On Two Concepts of Ultrafilter Extensions of First-Order Models and Their Generalizations

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Abstract. There exist two known concepts of ultrafilter extensions of first-order models, both in a certain sense canonical. One of them [1] comes from modal logic and universal algebra, and in fact goes back to [2]. Another one [3,4] comes from model theory and algebra of ultrafilters, with ultrafilter extensions of semigroups [5] as its main precursor. By a classical fact, the space of ultrafilters over a discrete space is its largest compactification. The main result of [3,4], which confirms a canonicity of this extension, generalizes this fact to discrete spaces endowed with a first-order structure. An analogous result for the former type of ultrafilter extensions was obtained in [6].

Here we offer a uniform approach to both types of extensions. It is based on the idea to extend the extension procedure itself. We propose a generalization of the standard concept of first-order models in which functional and relational symbols are interpreted rather by ultrafilters over sets of functions and relations than by functions and relations themselves. We provide two specific operations which turn generalized models into ordinary ones, and establish necessary and sufficient conditions under which the latter are the two canonical ultrafilter extensions of some models.

1. Fix a first-order language and consider an arbitrary model

$$\mathfrak{A} = (X, F, \dots, R, \dots)$$

with the universe X, operations F, \ldots , and relations R, \ldots . Let us define an abstract *ultrafilter extension* of \mathfrak{A} as a model \mathfrak{A}' (in the same language) of form

$$\mathfrak{A}' = (\beta X, F', \dots, R', \dots)$$

where βX is the set of ultrafilters over X (one lets $X \subseteq \beta X$ by identifying each $x \in X$ with the principal ultrafilter given by x), and operations F', \ldots and relations R', \ldots on βX extend F, \ldots and R, \ldots resp. There are essentially two

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known ways to extend relations by ultrafilters, and *one* to extend maps. Partial cases of these extensions were discovered by various authors in different time and different areas, typically, without a knowledge of parallel studies in adjacent areas.

Recall that βX carries a natural topology generated by basic open sets

$$\widetilde{A} = \{ \mathfrak{u} \in \beta X : A \in \mathfrak{u} \}$$

for all $A \subseteq X$. Easily, the sets are also closed, so the space βX is zerodimensional. In fact, βX is compact, Hausdorff, extremally disconnected (the closure of any open set is open), and the largest compactification of the discrete space X. This means that X is dense in βX and every (trivially continuous) map h of X into any compact Hausdorff space Y uniquely extends to a continuous map \tilde{h} of βX into Y:



The largest compactification of Tychonoff spaces was discovered independently by Čech [7] and Stone [8]; then Wallman [9] did the same for T_1 spaces (by using ultrafilters on lattices of closed sets); see [5,10,11] for more information.

The ultrafilter extensions of *unary* maps F and relations R are exactly \tilde{F} and \tilde{R} (for $F: X \to X$ let $Y = \beta X$); thus in the unary case the procedure gives classical objects known in 30s. As for mappings and relations of greater arities, several instances of their ultrafilter extensions were discovered only in 60s.

Studying ultraproducts, Kochen [12] and Frayne et al. [13] considered a "multiplication" of ultrafilters, which actually is the ultrafilter extension of the *n*-ary operation of taking *n*-tuples. They shown that the successive iteration of ultrapowers by ultrafilters u_1, \ldots, u_n is isomorphic to a single ultrapower by their "product". This has leaded to the general construction of iterated ultrapowers, invented by Gaifman and elaborated by Kunen, which has become common in model theory and set theory (see [14, 15]).

Ultrafilter extensions of semigroups appeared in 60s as subspaces of function spaces; the first explicit construction of the ultrafilter extension of a group is due to Ellis [16]. In 70s Galvin and Glazer applied them to give an easy proof of what now known as Hindman's Finite Sums Theorem; the key idea was to use idempotent ultrafilters. The method was developed then by Blass, van Douwen, Hindman, Protasov, Strauss, and many others, and gave numerous Ramsey-theoretic applications in number theory, algebra, topological dynamics, and ergodic theory. The book [5] is a comprehensive treatise of this area, with an historical information. This technique was applied also for obtaining analogous results for certain non-associative algebras (see [17, 18]).

Ultrafilter extensions of arbitrary *n*-ary maps have been introduced independently by Goranko [1] and Saveliev [3,4]. For $F: X_1 \times \ldots \times X_n \to Y$, the extended map $\tilde{F}: \beta X_1 \times \ldots \times \beta X_n \to \beta Y$ is defined by letting

$$\widetilde{F}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n) = \{A \subseteq Y : \{x_1 \in X_1 : \ldots \{x_n \in X_n : F(x_1,\ldots,x_n) \in A\} \in \mathfrak{u}_n \ldots\} \in \mathfrak{u}_1 \}.$$

One can simplify this cumbersome notation by introducing *ultrafilter quantifiers*: let $(\forall^{\mathfrak{u}}x) \varphi(x,...)$ means $\{x : \varphi(x,...)\} \in \mathfrak{u}$. In fact, this is a second-order quantifier: $(\forall^{\mathfrak{u}}x)$ is equivalent to $(\forall A \in \mathfrak{u})(\exists x \in A)$, and also (since \mathfrak{u} is ultra) to $(\exists A \in \mathfrak{u})(\forall x \in A)$. Such quantifiers are self-dual, i.e. $\forall^{\mathfrak{u}}$ and $\exists^{\mathfrak{u}}$ coincide, and generally do not commute with each other, i.e. $(\forall^{\mathfrak{u}}x)(\forall^{\mathfrak{v}}y)$ and $(\forall^{\mathfrak{v}}y)(\forall^{\mathfrak{u}}x)$ are not equivalent. Then the definition above is rewritten as follows:

 $\widetilde{F}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n)=\big\{A\subseteq Y:(\forall^{\mathfrak{u}_1}x_1)\ldots(\forall^{\mathfrak{u}_n}x_n)\ F(x_1,\ldots,x_n)\in A\big\}.$

The map \widetilde{F} can be also described as the composition of the ultrafilter extension of taking *n*-tuples, which maps $\beta X_1 \times \ldots \times \beta X_n$ into $\beta (X_1 \times \ldots \times X_n)$, and the continuous extension of F considered as a unary map, which maps $\beta (X_1 \times \ldots \times X_n)$ into βY .

One type of ultrafilter extensions of relations goes back to a seminal paper by Jónsson and Tarski [2] where they have been appeared implicitly, in terms of representations of Boolean algebras with operators. For binary relations, their representation theory was rediscovered in modal logic by Lemmon [19] who credited much of this work to Scott, see footnote 6 on p. 204 (see also [20]). Goldblatt and Thomason [21] used this to characterize modal definability (where Sect. 2 was entirely due to Goldblatt); the term "ultrafilter extension" has been introduced probably in the subsequent work by van Benthem [22] (for modal definability see also [23, 24]). Later Goldblatt [25] generalized the extension to n-ary relations.

Let us give an equivalent formulation: for $R \subseteq X_1 \times \ldots \times X_n$, the extended relation $R^* \subseteq \beta X_1 \times \ldots \times \beta X_n$ is defined by letting

$$R^*(\mathfrak{u}_1,\ldots,\mathfrak{u}_n) \text{ iff}$$
$$(\forall A_1 \in \mathfrak{u}_1)\ldots(\forall A_n \in \mathfrak{u}_n)(\exists x_1 \in A_1)\ldots(\exists x_n \in A_n) R(x_1,\ldots,x_n)$$

Another type of ultrafilter extensions of n-ary relations has been recently discovered in [3, 4]:

$$\widetilde{R}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n) \text{ iff} \\ \left\{ x_1 \in X_1 : \ldots \{ x_n \in X_n : R(x_1,\ldots,x_n) \} \in \mathfrak{u}_n \ldots \right\} \in \mathfrak{u}_1,$$

or rewritting this via ultrafilter quantifiers,

$$\widetilde{R}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n)$$
 iff $(\forall^{\mathfrak{u}_1}x_1)\ldots(\forall^{\mathfrak{u}_n}x_n) R(x_1,\ldots,x_n).$

Or else, by decoding ultrafilter quantifiers, this can be rewritten by

$$\widetilde{R}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n) \text{ iff} (\forall A_1 \in \mathfrak{u}_1)(\exists x_1 \in A_1)\ldots(\forall A_n \in \mathfrak{u}_n)(\exists x_n \in A_n) R(x_1,\ldots,x_n),$$

whence it is clear that $\widetilde{R} \subseteq R^*$. For unary R both extensions coincide with the basic open set given by R. If R is functional then R^* (but not \widetilde{R}) coincides with the above-defined extension of R as a map. An easy instance of \sim -extensions (where R are linear orders) is studied in [26].

A systematic comparative study of both extensions (for binary R) is undertaken in [6]. In particular, there is shown that the * - and the ~ -extensions have a dual character w.r.t. relation-algebraic operations: the former commutes with composition and inversion but not Boolean operations except for union, while the latter commutes with all Boolean operations but neither composition nor inversion. Also [6] contains topological characterizations of \tilde{R} and R^* in terms of appropriate closure operations and in terms of Vietoris-type topologies (regarding R as multi-valued maps).

Ultrafilter extensions of arbitrary first-order models were considered for the first time independently in [1] with *-extensions of relations, and in [3] with their \sim -extensions. We shall denote them by \mathfrak{A}^* and $\widetilde{\mathfrak{A}}$ resp. Thus for a model $\mathfrak{A} = (X, F, \ldots, R, \ldots)$ we let

$$\mathfrak{A}^* = (\beta X, \widetilde{F}, \dots, R^*, \dots) \text{ and } \widetilde{\mathfrak{A}} = (\beta X, \widetilde{F}, \dots, \widetilde{R}, \dots)$$

The following is the main result of [1]:

Theorem 1. If h is a homomorphism between models \mathfrak{A} and \mathfrak{B} , then the continuous extension \tilde{h} is a homomorphism between \mathfrak{A}^* and \mathfrak{B}^* :



A full analog of Theorem 1 for the \sim -extensions has been appeared in [3] (called the First Extension Theorem in [4]):

Theorem 2. If h is a homomorphism between models \mathfrak{A} and \mathfrak{B} , then the continuous extension \widetilde{h} is a homomorphism between $\widetilde{\mathfrak{A}}$ and $\widetilde{\mathfrak{B}}$:

Moreover, both theorems remain true for embeddings and some other model-theoretic interrelations (see [1,3,4]).

Theorem 2 is actually is a partial case of a much stronger result of [3] (called the Second Extension Theorem in [4]). To formulate this, we need the following concepts (introduced in [3]).

Let X_1, \ldots, X_n, Y be topological spaces, and let $A_1 \subseteq X_1, \ldots, A_{n-1} \subseteq X_{n-1}$. An *n*-ary function $F: X_1 \times \ldots \times X_n \to Y$ is right continuous w.r.t. A_1, \ldots, A_{n-1} iff for each $i, 1 \leq i \leq n$, and every $a_1 \in A_1, \ldots, a_{i-1} \in A_{i-1}$ and $x_{i+1} \in X_{i+1}, \ldots, x_n \in X_n$, the map

$$x \mapsto F(a_1, \ldots, a_{i-1}, x, x_{i+1}, \ldots, x_n)$$

of X_i into Y is continuous. An n-ary relation $R \subseteq X_1 \times \ldots \times X_n$ is right open (right closed, etc.) w.r.t. A_1, \ldots, A_{n-1} iff for each $i, 1 \leq i \leq n$, and every $a_1 \in A_1, \ldots, a_{i-1} \in A_{i-1}$ and $x_{i+1} \in X_{i+1}, \ldots, x_n \in X_n$, the set

$$\{x \in X_i : R(a_1, \dots, a_{i-1}, x, x_{i+1}, \dots, x_n)\}$$

is open (closed, etc.) in X_i .

Theorem 3 [3,4] characterizes topological properties of \sim -extensions, it is a base of Theorem 4 (the Second Extension Theorem of [4]).

Theorem 3. Let \mathfrak{A} be a model. In the extension \mathfrak{A} , all operations are right continuous and all relations right clopen w.r.t. the universe of \mathfrak{A} .

Theorem 4. Let \mathfrak{A} and \mathfrak{C} be two models, h a homomorphism of \mathfrak{A} into \mathfrak{C} , and let \mathfrak{C} carry a compact Hausdorff topology in which all operations are right continuous and all relations are right closed w.r.t. the image of the universe of \mathfrak{A} under h. Then \tilde{h} is a homomorphism of $\widetilde{\mathfrak{A}}$ into \mathfrak{C} :



Theorem 2 (for homomorphisms) easily follows: take \mathfrak{B} as \mathfrak{C} . The meaning of Theorem 4 is that it generalizes the classical Čech–Stone result to the case when the underlying discrete space X carries an arbitrary first-order structure.

A natural question is whether *-extensions also canonical in a similar sense. The answer is positive; two following theorems are counterparts of Theorems 3 and 4 resp. (essentially both have been proved in [6]).

Theorem 5. Let \mathfrak{A} be a model. In the extension \mathfrak{A}^* , all relations are closed (and all operations are right continuous w.r.t. the universe of \mathfrak{A}).

Theorem 6. Let \mathfrak{A} and \mathfrak{C} be two models, h a homomorphism of \mathfrak{A} into \mathfrak{C} , and let \mathfrak{C} carry a compact Hausdorff topology in which all operations are right continuous w.r.t. the image of the universe of \mathfrak{A} under h, and all relations are closed. Then \tilde{h} is a homomorphism of \mathfrak{A}^* into \mathfrak{C} .

$$\begin{array}{c} \mathfrak{A}^* \\ \uparrow & \searrow & \widetilde{h} \\ \mathfrak{A} & \xrightarrow{h} & \mathfrak{C} \end{array}$$

Similarly, Theorem 1 (for homomorphisms) follows from Theorem 6. The latter also generalizes the Čech–Stone result for discrete spaces to discrete models but with a narrow class of target models \mathfrak{C} : having relations rather closed than right closed in Theorem 4.

2. The immediate purpose of this section is to provide a uniform approach to both types of extensions. This approach will lead us to certain structures, called here generalized models, which generalize ultrafilter extensions of each of the two types.

First we shall show that the *-extension can be described in terms of the basic (cl)open sets and the continuous extension of maps. For this, let us consider the continuous extension of the continuous extension operation *itself*. To make notation easier, denote by ext the operation of continuous extension of maps; i.e. ext(f) is another notation for \tilde{f} :

$$\operatorname{ext}(f) = \widetilde{f}.$$

So if we consider maps of X into Y, then ext is a map of Y^X into $C(\beta X, \beta Y)$. Since $C(\beta X, \beta Y)$ with the standard (i.e. pointwise convergence) topology is a compact Hausdorff space, ext continuously extends to the map ext of $\beta(Y^X)$ into this space:

The extended map ext is surjective and non-injective.

Lemma 1. Let $R \subseteq Y^X$. Then ext maps the closure of R in the space $\beta(Y^X)$ onto the closure of R in the space $C(\beta X, \beta Y)$:

$$\left\{\widetilde{\operatorname{ext}}(\mathfrak{f}):\mathfrak{f}\in\operatorname{cl}_{\beta(Y^X)}R\right\}=\operatorname{cl}_{C(\beta X,\beta Y)}R.$$

For our purpose, let X = n. Then $\beta X = n$ and $C(\beta X, \beta Y)$ is $(\beta Y)^n$, which can be identified with $\beta Y \times \ldots \times \beta Y$ (*n* times). Now the required description of the *-extension follows from Theorem 5:

Theorem 7. Let $R \subseteq X \times \ldots \times X$. Then $R^* \subseteq \beta X \times \ldots \times \beta X$ is (identified with) the image of $cl_{\beta(X^n)}R$ under ext.

Using ultrafilters over maps leads to the following concept. Given a language, we define a generalized (or ultrafilter) interpretation (the term is ad hoc) as a map i that takes each *n*-ary functional symbol F to an ultrafilter over the set of *n*-ary operations on X, and each *n*-ary predicate symbol R to an ultrafilter over the set of *n*-ary relations on X; let also v be an ultrafilter valuation of variables, i.e. a valuation which takes each variable x to an ultrafilter over a given set X:

$$v(x) \in \beta X, \quad i(F) \in \beta(X^{X \times \dots \times X}), \quad i(R) \in \beta \mathcal{P}(X \times \dots \times X).$$

The set $(\beta X, \iota(F), \ldots, \iota(R), \ldots)$ is a generalized model. Now we are going to define the satisfiability relation in generalized models, which will be denoted by the symbol \models .

First, given an interpretation i of non-logical symbols, we expand any valuation v of variables to the map v_i defined on all terms as follows. Let app : $X_1 \times \ldots \times X_n \times Y^{X_1 \times \ldots \times X_n} \to Y$ be the *application* operation:

$$\operatorname{app}(a_1,\ldots,a_n,f) = f(a_1,\ldots,a_n)$$

Extend it to the map $\widetilde{\text{app}} : \beta X_1 \times \ldots \times \beta X_n \times \beta (Y^{X_1 \times \ldots \times X_n}) \to \beta Y$ right continuous w.r.t. the principal ultrafilters, in the usual way:

$$\beta X_1 \times \ldots \times \beta X_n \times \beta (Y^{X_1 \times \ldots \times X_n}) \xrightarrow{\operatorname{app}} \beta Y$$

$$\uparrow \qquad \qquad \uparrow \qquad \qquad \uparrow$$

$$X_1 \times \ldots \times X_n \times Y^{X_1 \times \ldots \times X_n} \xrightarrow{\operatorname{app}} Y$$

Let v_i coincide with v on variables, and if v_i has been already defined on terms t_1, \ldots, t_n , we let

$$v_{\iota}(F(t_1,\ldots,t_n)) = \widetilde{\operatorname{app}}(v_{\iota}(t_1),\ldots,v_{\iota}(t_n),\iota(F)).$$

Further, given a generalized model $\mathfrak{A} = (\beta X, \iota(F), \ldots, \iota(R), \ldots)$, define the satisfiability in \mathfrak{A} as follows. Let in $\subseteq X_1 \times \ldots \times X_n \times \mathcal{P}(X_1 \times \ldots \times X_n)$ be the *membership* predicate:

in
$$(a_1, \ldots, a_n, R)$$
 iff $(a_1, \ldots, a_n) \in R$.

Extend it to the relation $\widetilde{in} \subseteq \beta X_1 \times \ldots \times \beta X_n \times \beta \mathcal{P}(X_1 \times \ldots \times X_n)$ right clopen w.r.t. principal ultrafilters. Let

$$\mathfrak{A} \models t_1 = t_2 [v] \text{ iff } v_i(t_1) = v_i(t_2).$$

If $R(t_1, \ldots, t_n)$ is an atomic formula in which R is not the equality predicate, we let

$$\mathfrak{A} \models R(t_1, \dots, t_n) [v] \text{ iff } \widetilde{\mathrm{in}} (v_i(t_1), \dots, v_i(t_n), i(P))$$

(Equivalently, we could define the satisfiability of atomic formulas by identifying predicates with their characteristic functions and using the satisfiability of equalities of the resulting terms.) Finally, if $\varphi(t_1, \ldots, t_n)$ is obtained by negation, conjunction, or quantification from formulas for which \models has been already defined, we define $\mathfrak{A} \models \varphi[v]$ in the standard way.

When needed, we shall use variants of notation commonly used for ordinary models and satisfiability, for the generalized ones. E.g. for a generalized model \mathfrak{A} with the universe βX , a formula $\varphi(x_1, \ldots, x_n)$, and elements $\mathfrak{u}_1, \ldots, \mathfrak{u}_n$ of βX , the notation $\mathfrak{A} \models \varphi[\mathfrak{u}_1, \ldots, \mathfrak{u}_n]$ means that φ is satisfied in \mathfrak{A} under a valuation taking the variables x_1, \ldots, x_n to the ultrafilters $\mathfrak{u}_1, \ldots, \mathfrak{u}_n$.

Generalized models actually generalize not all ordinary models but those that are ultrafilter extensions of some models. It is worth also pointing out that whenever a generalized interpretation is *principal*, i.e. all non-logical symbols are interpreted by principal ultrafilters, we naturally identify it with the obvious ordinary interpretation with the same universe βX ; however, not every ordinary interpretation with the universe βX is of this form. Precise relationships between generalized models, ordinary models, and ultrafilter extensions will be described in Theorems 9 and 10.

An ultrafilter valuation v is *principal* iff it takes any variable to a principal ultrafilter.

Lemma 2. Let two generalized models $\mathfrak{A} = (\beta X, \iota(F), \ldots, \iota(R), \ldots)$ and $\mathfrak{B} = (\beta X, \jmath(F), \ldots, \jmath(R), \ldots)$ have the same universe βX . If for all functional symbols F, predicate symbols R, variables x_1, \ldots, x_n , and principal valuations v,

$$\widetilde{\operatorname{app}}(v(x_1),\ldots,v(x_n),\iota(F)) = \widetilde{\operatorname{app}}(v(x_1),\ldots,v(x_n),\mathfrak{g}(F)),$$

$$\widetilde{\operatorname{in}}(v(x_1),\ldots,v(x_n),\iota(R)) \quad iff \quad \widetilde{\operatorname{in}}(v(x_1),\ldots,v(x_n),\mathfrak{g}(R)),$$

then for all formulas φ , terms t_1, \ldots, t_n , and valuations v,

$$\mathfrak{A} \models \varphi(t_1, \ldots, t_n)[v] \text{ iff } \mathfrak{B} \models \varphi(t_1, \ldots, t_n)[v].$$

Corollary 1. Let $\mathfrak{A} = (\beta X, \imath(F), \ldots, \imath(R), \ldots)$ be a generalized model and $\mathfrak{B} = (\beta X, \jmath(F), \ldots, \jmath(R), \ldots)$ the generalized model having the same universe βX and such that \jmath coincides with \imath on functional symbols and for each predicate symbol $R, \jmath(R)$ is the principal ultrafilter given by

$$\{(a_1,\ldots,a_n)\in X^n: in(a_1,\ldots,a_n,i(R))\}.$$

Then for all valuations v, formulas φ , and terms t_1, \ldots, t_n ,

$$\mathfrak{A} \models \varphi(t_1, \ldots, t_n) [v] \text{ iff } \mathfrak{B} \models \varphi(t_1, \ldots, t_n) [v].$$

Let us say that an ultrafilter \mathfrak{f} over functions is *pseudo-principal* iff $\widetilde{\text{app}}$ takes any tuple consisting of principal ultrafilters together with \mathfrak{f} to a principal ultrafilter, i.e. for $\mathfrak{f} \in \beta(Y^{X_1 \times \ldots \times X_n})$,

 $a_1 \in X_1, \ldots, a_n \in X_n$ implies $\widetilde{\operatorname{app}}(a_1, \ldots, a_n, \mathfrak{f}) \in Y$.

Every principal f is pseudo-principal, and there exist pseudo-principal ultrafilters that are not principal as well as ultrafilters that are not pseudo-principal. A generalized interpretation i is *pseudo-principal on functional symbols* iff i(F) is a pseudo-principal ultrafilter for each functional symbol F (and then, for each term t).

Corollary 2. Let $\mathfrak{A} = (\beta X, \iota(F), \ldots, \iota(R), \ldots)$ be a generalized model with ι pseudo-principal on functional symbols. Let $\mathfrak{B} = (\beta X, \jmath(F), \ldots, \jmath(R), \ldots)$ be the generalized model having the same universe βX and such that \jmath coincides with ι on predicate symbols and for each functional symbol F, $\jmath(F)$ is the principal ultrafilter given by $f: X^n \to X$ defined by letting

$$f(a_1,\ldots,a_n) = \widetilde{\operatorname{app}}(a_1,\ldots,a_n,\imath(F)).$$

Then for all valuations v, formulas φ , and terms t_1, \ldots, t_n ,

$$\mathfrak{A} \models \varphi(t_1, \ldots, t_n) [v] \text{ iff } \mathfrak{B} \models \varphi(t_1, \ldots, t_n) [v].$$

It follows that for any generalized model \mathfrak{A} whose interpretation is pseudoprincipal on functional symbols, by replacing its relations as in Corollary 1 and its operations as in Corollary 2, one obtains an ordinary model \mathfrak{B} with the same universe such that for all formulas φ and elements $\mathfrak{u}_1, \ldots, \mathfrak{u}_n$ of the universe, $\mathfrak{A} \models \varphi [\mathfrak{u}_1, \ldots, \mathfrak{u}_n]$ iff $\mathfrak{B} \models \varphi [\mathfrak{u}_1, \ldots, \mathfrak{u}_n]$.

We do not formulate this fact as a separate theorem since we shall be able to establish stronger facts soon. In Theorem 8, we shall establish that for any generalized model \mathfrak{A} , not only one with a pseudo-principal interpretation, one can construct a certain ordinary model $e(\mathfrak{A})$ satisfying the same formulas; and then, in Theorem 9, that whenever \mathfrak{A} has a pseudo-principal interpretation, $e(\mathfrak{A})$ is nothing but the $\tilde{}$ -extension of some model. In fact, in the latter case, $e(\mathfrak{A})$ coincides with \mathfrak{B} from the previous paragraph.

Now we provide two operations, e and E, which turn generalized models into certain ordinary models that generalize * - and ~ -extensions. Both operations are surjective and non-injective.

Define a map e on ultrafilters over functions to functions over ultrafilters,

$$e:\beta(Y^{X_1\times\ldots\times X_n})\to\beta Y^{\beta X_1\times\ldots\times\beta X_n},$$

by induction on *n*. For n = 1, let *e* coincide with ext. Assume that *e* has been already defined for *n*. First we identify $Y^{X_1 \times X_2 \times \ldots \times X_{n+1}}$ with $(Y^{X_2 \times \ldots \times X_{n+1}})^{X_1}$ (by the so-called evaluation map, or carrying). Under this identification, each $\mathfrak{f} \in \beta(Y^{X_1 \times X_2 \times \ldots \times X_{n+1}})$ corresponds to a certain $\mathfrak{f}' \in \beta((Y^{X_2 \times \ldots \times X_{n+1}})^{X_1})$. Now we define $e(\mathfrak{f})$ by letting

$$e(\mathfrak{f})(\mathfrak{u}_1,\mathfrak{u}_2,\ldots,\mathfrak{u}_{n+1})=e(e(\mathfrak{f}')(\mathfrak{u}_1))(\mathfrak{u}_2,\ldots,\mathfrak{u}_{n+1})$$

(since e has been already defined on \mathfrak{f}' and $e(\mathfrak{f}')(\mathfrak{u}_1)$ by induction hypothesis).

Alternatively, we can define e as follows. Expand the domain of ext by letting

$$\operatorname{ext}(f) = \widetilde{f}$$

for *n*-ary functions f with any n, not only unary ones. Thus, if we consider functions of $X_1 \times \ldots \times X_n$ into Y, then ext maps $Y^{X_1 \times \ldots \times X_n}$ into $RC_{X_1,\ldots,X_{n-1}}(\beta X_1 \times \ldots \times \beta X_n,\beta Y)$, the set of all functions of $\beta X_1 \times \ldots \times \beta X_n$ into βY that are right continuous w.r.t. X_1,\ldots,X_{n-1} . It can be shown that the latter set forms a closed subspace in the compact Hausdorff space $\beta Y^{\beta X_1 \times \ldots \times \beta X_n}$ of all functions of $\beta X_1 \times \ldots \times \beta X_n$ into βY with the standard (i.e. pointwise convergence) topology, and hence, is compact Hausdorff too. Therefore, ext continuously extends to the map ext of $\beta (Y^{X_1 \times \ldots \times X_n})$ into it:

Now we can identify e with ext in this expanded meaning.

By identifying relations with their characteristic functions, we can also let that e takes ultrafilters over relations to relations over ultrafilters:

$$e:\beta \mathcal{P}(X_1 \times \ldots \times X_n) \to \mathcal{P}(\beta X_1 \times \ldots \times \beta X_n).$$

In fact, e and $\widetilde{\text{app}}$ (or in) are expressed via each other:

Lemma 3. For all $\mathfrak{f} \in \beta(Y^{X_1 \times \ldots \times X_n})$, $\mathfrak{r} \in \beta \mathcal{P}(X_1 \times \ldots \times X_n)$, and $\mathfrak{u}_1 \in \beta X_1, \ldots, \mathfrak{u}_n \in \beta X_n$,

$$e(\mathfrak{f})(\mathfrak{u}_1,\ldots,\mathfrak{u}_n) = \widetilde{\operatorname{app}}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n,\mathfrak{f}),$$
$$e(\mathfrak{r})(\mathfrak{u}_1,\ldots,\mathfrak{u}_n) \quad iff \quad i\widetilde{n}(\mathfrak{u}_1,\ldots,\mathfrak{u}_n,\mathfrak{r}).$$

In other words,

$$e(\mathfrak{f}) = \big\{ (\mathfrak{u}_1, \dots, \mathfrak{u}_n, \mathfrak{v}) \in \beta X_1 \times \dots \times \beta X_n \times \beta Y : \widetilde{\operatorname{app}}(\mathfrak{u}_1, \dots, \mathfrak{u}_n, \mathfrak{f}) = \mathfrak{v} \big\},\\ e(\mathfrak{r}) = \big\{ (\mathfrak{u}_1, \dots, \mathfrak{u}_n) \in \beta X_1 \times \dots \times \beta X_n : \widetilde{\operatorname{in}}(\mathfrak{u}_1, \dots, \mathfrak{u}_n, \mathfrak{r}) \big\}.$$

Corollary 3. For all generalized models $\mathfrak{A} = (\beta X, \imath(F), \ldots, \imath(R), \ldots)$ and valuations v,

$$v_{\iota}(F(t_1,\ldots,t_n)) = e(\iota(F))(v_{\iota}(t_1),\ldots,v_{\iota}(t_n)),$$

$$\mathfrak{A} \models R(t_1,\ldots,t_n) [v] \quad iff \quad e(\iota(R))(v_{\iota}(t_1),\ldots,v_{\iota}(t_n)).$$

For a generalized model $\mathfrak{B} = (\beta X, \mathfrak{f}, \dots, \mathfrak{r}, \dots)$, let

$$e(\mathfrak{B}) = (\beta X, e(\mathfrak{f}), \dots, e(\mathfrak{r}), \dots).$$

Note that $e(\mathfrak{B})$ is an ordinary model.

Theorem 8. If \mathfrak{A} is a generalized model, then for all formulas φ and elements $\mathfrak{u}_1, \ldots, \mathfrak{u}_n$ of the universe of \mathfrak{A} ,

$$\mathfrak{A} \models \varphi [\mathfrak{u}_1, \dots, \mathfrak{u}_n] \quad iff \ e(\mathfrak{A}) \models \varphi [\mathfrak{u}_1, \dots, \mathfrak{u}_n].$$

Define a map E, with the same domain and range that the map e has, as follows: E and e coincide on $\beta(Y^{X_1 \times \ldots \times X_n})$, and if $\mathfrak{r} \in \beta \mathcal{P}(X_1 \times \ldots \times X_n)$ then

$$E(\mathfrak{r}) = \{ \widetilde{\operatorname{ext}}(\mathfrak{q}) : \mathfrak{q} \in \widetilde{\operatorname{ext}}(\mathfrak{r}) \}.$$

Here $\widetilde{\operatorname{ext}}(\mathfrak{r})$ is a clopen subset of $\beta(X_1 \times \ldots \times X_n)$, if $\mathfrak{q} \in \widetilde{\operatorname{ext}}(\mathfrak{r})$ then $\widetilde{\operatorname{ext}}(\mathfrak{q})$ is identified with an element of the space $\beta X_1 \times \ldots \times \beta X_n$ (as in Theorem 7), and the resulting $E(\mathfrak{r})$ is closed in the space.

Lemma 4. Let $\mathfrak{r} \in \beta \mathcal{P}(X_1 \times \ldots \times X_n)$. Then

$$e(\mathbf{r}) = \widetilde{R} \text{ and } E(\mathbf{r}) = R^*$$

for $R = e(\mathfrak{r}) \cap (X_1 \times \ldots \times X_n) = E(\mathfrak{r}) \cap (X_1 \times \ldots \times X_n) = \bigcap_{S \in \mathfrak{r}} \bigcup S.$

One may write up this R more explicitly:

$$R = \{(a_1, \ldots, a_n) \in X_1 \times \ldots \times X_n : (\forall S \in \mathfrak{r}) \ (\exists Q \in S) \ Q(a_1, \ldots, a_n)\}.$$

For a generalized model $\mathfrak{B} = (\beta X, \mathfrak{f}, \dots, \mathfrak{r}, \dots)$, let

$$E(\mathfrak{B}) = (\beta X, E(\mathfrak{f}), \dots, E(\mathfrak{r}), \dots).$$

Then $E(\mathfrak{B})$, like $e(\mathfrak{B})$, is an ordinary model.

By Lemma 4, relations of the model $e(\mathfrak{B})$ are \sim -extensions of some relations on X, while relations of the model $E(\mathfrak{B})$ are *-extensions of the same relations. Whether the whole models $e(\mathfrak{B})$ and $E(\mathfrak{B})$ are ultrafilter extensions of some models depends only on the (generalized) interpretation of functional symbols in \mathfrak{B} :

Theorem 9. Let \mathfrak{B} be a generalized model with the universe βX . The following are equivalent:

- (i) $e(\mathfrak{B}) = \widetilde{\mathfrak{A}}$ for a model \mathfrak{A} with the universe X,
- (ii) $E(\mathfrak{B}) = \mathfrak{A}^*$ for a model \mathfrak{A} with the universe X,
- (iii) The interpretation in \mathfrak{B} is pseudo-principal on functional symbols.

Moreover, the model \mathfrak{A} in (i) and (ii) is the same.

Finally, we point out that the fact whether an ordinary model with the universe βX is of form $e(\mathfrak{B})$, and whether it is of form $E(\mathfrak{B})$, for some generalized model \mathfrak{B} (clearly, with the same universe βX) depends only on its topological properties:

Theorem 10. Let \mathfrak{A} be a model with the universe βX . Then:

- (i) $\mathfrak{A} = e(\mathfrak{B})$ for a generalized model \mathfrak{B} iff in \mathfrak{A} all operations are right continuous and all relations right clopen w.r.t. X,
- (ii) $\mathfrak{A} = E(\mathfrak{B})$ for a generalized model \mathfrak{B} iff in \mathfrak{A} all operations are right continuous w.r.t. X and all relations closed.

Since by Theorem 9, e and E applied to generalized models with pseudoprincipal interpretations give the \sim - and *-extensions of ordinary models, Theorem 10 can be considered as a generalization of Theorems 3 and 5.

In conclusion, let us mention that various characterizations of both types of ultrafilter extensions lead to a spectrum of similar extensions as proposed at the end of [6]; so natural tasks are to study all of the spectrum as well as to isolate special features of the two canonical extensions among others.

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