

Direct Net Present Value Open Pit Optimisation with Probabilistic Models

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Abstract Traditional implementations of open pit optimisation algorithms are designed simply to find a set of nested open pit limits that maximise the undiscounted financial pay-off for a series of commodity prices using a single ‘estimated’ orebody model. Then, the maximum Net Present Value (NPV) open pit limit is derived by considering alternate (usually only best and worst-case) mining schedules for each open pit limit. Divorcing the open pit limit delineation from the NPV calculation in this two-step approach does not guarantee that an optimal NPV open pit solution will be found. A new open pit optimisation algorithm that considers the mining schedule is proposed. As a consequence, it can also account explicitly for commodity price cycles and uncertainty that can be modelled by stochastic simulation techniques. This state-of-the-art algorithm integrates Monte Carlo-based simulation and heuristic optimisation techniques into a global system that directly provides NPV optimal pit outlines. This new approach to open pit optimisation is demonstrated for a large copper deposit using multiple orebody models.

Introduction

Several open pit optimisation techniques such the Lerchs–Grossman algorithm (Lerchs and Grossman 1965), network flow (Johnson 1968), pseudoflow network models (Hochbaum and Chan 2000) and others, involve a 3D grid of regular blocks that is converted *a priori* into a pay-off matrix by considering a 3D block model of mineral grades and economic and mining parameters. These algorithms rely on the block pay-offs averaging linearly, as is the case when undiscounted block pay-offs are considered. However, the Net Present Value (NPV) of the block pay-offs is a non-linear function of the undiscounted block pay-offs that depends explicitly on the discount to be applied to the individual blocks, which in turn depends on the

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30 block mining schedule. To overcome the issue of discounting block pay-offs, tradi-
31 tional implementations of open pit optimisation algorithms are designed simply to
32 find a set of nested open pit limits that maximise the undiscounted financial pay-off
33 for a series of constant commodity prices using a single ‘estimated’ orebody model.
34 Then, the maximum NPV open pit limit is derived by considering alternate (usually
35 only best and worst-case) mining schedules for each open pit limit. This two-step
36 approach to finding the maximum NPV open pit limit raises three significant issues:

- 37 1. Divorcing the open pit limit delineation from the NPV calculation does not
38 guarantee that an optimal (maximum) NPV open pit solution will be found;
- 39 2. NPV calculations are based on a constant commodity price that fails to consider
40 its time-dependant and uncertain nature; and
- 41 3. The single ‘estimated’ orebody model is invariably smoothed, thus it fails to
42 consider short-scale grade variations.
43

44 Consequently, the block model does not accurately reflect the grade and tonnage
45 of ore that will be extracted and processed during mining.

46 To overcome the inadequacy of undiscounted pay-offs in commonly used
47 algorithms for open pit optimisation, it is proposed to embed a scheduling heuristic
48 within an open pit optimisation algorithm. This may be seen as an alternative
49 avenue to that taken by mixed integer programming approaches (eg. Caccetta and
50 Hill 2003; Ramazan 2007; Stone et al. 2017; Menabde et al. 2007) that may become
51 numerically demanding in the case of large deposits. As a consequence, uncertain
52 and time-dependent variables such as commodity prices can also be incorporated
53 stochastically into the optimisation process. This permits strategic options for
54 project timing and staging to be assessed as discrete optimisation problems and
55 compared quantitatively and is more advanced than other recent approaches
56 (Monkhouse and Yates 2017 in this volume; Dimitrakopoulos and Abdel Sabour
57 2007). It is also proposed to consider multiple conditional simulations in the
58 optimisation process such that the mining and financial implications related to
59 small-scale grade variations are honoured (Menabde et al. 2017 in this volume;
60 Ramazan and Dimitrakopoulos 2013, 2017 in this volume; Leite and
61 Dimitrakopoulos 2007; Godoy and Dimitrakopoulos 2004; Ravenscroft 1992). By
62 considering discounted block pay-offs, stochastic models of commodity prices and
63 short-scale grade variations a more accurate discounted pay-off matrix (revenue
64 block model) is generated, which in turn will yield an open pit limit that will be
65 closer to the true optimum.

66 NPV Calculations with Uncertain Variables

67 Calculation of the NPV for a given open pit limit relies on estimates of numerous
68 parameters, including (but not restricted to) the mineral grades, extraction sequence
69 and timing, mineral recovery, prevailing commodity price and capital and operating

costs. All of these parameters are uncertain and should be modelled stochastically. For example, mineral grade values by geostatistical simulations, operating costs with growth functions and commodity prices using long-term mean reverting models that account for periodicity. Consequently, the cumulative distribution of total financial pay-offs for an open pit limit can be derived from the combination of a series of stochastic models of mineral grades, costs, prices, recoveries, etc.

Given L potential NPV outcomes for a block (related to L realisations of grade values, commodity prices, etc), we can calculate the NPV for any realisation l :

$$NPV_l = \sum_{j=1}^B d^l(b_j) i_j \quad (1)$$

and the expected NPV for L realisations:

$$NPV_L = \frac{1}{L} \left\{ \sum_{l=1}^L NPV_l \right\} \quad (2)$$

where:

- B is the number of blocks under consideration
- $d^l(b_j)$ is the discounted value for block b_j for the l th realisation
- $i_j = 1$ if b_j falls within the open pit limit and 0 otherwise

The idea being to find the open pit limit that maximises NPV_L . Additional financial goals, for example minimising downside risk (Richmond 2004a) could also be considered, but are outside the scope of this paper.

Accounting for Multiple Orebody Models

Pit optimisation algorithms found in the literature invariably consider an orebody block model with a single grade value for each block (or parcel). In such an approach, a simple decision rule is used where block b_j is processed using option k if $g_k \leq z^*(b_j) < g_{k+1}$, where:

g_k is the cut-off grade for processing option k (by convention $g_1 = 0$ and $k = 1$ indicates waste)

z^* is the estimated grade value

To account for grade uncertainty in open pit optimisation, Richmond (2004a) proposed incorporating L grade values for each block. In this approach, multiple grade values $z^l(b_j)$, $l=1, \dots, L$ were generated by conditional simulation and a processing option $k^l(b_j)$ was determined for each realisation. Alternatively, conditional simulation provides short-scale grade variations that permit local ore loss and mining dilution to be readily accounted for in open pit optimisation by (Richmond 2004a):

- Generating geometrically irregular dig-lines (that separate ore and waste) based on small-scale grade simulations with a floating circle algorithm, and
- Assimilating the dig-lines into large-scale geometrically regular blocks by a novel re-blocking method.

This two-step approach accounts for short-scale grade variation, but also provides ‘recoverable’ grade and tonnage information for large regular blocks suitable for open pit optimisation. In other words, the simulated grade models are compressed without loss of accuracy so that optimisation is computationally tractable.

An NPV Open Pit Optimisation Algorithm

For the vast majority of open pit optimisation techniques a directed graph is superimposed onto the pay-off matrix to identify the blocks that constitute an optimal open pit limit. To paraphrase Dowd and Onur (1993)—each block in the grid, represented by a vertex, is assigned a mass equal to its net expected revenue. The vertices are connected by arcs in such a way that the connections leading from a particular vertex to the surface define the set of vertices (blocks) that must be removed if that vertex (block) is to be mined. A simple 2D example is shown in Fig. 1. Blocks connected by an arc pointing away from the vertex of a block are termed successors of that block, ie. b_i is a successor of b_j if there exists an arc directed from b_i to b_j . In this paper, the set of all successors of b_j will be denoted as Γ_j . For example, in Fig. 1, $\Gamma_8 = \{2, 3, 4\}$. A closure of a directed graph, which consists of a set of blocks B , is a set of blocks $B_p \subset B$ such that if $b_j \in B_p$ then $\Gamma_j \in B_p$. For example, in Fig. 1, $B_p = \{1-5, 7-9, 13\}$ is a closure of the directed graph. The value of a closure is the sum of the pay-offs of the vertices in the closure. As each closure defines a possible open pit limit, the closure with the maximum value defines the optimal open pit limit.

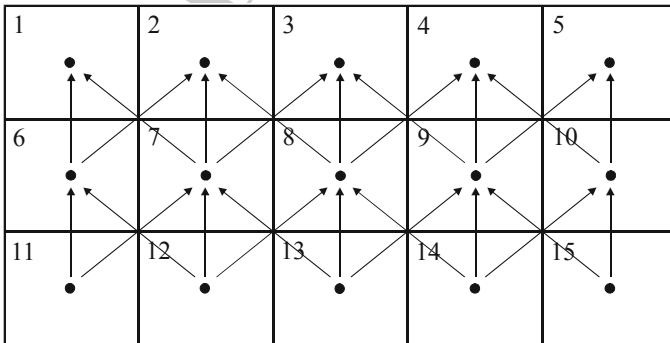


Fig. 1 Directed graph representing 2D vertical orebody model

For simplicity of notation, the algorithm proposed in this paper is described for a single orebody model. The undiscounted pay-off matrix $\{w(b), b \in B\}$ typically used for open pit optimisation is calculated as:

$$w(b) = ton_b(vz(b)r_k - c_k) \quad (3)$$

where:

ton_b represents the tonnage of block b

v is the commodity (attribute z) value per concentration unit

r_k is the proportion of the mineral recovered using processing option k

c_k is the mining and processing cost for k (\$/ton)

In practice, r_k and c_k commonly vary spatially and v and c_k temporally. The discounted pay-off matrix $\{d(b|S), b \in B\}$, conditional to a mining schedule S , that is required for NPV open pit optimisation is calculated as:

$$d(b|S) = [ton_b(v_t z(b)r_k - c_{k,t})]/(1 + DR)^t \quad (4)$$

where:

t is the time period in which block b is scheduled for extraction and processing

$v_t, c_{k,t}$ are the prevailing commodity price and operating cost at time t

DR is the discount rate

In Eq. 4, discounted pay-offs are conditional to the mining schedule as alternate schedules can be derived for the same open pit closure. It is also important to note that, cut-off grades and consequently the processing option k , may change in response to commodity price and operating cost fluctuations over time. Does not imply that the discounted value for is positive.

The traditional floating cone algorithm decomposes the full directed graph problem into a series of independent evaluations of individual Γ_j and if the sum of the pay-offs associated with Γ_j is positive, then b_j is added to B_p . However, a positive undiscounted value for Γ_j , does not imply that the discounted value for Γ_j is positive. In other words, negatively-valued successors b_j or block b_j that may be mined significantly earlier in the mining schedule and receive substantially less discounting may not be carried by a more heavily discounted positively-valued b_j . Furthermore, the modified schedule may have shifted more profitable b_j into later periods and additional wasted blocks into earlier periods, reducing the discounted value of the pit. As traditional independent evaluation of locally decomposed Γ_j .

To allow for discounting, it is proposed that a Direct NPV Floating Cone algorithm (DFC) proceeds as follows:

1. Select the time for initial investment (start of construction) t_i ;
2. Define a cone that satisfies the physical constraints of the desired open pit slope angles;

- 176
- 177 3. Define an ordered sequence of visiting blocks $[1, 2, \dots, \# < B]$ with positive $w(b)$,
- 178 by ordering the blocks b_i firstly on decreasing elevation, and then for blocks
- 179 with identical elevations on decreasing value in $w(b_i)$;
- 180 4. Set the open pit closure counter $n = 0$, the initial open pit closure B_p^n to a null set
- 181 of blocks, and the Net Present Value of initial open pit closure $NPV^n = 0$;
- 182 5. Set $j = 0$;
- 183 6. Set $j = j + 1$;
- 184 7. Float the cone to b_j to create a new closure $B_p^{n+1} = B_p^n + \Gamma_j$ (excluding from Γ_j
- 185 any block that currently belongs to B_p^n);
- 186 8. Determine the schedule S for the new closure B_p^{n+1} ;
- 187 9. Calculate the discounted pay-off matrix $\{d(b|S), b \in B_p^{n+1}\}$ using Eq. 4 and the
- 188 Net Present Value of the new closure using Eq. 1;
- 189 10. Accept the new closure if $NPV^{n+1} - NPV^n > 0$, whereupon the current closure
- 190 is updated into a new optimal closure, ie. $n = n + 1$ and go to step 5; and
- 191 11. if $j < \#$, the number of blocks with positive pay-offs $w(b)$, then go to step 6.
- 192

193 The deterministic floating cone algorithm presented above is heuristic in nature
194 and not be optimal. Alternate B_p can be generated by varying the initial investment
195 timing (step 1), the ordered path (step 3) and/or the mining schedule (step 8).

196 Investment timing to satisfy corporate constraints or to take advantage of
197 cyclical commodity prices can be investigated as mutually exclusive opportunities
198 by varying t_1 , which modifies implicitly the mining schedule in step 8 above. For
199 example, given a schedule S commencing at $t = 0$, the modified schedule $t' = t + t_1$.
200 For delayed investment, the NPV for many potential production assets will typically
201 be reduced unless maximum production/grade happens to coincide with the peak in
202 cyclical commodity prices. However, for a risk averse and capital constrained
203 company, the shift of the capital cost into future years may be strategically
204 advantageous when considered in conjunction with other mining assets.
205 Re-initiating the test sequence from the top of the mineral deposit each time a
206 positively-valued cone is found and added to the closure is generally regarded to
207 estimate the heuristic maximum undiscounted pay-off solution (Lemieux 1979).
208 Computational experimentation on the ordering of blocks in step 3 above suggested
209 that this also holds true for the discounted case when t_1 is fixed. Note that, due to
210 re-initiation of the test sequence it is p common for $B_p^{n+1} = B_p^n$ in step 7 above. For
211 such instances, steps 8–10 above are ignored.

212 It is well known that the floating cone algorithm may not return the maximum
213 undiscounted pay-off solution. However, it is used in the algorithm presented above
214 to generate physically feasible solutions. The author has not investigated whether
215 the Lerchs–Grossman and network flow algorithms could be substituted for the
216 floating cone algorithm, but the non-linearity of the proposed objective function
217 may present some difficulty. The computational efficiency of the proposed algo-
218 rithm is enhanced significantly when a simple scheduling algorithm in step 8 above
219 is employed. However, more complex risk-based scheduling algorithms to account

220 for multiple orebody models and production goals (eg. Godoy 2002) could be
 221 considered.

222 Application to a Copper Deposit

223 This section demonstrates the proposed concepts for a large subvertical copper
 224 deposit. The geometry and contained copper per level are variable, but there is no
 225 strong trend. The options considered in this study were:

- 226 • Two processing options (ore and waste), ie. $K = 2$;
- 227 • 60 Mt/year mill constraint;
- 228 • 25 realisations of copper grades by Sequential Gaussian Simulation (SGS);
- 229 • 25 stochastic simulations of future copper prices with a two factor Pilipovic
 230 model that was modified to account for periodicity and cap and collar aversion
 231 (Fig. 2);
- 232 • 25 stochastic simulations of operating costs with a growth model (Fig. 3);
- 233 • Monthly copper recoveries randomly drawn from normal distribution with mean
 234 of 80% and a standard deviation of 1%²;
- 235 • A fixed annual discount rate of 10%; and
- 236 • Initial investment timings at discrete yearly intervals for five years.

238 Figure 2 shows 25 stochastic simulations of future copper prices. The assump-
 239 tions in this study were:

- 240 • A long-term copper price of \$1.30/lb,

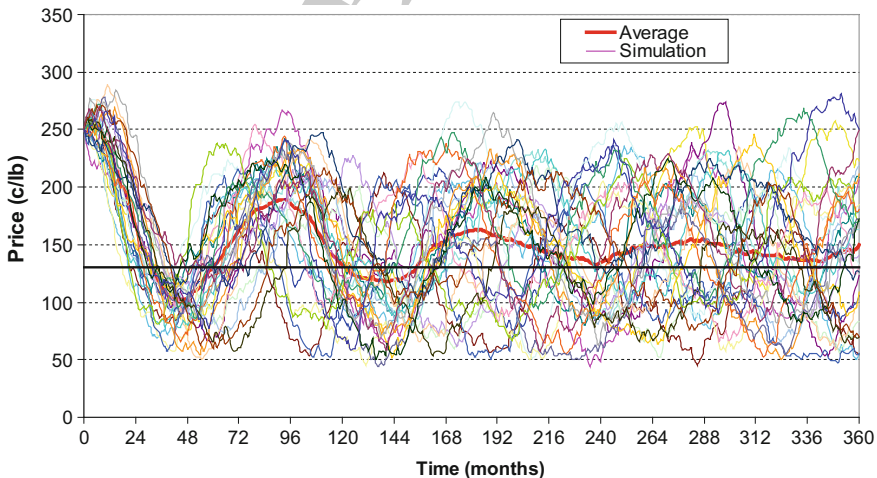


Fig. 2 Thirty year future copper price simulations with mean reversion and collar and cap aversion

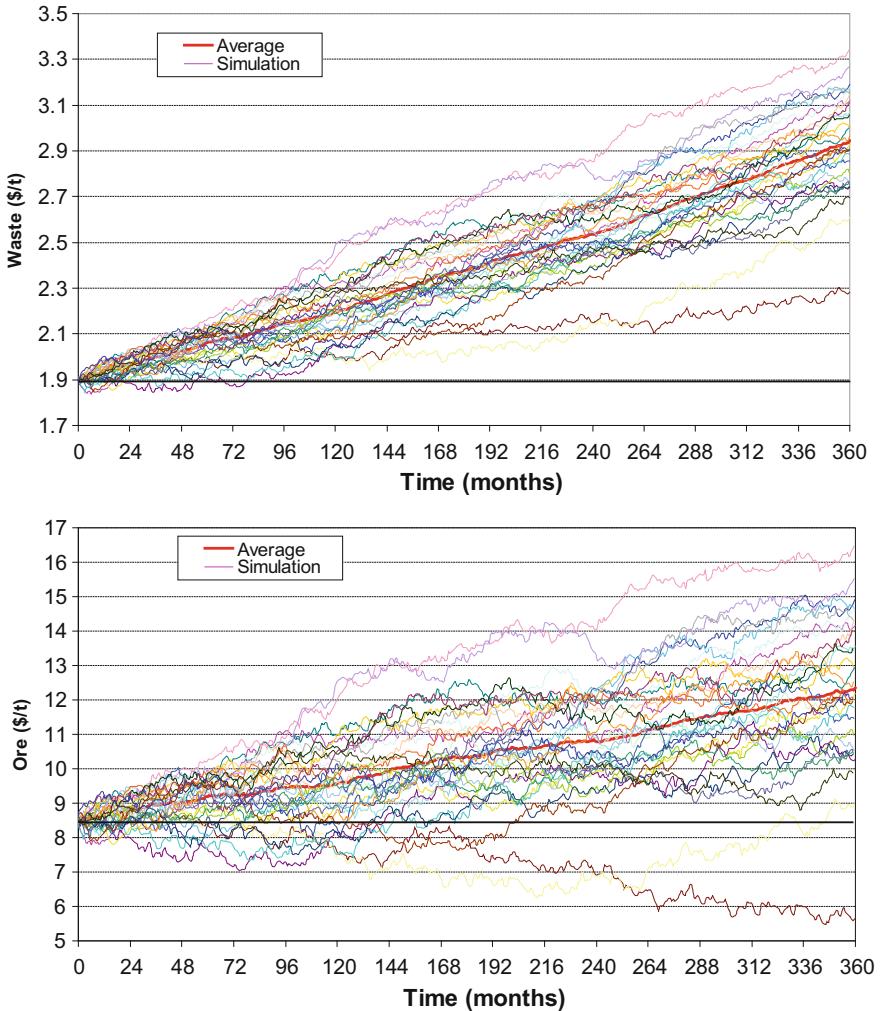


Fig. 3 Thirty year waste and ore processing cost simulations

- 241
- The present time (\$2.50/lb) was near the peak of the price cycle,
 - An average eight year copper price cycle, and
 - \$0.50/lb and \$3.00/lb lower and upper aversion values.
- 242
243
244

245 Note that, as time increases uncertainty in the simulated copper price increases
246 and the deviation of the average simulated value to the long-term price decreases.
247 The average copper price does not fluctuate symmetrically around the long-term
248 copper price due to the asymmetrical aversion limits. Figure 3 shows 25 stochastic
249 simulations of waste and ore processing costs.

250 To assess the potential improvement in NPV against the traditional two-stage pit
 251 optimisation approach a base case scenario (\$1.30/lb—80% recovery, \$1.90/t waste
 252 cost and \$8.50/t milling cost) was run to generate a series of nested pits using a
 253 FCA. The E-type (or average) of the 25 SGS realisations was adopted as the single
 254 grade model as it is known to be smoothed. The NPV for this series of pits using the
 255 base case assumptions are shown in Fig. 4 as crosses. The maximum NPV under
 256 the base case scenario is associated with a pit closure of 26,402 blocks. Note that,
 257 the capital cost, which could also be modelled stochastically, was not included in
 258 this study.

259 The NPV for the FCA nested pits were also calculated using the simulated
 260 grades, metal prices, costs and recoveries for the six annual investment timings,
 261 shown in Fig. 4. Note that:

- 262 • These curves vary substantially from the base case.
- 263 • In all instances the maximum NPV pit is significantly larger (49,239–85,093
 264 blocks) than the base case and the maximum NPV is higher than for the base
 265 case.
- 266 • Delaying the investment from Year 3 to Year 5 results in a higher NPV (\$3.02
 267 billion versus \$2.88 billion). At first this relationship appears counter-intuitive
 268 as costs are greater and discounting greater. However it is related to higher Cu
 269 prices in key production periods.
 270

271 The NPV of the proposed DFC approach for the six annual investment timings
 272 are also shown in Fig. 4. Note that, considering the mining schedule explicitly in
 273 the optimisation process was successful in finding the maximum NPV pit in a single
 274 run. Whilst the improvement over the maximum NPV pit from the two-step

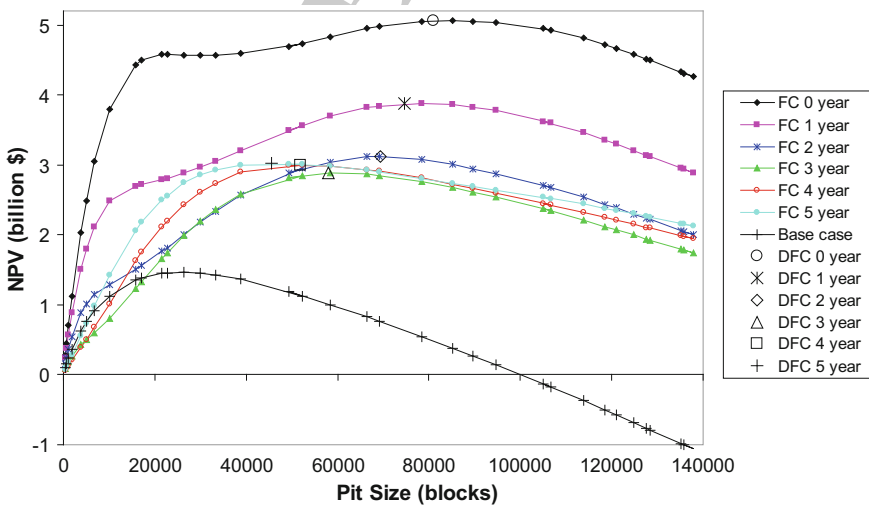


Fig. 4 Pit size versus NPV (FCA = floating cone algorithm; DFC = proposed direct NPV FCA)

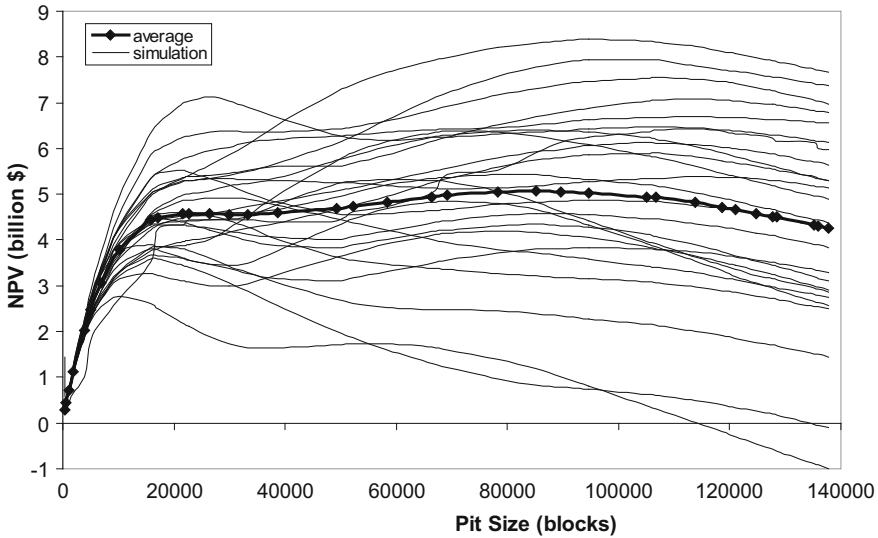


Fig. 5 Pit size versus NPV distribution

275 approach that considered the stochastic inputs was limited (usually $<0.5\%$ in NPV),
 276 there was often some difference in the pit dimension. It is likely that these differ-
 277 ences would be reduced further if additional pit closures had been generated for
 278 evaluation in the two-step approach. Computationally, it was more efficient to post
 279 process a finite series of pit closures than embed the scheduler in the pit optimi-
 280 sation process. In the example shown, the DFC approach that generated a single pit
 281 required around the same computational time as that required in generating 36
 282 nested pits by a simple FC approach.

283 Figure 5 shows the distribution of potential NPVs for the set of nested FCA pits
 284 without any investment delay. As expected, the uncertainty increases with pit size
 285 with some possibility of negative NPVs for large pit closures. If minimising
 286 downside financial risk is of greater importance than maximising the NPV then the
 287 financially efficient set (frontier) of open pit limits could be determined under a
 288 stochastic framework (Richmond 2004a).

289 Conclusions

290 A novel method for working with discounted pay-off matrices during open pit
 291 optimisation was proposed. The approach used in this study embedded a simple ore
 292 scheduler in a floating cone-based heuristic algorithm. It was a trivial exercise to
 293 further consider multiple orebody models, local ore loss and mining dilution,
 294 time-dependent commodity prices and costs and variable metal recoveries during



295 optimisation. As a consequence, alternate project development timings could be
296 strategically assessed. Traditional evaluation of a set of nested pit shells with
297 constant metal prices and operating costs failed to determine the maximum NPV pit
298 under uncertain conditions. However, provided that sufficient pit shells were gen-
299 erated and evaluated with the same stochastic price and cost input as for the
300 proposed algorithm there was little difference in the maximum NPV shell derived.
301 Further experimentation should be undertaken to determine whether this observa-
302 tion holds for more complex mining schedule algorithms and geometrically irreg-
303 ular orebodies, as well as when a smoothed block model other than the E-type of
304 the stochastic grade model is used to generate a series of nested closures.

305 This study demonstrated that uncertainty in future metal prices and operating
306 costs cannot be adequately captured in open pit optimisation by simply
307 post-processing a series of nested pit closures with constant values. Stochastic
308 modelling of mineral grades, mineral recovery, commodity prices and capital and
309 operating costs provide an ideal platform to:

- 310 • Generate an optimal pit to maximise the overall project NPV considering
311 geological and market uncertainty,
- 312 • Determine the optimum investment and project start up timing, and
- 313 • Quantify the multiple aspects of uncertainty in a mine plan.
314

315 The example studied in this paper indicates periods of potential financial
316 weakness that could benefit from management focus (eg. forward selling strategies
317 and placing the mine on care and maintenance) prior to difficulties arising.

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