $|L_a|$ and $||L_a|$ rules, with the following side conditions: $X_1 \neq \epsilon$ or $X_2 \neq \epsilon$ in $[||]L_a; Z_3 \neq \epsilon$ in $\{|\}L_a$.¹⁵

(i) * John₁ expects Mary to like himself₁.

ĺ

(ii) * John₁ believes Mary to expect Susan to like himself₁.

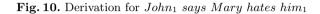
In addition, the side conditions on the $[||]L_a$ rule inadequately license pronouns to be bound by an antecedent within a conjunctive nominal phrase, as exemplified below. Indeed, *Mary* is taken as an argument of the functional type commonly assigned to *and*:

- (iii) * John and Mary₁ praised her₁.
- (iv) John₁ and Mary talked about him_1 .

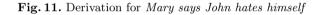
It appears that the unbracketed type assigned to the conjunction and has to be differentiated from the (bracketed) functional types assigned, for example, to of— $(n\backslash n)/\langle n\rangle$ and s— $\langle n\rangle \backslash (n/cn)$. A distinction between a collective and a distributive type for and also seems to be relevant: roughly, $X \backslash X/X$ and $\langle X \rangle \backslash X/\langle X \rangle$, for example. For reasons of space, and since judgments seem to vary among speakers and sentences, we defer this problem to future research.

¹⁵ As an anonymous reviewer pointed out, if the type $(\langle n \rangle \backslash s)/s$ were assigned to ECM verbs to differentiate them from propositional verbs, it would allow for ungrammatical sentences like (i–ii) below. To block binding of a reflexive in an object position by a non-local antecedent it seems we would have to impose some condition on the sequence Z_1 in []]L_a. We plan to address the challenge posed by ECM constructions in future investigations.

$$\frac{-\frac{n \Rightarrow n}{[n] \Rightarrow \langle n \rangle} \langle \ \rangle \mathbf{R} \quad \frac{\frac{n \Rightarrow n}{[n] \Rightarrow \langle n \rangle} \langle \ \rangle \mathbf{R} \quad \frac{s \Rightarrow s}{\{s\} \Rightarrow [s]} \left[\] \mathbf{R} \\ \frac{[n] \Rightarrow \langle n \rangle}{\{[n], \langle n \rangle \setminus s\} \Rightarrow [s]} \left\{ \\ \frac{\{[n], \langle (n \rangle \setminus s) / \langle n \rangle, [n]\} \Rightarrow [s]}{\left[\frac{[n], \langle (n \rangle \setminus s) / \langle n \rangle, [n]\} \Rightarrow s}{[n], \langle (n \rangle \setminus s) / \langle n \rangle, [n]\} \Rightarrow s} \left[\frac{[n], \langle (n \rangle \setminus s) / [s], \{[n], \langle (n \rangle \setminus s) / \langle n \rangle, [n]\} \Rightarrow s}{[n], \langle (n \rangle \setminus s) / [s], \{[n], \langle (n \rangle \setminus s) / \langle n \rangle, [n]\|n]\} \Rightarrow s} \left\{ \| \} \mathbf{L}_{a} \right\}$$



$$\begin{array}{c} \vdots \\ \hline [n], (\langle n \rangle \setminus s) / \langle n \rangle, [n] \Rightarrow s \\ \hline [n], (\langle n \rangle \setminus s) / \langle n \rangle, [n|n] \Rightarrow s \\ \hline [n], (\langle n \rangle \setminus s) / \langle n \rangle, [n|n] \} \Rightarrow \lceil s \\ \hline [n], \langle n \rangle \setminus [s], \{[n], (\langle n \rangle \setminus s) / \langle n \rangle, [n|n] \} \Rightarrow s \\ \end{array} \right] / L$$



$$\begin{array}{c} \vdots \\ \left\{ [n], \left(\langle n \rangle \setminus s \right) / \langle n \rangle, [n] \right\} \Rightarrow \left\lceil s \right\rfloor \\ [n], \left\langle n \rangle \setminus s \Rightarrow s \\ \hline \frac{[n], \left(\langle n \rangle \setminus s \right) / \left\lceil s \right\rfloor, \left\{ [n], \left(\langle n \rangle \setminus s \right) / \langle n \rangle, [n] \right\} \Rightarrow s \\ [n], \left(\langle n \rangle \setminus s \right) / \left\lceil s \right\rfloor, \left\{ [n], \left(\langle n \rangle \setminus s \right) / \langle n \rangle, [n|n] \right\} \Rightarrow s \\ \end{array} \right\} / L$$

Fig. 12. Illicit derivation for $John_1$ says Peter hates himself₁

The rule [|]L_a in Fig. 8 preserves the prominence condition on the binder for reflexives: given that the reflexive within an argument domain [] takes [A] as its binder, the binder itself is not part of the subject (i.e. higher) argument. Conversely, the side conditions on the sequences X_1 and X_2 in [||]L_a impede binding of a pronoun in an argument position if the binder is not part of the higher subject argument.

Note that there is no condition on the sequences X_1 and X_2 in the rule for pronouns $\{\|\}L_a$ in Fig. 9. Thus, a pronoun within a propositional complement can be bound by a matrix subject (see Fig. 10). The side condition on the reflexive $\{|\}L_a$ rule ensures that the reflexive stands in an embedded position within the propositional complement clause (contrast Figs. 11 and 12).

Nominal Complements. In nominal complements complementary distribution is fully verified: where the anaphora and its antecedent are co-arguments of the same function, a bound pronoun is ruled out. Nevertheless, a pronoun in the object complement position can still be bound provided that the binder itself is an argument of another functional type, as exemplified below:

- (23) $[John_1's father]_2 loves him_1/himself_2.$
- (24) [The father of $John_1$]₂ loves $him_1/himself_2$.

From the following (bracketed) lexicon we may obtain the correct binding relations for reflexives and pronouns within the (direct) object argument position, as exemplified in (23–24) above and in (25–28) below (Fig. 13):

see/like : $(\langle n \rangle \langle s \rangle) / \langle n \rangle$ john/mary : nfather/picture : cnthe/a : n/cnof : $(n \backslash n)/n$'s : $n \backslash (n/cn)$

- (25) John likes himself/*him.
- (26) John takes a picture of himself/*him.
- (27) * John saw himself's mother.
- (28) John believes himself/*him to kiss Mary.

$$\begin{array}{c} \vdots \\ \underline{[n], (\langle n \rangle \setminus s) / \langle n \rangle, [n] \Rightarrow s} \\ \overline{[n], (\langle n \rangle \setminus s) / \langle n \rangle, [n|n] \Rightarrow s} \end{array} []] \mathcal{L}_a / []] \mathcal{L}_a *$$

Fig. 13. Derivation for $John_1$ likes $himself_1 / *him_1$

Prepositional Phrases. In general terms, scholars agree that prepositional phrases (PPs) selected by a verb can only contain a reflexive but not a bound pronoun, while prepositional phrases operating as adjuncts allow both a reflexive and a bound pronoun (cf. [4]).

Since our proposal strongly depends on the syntactic types assigned to the lexical items into the lexicon, the correctness of our proposal for anaphoric items within prepositional phrases mainly rests on the type assigned to the different classes of verbs.

Unfortunately, the distinction between complement prepositional phrases and adjunct phrases is not so pure in some cases. As claimed in [10], locative PPs, including those selected by a verb, must be distinguished from other PPs. Those verbs that select a PP bearing a locative role like *put* and *sit*, allow several locative prepositions, such as *in*, *on*, *near*, *into*, *next*, *in front of*. In this sense, locative PPs resemble adjunct PPs. By contrast verbs like *relies*, despite selecting a PP as complement, also select some specific preposition. The PP headed by *on/upon* in *relies on/upon* does not bear a locative role. Given this, it seems clear that we need to set a distinction between the PP selected by verbs like *put* and

the PP selected by verbs like *relies*. In other terms, we need to set a bipartition into the set of PP complements: locative PPs and non-locative PPs. By using the bracket modality $\langle \rangle$ we mark the non-locative PP complement position into the lexical entry of the corresponding verbs and, taking into consideration the similarity between adjunct PPs and locative PP complements, we leave the PP position for the locative complement unmarked.¹⁶

Given that reflexives and anaphoric pronouns can occur within an unmarked position, we assume the following rules to process them (Fig. 14):

$$\frac{[X_1, A, X_2], Z, B, W \Rightarrow C}{[X_1, A, X_2], Z, B | A, W \Rightarrow C} \mid \mathbf{L}_a \quad \frac{[X_1, A, X_2], Z, B, W \Rightarrow C}{[X_1, A, X_2], Z, B ||A, W \Rightarrow C} \mid \mathbf{L}_a$$

Fig. 14. Rule for reflexives and anaphoric pronouns out of bracketed domains

Thus, assuming the following lexicon, we obtain the correct binding relation in different prepositional phrases, as exemplified in (29-32):

put/see : $((\langle n \rangle \backslash s)/pp) / \langle n \rangle$ **glance** : $(\langle n \rangle \backslash s)/pp$ **rely** : $(\langle n \rangle \backslash s) / \langle pp \rangle$ **on/upon/behind/next** : pp/n

- (29) John₁ relies on himself₁/*him₁.
- (30) John₁ glanced behind $himself_1/him_1$.
- (31) John₁ put the gun near/under/on $himself_1/him_1$.
- (32) John₁ saw a gun near himself₁/him₁.

4.2 Object-Oriented Reflexives and Bound Pronouns

Nominal Complements. Verbs like *show, give, send, promise, introduce* may select two nominal phrases as complements, and thus give rise to double-object constructions. These structures allow then for another pattern of reflexivization: reflexives bound by a nominal within a verb complement position. In other terms, besides subject-oriented reflexives, double-object constructions also allow for object-oriented ones. Double-object constructions alternate with oblique dative structures:

- (33) Mary showed/gave/sent/promised John a gift.
- (34) Mary showed/gave/sent/promised a gift to John.

¹⁶ Alternatively, Reinhart and Reuland [22] consider that *relies on* forms a complex (semantic and syntactic) unit selecting a nominal complement, whilst *put* selects a prepositional complement. In view of this fact, we would assign the type $(\langle n \rangle \backslash s) / \langle n \rangle$ to *relies on/upon*.

Although the two structures display a different linear order, both reveal the same behavior when licensing anaphors. They show an asymmetry with respect to the licensing of object-oriented reflexives. As shown in the examples below, the correct binding relation for object-oriented reflexives in both structures depends on an ordering of the complements.

- (35) Mary showed/presented John₁ himself₁.
- (36) Mary showed/presented John₁ to himself₁.

Thus, these structures indicate that a hierarchical order between the two objects has to be imposed. To deal with double-object constructions we extend the calculus by adding a new product type-constructor.¹⁷ We also present a different inference rule for object-oriented reflexives, where Z_1 does not contain a subtype s (Figs. 15 and 17).¹⁸

$$\frac{[X] \Rightarrow A \quad [Y] \Rightarrow B}{[X, [Y]] \Rightarrow A \otimes B} \ [\odot] \mathbf{R}$$

Fig. 15. Right rule for non-commutative \otimes asymmetrical product

The correct binding relation is ensured by the following lexical assignment:

show/give/present/send : $(\langle n \rangle \setminus s)/(\langle n \rangle \otimes \langle n \rangle)$ **show/give/present/send** : $(\langle n \rangle \setminus s)/(\langle n \rangle \otimes \langle pp \rangle)$

(i) I gave to the students presents that I had brought back from Spain.

To deal with double-object structures, Hepple [6] extends the **L** calculus by adding a new slash type-constructor \oint and a modality \triangleright . Since the slash type-constructor lacks introduction rules, it may encode the hierarchical ordering of the nominal complements; the modality allows the nominal complements to be reordered to obtain the correct surface word-order.

- ¹⁸ We note that a slightly modified version of the rule in Fig. 16 may also be used for anaphors in a complement of possessives, which are not either subject nor objectoriented. Once again, it appears that a distinction between the functional type assigned to of or 's and and has to be made to prevent He and himself from assigning the type n.
 - (i) John₁'s description of himself₁ is lovely.
 - (ii) Lisa burned Andy Warhol₁'s portrait of himself₁.

It seems that it could be possible to also encode a hierarchical ordering into the rules for the Lambek slash type-constructors. In this case, it would be possible to deal with subject- and object-oriented anaphors in a uniform way.

¹⁷ Note that the product ⊗ is not a discontinuous (or wrapping) type-constructor, unlike that of [1] or [19]. Since ⊗ is non-commutative, we would not be able to derive cases of "heavy" NP, as exemplified below. Nevertheless, in the following section we shall adopt a commutative product-type for the treatment of prepositional phrases.

$$\frac{X_1, [X_2, A, Z_1, [Z_2, B]], W \Rightarrow C}{X_1, [X_2, A, Z_1, [Z_2, B|A]], W \Rightarrow C} \ []] \mathcal{L}_a$$

Fig. 16. Rule for object-oriented reflexives

Fig. 17. Derivation for Mary presents John himself

Prepositional Phrases: The *About-Phrase.* Verbs selecting two prepositional phrases also challenge several binding theories. In this case, there is also no complete agreement among scholars with respect to their syntactic status.¹⁹ As is known, two prepositional complements may appear in either order:

- (37) John talked to Mary about Bill.
- (38) John talked about Bill to Mary.

Despite the free word-order, the occurrence of a reflexive within a prepositional phrase, such as in double-object structures, indicates that a structural ordering between the *about*-phrase and the *to*-phrase has to be imposed.

- (39) John talked to $Mary_1$ about herself₁.
- (40) * John talked about $Mary_1$ to herself₁.

(ii) * Mary talked about himself₁ to John₁.

¹⁹ In some generative theories, the *about*-phrase is evaluated as an adjunct phrase and thus is separated from the *to*-phrase or *with*-phrase complement (cf. [22]). This would explain the ungrammaticality of (40), but not the ungrammaticality of (i) below.

⁽i) * Mary talked to $John_1$ about him_1 .

In other theories, the *about*-phrase, as well as the *to*-phrase, is considered a verb complement; the difference between these PPs is made by assuming an ordering with respect to their relative obliqueness: the *about*-phrase is more oblique than the *to*-phrase (cf. [20]). Since the anaphor has to be bound by a less oblique co-argument, the relationship of relative obliqueness would account for (i) above, but not for (ii) below, where the linear word-order seems to be also relevant.

In addition, [3] suggests an approach in which the verb talk (and also speak) and the preposition to are reanalyzed as one verb taking a nominal object (and a prepositional complement) (cf. also [23]). Thus, talk would be analogous to (one of the forms of) tell. To formalize this proposal, besides encoding free linear word-order and relative obliqueness, the syntactic functional type assigned to the talk to-phrase would have to encode discontinuity as well.

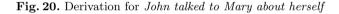
Binding Domains for Anaphors in CG 15

$$\frac{[X] \Rightarrow A \qquad [Y] \Rightarrow B}{[X, [Y]] \Rightarrow A \otimes B} \ [\heartsuit] \mathbf{R}_1 \qquad \qquad \frac{[X] \Rightarrow B \qquad [Y] \Rightarrow A}{[[X], Y] \Rightarrow A \otimes B} \ [\heartsuit] \mathbf{R}_2$$

Fig. 18. Right rules for commutative asymmetrical product ⊗

Fig. 19. Illicit derivation for *John talked about Mary to herself

$$\begin{array}{c} \vdots & \vdots \\ \hline \left[PP_{to}/n, n \right] \Rightarrow \langle PP_{to} \rangle & \left[pp_{about}/n, n \right] \Rightarrow \langle pp_{about} \rangle \\ \hline \left[\frac{\left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle }{\left[n \right], \langle n \rangle \setminus s \Rightarrow s} \right] \\ \hline \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / (\langle pp_{to} \rangle \otimes \langle pp_{about} \rangle), \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow s}{\left[n \right], \langle n \rangle \setminus s \rangle / (\langle pp_{to} \rangle \otimes \langle pp_{about} \rangle), \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow s} \right] \\ \hline \end{array} \right] \\ \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / (\langle pp_{to} \rangle \otimes \langle pp_{about} \rangle), \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow s}{\left[n \right], \langle n \rangle \setminus s \rangle / \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow s} \right] \\ \end{array} \right] \\ \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow s}{\left[n \right] \left[n \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \Rightarrow s} \right] \\ \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \setminus s \rangle / \langle pp_{to} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s / \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s / \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s / \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle, \left[pp_{to}/n, n, \left[pp_{about}/n, n \right] \right] \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s / \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle \otimes s \rangle \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s / \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle \otimes \langle pp_{about} \rangle \otimes s \rangle \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s }{\left[n \right], \langle n \rangle \otimes s \rangle \otimes s \rangle } \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s }{\left[n \right], \langle n \rangle \otimes s \rangle \otimes s \rangle \otimes s \rangle \otimes s \rangle } \right] \\ \left[\frac{\left[n \right], \langle n \rangle \otimes s }{\left[n \right], \langle n \rangle \otimes s \rangle$$



In a categorial framework, it is the functional type assigned to a verb like *talk* which has to express the different syntactic relation that these two PPs maintain with the verb. In [16], for example, the type assigned to *talk* is $((n \ s)/pp)/pp$, while in [19] it is the type $(n \ s)/(pp_{to} \otimes pp_{about})$, where \otimes is the nondeterministic continuous product of the Displacement Calculus **D**. Thus, this last type captures the alternative surface word-order. This notwithstanding, by using either the former or the latter type, prepositional phrases both get the same syntactic non-hierarchical status of verb complements.²⁰

Hence, the different hierarchical relation the PPs complements maintain with the verb seems to call for a new type-constructor that is analogous to that we have used to deal with double-object constructions, but which encodes commutativity as well (Figs. 18, 19 and 20).

With this type-constructor at hand, we then propose the following lexical assignment:

talk : $(\langle n \rangle \backslash s) / (\langle pp_{to} \rangle \otimes \langle pp_{about} \rangle)$

4.3 Long-Distance Anaphors

Anaphors in Icelandic are necessarily subject-oriented and do not respect Principle A for anaphors, as they can be bound by a long-distance antecedent, provided

²⁰ Since the calculus **D** also contains a nondeterministic discontinuous product \odot , the type $(n \setminus s)/(pp \odot pp)$ would take the structural ordering into account if the premisses of the right rule were bracketed sequences.

that the anaphora stands in a subjunctive clause. In this sense, long-distance anaphors resemble anaphoric pronouns in propositional (finite) complements from English. In addition, the subjunctive mood in Icelandic may be propagated down through embedded complements (this is the so-called *domino effect*). Given that bracket modalities have been applied in Type-Logical Grammar to delimit syntactic domains, we suggest using the bracket $\{ \}$ to simulate the domino effect of the subjunctive mood generated by some verbs (e.g. *segir* 'say' vs. *vita* 'know') and the bracket [] to ensure binding only by the subject (that is, the subject condition; cf. [18]). Since the licensing of a long-distance anaphor in this language also depends on the case properties of the binder and bindee, we merely sketch an analysis here. The left rule for long distance anaphors is given in Fig. 21.

segir 'say': $(\langle n \rangle \langle s \rangle / \lceil s \rfloor)$ víta 'know': $(\langle n \rangle \langle s \rangle / s)$ elskar 'love': $(\langle n \rangle \langle s \rangle / n)$

- (41) Jón₁ segir að María₂ elski sig₁.
 John say that Maria love.SUBJ SE-ANAPHOR.ACC
 'John says that Mary loves him.'
- (42) ?Jón₁ veit að María elskar sig₁.
 John know that Maria love.IND SE-ANAPHOR.ACC
 'John knows that Mary loves him.'

 $\frac{X, [A], Z_1, \{Z_2, B, W\} \Rightarrow C}{X, [A], Z_1, \{Z_2, B|A, W\} \Rightarrow C} |\mathcal{L}_{lg}$

Fig. 21. Rule for long-distance anaphor

5 Conclusions

In this paper we have proposed different rules to deal with anaphoric and pronominal pronouns occurring in several syntactic domains. Although both the type assignment for pronouns and our initial idea for the pronominal rules come from Jäeger [8], we have proposed a different type assignment for reflexives and pronominal pronouns and we have modified the rules of **LLC**. The inspiration for lexical entries encoding marked argument positions comes from [18]. By adopting bracket modalities we have identified different syntactic domains; in light of the latter, we have encoded binding restrictions into the left anaphoric rule of **LLBE**. The right pronominal rule of **LLBE**, in turn, evidences that despite the fact that an antecedent A could occur in the local syntactic domain [], a free pronoun is derived by assuming an antecedent in a non-local domain [].

The rules of **LLBE** reveal then that free and anaphoric pronouns on the one hand, and bound pronouns on the other, are generally processed in different steps in a proof: if $X_1 = X_2 = \epsilon$ and so the antecedent A is left-peripheral, free pronouns and reflexives, but not bound pronouns, can be inserted into a derivation. Our proposal preserves the prominence condition on the binder for reflexives: the binder may not be an argument lower in the hierarchy and neither may it be part of an argument higher in the hierarchy. In addition, we have incorporated the previous modal categorial analysis for *say* in terms of structural modalities, in accordance with our overall proposal. Further, we have suggested how this proposal can be used to deal with long-distance anaphors in Icelandic. Our rule for object-oriented reflexives, such as anaphors in possessive complements.

In future work, we propose to investigate how to impose structural conditions upon the sequences of the left rules for the customized slash type-constructors in order to reduce the number of pronominal rules and thus to deal with subject and object anaphoric pronouns in a more uniform way. We also plan to explore how to deal with ECM constructions.

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Appendix

The proof for the Cut elimination theorem requires the use of the following bracketed versions of the structural rules of Permutation and Expansion (Fig. 22). In order to prove Cut Elimination for **LLBE** we have to consider two more cases for principal Cut: the left premise of Cut is the conclusion of $||L_a|$ or that of $||L_p|$ and the right premise is the conclusion of ||R. These two configurations are given schematically in Figs. 23 and 24. In both cases, the principal Cut is replaced by a Cut of lower degree. Since no rule introduces a formula [|A|] into the right side of a sequent (i.e. there are only antecedent occurrences of the formula [|A|]), the Cut formula could not have been derived by applying either of the bracketed structural rules.

$$\frac{X, A, Y, Z \Rightarrow C}{X, A, Y, \llbracket A \rVert, Z \Rightarrow C} \llbracket E \rrbracket \qquad \qquad \frac{X, \llbracket A \rrbracket, Y, Z \Rightarrow C}{X, Y, \llbracket A \rVert, Z \Rightarrow C} \llbracket P \rrbracket$$

Fig. 22. Bracketed structural rules

$$\frac{Z_1 \Rightarrow C}{\llbracket A \rrbracket, Z_1 \Rightarrow C \Vert A} \Vert R \qquad \frac{X_2, A, Z_2, C, W_2 \Rightarrow D}{X_2, A, Z_2, C \Vert A, W_2 \Rightarrow D} \Vert L_a \\
\xrightarrow{X_2, A, Z_2, \llbracket A \rrbracket, Z_1, W_2 \Rightarrow D} Cut \\
\xrightarrow{\sim} \\
\frac{Z_1 \Rightarrow C \qquad X_2, A, Z_2, C, W_2 \Rightarrow D}{X_2, A, Z_2, Z_1, W_2 \Rightarrow D} Cut \\
\xrightarrow{X_2, A, Z_2, \llbracket A \rrbracket, Z_1, W_2 \Rightarrow D} \Vert E \rfloor$$

Fig. 23. Principal cut for $\parallel : \parallel L_a$

$$\frac{Z_1 \Rightarrow C}{\llbracket A \rrbracket, Z_1 \Rightarrow C \Vert A} \Vert \mathbf{R} \qquad \frac{X_2, \llbracket A \rrbracket, Z_2, C, W_2 \Rightarrow D}{X_2, Z_2, C \Vert A, W_2 \Rightarrow D} \Vert \mathbf{L}_p$$
$$\frac{X_2, Z_2, \llbracket A \rrbracket, Z_1, W_2 \Rightarrow D}{X_2, Z_2, \llbracket A \rrbracket, Z_1, W_2 \Rightarrow D} Cut$$

$$\frac{Z_1 \Rightarrow C \qquad X_2, [\llbracket A \rrbracket, Z_2, C, W_2 \Rightarrow D}{X_2, \llbracket A \rrbracket, Z_2, Z_1, W_2 \Rightarrow D} Cut$$

$$\frac{X_2, \llbracket A \rrbracket, Z_2, \llbracket A \rrbracket, Z_1, W_2 \Rightarrow D}{X_2, Z_2, \llbracket A \rrbracket, Z_1, W_2 \Rightarrow D} [\llbracket P \rrbracket$$

Fig. 24. Principal cut for $\parallel : \parallel \mathbf{L}_p$

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