Extending the Soft Constraint Based Mining Paradigm

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Abstract. The paradigm of pattern discovery based on constraints has been recognized a[s](#page-16-0) a core technique in inductive querying: constraints provide to the user a tool to drive the discovery process towards potentially interesting patterns, with the positive side effect of achieving a more efficient computation. So far the research on this paradigm has mainly focussed on the latter aspect: the development of efficient algorithms for the evaluation of constraint-based mining queries. Due to the lack of research on methodological issues, the constraint-based pattern mining framework still suffers from many problems which limit its practical relevance. In our previous work [5], we analyzed such limitations and showed how they flow out from the same source: the fact that in the classical constraint-based mining, a constraint is a rigid boolean function which returns either true or false. To overcome such limitations we introduced the new paradigm of pattern discovery based on Soft Constraints, and instantiated our idea to the fuzzy soft constraints. In this paper we extend the framework to deal with probabilistic and weighted soft constraints: we provide theoretical basis and detailed experimental analysis. We also discuss a straightforward solution to deal with $top-k$ queries. Finally we show how the ideas presented in this paper have been implemented in a real Inductive Database system.

1 Introduction

[Th](#page-16-1)e paradigm of pattern discovery based on c[ons](#page-17-0)[tra](#page-17-1)ints was introduced with the aim of providing to the user a tool to drive the discovery process towards potentially interesting patterns, with the positive side effect of achieving a more efficient computation. So far the research on this paradigm has mainly focused on the latter aspect: the study [of](#page-17-2) constraint properties and, on the basis of these properties, the development of efficient algorithms for the evaluation of constraint-based mining queries. Despite such algorithmic research effort, and regardless some successful applications, e.g., in medical domains [13,18], or in biological domains [4], the constraint-based pattern mining framework still suffers from many problems which limit its practical relevance. In our previous work [5], we analyzed such limitations and showed how they flow out from the

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same source: the fact that in the classical constraint-based mining, a constraint is a rigid boolean function which returns either true or false. Indeed, interestingness is not a dichotomy. Following this consideration, we introduced in [5] the new paradigm of pattern discovery based on Soft Constraints, where constraints are no longer rigid boolean functions. In particular we adopted a definition of soft constraints based on the mathematical concept of semiring. Albeit based on a simple idea, our proposal has the merit of providing a rigorous theoretical framework, which is very general (having the classical paradigm as a particular instance), and which overcomes all the major methodological drawbacks of the classical constraint-based paradigm, representing a step further towards practical pattern discovery.

While in our previous paper we instantiated the framework to the *fuzzy* semiring, in this paper we ex[ten](#page-17-3)d the framework to deal with the probabilistic and the weighted semirings: these different constraints instances can be used to model different situations, depending on the application at hand. We provide the formal problemdefinition and the theoreticalbasisto develop concretesolversforthe mining problems we defined. In particular, we will show how to build a concrete softconstraint based pattern discovery system, by means of a set of appropriate wrappers around a crisp constraint pattern mining system. The mining system for classical constraint-based pattern discover that we adopted is CONQUEST, a system which we have developed at Pisa KDD Laboratory[8]. Such a system is based on a mining engine which is a general Apriori-like algorithm which, by means of data reduction and search space pruning, is able to push a wide variety of constraints (practically all possible kinds of constraints which have been studied and characterized) into the frequent itemsets computation. Finally, we discuss how to answer to top-k queries.

2 Soft Constraint Based Pattern Mining

Classical constraint (or crisp constraints) are use[d](#page-16-2) to discriminate admissible and/or non-admissible values for a specific (set of) variable. However, sometimes this discrimination does not help to select a set of assignments for the variable (consider for instance overconstrained problems, or not discriminating enough constraints). In this case is preferable to use soft constraints where a specific cost/preference is assigned to each variable assignments and the best solution is selected by looking for the less expensive/more preferable complete assignment.

Several formalizations of the concept of soft constraints are currently available. In the following, we refer to the formalization based on c-semirings [7]. Using this framework, classical/crisp constraints are represented by using the boolean true and false representing the admissible and/or non-admissible values; when cost or preference are used, the values are instead instantiations over a partial order set (for instance, the reals, or the interval [0,1]).

Moreover the formalism must provide suitable operations for combination (\times) of constraints satisfaction level, and comparison $(+)$ of patterns under a combination of constraints. This is why this formalization is based on the mathematical concept of semiring.

Definition 1 (c-semirings [7]). A semiring is a tuple $\langle A, +, \times, 0, 1 \rangle$ such that: A is a set and $0, 1 \in A$; $+$ is commutative, associative and 0 is its unit element; \times is associative, distributes over $+$, **1** is its unit element and **0** is its absorbing element. A c-semiring ("c" stands for "constraint-based") is a semiring $\langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$ such that $+$ is idempotent with **1** as its absorbing element and \times is commutative.

Definition 2 (soft constraint on c-semiring [7]). Given a c-semiring $S =$ $\langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$ and an ordered set of variables V over a finite domain D, a constraint is a function which, given an assignment $\eta: V \to D$ of the variables, returns a value of the c-semiring. By using this notation we define $\mathcal{C} = \eta \rightarrow A$ as the set of all possible constraints that can be built starting from S , D and V .

In the following we will always use the word semiring as standing for c-semiring.

Example 1. The following example illustrates the definition of soft constraints based on semiring, using the example mining query:

 $\mathcal{Q}: supp_{\mathcal{D}}(X) \geq 1500 \land avg(X.weight) \leq 5 \land sum(X.price) \geq 20$

which requires to mine, from database D , all patterns which are frequent (have a support at least 1500), have average weight at most 5 and a sum of prices at least 20. In this context, we have that the ordered set of variables V is $\langle supp_{\mathcal{D}}(X), avg(X, weight), sum(X, price) \rangle$; the domain D is: $D(supp_{\mathcal{D}}(X)) =$ $\mathbb{N}, D(avg(X.weight)) = \mathbb{R}^+, \text{ and } D(sum(X.price)) = \mathbb{N}.$ If we consider the classical crisp framework (i.e., hard constraints) we are on the boolean semiring: $S_{Bool} = \langle \{true, false\}, \vee, \wedge, false, true \rangle$. A soft constraint C is a function $V \to D \to A$; e.g., $supp_{\mathcal{D}}(X) \to 1700 \to true$.

The + operator is what we use to compare the level of constraints satisfaction for various patterns. Let us consider the relation \leq_S (where S stands for the specified semiring) over A such that $a \leq_S b$ iff $a + b = b$. It is possible to prove that: \leq_S is a partial order; + and \times are monotone on \leq_S ; **0** is its minimum and **1** its maximum, and $\langle A, \leq_S \rangle$ is a complete lattice with least upper bound operator +. In the context of pattern discovery $a \leq_S b$ means that the pattern b is *more interesting* than a , where interestingness is defined by a combination of soft constraints. When using (soft) constraints it is necessary to specify, via suitable combination operators, how the level of interest of a combination of constraints is obtai[ne](#page-2-0)d from the interest level of each constraint. The combined weight (or interest) of a combination of constraints is computed by using the operator \otimes : $\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ defined as $(C_1 \otimes C_2)\eta =$ $C_1\eta \times_S C_2\eta$.

Example 2. In this example, and in the rest of the paper, we use for the patterns the notation $p: \langle v_1, v_2, v_3 \rangle$, where p is an itemset, and $\langle v_1, v_2, v_3 \rangle$ denote the three values $\langle supp_{\mathcal{D}}(p), avg(p.weight), sum(p.price) \rangle$ corresponding to the three constraints in the conjunction in the query Q of Example 1. Consider, for instance, the following three patterns: p_1 : $\langle 1700, 0.8, 19 \rangle$, p_2 : $\langle 1550, 4.8, 54 \rangle$, p_3 :

 $(1550, 2.2, 26)$. If we adopt the classical crisp framework, in the mining query $\mathcal Q$ we have to combine the three constraints using the \land operator (which is the \times in the boolean semiring S_{Bool}). Consider for instance the pattern p_1 : $\langle 1700, 0.8, 19 \rangle$ for the ordered set of variables $V = \langle supp_{\mathcal{D}}(X), avg(X,weight), sum(X,price) \rangle$. The first and the second constraint are satisfied leading to the semiring level true, while the third one is not satisfied and has associated level *false*. Combining the three values with \land we [ob](#page-16-0)tain true \land true \land false = false and we can conclude that the pattern $\langle 1700, 0.8, 19 \rangle$ is not interesting w.r.t. our purposes. Similarly, we can instead compute level *true* for pattern p_3 : $\langle 1550, 2.2, 26 \rangle$ corresponding to an interest w.r.t. our goals.

However, dividing patterns in interesting and non-interesting is sometimes not meaningful nor useful. Most of the times we want to say that each pattern is interesting with a specific level of preference. This idea is at the basis of the soft constraint based pattern mining paradigm [5].

Definition 3 (Soft Constraint Based Pattern Mining). Let P denote the domain of possible patterns. A soft constraint on patterns is a function $C: \mathcal{P} \to A$ where A is the carrier set of a semiring $S = \langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$. Given a combination of soft constraints $\otimes \mathcal{C}$, i.e., a description of what is considered by the user an interesting pattern, we define two different problems:

- λ **-interesting:** given a minimum interest threshold $\lambda \in A$, it is required to mine the set of all λ -interesting patterns, i.e., $\{p \in \mathcal{P} \mid \otimes \mathcal{C}(p) \geq_S \lambda\}.$
- **top-k:** given a threshold $k \in \mathbb{N}$, it is required to mine the top-k patterns $p \in \mathcal{P}$ w.r.t. the order \leq_S .

In the rest of the paper we adopt the notation $int_S^p(\lambda)$ to denote the problem of mining λ -interesting patterns (from pattern domain \mathcal{P}) on the semiring S, and similarly $top_S^p(k)$, for the corresponding top-k mining problem. Note that the Soft Constraint Based P[att](#page-3-0)ern Mining paradigm just defined, has many degrees of freedom. In particular, it can be instantiated:

- 1. on the domain of patterns P in analysis (e.g., itemsets, sequences, trees or graphs),
- 2. on the semiring $S = \langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$ $S = \langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$ $S = \langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$ (e.g., boolean, fuzzy, weighted or probabilistic), and
- 3. on one of the two possible mining problems, i.e., λ -interesting or top-k mining.

In other terms, by means of Definition 3, we have defined many different mining problems: it is worth noting that the classical constraint based frequent itemsets mining, is just a particular instance of our framework. In particular, it corresponds to the mining of λ -interesting itemsets on the boolean semiring, where $\lambda = true$, i.e., $int_b^L (true)$. In our previous paper [5] we have shown how to deal with the mining problem $int_f^1(\lambda)$ (i.e., λ -interesting Itemsets on the Fuzzy Semiring), in this paper we show how to extend our framework to deal with (i) $int_p^{\mathcal{I}}(\lambda)$ (i.e., λ -interesting Itemsets on the Probabilistic Semiring), (*ii*) $int_w^{\mathcal{I}}(\lambda)$

(i.e., λ -i[nte](#page-16-2)resting Itemsets on the Weighted Semiring), and *(iii)* mining top-k itemsets on any semiring.

The methodology we adopt is based on the property that in a c-semiring $S = \langle A, +, \times, \mathbf{0}, \mathbf{1} \rangle$ the \times -operator is extensive [7], i.e, $a \times b \leq_{S} a$ for all $a, b \in A$. Thanks to this property, we can easily prune away some patterns from the set of possibly interesting ones. In particular this result directly applies when we want to solve a λ -interesting problem. In fact for any semiring (fuzzy, weighted, probabilistic) we have that [7]:

Proposition 1. Given a combination of soft constraints $\otimes \mathcal{C} = C_1 \otimes \ldots \otimes C_n$ based on a semiring S, for any pattern $p \in \mathcal{P}$:

 $\otimes \mathcal{C}(p) \geq_S \lambda \Rightarrow \forall i \in \{1,\ldots,n\} : C_i(p) \geq_S \lambda.$

Proof. Straightforward from the extensivity of \times .

Therefore, computing all the λ -interesting patterns can be done by solving a crisp problem where all the constraint instances with semiring level lower than λ have been assigned level *false*, and all the instances with semiring level greater or equal to λ have been assigned level *true*. In fact, if a pattern does not satisfy such conjunction of crisp constraints, it will not be neither interesting w.r.t. the soft constraints. Using this theoretical result, and some simple arithmetic we can transform each soft constraint in a corresponding crisp constraint, push the crisp constraint in the mining computation to prune uninteresting patterns, and when needed, post-process the solution of the crisp problem, to remove uninteresting patterns from it.

3 Mining $int_p^{\mathcal{I}}(\lambda)$ (λ -Interesting Itemsets on the **[Pr](#page-17-4)obabilistic Semiring)**

Probabilistic CSPs (Prob-CSPs) were introduced to model those situations where each constraint c has a certain probability $p(c)$, independent from the probability of the other constraints, to be part of the given [pro](#page-5-0)blem (actually, the probability is not of the constraint, but of the situation which corresponds to the constraint: saying that c has probability p means that the situation corresponding to c has probability p of occurring in the real-life problem). Using the probabilistic constraints framework [14] we suppose each constraint to have an independent probability law, and combination is computed performing the product of the semiring value of each constraint instantiations. As a result, the semiring corresponding to the probabilistic framework is $S_P = \langle [0, 1], max, \times, 0, 1 \rangle$.

Consider the constraints graphical representations in Figure 1, where the semiring values between 0 and 1 are this time interpreted as probabilities. In this situation for the pattern $p_1 = \langle 1700, 0.8, 19 \rangle$ we obtain that: $C_1(p_1) =$ 0.83, $C_2(p_1) = 1$ and $C_3(p_1) = 0.45$. Since in the probabilistic semiring the combination operator \times is the arithmetic multiplication, we got that the interest level of p_1 is 0.37. Similarly for p_2 and p_3 :

- $p_1 : C_1 \otimes C_2 \otimes C_3(1700, 0.8, 19) = \times (0.83, 1, 0.45) = 0.37$
- $-p_2$: $C_1 \otimes C_2 \otimes C_3(1550, 4.8, 54) = \times (0.58, 0.6, 1) = 0.35$
- $-p_3$: $C_1 \otimes C_2 \otimes C_3(1550, 2.2, 26) = \times (0.58, 1, 0.8) = 0.46$

Therefore, with this particular instance we got that $p_2 <_{S_P} p_1 <_{S_P} p_3$, i.e., p_3 is the most interesting pattern among the three. Dealing with the probabilistic semiring, we can readapt most of the framework developed for the fuzzy semiring. In fact the two semirings are based on the same set $[0, 1]$ and on the same $+$ operator which is max. The only distinguishing element is the \times operator which is *min* for the fuzzy semiring, while it is the arithmetic *times* for the probabilistic semiring. This means that we can straightforwardly readapt the problem definition, the way of defining the behaviour of soft constraints, and the crisp translation.

Definition 4. Let $\mathcal{I} = \{x_1, ..., x_n\}$ be a set of items, where an item is an object with some predefined attributes (e.g., price, type, etc.). A soft constraint on itemsets, based on the probabilistic semiring, is a function $C: 2^{\mathcal{I}} \to [0,1]$. Given a combination of such soft constraints $\mathcal{C} \equiv \mathcal{C}_1 \otimes \ldots \otimes \mathcal{C}_n$, we define the interest level of an itemset $X \in 2^{\mathcal{I}}$ as $\otimes \mathcal{C}(X) = \mathcal{C}_1(X) \times \cdots \times \mathcal{C}_n(X)$. Given a minimum interest threshold $\lambda \in]0,1]$, the λ -interesting itemsets mining problem, requires to compute $int_p^{\mathcal{I}}(\lambda) = \{X \in 2^{\mathcal{I}} | \otimes \mathcal{C}(X) \geq \lambda\}.$

Definition 5. A soft constraint C on itemsets, based on the probabilistic semiring, is defined by a quintuple $\langle Agg, Att, \theta, t, \alpha \rangle$, where:

- **–** Agg ∈ {supp, min, max, count, sum, range, avg, var, median, std, md};
- **–** Att is the name of the attribute on which the aggregate agg is computed (or the transaction database, in the case of the frequency constraint);
- $-\theta \in \{\leq, \geq\};$
- **–** t ∈ R corresponds to the center of the interval and it is associated to the semiring value 0.5;
- $-\alpha \in \mathbb{R}^+$ is the softness parameter, which defines the inclination of the preference function (and thus the width of the interval).

Fig. 1. Graphical representation of possible probabilistic instance of the constraints in the mining query Q in Example 1

In particular, if $\theta = \leq$ (as in Figure 1(C₂)) then $\mathcal{C}(X)$ is 1 for $X \leq (t - \alpha t)$, is 0 for $X \ge (t + \alpha t)$, and is linearly decreasing from 1 to 0 within the interval $[t-\alpha t, t+\alpha t]$. The other way around if $\theta = \geq$ (as, for instance, in Figure 1(C₃)). Note that if the softness parameter α is 0, then we obtain the crisp (or hard) version of the constraint.

Example 3. Consider again the query \mathcal{Q} given in Example 1, and its probabilistic instance graphically described by Figure 1. Such query can be expressed in our constraint language as:

 $\langle supp, D \rangle \geq$, 1500, 0.2 \rangle , $\langle avg, weight, \leq$, 5, 0.2 \rangle , $\langle sum, price, \geq$, 20, 0.5 \rangle

Definition 6. Given a probabilistic soft constraint $\mathcal{C} \equiv \langle Agg, Att, \theta, t, \alpha \rangle$, and a minimum interest threshold λ , we define the crisp translation of C w.r.t. λ as:

$$
\mathcal{C}_{crisp}^{\lambda} \equiv \begin{cases} Agg(Att) \ge t - \alpha t + 2\lambda \alpha t, & \text{if } \theta = \ge \\ Agg(Att) \le t + \alpha t - 2\lambda \alpha t, & \text{if } \theta = \le t \end{cases}
$$

In [5] we proved that, on the fuzzy semiring, given a combination of soft constraints $\otimes \mathcal{C} \equiv \mathcal{C}_1 \otimes \ldots \otimes \mathcal{C}_n$, and a minimum interest threshold λ , if we consider the conjunction of crisp constraints obtained by conjoining the crisp translation of each constraint in $\mathcal{C}\mathcal{C}$ w.r.t. λ (i.e., $\mathcal{C}' \equiv \mathcal{C}_1 \overline{\lambda}_{crisp} \wedge \ldots \wedge \mathcal{C}_n \overline{\lambda}_{crisp}$), it holds that

$$
int_f^{\mathcal{I}}(\lambda) = \{ X \in 2^{\mathcal{I}} | \otimes \mathcal{C}(X) \ge \lambda \} = Th(\mathcal{C}')
$$

Similarly, the following property holds:

Proposition 2. Given the vocabulary of items I , a combination of soft constraints $\otimes \mathcal{C} \equiv \mathcal{C}1 \otimes \ldots \otimes \mathcal{C}n$, and a minimum interest threshold λ . It holds that:

$$
int_p^{\mathcal{I}}(\lambda) \subseteq int_f^{\mathcal{I}}(\lambda)
$$

Proof. Consider two real numbers x_1, x_2 in the interval [0, 1]. It holds that $x_1 \times$ $x_2 \leq min(x_1, x_2)$. Therefore, for a given pattern *i*, if in the probabilistic semiring $\otimes \mathcal{C}(i) \geq p \lambda$, then also in the fuzzy semiring $\otimes \mathcal{C}(i) \geq f \lambda$.

	\mathcal{D}	t	α		α		α
\mathcal{Q}_1	RETAIL	20	0.8	10000	0.5	20000	0.5
\mathcal{Q}_2	RETAIL	20	0.5	10000	0.5	20000	0.5
$\overline{\mathcal{Q}_3}$	RETAIL	20	0.2	10000	0.5	20000	0.5
$\overline{\mathcal{Q}_4}$	RETAIL	20	0.8	5000	0.2	20000	0.5
\mathcal{Q}_5	RETAIL	20	0.8	5000	0.8	20000	0.5
	Q_6 T40I10D100K 800 0.75 15000				0.2	100000	0.5
\mathcal{Q}_7	T40I10D100K 800 0.75 15000				0.9	100000	0.5
	Q_8 T40I10D100K 800 0.25 15000				0.2	100000	0.2

 $\langle supp, \mathcal{D}, \geq, t, \alpha \rangle$ $\langle avg, weight, \leq, t, \alpha \rangle$ $\langle sum, price, \geq, t, \alpha \rangle$

Fig. 2. Description of queries experimented

Fig. 3. Experimental results on the RETAIL dataset with λ ranging in [0, 1] in the probabilistic semiring: number of solutions (a), and ratio with the number of solutions in the fuzzy semiring (b)

When dealing with the probabilistic semiring, [w](#page-6-0)e translate the given query to a crisp one. But afterwards, we need a post-processing step in which we select, among the solutions to the crisp query, the λ -interesting patterns. It is natural to ask ourselves how much selective is this post-processing. This could provide a measure of the kind of improvement that one could get by studying and developing ad-hoc techniques, to push probabilistic soft constraints into the pattern extraction computation.

In Figure 3, for the RETAIL dataset and the queries of Figure 2, we report: in (a), the number of λ -interesting patterns in the probabilistic semiring, while in (b) the ratio of this number with the number of solutions in the fuzzy semiring, i.e., $|int_p^1(\lambda)|/|int_f^1(\lambda)|$. The execution time of the post-processing is not reported in the plots, because in all the experiments conducted, it was always in the order of few milliseconds, thus n[egl](#page-8-0)igible w.r.t. the mining time. Observing the ratio we can note that it is always equals to 1 for $\lambda = 0$ and $\lambda = 1$. In fact a pattern having at least a constraint for which it returns 0, will receive a semiring value of 0 in both the fuzzy semiring (min combination operator), and the probabilistic semiring (\times combination operator). Similarly, for $\lambda = 1$, to be a solution a pattern must return a value of 1 for all the constraints in the combination, in both the semirings. Then we can observe that this ratio is quite high, always larger than 0.7 in the RETAIL dataset. This is no longer true for the queries on the T40I10D100K dataset, reported in Figure 4 (a) and (b): the ratio reach a minimum value of 0.244 for query \mathcal{Q}_7 when $\lambda = 0.2$.

What we can observe is that the ratio does not depend neither on the number of solutions nor on λ (apart the extreme cases 0 and 1). The ratio depends on the softness of the query: the softer the query the lower the ratio, i.e., more patterns discarded by the post-processing. This can be observed in both Figure 3(b) and 4(b): for instance, among the first three queries \mathcal{Q}_1 is softer than \mathcal{Q}_2 which in turns is softer than \mathcal{Q}_3 , and this is reflected in the ratio

Fig. 4. Experimental results on the T40I10D100K dataset with λ ranging in [0, 1] in the probabilistic semiring: number of solutions (a), and ratio with the number of solutions in the fuzzy semiring (b)

which is lower for Q_1 ; similarly Q_5 is softer than Q_4 and its ratio is lower; in $4(b)$ \mathcal{Q}_8 is the least soft while \mathcal{Q}_7 is the most soft, and accordingly behaves the ratio.

4 Mining $int_w^{\mathcal{I}}(\lambda)$ (λ -Interesting Itemsets on the **Weighted Semiring)**

While in the fuzzy semi[rin](#page-9-0)g each pattern has an associated level of preference (or interestingness) for each constraint, and in the probabilistic semiring a value which represents a probability, in the weighted semiring they have an associated cost. Therefore, in the weighted semiring the cost function is defined by summing up the costs of all constraints. According to the informal description given above, the weighted semiring is $S_W = \langle \mathbb{R}^+, min, sum, +\infty, 0 \rangle$.

Example 4. Consider the following weighted instance for the constraints in the query Q (graphically represented in Figure 5):

$$
-C_1(supp_{\mathcal{D}}(X)) = \begin{cases} 1750 - supp_{\mathcal{D}}(X), & \text{if } supp_{\mathcal{D}}(X) < 1750 \\ 0, & otherwise. \end{cases}
$$

 $- C_2(avg(X.weight)) = 25 * avg(X.weight)$

$$
- C_3(sum(X-price)) = \begin{cases} 5*(60 - sum(X-price)), & \text{if } sum(X-price) < 60 \\ 0, & \text{otherwise.} \end{cases}
$$

Note how the soft version of the constraints are defined in the weighted framework: C_1 for instance, since bigger support is better, gives a cost of 0 when the support is greater than 1750 and an increasing cost as the support decreases. Similarly for constraint C_3 : we assign a cost 0 when the sum of prices is at least 60, while the cost increases linearly as the sum of prices shrinks. Constraint C_2 instead aims to have an average weight as lower as possible, and thus larger average weight will produce larger (worse) cost. In this situation we got that:

 $– p_1 : C_1 \otimes C_2 \otimes C_3(1700, 0.8, 19) = sum(50, 20, 205) = 275$

 $– p_2 : C_1 \otimes C_2 \otimes C_3(1550, 4.8, 54) = sum(200, 120, 30) = 350$ $– p_3 : C_1 \otimes C_2 \otimes C_3(1550, 2.2, 26) = sum(200, 55, 170) = 425$

Therefore, with this particular instance we got that $p_3 <_{S_W} p_2 <_{S_W} p_1$ (remember that the order \leq_{S_W} correspond to the \geq on real numbers). In other terms, p_1 is the most interesting pattern w.r.t. this constraints instance.

Since in the weighted semiring, the values correspond to costs, instead of looking for patterns with an interest level larger than λ , we seek for patterns with a cost smaller than λ .

Definition 7. Let $\mathcal{I} = \{x_1, ..., x_n\}$ be a set of items, where an item is an object with some predefined attributes (e.g., price, type, etc.). A soft constraint on itemsets, based on the weighted semiring, is a function $\mathcal{C}: 2^{\mathcal{I}} \to \mathbb{R}^+$. Given a combination of such soft constraints $\mathcal{C} \equiv \mathcal{C}_1 \otimes \ldots \otimes \mathcal{C}_n$, we define the interest level of an itemset $X \in 2^{\mathcal{I}}$ as $\otimes \mathcal{C}(X) = \sum_{i=1,...,n} \mathcal{C}_i(X)$. Given a maximum cost threshold $\lambda \in \mathbb{R}^+$, the λ -interesting itemsets mining problem, requires to compute $int_w^{\mathcal{I}}(\lambda) = \{ X \in 2^{\mathcal{I}} | \otimes \mathcal{C}(X) \leq \lambda \}.$

Fig. 5. Graphical representation of possible weighted instances of the constraints in in the mining query Q in Example 1

For sake of simplicity, we restrict to weighted constraints with a linear behavior as those ones described in Figure 5. To describe such simple behavior, we need a new parameter $\beta \in \mathbb{R}^+$ $\beta \in \mathbb{R}^+$ $\beta \in \mathbb{R}^+$ that represents the semiring value associated to the t point (playing the role of the implicitly given 0.5 value for the fuzzy and probabilistic semiring). In other words we provide two points to describe the straight line passing through them: the point (t, β) and the point $(t - \alpha t, 0)$ for $\theta = \leq$ or $(t + \alpha t, 0)$ for $\theta = \geq$. Note that α still plays the role of the softness knob.

Definition 8. A soft constraint C on itemsets, based on the weighted semiring, is defined by a sextuple $\langle Agg, Att, \theta, t, \beta, \alpha \rangle$, where: Agg, Att, θ and α are defined as for the fuzzy/probabilistic case (Definition 5), t is a point in the carrier set of the weighted semiring, i.e., $t \in \mathbb{R}^+$, and β represents the semiring value associated to t.

Example 5. Consider again the query Q given in Example 1, and its weighted instance graphically described by Figure 5. Such query can be expressed in our constraint language as:

$$
\langle supp, \mathcal{D}, \geq, 1500, 250, \frac{1}{6}\rangle, \langle avg, weight, \leq, 5, 125, 1\rangle, \langle sum, price, \geq, 20, 200, 1\rangle
$$

For the weighted semiring we can still rely on Proposition 1, which states that a pattern in order to be λ -interesting, must return a semiring value smaller than λ (we are dealing this time with costs; i.e., \geq_W is \leq) for each single constraint in the query: this assures us that if a pattern does not satisfy the crisp translation of the given query, it will not be λ -interesting neither in the weighted semiring. In other words we can always use the same methodology described for the probabilistic semiring: translate the query to a crisp one, evaluate it, post-process the result to select the exact solution set.

Definition 9. Given a weighted soft constraint $\mathcal{C} \equiv \langle Agg, Att, \theta, t, \beta, \alpha \rangle$, and a maximum cost threshold λ , we define the crisp translation of C w.r.t. λ as:

$$
\mathcal{C}_{crisp}^{\lambda} \equiv \begin{cases} Agg(Att) \le t - \alpha t + \frac{1}{\beta} \lambda \alpha t, & \text{if } \theta = \le \\ Agg(Att) \ge t + \alpha t - \frac{1}{\beta} \lambda \alpha t, & \text{if } \theta = \ge \end{cases}
$$

Example 6. Given the weighted soft constraint $\langle sum, price, \geq, 20, 200, 1 \rangle$, its crisp translation is $sum(X,price) \ge 24$ for $\lambda = 180$, it is $sum(X,price) \ge 10$ for $\lambda = 250.$

Proposition 3. Given the vocabulary of items I , a combination of weighted soft constraints $\otimes \mathcal{C} \equiv \mathcal{C}_1 \otimes \ldots \otimes \mathcal{C}_n$, and a maximum interest threshold λ . Let \mathcal{C}' be the conjunction of crisp constraints obtained by conjoining the crisp translation of each constraint in \mathcal{LC} w.r.t. $\lambda: \mathcal{C}' \equiv \mathcal{C}_1 \lambda_{crisp} \wedge \ldots \wedge \mathcal{C}_n \lambda_{crisp}$. It holds that:

$$
int_w^{\mathcal{I}}(\lambda) \subseteq \{ X \in 2^{\mathcal{I}} | \otimes \mathcal{C}(X) \le \lambda \} = Th(\mathcal{C}')
$$

where $Th(\mathcal{C}')$ is the solution set for the crisp problem, according to the notation introduced in Definition 2.

In the following we report the result[s](#page-11-0) of some ex[per](#page-11-1)iments that we have conducted on the same datasets used before for the fuzzy and the probabilistic semirings. We have compared 8 different instances (described in Figure 6) of the query \mathcal{Q} :

$$
\langle supp, \mathcal{D}, \geq, t, \beta, \alpha \rangle \langle avg, weight, \leq, t, \beta, \alpha \rangle, \langle sum, price, \geq, t, \beta, \alpha \rangle
$$

The results of the experiments are reported in Figure 7 and Figure 8. A first observation is that, on the contrary of what happening in the probabilistic and fuzzy semiring, here the larger is λ the larger is the number of solutions. This is trivially because the order of the weighted semiring says that smaller is better.

	$\langle supp, \mathcal{D}, \geq, t, \beta, \alpha \rangle$						$\langle avg, weight, \leq, t, \beta, \alpha \rangle$	$\langle sum, price, \geq, t, \beta, \alpha \rangle$		
	$\mathcal{D}% _{T}=\mathcal{D}_{T}\!\left(a,b\right) ,\ \mathcal{D}_{T}=\mathcal{D}_{T}\!\left(a,b\right) ,$		з	α			α			α
\mathcal{Q}_9	RETAIL	20	600 0.8		5000	100	0.2	20000	250	0.5
\mathcal{Q}_{10}	RETAIL	20	600 0.2		5000	100	0.2	20000	250	0.5
\mathcal{Q}_{11}	RETAIL	20			600 0.8 5000	100	0.8	20000	250	0.5
Q_{12}	RETAIL	20			600 0.8 5000	500	0.2	20000	250	0.5
Q_{13}	RETAIL	20			600 0.8 5000	1000	0.2	20000	500	0.5
	Q_{14} T40I10D100K				800 500 0.8 5000	200	0.5	80000	400	0.8
	Q_{15} T40I10D100K				600 600 0.8 15000	500	0.5	80000	400	0.8
	Q_{16} T40I10D100K 1000 500 0.5 15000 500						0.5	100000 600		0.9

 $supp, \mathcal{D}, \geq, t, \beta, \alpha$ < $avg,weight, \leq, t, \beta, \alpha \rangle$ $sum, price, \geq, t, \beta, \alpha \rangle$

Fig. 6. Description of queries experimented

Fig. 7. Experimental results on the RETAIL dataset with λ ranging in [0, 1000] in the weighted semiring: number of solutions (a), and ratio with the number of solutions of the crisp translation (b)

Fig. 8. Experimental results on the T40I10D100K dataset with λ ranging in [0, 1000] in the weighted semiring: number of solutions (a), and ratio with the number of solutions of the crisp translation (b)

In Figure 7(a) we can observe that queries \mathcal{Q}_{12} and \mathcal{Q}_{13} always return a small number of solutions: this is due to the high values of β in the constraints, which means high costs, making difficult for patterns to produce a total cost smaller than λ . In Figure 7(b) and Figure 8(b[\) w](#page-11-0)e report the ratio of the number of solution with the cardinality of the theory corresponding to the crisp translation of the queries, i.e., $|int_w^{\mathcal{I}}(\lambda)| / |Th(\mathcal{C}')|$. This gives a measure of how good is the approximation of the crisp translation, or in other terms, the amount of post-processing needed (which, however, has negligible computational cost). The approximation we obtain using our crisp solver is still quite good but, as we expected, not as good as in the probabilistic semiring. Also in this case, the softer the query the lower the ratio, i.e., the crisp approximation is better for harder constra[in](#page-16-3)ts (closer to crisp). For instance in Figure 7(b) we can observe that Q_{10} , which is the query with smaller values for the softness parameter α , always present a very high ratio.

5 Mining Top-*k* **Itemsets**

For sake of completeness, in this section we sketch a simple methodology to deal with $top-k$ queries, according to [6]. In the following we do not distinguish between the possible semiring instances, we just describe the general methodology.

The main difficult to solve *top-k* queries is that we can know the number of solutions only after the evaluation of a query. Therefore, given k , the simple idea is to repeatedly run λ -interesting queries with different λ thresholds: we start from extremely selective λ (fast mining) decreasing in selectivity, until we do not extract a solution set which is large enough (more than k).

Considering for instance the fuzzy semiring, where the best semiring value is 1: we could start by performing a 0.95-interesting query, and if the query results in a solution set of cardinality larger than k , then we sort the solution according to their semiring value and return the best k , otherwise we slowly decrease the threshold, for instance $\lambda = 0.9$, and so on. Notice that is important to start from a very high threshold in order to perform fast mining extractions with small solution sets, and only if needed decrease the threshold to get more solutions at the cost of longer computations.

6 Soft Constraints in ConQueSt

In this section we describe how the ideas presented in this paper have been integrated within the CONQUEST inductive database system. CONQUEST is a constraint-based querying system devised with the aim of supporting the intrinsically exploratory nature of pattern discovery. It provides users with an expressive constraint-based query language (named SPQL) which allows the discovery process to be effectively driven toward potentially interesting patterns. The system is built around an efficient constraint-based mining engine which entails several data and search space reduction techniques, and allows new user-defined constraints to be easily added (for deeper details on the CONQUEST system, see also other paper in this volume [10]).

In order to integrate the soft constraint based pattern mining paradigm within CONQUEST, we first extended the $SPQL$ query language to allow definition of soft constraints.

Example 7. In this example we show a complex SPQL query exploiting the soft constraint paradigm. In particular it requires to mine, in the probabilistic semiring, the top 5 patterns w.r.t. a given combination of 3 soft constraint: the frequency constraint, support larger than 5 with 0.4 softness, plus two aggregate soft constraints defined over the attributes product.gross weight and product.units per case. This is a true mining query, defined within Con-QueSt on the famous foodmart2000 datamart.

Fig. 9. CONQUEST window for the definition of a soft constraint (a), and another window with the graphical representation of the soft constraint defined (b)

In line 1. we got the soft constraint query type definition (i.e., if top-k or λ interesting with the appropriate threshold) and the semiring in which the query must be evaluated. In line 2 a minimum frequency constraint is defined with

threshold 5 and 0.4 softness level. From line 3 to 5 we got a typical SQL selectfrom-where statement defining the data source for the query. Lines from 6 to 8 contains the mining view definition, or in other terms, how transactions must be built from the source data (pre-processing). Line 9 contains the two other constraints with their associated softness parameters.

Fig. 10. CONQUEST soft constraints query definition

[T](#page-15-0)his query seems quite complex to be written, but CONQUEST offers simple mechanisms to facilitate the definition of a query. In figure 9 we show the window for the definition of a soft constraint, and the window with the graphical representation of the soft constraint defined.

In Figure 10 we show CONQUEST's constraint definition module, where all the three constraints of the query in Example 7 are reported. Note the dropdown menus to choose among top-k or λ -interesting, and to choose the semiring.

Finally, in Figure 11 we show CONQUEST global view with the query in Example 7 ready to be run, then the resulting top 5 patterns with two different possible views (that can be chosen from the menu): with the actu[al](#page-17-5) [va](#page-17-5)lue of each pattern for each aggregate in a con[stra](#page-17-6)int, or with the respective interestingness value.

7 Related Work

Since in this paper we extend a novel paradigm that we introduced last year, there are not many related works in a strict sense. In a larger sense, all the work done on interestingness of extracted patterns can be considered related. In [22] all [thes](#page-17-7)[e](#page-17-8) works are divided in four classes: objective interestingness measures [12,3,21,15], visualization-based approaches [17], subjective domain-dependent measures of interest [20], and constraint-based approaches. Our proposal clearly collocates within the last class. As already stated in the introduction, a lot of work has been done on constraint-based pattern discovery, but almost all has been done on the development of efficient constraint-pushing algorithms. Entering in the details of these computational techniques, for which we have provided references in the introduction, is beyond the scope of this paper. The reader should refer to [11,9] for un updated state-of-the-art. What we can say here is that most of these techniques have been adopted to build CONQUEST's mining engine [8].

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Fig. 11. CONQUEST global view with the query in Example 7 ready to be run(a); the pattern browser showing top 5 patterns (b); the pattern browser where for each pattern the interestingness level for each constraint is shown

To the best of our knowledge only few works [16,2] have studied the constraintbased paradigm by a methodological point of view, mainly criticizing some of its weak points. To overcome these weak points in this paper we have introduced the use of soft-constraints. [A](#page-16-4) similar approach, based on relaxation of constraints, has been adopted in [1] but for sequential patterns. In the context of sequential patterns, constraints are usually defined by means of regular languages: a pattern is a solution to the query only if it is frequent and it is accepted by the regular language. In this case, constraint-based techniques ado[pt a](#page-17-9) deterministic finite automaton to define the regular language.

The use of regular languages transforms the pattern mining process into the verification of which of the sequences of the language are frequent, completely blocking the discovery of novel patterns. In [1] the authors propose a new mining methodology based on the use of constraint relaxations, which assumes that the user is responsible for choosing the strength of the restriction used to constrain the mining process. A hierarchy of constraint relaxations is developed.

Another recent work using softness in a inductive database context is [19]. In this paper the softness issue addressed, is mostly related to the frequency constraint, i.e., avoiding the exact match between candidate patterns and data instances. The work is developed for substring patterns.

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