# Mining Schedule Optimisation for Conditionally Simulated Orebodies

#### M. Menabde, G. Froyland, P. Stone and G.A. Yeates

Abstract Traditionally the process of mine development, pit design and long-term 4 scheduling is based on a single deterministic orebody model built by the interpo-5 lation of drill hole data using some form of spatial interpolation procedure, e.g. 6 kriging. Typical steps in mine design would include the determination of the ultimate pit, the development of a number of mining phases (pushbacks) and then the 8 development of a life-of-mine schedule. All of these steps would have the aim of q maximising the mine's net present value (NPV), along with meeting numerous 10 other business and physical constraints. There are a number of software packages 11 commercially available and widely used in the mining industry that deal with some 12 or all of these issues. The methods employed by all of these packages treat the 13 process described above in a strictly deterministic way. In reality, given the sparse 14 drill hole data, there is usually significant and variable uncertainty associated with a 15 single or unique deterministic block model. This uncertainty is not captured or used 16 in the planning process. This paper describes work undertaken by the Exploration 17 and Mining Technology Group within BHP Billiton to develop a new mathematical 18 algorithm for mine optimisation under orebody uncertainty. This uncertainty is 19 expressed as a number of conditionally simulated orebody models. This optimi-20 sation algorithm is implemented in a new software package. The software uses a 21 number of proprietary algorithms along with the commercially available mixed 22 integer-programming package ILOG CPLEX. The development targets all phases 23 of mine optimisation, including the NPV optimal block extraction sequence, 24 pushback design, and simultaneous cut-off grade and mining schedule optimisation. 26

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#### Introduction 27

This paper describes the development and implementation of a new software 28 package for open pit mine development and scheduling optimisation under con-29 ditions of orebody uncertainty and is based on the mixed integer programming 30 method. The approach uses multiple conditionally simulated realisations of the 31 orebody as input to characterise the orebody along with the uncertainty in the 32 estimate. 33

Traditionally open pit mine planning, pit design and long-term scheduling is 34 based on a block model of the orebody built by interpolation techniques such as 35 kriging from the drill hole sample data. This single model is assumed to be a fair 36 representation of reality and is used for mine design and optimisation. The design 37 process consists of four main steps: 38

- 1. Determining the ultimate pit shell to define the scheduling universe. 39
- 2. Finding the block extraction sequence which produces the best net present value 40 (NPV) whilst satisfying the geotechnical slope constraints. 41
- 3. Designing the practically minable mine phases (pushbacks) which are roughly 42 based on the optimal block sequence. 43
- 4. Optimising the mining schedule and cut-off grades (COG) within a set of 44 business and operational constraints. The NPV of this 'optimal' schedule is 45 considered as a main criterion of the economical viability of the project. 47

In reality, there are many uncertainties in the models and parameters used in 48 optimisation. Thus, the adoption of a single economic criterion for a project can be 49 very questionable. One of the most important sources of uncertainty is the block 50 model itself. The drill hole data for a mining project is typically sparse, particularly 51 at the scale of the selective mining unit and could support a range of possible 52 outcomes for the orebody. A unique deterministic block model will often be a good 53 representation of the global resource, but will not be representative of the potential 54 local variability or the uncertainty in the estimate. An approach that quantifies both 55 the local variability and the potential uncertainty is to use multiple conditional 56 simulation realisation to represent the orebody (see Dimitrakopoulos 1998). This 57 approach allows the generation of a number of equally probable realisations of the 58 block model, at the selective mining unit (SMU) scale, with all of them honouring 59 the drill hole data along with the first and second order statistics of the orebody 60 represented, respectively, by the probability distribution and variogram (e.g. Isaaks 61 and Srivastava 1989). 62

The simplest and most straightforward use of this set of orebody realisations is to 63 estimate the variability in the project NPV associated with the orebody uncertainty 64 by valuing the 'optimal' schedule obtained from the kriged deterministic model 65 through each of the conditionally simulated realisations. 66

The more interesting question is whether it is possible to use the set of condi-67 tionally simulated realisations to produce a better mine design and production 68 schedule. By 'better' we mean here a higher expected NPV (which becomes a 69

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random variable in case of multiple realisations of the orebody model) and/or less
 variability from one realisation to other (i.e. lower variance of NPV). A new
 promising approach to this problem is presented in Ramazan and Dimitrakopoulos
 (2007, this volume); Jewbali (2006).

In this paper we address one particular aspect of the optimisation under uncer-74 tainty, namely the simultaneous optimisation of the extraction sequence and COG. 75 The use and importance of optimal (variable) COG to mining projects has been 76 known for a long time (e.g. Lane 1988). It will be demonstrated here that the use of 77 variable COG optimised under uncertainty, using the set of equi-probable realisa-78 tions of the orebody can provide a substantial improvement in terms of expected 79 NPV. The approach based on mixed integer programming techniques can provide a 80 truly optimal schedule, as opposed to various heuristic methods used in most of the 81 commercially available mining optimisation software packages. 82

# <sup>83</sup> Mining Schedule Optimisation as a Mixed Integer

# 84 **Programming Model**

Typically, the orebody block model contains between 50,000 and 5,000,000 blocks, which must be scheduled over a period of say 5–25 years. The objective of any scheduling procedure is to find the block extraction sequence, which produces the maximum possible net present value (NPV) and obeys a number of constraints. The latter include:

- geotechnical slope constraints, which are modelled by a set of precedence arcs
   between individual blocks;
- mining constraints, i.e. total maximum amount of rock which can be mined in
   one time period (usually one year);
- <sup>94</sup> 3. processing constraints, i.e. maximum amount of ore which can be processed
   <sup>95</sup> through a given processing plant in one time period; and
- the market constraints, i.e. the maximum amount of metal that can be sold in one time period.

The mathematical formulation of the scheduling procedure in terms of binary decision variables describing in which period the particular block is extracted and what its destination is (either processing plant, stockpile or waste dump), is quite straightforward. The size of the problem is, however, prohibitively large. Apart from the computational difficulties, the hypothetical optimal block extraction sequence may be completely impractical due to the requirements for the mining equipment access and relocation.

Because of these problems the mine scheduling is done using much bigger elementary units that are typically aggregations of hundreds or even thousands of blocks. The aggregation of blocks is a nontrivial problem. For example, simply combining rectangular blocks into a larger rectangular block with dimensions multiples of that of individual blocks can effectively reduce the size of the problem

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but will provide a very poor approximation for the geotechnical slopes. An inter-111 esting approach to block aggregation based on the concept of 'fundamental trees' 112 has been recently developed by Ramazan (2007, this volume). In this method the 113 aggregations of blocks-fundamental trees-obey the slope constraints and can 114 substantially reduce the number of integer variables required for the scheduling 115 model. However, the number of these aggregations is not user controllable and in 116 many cases the problem can be still too big to be solved by a direct application of 117 the mixed integer programming techniques. 118

We have recently developed a new algorithm for block aggregation, which preserves the slope constraints, and is very flexible allowing the user to fully control the size and shape of these aggregations. The details of this algorithm will not be discussed here. The optimisation procedure, however, can be applied to any aggregation of blocks with a set of precedence arcs, prescribing which blocks should be extracted before the given one. As an example we consider here the scheduling of mining phases.

In practice, the open pit mine is divided into a number of mining phases, which 126 are mined bench by bench, each bench represented by a horizontal layer of blocks 127 within the given mining phase and having the same elevation. A bench within a 128 mining phase is sometimes referred to as a "panel". The mining phases can be 129 mined one by one from top to bottom, however this kind of schedule is usually 130 suboptimal. Mining several phases simultaneously and applying variable COG can 131 produce much better results in terms of NPV. There are several commercially 132 available packages, which use proprietary (and undisclosed) heuristics to optimise 133 the schedule and COG. It is difficult to estimate their effectiveness, as the upper 134 theoretical limit on NPV remains unknown. Moreover, these methods can only be 135 used on a single orebody representation and cannot be directly used on a set of 136 conditionally simulated orebody realisations. 137

The standard optimisation technique widely used in many industrial applications 138 is the leanear and integer programming (e.g. industrial applications is the linear and 139 integer programming (e.g. Padberg 1995). The main difficulty in its application to 140 mining scheduling is that the optimisation with variable COG in its direct formu-141 lation leads to a non-linear problem, which is much harder to solve. Our approach 142 provides an effective linearsation of this problem, making it possible to use a mixed 143 integer programming (MIP) formulation for a simultaneous optimization of the 144 extraction sequence and COG for a number of conditionally simulated orebody 145 models. The MIP formulation we use here is similar to the one used by Caccetta and 146 Hill (2003) but is generalised to include the multiple realisations of conditional 147 simulations and variable cut-off grades. This approach also allows one to estimate 148 the gap between the obtained solution and the upper theoretical limit. 149

We consider the simplest case when we have one rock type containing one metal type, which can be processed through one processing plant. Generalisation to the case of multiple rock types, metals and processing streams is cumbersome but straightforward. For simplicity we consider here only the case of a discrete set of COGs, though it is possible to generalise the results to the continuous COG case. We use the following notations:

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- $T_{156}$  T is the number of scheduling periods
- $_{157}$  N is the number of simulations
- <sup>158</sup> P is the total number of panels
- G is the number of all possible cut-off grades
- $R_i^n$  is the total rock in the panel *i* in simulations *n*
- $Q_{ij}^n$  is the total ore in the panel *i*, simulation *n*, when mined with the COG *j*
- $V_{ij}^n$  is the value of the panel *i*, simulation *n*, when mined and processed with the COG *j*
- $_{163}$   $R_t^0$  is the maximum mining capacity in period t
- $_{164}$   $Q_t^0$  is the maximum processing rate in period t
- $_{165}$  Si is the set of panels that must be removed before starting the panel i
- $d^t$  is the time discount factor
- 167 Xijt is the fraction of the panel i is extracted with the COG j in period t
- <sup>168</sup> *Yit* is a binary variable equal to 1 if the extraction of the panel *I* has started in periods 1 to t, and equal to 0 otherwise;
- $\delta jt$  is a binary variable controlling the selection of the COG applied in period t
- The MIP formulation is:

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Maximise  $\left(\frac{1}{N}\sum_{n=1}^{N}\sum_{i=1}^{P}\sum_{i=1}^{G}\sum_{t=1}^{T}V_{ij}^{n}x_{ijt}d^{t}\right)$ (1)

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<sup>175</sup> Subject to the following constraints:

$$\frac{1}{N} \sum_{n=1}^{N} \sum_{i=1}^{P} \sum_{j=1}^{G} R_{i}^{n} x_{ijt} \le R_{t, for all t}^{0}$$
(2)

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 $\frac{1}{N}\sum_{n=1}^{N}\sum_{i=1}^{P}\sum_{j=1}^{G}Q_{ij}^{n}x_{ijt} \le Q_{t, \text{ for all }t}^{0}$ (3)

$$\mathbf{y}_{i, t-1} \leq \mathbf{y}_{it, for all \, i \, and \, t} \tag{4}$$

$$\sum_{\tau=1}^{t} \sum_{j=1}^{G} x_{ij\tau} \le y_{it, \text{ for all}}$$
(5)

$$y_{it \leq \sum_{j=1}^{G} \sum_{\tau=1}^{T} x_{kj\tau, \text{ for all } i, \tau \text{ and } k \subset S_i}$$

$$(6)$$

$$\sum_{j=1}^{G} \delta_{jt=1, \quad for \, all \, t} \tag{7}$$

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$$x_{ijt} \leq \delta_{jt}$$
, for all  $i, j$  and  $t$ 

(8)

The objective function (1) represents the discounted cash flow. Constraints (2) and (3) enforce the mining and processing limits on average. Constraints (4)–(6) enforce the panel extraction precedence constraints, and constraints (7) and (8) ensure that the same COG is applied to all panels extracted in any given time period.

This MIP formulation is solved by the commercially available software package CPLEX version 9.0, by ILOG Inc.

### 198 Case Study

To test the algorithm we have chosen ten conditional simulations of a block model 199 containing one type of metal and using one processing plant. Because of confi-200 dentiality requirements, all the economic parameters were rescaled and do not 201 represent reality. All of the relative characteristics which demonstrate the potential 202 of this new method are not affected by this rescaling. The ultimate pit for the design 203 is chosen by applying the Lersch-Grossmann algorithm (Lersch and Grossmann 204 1965) in a procedure similar to that used in Whittle Four-X software. The ultimate 205 pit contains 191 million tonnes of rock and  $62.9 \pm 2.7$  million tonne of ore (above 206 the marginal COG = 0.6%). The undiscounted value in the ultimate pit (if pro-207 cessed with the marginal COG) is  $(1316 \pm 99)$  million. It was divided into six 208 mining phases and scheduled over 12 years. The mining rate was set to 30 Mtpa 209 and the processing rate to 5 Mtpa. The initial capital investment was assumed to be 210 \$300 million, and the discount rate 10%. The base case optimisation was done 211 using the marginal COG applied individually to all conditional simulation. 212 The NPV for this case was  $(404 \pm 31)$  million. The mining schedule is shown in 213 Fig. 1. The second optimisation was done using the variable COG, but was based 214 on the mean grade block model, i.e. it was similar to an optimisation generated by 215 using a single deterministic model. This schedule was evaluated against all ten 216 realisations of orebody model and produced the NPV =  $(485 \pm 40)$  million, an 217 increase of 20% over the base case. This mining schedule is shown in Fig. 2. The 218 third optimisation was done using the algorithm described in earlier, using all 219 realisations as input to the optimisation and produced the orebody 220 NPV =  $(505 \pm 43)$  million, a further increase of 4.1% over the case of mean 221 grade based optimisation. This mining schedule is shown in Fig. 3. The relative 222 variability of NPV in all cases was roughly the same, about 8%. The cumulative 223 NPV graphs for the three different schedules are shown in Fig. 4, and the com-224 parison between expected NPVs and their variability is shown in Fig. 5. Another 225 important result of the variable COG policy is that the pay-back period (defined 226 here as the time when the cumulative NPV becomes equal to zero) is decreased 227 from five to three years (see Fig. 4). 228

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Fig. 1 Mining schedule optimised with the marginal COG



Fig. 2 Mining schedule optimised with the mean grade model

The increase of 4.1% in NPV may be not seen as a very substantial, but it should be mentioned that the block model considered does not have a high variability. The relative variance in the undiscounted value of the ultimate pit is only 7.6%. There are many deposits that have variability of the order of 20–30%. For these kind of deposits the potential improvement in the expected NPV may be substantially higher.

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Fig. 3 Mining schedule optimised with the set of conditional simulations



Fig. 4 Cumulative NPV for different missing schedules (solid line-variable COG on conditional simulations; dashed line-variable-COG on the mean grade model; dotted line-marginal COG)



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Fig. 5 Comparison of expected NPVs and their variability for different mining schedules (circle-variable COG on the conditional simulations; square-variable COG on the mean grade model; traingle-marginal COG)

#### 235 Conclusions

A new method for simultaneous optimisation of the extraction sequence and cut-off 236 grade policy for a set of conditionally simulated orebody realisations has been 237 developed and demonstrated. This method is based on the mixed integer pro-238 gramming model and uses the commercially available software package CPLEX by 239 ILOG Inc. The goal of the optimisation is to find the extraction sequence and cut-off 240 grade policy, which, when evaluated through the whole set of conditionally sim-241 ulated orebodies (representing the range of possible outcomes), will produce the 242 best possible expected NPV. The degree of accuracy of this optimised schedule can 243 be estimated precisely, in contrast to a number of heuristic routines used in current 244 commercially available mining optimisation software packages. A fully functional 245 software prototype that uses the new optimisation method has been developed. 246

In this study, we were using the expected NPV as the objective function and the 247 mining and processing constraints were applied to the mean rock and ore tonnages. 248 Some of the possible extensions of this method may include some kind of penalty 249 functions in the objective function in order to find a schedule with a reduced 250 variability in NPV, defining hard constraints bounding the NPV from below, or 251 defining a lower bound on the annual cash flows. Another very interesting gener-252 alisation may include a stochastic price model for metals and adjustable cut-off 253 grade policy. 254



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