

# Dynamic *JChoc*: A Distributed Constraints Reasoning Platform for Dynamically Changing Environments

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**Abstract.** In Artificial Intelligence, a large number of problems (i.e. distributed resource management, distributed air traffic management, Distributed Sensor Network [1]) can be modeled and solved as Distributed Constraint Satisfaction Problems (DisCSPs). As many real world problems change continuously and incessantly over time, some methods have been developed (e.g. DynABT), for solving problems which exhibit this dynamic behavior. Meanwhile, there was no available framework that helped users to develop intelligent multi-agent systems based on Dynamic and Distributed Constraints Reasoning (DCR) techniques.

In this paper, we propose a new platform, called *JChoc*, supporting the dynamic aspect for DisCSPs. *JChoc* is an easy to use platform, based on an elegant Multi-agent communication sub-platform (e.i. JADE). It deals with agents with local complex problems and allows a realistic use of agents on a real distributed and dynamic framework.

A real distributed problem is addressed to illustrate how the platform can be used to solve dynamically changing problems. However, the experimental results show the defectiveness of our platform.

**Keywords:** Dynamically changing environments · Constraint programming (CP) · Multi-agent systems · Distributed problem solving · Agent models and architectures · Distributed constraints reasoning · Realistic use · Constraint satisfaction problem (CSP) · Distributed CSP (DisCSP)

## 1 Introduction

Since the onset of real time electronic devices, mobiles, ubiquitous, and intelligent computing, new combinatorial problems have emerged in the AI community such as: distributed resource management, distributed air traffic management, Distributed Sensor Network [1], disaster rescue [2] and distributed Meeting Scheduling Problems (SMP), for which it is not suitable to collect all data of problem in one site, to solve it by a centralized algorithm. The reasons are communication time and cost of translation of each sub-problem in a

common format. In addition, to give a single agent all data of the problem can also be excluded for reasons of security and confidentiality. Therefore, some of the AI communities are motivated to take an interest in Distributed Constraint Reasoning (DCR), giving birth to other distributed formalism [5], whose work focused on developing techniques for modeling and solving distributed combinatorial problems with or without optimization criterion. Distributed Constraint Satisfaction Problems (DisCSP), Distributed Constraint Optimization Problems (DCOP) and Dynamic Distributed Constraints Satisfaction Problems provide a useful framework of multiagent systems for distributed and dynamic resolution of combinatorial problems [3–5, 16, 17].

In this context, an agent must have a communication platform that allows the exchange of information or dialogue to coordinate their decision-making. This reliable communication tool allows agents to send and receive messages according to a given distributed protocol. However, various sophisticated solvers have been developed: DisChoco [18], Disolver [6], MELY [7], Frodo [8]. Those solvers rely on several algorithms for solving DisCSP problems such as Asynchronous Backtracking (ABT [4], ABT Family [9]), Asynchronous Forward Checking (AFC) [10] and Nogood-based Asynchronous Forward-Checking (AFC-ng) [11]. Asynchronous Distributed Constraints Optimization (ADOPT) [12], Asynchronous Forward Bounding (AFB) [13], Asynchronous Branch-and-Bound (BnB-ADOPT) [14] and Dynamic Backtracking for distributed constraint optimization (DyBop) [15] were developed to solve DCOP problems. As well as the authors recognise that most of these tools are specially developed for simulation context. This fact can be clearly observed from its experimental setups. Given the difficulty that researchers are facing, they often make many simplifying assumptions (i.e. simple agent (one variable per agent), agents as multi-thread, single physical platform, communication via simulated perfect FIFO channels, etc.) about the underlying distributed problem, which might affect the predictions obtained from the simulation in non-trivial ways. Switching from the simulation to the actual development practice often leads to loss of accuracy. Hence, bridging the gap between simulation and actual development and deployment within distributed constraints solvers and include dynamic aspect are the motivations for presenting the different ideas discussed in the present paper.

In this paper we focus on the development of a Multi-agent platform for Distributed Constraint Reasoning and Dynamic Distributed Constraints Problems, namely *JChoc DisSolver*. This proposed platform allows non-expert user to address and solve easily, not only distributed constraint satisfaction problems, but also real Dynamic Distributed Constraint Satisfaction Problems.

This document is organized as follows. Section 2 presents a brief definition of Distributed Constraint Satisfaction Problem (DisCSP) and Dynamic and Distributed Constraint Satisfaction Problem (DDisCSP) with an example. In Sect. 3, we present related work. Section 4 presents the global architecture of *JChoc* platform. In Sect. 5, we show how use this platform in a distributed environment even if it changes dynamically. And finally, in Sect. 6 we conclude the paper by experiment this platform within a real Distributed and Dynamic Constraints Satisfaction Problems.

## 2 Preliminaries

### 2.1 Distributed Constraint Satisfaction Problems

Constraint Programming distinguishes between the description of the constraints involved in a problem on the one hand, and the algorithms and heuristics used to solve the problem on the other hand. Modeling and solving problems is through a very elegant mathematical formalism, called the Constraint Satisfaction Problems CSPs.

The Distributed Constraint Satisfaction Problem (DisCSP) is represented by a constraint network where variables and constraints are distributed among multiple automated agents.

**Definition:** A finite DisCSP is defined by a 5-tuple  $(A, X, D, C, \psi)$ , where:

- $A = \{A_1, \dots, A_p\}$  is a set of  $p$  agents.
- $X = \{x_1, \dots, x_n\}$  is a set of  $n$  variables such that each variable  $x_i$  is controlled by one agent in  $A$ .
- $D = \{D(x_1), \dots, D(x_n)\}$  is a set of current domains, where  $D(x_i)$  is a finite set of possible values for variable  $x_i$ .
- $C = \{c_1, \dots, c_m\}$  is a set of  $m$  constraints that specify the combinations of values allowed for the variables they involve. We note that the constraints are distributed among the automated agents. Hence, constraints divide into two broad classes: inter-agent and intra-agent.
- $\psi : X \rightarrow A$  is a function that maps each variable to its agent.

A solution to a DisCSP is an assignment of a value from its domain to every variable of the distributed constraint network, in such a way that every constraint is satisfied. Solutions to DisCSPs can be found by searching through the possible assignments of values to variables such as ABT algorithm [4].

Many hard practical problems can be seen as DisCSPs. Most Distributed Reasoning platform however assume that problem are static. This has a limitation for dynamic problems that change over time, for example timetabling shifts in a large hospital where availability staff change over time. Also in a dynamic environment a DisCSP may change over time, these changes could be due to addition/deletion of variables, constraints, or agents. The Distributed and Dynamic Constraint Satisfaction Problems (DDisCSPs) can be described as a five tuple  $(A, X, D, C, \delta)$  where:

- $A, X, D$  and  $C$  remain as described in DisCSP.
- $\delta$  is the change function which introduces changes.

Many DDisCSPs approaches (e.i : DynABT [25], DynBDA [26]) are proposed to solve such problems, and can be easily implemented in This platform.

### 2.2 Meeting Scheduling Problem as a DisCSP

The Distributed Meeting Scheduling Problem (MSP) is a real distributed problem where agents may not desire to deliver their personal information to a centralized agent to solve the whole problem [20, 21].

The MSP involves a set of  $n$  agents having a personal private calendar and a set of  $m$  meetings each taking place in a specified location. Each agent,  $A_i \in A$ , knows the set of the  $k_i$  among  $m$  meetings he/she must attend. It is assumed that each agent knows the traveling time between the locations where his/her meetings will be held. The traveling time between locations where two meetings  $m_i$  and  $m_j$  will be held is denoted by  $TravellingTime(m_i, m_j)$ . Solving the problem consists in satisfying the following constraints: (i) all agents attending a meeting must agree on when it will occur, (ii) an agent cannot attend two meetings at same time, (iii) an agent must have enough time to travel from the location where he/she is to the location where the next meeting will be held.

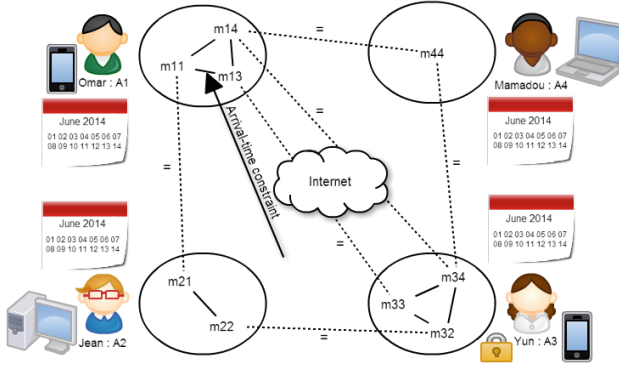
We illustrate in Fig. 1 the encoding of the instance of the meeting scheduling problem in the distributed constraint network formalism. This figure shows 4 agents where each agent has a personal private calendar and a set of meetings each taking place in a specified location. Thus, we get the following DisCSP:

- $A = \{A_1, A_2, A_3, A_4\}$  each agent  $A_i$  corresponds to a real agent,
- For each agent  $A_i \in A$  there is a variable  $m_{ik}$ , for every meeting  $m_k$  that  $A_i$  attends,
- $X = \{m_{11}, m_{13}, m_{14}, m_{21}, m_{22}, m_{32}, m_{33}, m_{34}, m_{44}\}$ .
- $D = \{D(m_{ik}) | m_{ik} \in X\}$  where,
  - \*  $D(m_{11}) = D(m_{13}) = D(m_{14}) = \{s \mid s \text{ is a slot in calendar}(A_1)\}$ .
  - \*  $D(m_{21}) = D(m_{22}) = \{s \mid s \text{ is a slot in calendar}(A_2)\}$ .
  - \*  $D(m_{32}) = D(m_{33}) = D(m_{34}) = \{s \mid s \text{ is a slot in calendar}(A_3)\}$ .
  - \*  $D(m_{44}) = \{s \mid s \text{ is a slot in calendar}(A_4)\}$ .
- For each agent  $A_i$ , there is a private arrival-time constraint ( $c_{kl}^i$  intra-agent constraint) between every pair of its local variables ( $m_{ik}, m_{il}$ ) (e.g. Omar must attend tree meetings  $m_1, m_2$  and  $m_3$ ). For each two agents  $A_i, A_j$  that attend the same meeting  $m_k$  there is an equality inter-agent constraint ( $c_k^{ij}$ ) between the variables  $m_{ik}$  and  $m_{jk}$ , corresponding to the meeting  $m_k$  on agent  $A_i$  and  $A_j$  (e.g. Omar and Jean participate in the same meeting  $m_1$ ). Then,  $C = \{c_{kl}^i, c_k^{ij}\}$ .

Given this example, a Distributed Constraint Reasoning (DCR) platform must allow agents to have a reliable communication tool that allows sending and receiving messages, in order to find the feasible solutions.

### 3 Related Work

Recently, B. Lutati and et al. [23] have proposed a MAS platform, called AgentZero. This tool can be considered as a new addition to the available MAS tools in general and to the DCR research field in particular. The authors claim that AgentZero is generic and applicable to many domains, specifically introducing benefits for the DCR simulation domain. However, the platform has been designed only for simulation use and used only by researchers in Distributed Constraint Reasoning. So developing and setting computer software for real problems based on DCR is not simple and remains a difficult task for users in general.



**Fig. 1.** Meeting scheduling problem modeled as DisCSP.

In [8] A. Petcu. Proposes a Framework for Open Distributed Optimization (FRODO). The framework is implemented in Java, and simulates a multiagent environment in a single Java virtual machine. Each agent in the environment is executed asynchronously in a separate execution thread, and communicates with its peers through message exchange. FRODO comes with several built in algorithms and a suite of problem generators for benchmarking.

The authors of [24] proposed an open-source tool for solving DCR, called DCOPolis. DCOPolis is an open-source framework designed to abstract algorithm implementation from the underlying platform (i.e. hardware, network, operating system). This allows a single implementation of an algorithm to be run in simulation (i.e. on top of the NS2 network simulator with AgentJ). DCOPolis differs from existing DCR frameworks and simulators, however, it supports a novel type of simulation in which the runtime of any distributed algorithm can be accurately estimated on a single physical computer.

Researchers in DCR are concerned with developing new algorithms, and comparing their performance with existing algorithms. Therefore, in [18] the authors present an open source Java library, called DisChoco which aims at implementing DCR algorithms from an abstract model of agent. DisChoco allows to represent both DisCSPs and DCOPs, as opposed to other platforms. A single implementation of a DCR algorithm can run as simulation on a single machine. DisChoco is an elegant platform, but all the different issues of realistic uses and actual deployment have not been addressed.

Developing intelligent software applications based on DCR algorithms is a difficult task, because the programmer must explicitly juggle between many very different concerns, including centralized programming, distributed programming, asynchronous and concurrent management of distributed structures, communication concerns and others. In addition, there are very few open-source tools for solving DCR problems in a physically distributed environment. In this paper we have been looking for a singular platform that would possess not only simulation qualities, but especially designed for realistic and actual deployment. JChoc

platform is a new added value which allows bridging the gap between simulation and realistic use. To our knowledge, this is the first DCR platform respecting FIPA standards and specifications.

## 4 JChoc Platform

### 4.1 JChoc Description

The best way to prove the effectiveness of a proposed distributed constraint reasoning algorithm, is to use it in a realistic multi-platform agent. This is how we can reduce the gap between theory and practice. JChoc is a distributed constraints multiagent platform proposed for solving combinatorial problems within a specific distributed environment. It can also be used to analyze and test the algorithms proposed by constraints programming community. This platform is presented in the form of programming environment (API) and applications to help different types of users. Hence, JChoc can be easily appropriated by two main actors:

- Developers to design and develop applications (e.i. client application, web application, mobile application, etc.) within distributed constraints programming based on JChoc API;
- Non-expert user to interact directly with applications based on distributed constraints programming.

This proposed platform has several advantages:

- A distributed constraints problem can be easily addressed and solved in a realistic environment by unsophisticated users;
- The performances of the proposed protocols (i.e. ABT, AFC, Adopt, etc.) can be actually tested and proved in a realistic communication channel (i.e. WLAN WPAN WMAN WWAN);
- It offers a modular software architecture which accepts extensions easily (i.e. security, confidentiality, cryptography, etc.);
- Thanks to the extensibility of JADE communication model [19], JChoc allows the development of multiagent systems and applications consistent with Foundation for Intelligent Physical Agents (FIPA)<sup>1</sup> standards and specifications;
- Thanks to the robustness of Choco platform [22], complex agent (i.e. multiple variables per agent) can easily address and solve its local sub-problem and use solutions as a compiled domain.

This platform consists of several modules presented as services. The main constraint programming services offered are based Distributed Constraint Reasoning Protocols (DCRP) and Choco Solver (CS). Choco is a platform for research in centralized constraint programming and combinatorial optimization. This choice of Choco enabled us to benefit from the modules already implemented in it. In the next section, we will study the different elements of JChoc platform.

<sup>1</sup> <http://www.fipa.org/>.

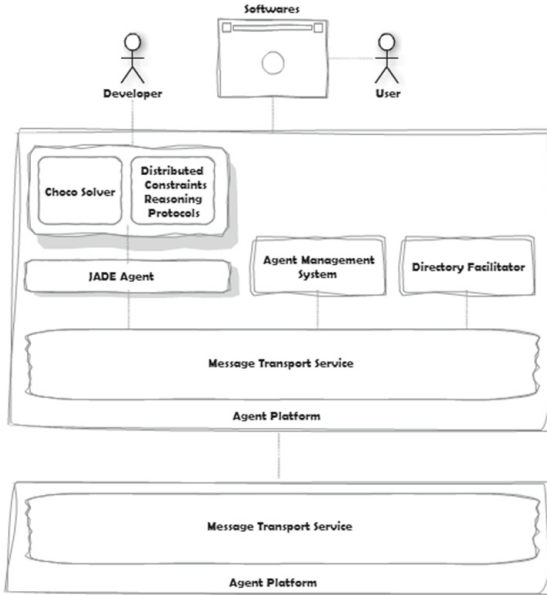


Fig. 2. The JChoc architecture.

## 4.2 JChoc Architecture

JChoc architecture is motivated by FIPA specifications, it allows the development of multiagent systems and applications conforming to MAS standards. It is implemented in JAVA and provides classes that implement and inherit from JADE and Choco platforms to define the behavior of specific agents. Figure 2 represents the main JChoc architectural elements. This platform has five main modules.

- DCRP «Distributed Constraint Reasoning Protocols» provides distributed constraints protocols as service. This element defines new types of messages and implements the behavior of the agent when receiving and sending a specific type of information (e.i. ABT, AFC, Adopt, etc.);
- CS «Choco Solver» provides the ability to address and resolve local CSP sub-problem;
- DF «Director Facilitator» provides a service of “yellow pages” to the platform;
- ACC «Agent Communication Channel» manages the communication between agents;
- AMS «Agent Management System» oversees the registration of agents, their authentication, their access and the use of the system.

These five modules are activated at each time the platform is started.

The JADE agent is also a key player in our platform. Thanks to JADE an Agent Identifier (AID) identifies an agent uniquely.

JChoc uses extensively a sniffing tool for debugging, or simply documenting conversations between agents. The sniffer subscribes to AMS agent to be notified of all platform events and of all message exchanges between a set of specified agents. When the user decides to monitor an agent or a group of agents, every message directed to, or coming from, that agent/group is tracked and displayed in the sniffer GUI. The user can select and view the details of every individual message, save the message or serialize an entire conversation as a binary file.

## 5 Using Dynamic JChoc

### 5.1 Using JChoc in Distributed Environment

In this section we present how to use the JChoc platform in real distributed environment. The MSP problem depicted in Fig. 1 is used to illustrate this proposed platform. Initially we generate a sub-problem for each agent involved in the global DisCSP problem, modeled by an expert as an XML file, which allows standardizing the syntactic structure of the sub-problems. A sub-problem containing only the information necessary for a single agent, so he can participate in solving the global problem in a real distributed environment.

Figure 3 shows an example of representation of the MSP sub-problem defined above in the XDisCSP format. Each variable has a unique ID, which is the concatenation of the ID of its owner agent and index of the variable in the agent. This is necessary when defining constraints (scope of constraints). For constraints, we used two types of constraints: TKC for Totally Known Constraint and PKC for Partially Known Constraint. Constraints can be defined in extension or as a Boolean function. Different types of constraints are pre-defined: equal to  $eq(M_i, M_j)$ , different from  $ne(M_i, M_j)$ , greater than or equal  $ge(M_i, M_j)$ , greater than  $gt(M_i, M_j)$ , etc. In this sub-problem there is 1 complex agent  $A_3$  which controls exactly 3 variables. The domain of  $A_3$  contain 14 values  $D_3 = \{1...14\}$ . There are three constraints of Arrival time  $ge(abs(sub(M_i, M_j))$ : the first constraint is between  $M_{3,2}$  and  $M_{3,3}$  the second one is between  $M_{3,3}$  and  $M_{3,4}$  and the third is between  $M_{3,2}$  and  $M_{3,4}$ , three constraints of equality  $eq(M_i, M_j)$ : between  $M_{1,4}$  and  $M_{3,4}$ , between  $M_{1,3}$  and  $M_{3,3}$ , between  $M_{2,2}$  and  $M_{3,2}$  after defining our sub-problem we can configure our solver.

Once the sub-problem is generated, we can test the functioning of the platform in a physically distributed environment. So we chose machines that simulate the different agents of the problem, and filed each sub-problem in a machine, before launching it.

Figure 4 shows how the master launches its communication interface listening on the network. We start with instantiate the dissolver object (line 7), This class models the distributed problem when JChoc is used to solve a problem in a real distributed environment. All information on distributed problem is encapsulated in this object (identities of agents, inter-agent constraints, protocol, etc.). Then, we define the type of master (line 8) (ABT in this case). Finally, we trigger the container and we launch the master (lines 10–11).



```

1 <?xml version="1.0" encoding="UTF-8"?>
2 <instance>
3   <presentation name="MSP" type="DisCSP"
4     model="Complex" constraintModel="TKC" format="XDisCSP 1.0" />
5   <domains nbDomains="2">
6     <domain name="D1" nbValues="7">1..7</domain>
7   </domains>
8   <variables nbVariables="3">
9     <variable name="M3.2" id="1" domain="D1" description="M.2" />
10    <variable name="M3.3" id="2" domain="D1" description="M.3" />
11    <variable name="M3.4" id="1" domain="D1" description="M.4" />
12  </variables>
13  <constraints nbConstraints="3">
14    <constraint model="TKC" name="C0" reference="ArrivalTime" scope="M3.2 M3.3 2" arity="2">
15      <parameters>M3.2 M3.3 2</parameters>
16    </constraint>
17    <constraint model="TKC" name="C1" reference="ArrivalTime"
18      scope="M3.3 M3.4 2" arity="2">
19      <parameters>M3.3 M3.4 2</parameters>
20    </constraint>
21    <constraint model="TKC" name="C2" reference="ArrivalTime"
22      scope="M3.2 M3.4 2" arity="2">
23      <parameters>M3.2 M3.4 2</parameters>
24    </constraint>
25  </constraints>
26  <predicates nbPredicates="2">
27    <predicate name="ArrivalTime">
28      <parameters>int Mi,int Mj,int cte</parameters>
29      <expression>
30        <functional>ge(abs(sub(Mi,Mj)), cte)</functional>
31      </expression>
32    </predicate>
33    <predicate name="eq">
34      <parameters>int Mi,int Mj</parameters>
35      <expression>
36        <functional>eq(Mi,Mj)</functional>
37      </expression>
38    </predicate>
39  </predicates>
40  <agents.neighbours>
41    <agent name="A1">
42      <constraints nbConstraints="2">
43        <constraint model="TKC" name="C0" reference="eq" scope="M1.4 M3.4" arity="2">
44          <parameters>M1.4 M3.4</parameters>
45        </constraint>
46        <constraint model="TKC" name="C1" reference="eq" scope="M1.3 M3.3" arity="2">
47          <parameters>M1.3 M3.3</parameters>
48        </constraint>
49      </constraints>
50    </agent>
51    <agent name="A2">
52      <constraints nbConstraints="1">
53        <constraint model="TKC" name="C0" reference="eq" scope="M2.2 M3.2" arity="2">
54          <parameters>M2.2 M3.2</parameters>
55        </constraint>
56      </constraints>
57    </agent>
58  </agents.neighbours>
59  <agents.children>
60    <agent name="A4" id="5" variable="M3.4" />
61  </agents.children>
62  </agents.neighbours>
63 </instance>

```

**Fig. 3.** Definition of DMS sub-problem in XDisCSP format.

Figures 5, 6, 7 and 8 show how to launch JChoc agents. We start with instantiate the DisSolver object (line 7), followed by the agent and distributed sub-problem declaration which specifies the resolution algorithm to be used (line 8–9). Next, the declaration of the container containing the master with its IP address (line 10). Eventually, we launch the agent (line 11).

The master waits for the confirmation of creation of all agents before ordering the start of the search. Thus, the problem can be solved using a specified implemented protocol (ABT for example).

```

1 import JChoc.DisSolver;
2
3 public class Master
4 {
5     public static void main(String[] args)
6     {
7         DisSolver js = new DisSolver();
8         js.setType("MasterABT");
9         js.setGui(true);
10        js.setNumberOfAgents(4);
11        js.run();
12    }
13 }
14 }

```

**Fig. 4.** How the master launches its communication interface.

```

1 import JChoc.DisSolver;
2
3 public class Omar
4 {
5     public static void main(String[] args)
6     {
7         DisSolver js1 = new DisSolver();
8         js1.setType("AgentABT");
9         js1.addAgent("A1", "Problem1.xml", true, true);
10        js1.setContainer("192.168.1.8");
11        js1.run();
12    }
13 }

```

**Fig. 5.** How to implement and launch JChoc DisSolver in Omar agent (A1).

```

1 import JChoc.DisSolver;
2
3 public class Jean
4 {
5     public static void main(String[] args) {
6         DisSolver js1 = new DisSolver();
7         js1.setType("AgentABT");
8         js1.addAgent("A2", "Problem2.xml", true, true);
9         js1.setContainer("192.168.1.8");
10        js1.run();
11    }
12 }

```

**Fig. 6.** How to implement and launch JChoc DisSolver in Jean agent (A2).

```

1 import JChoc.DisSolver;
2
3 public class Yun
4 {
5     public static void main(String[] args) {
6         DisSolver js1 = new DisSolver();
7         js1.setType("AgentABT");
8         js1.addAgent("A3", "Problem3.xml", true, true);
9         js1.setContainer("192.168.1.8");
10        js1.run();
11    }
12 }

```

**Fig. 7.** How to implement and launch JChoc DisSolver in Yun agent (A3).

```

1 import JChoc.DisSolver;
2
3 public class Mamadou
4 {
5     public static void main(String[] args) {
6         DisSolver js1 = new DisSolver();
7         js1.setType("AgentABT");
8         js1.addAgent("A4", "Problem4.xml", true, true);
9         js1.setContainer("192.168.1.8");
10        js1.run();
11    }
12 }

```

Fig. 8. How to implement and launch JChoc DisSolver in Mamadou agent (A4).

```

1 <?xml version="1.0" encoding="UTF-8"?>
2 <instance>
3   <presentation name="MSP" type="DisCSP"
4     model="Complex" constraintModel="TKC"
5     format="XDisCSP 1.0" />
6   <domains nbDomains="1">
7     <domain name="D1" nbValues="20">1..20</domain>
8   </domains>
9   <variables nbVariables="1">
10    <variable name="X4.1" id="1" domain="D1" description="X4" />
11  </variables>
12  <predicates nbPredicates="0">
13  </predicates>
14  <constraints>
15  </constraints>
16  <agents.neighbours>
17    <agents.parent>
18      <agent name = "A1">
19        <constraints nbConstraints="1">
20          <constraint model="TKC" name="C1" reference="R1"
21            scope="X1.1 X4.1" arity="2" change = "add_4">
22            <parameters>X1.1 X4.1</parameters>
23          </constraint>
24        </constraints>
25      </agent>
26      <agent name = "A2">
27        <constraints nbConstraints="1">
28          <constraint model="TKC" name="C1" reference="R2"
29            scope="X2.1 X4.1" arity="2" change = "no">
30            <parameters>X2.1 X4.1</parameters>
31          </constraint>
32        </constraints>
33      </agent>
34      <agent name = "A3">
35        <constraints nbConstraints="1">
36          <constraint model="TKC" name="C1" reference="R3"
37            scope="X3.1 X4.1" arity="2" change = "remove_3">
38            <parameters>X3.1 X4.1</parameters>
39          </constraint>
40        </constraints>
41      </agent>
42    </agents.parent>
43    <agents.children>
44      <agent name="A5" id="1" variable="X4.1" />
45    </agents.children>
46  </agents.neighbours>
47  <relations>
48    <relation name="R1" semantics="conflicts">2 8|8 20|5 7|9 20|17 14|
49      5 16|16 7|9 10|19 2|4 13|18 1|1 19|20 6|16 2|
50      5 18|11 14|10 2|19 19|19 18</relation>
51    <relation name="R2" semantics="conflicts">6 16|14 2|19 15|20 2|8 2|
52      4 2|17 20|18 6|7 7|7 10|3 18|18 10|13 15|9 18|
53      14 16|19 6|13 18|3 14|14 20</relation>
54    <relation name="R3" semantics="conflicts">16 12|20 6|8 17|17 5|4 18|
55      12 18|19 5|20 17|15 13|6 5|17 18|3 1|17 12|1 16|
56      2 1|8 5|13 3|17 10|6 20|7 9</relation>
57    <relation name="R4" semantics="conflicts">1 6|5 2|7 19|15 1|17 15|
58      13 7|2 5|8 5|1 14|2 16|4 14|12 14|9 19|3 4|
59      19 18|10 8|9 4|1 2</relation>
60  </relations>
61 </instance>

```

Fig. 9. Definition of dynamic sub-problem in XDisCSP format.

## 5.2 Using JChoc in Dynamic Distributed Environment

The use of JChoc platform in a dynamic environment is not very different to that in the case of distributed static problems. The difference is seen in the xml file that defines the sub-problem of each agent.

To see the platform's exploitation in the Dynamic case of Distributed Satisfaction Problems, we take a random example composed of five agents, each agent has one variable. Figure 9, shows a model of representation of a dynamic sub-problem of an agent that has two constraints with two other agents, 3s after launching, one of the constraints is going to be removed, then after 4s, another link with a third agent will be added.

In addition of the definition of variables, domains and constraints, we define the constraints that will be either added or removed.

After the generation of the dynamic sub-problem, we can launch the resolution following the same approach as before, but instead to insert the name of an XML file of a static sub-problem as argument, we insert the name of dynamic sub-problem XML.

## 6 Experimental Results

### 6.1 Configuration Example

To experiment the JChoc platform in a physically distributed environment, we chose five machines with features **2.93 GHz, CORE(TM) 2 duo** with **2 GB RAM** that simulate agents. These machines are connected via the **WLAN** of our laboratory. We also chose *ABT* algorithm to solve Meeting Scheduling problems (MSP). In Fig. 1 above, we depict an example of problem solved by this platform in a live distributed environment network. This figure illustrates an instance of MSP viewed as DisCSP where each agent has a personal private calendar and a set of meetings each taking place in a specified location. In that example, there are four agents,  $A_1$ ,  $A_2$ ,  $A_3$  and  $A_4$ , and four meetings,  $m_1$ ,  $m_2$ ,  $m_3$  and  $m_4$ . Each agent has its own calendar divided into **14** slots. The time required for traveling among places where meetings can be scheduled is **2 slots**.

We have intentionally limited the number of agents to 4 for this problem needs, but the number of the agents can be easily extended to  $N \gg 4$  for the neediest problems.

Figures 10 and 11 show the GUI of the sniffer agent at the start and the end of ABT resolution. The canvas provides a graphical representation of the messages exchanged between sniffed **ABTagents**, where each arrow represents a message and each color identifies a type of conversation. For example agent  $A_1$  sends an **OK?** message to informs  $A_2$  that he has done a new assignment  $m_{1,1}:\mathbf{1}$  (line 5).

If no new consistent value is found (line 10),  $A_3$  generates a new **nogood**  $m_{1,3}:\mathbf{3} \wedge m_{1,4}:\mathbf{5} \Rightarrow m_{2,2} \neq 5$  by the resolution of existing nogoods. Eventually, the system can stabilize in a state where each agent has a value and no constraint is violated. This state is a global solution and the network has reached quiescence,

meaning that no message is traveling through it (lines 37, 40, 43, 46). Once the solution is found, the master should be advised to spread the stop order to all agents (lines 49–52) (Fig. 11).

A solution to this example is:

$$A_1 \rightarrow (m_{1,1} : 3; m_{1,3} : 7; m_{1,4} : 1), A_2 \rightarrow (m_{2,1} : 3; m_{2,2} : 5), A_3 \rightarrow (m_{3,2} : 5; m_{3,3} : 7; m_{3,4} : 1), A_4 \rightarrow (m_{4,4} : 1).$$

### 6.2 Platform Scalability

The scalability of JChoc is the ability of the system, network, and process to handle a growing amount of work in a capable manner and its ability to be enlarged to accommodate that growth. In order to experiment our platform, we consider a large number of MSP instances. These Meeting Scheduling Problem are characterized by  $\langle m, p, n, d, h, t, a \rangle$ , where  $m$  is the number of meetings,  $p$  is the number of participants,  $n$  is the number of inter-agent constraints  $d$  determines the number of days. Different time slots are available for each meeting, and  $h$  is the number of hours per day,  $t$  is a duration of the meeting and  $a$  is the percentage of availability for each participant. We present our results for the class  $\langle m, p, n, 5, 10, 1, 90\% \rangle$  and we vary three parameters:  $m, p, n$  (each agent has 2 meetings):

As shown in experimental results, in Fig. 7, the performance of our platform is measured in terms of network load (number of messages) and run-time execution. From these preliminary results we see that JChoc platform performs rapidly in small instances ( $\#p \in [4, 14]$ ). The number of messages increases for  $\#p \in [15, 18]$  and reduces for  $\#p > 18$ . This scalability behavior is due to complexity of MSP problems. When the instance is hard the problem can be solved rapidly (Fig. 12).

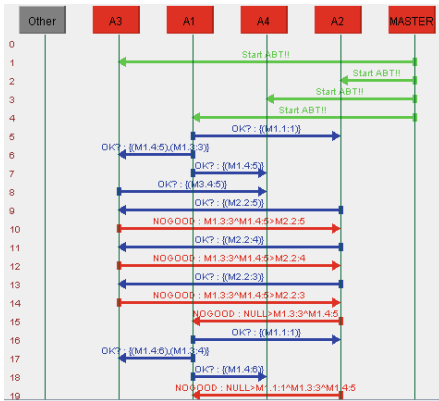


Fig. 10. The start on sniffer agent GUI.

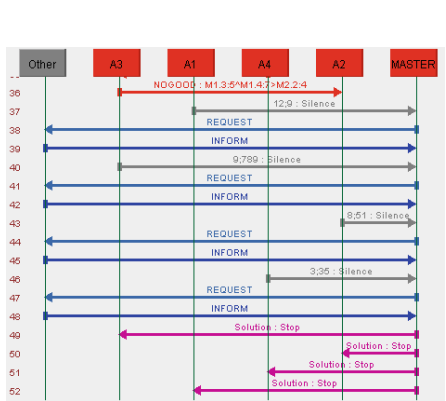
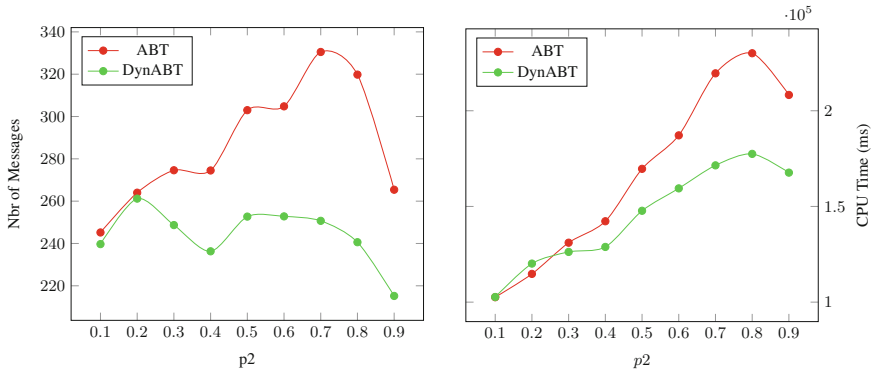


Fig. 11. The finish on sniffer agent GUI.

#p	#m	#n	#messages	Time (ms)
4	8	3	11	17070
5	10	5	11	17204
6	12	6	14	16144
7	14	7	14	17073
8	16	8	19	19180
9	18	9	24	20210
10	20	10	22	18294
11	22	11	32	20197
12	24	12	27	18516
13	26	15	30	20370
14	28	33	51	26073
15	30	35	105	31103
16	32	29	69	28914
17	34	33	175	38324
18	36	35	139	43172
19	38	38	141	37121
20	40	43	94	33457

**Fig. 12.** Performance of JChoc platform using ABT protocol on the Meeting Scheduling Problem (MSP).



**Fig. 13.** ABT vs DynABT.

### 6.3 Platform Scalability in a Dynamic Changed Environment

To compare the performance of the DDisCSPs with a platform that supports dynamic aspect and an other that doesn't. We made our experiments using ABT that cant solve such problem dynamically and resolve the problem when changes are available, and DynABT that can adapt changes and continuous problem's solving. We have introduced a rate change  $\delta$  as a percentage of the total constraints in the problem ( $\delta = 20\%$ ). In these experiments we generated problems randomly, with parameters  $(a, i, n, d, p1, p2)$  using the platform generator, where:  $a$  is the number of agents = 20,  $i$ : the number of instances = 10,  $n$ : the number

of variables = 20,  $p1$ : the density of constraints = 20%, and  $p2$ : the tightness of constraints with value 10%–90% step 10%, the range of tightness 10%–40% contains solvable problems, 50% contains both solvable and unsolvable problems, and 60%–90% problems are unsolvable.

The Fig. 13 shows the number of messages sent and CPU Time, measured for both ABT and DynABT implemented on JChoc platform and using our laboratory's wireless network, that allows the communication between Agents in the same environment and conditions. All results obtained show that DynABT significantly outperforms ABT in a dynamic changed environment. This comparison shows the benefits of solving dynamic distributed problems in a real distributed changed environment with an algorithm that support dynamic aspect implemented in a suitable Platform. The platform is user friendly and lets users implement their Multi-agents applications for dynamic environment.

## 7 Conclusion

In this paper, we have proposed a modular, reliable, deployable and distributed software architecture, called JChoc DisSolver, which can be used easily for several real dynamic combinatorial problems. The main purpose of our platform is to break down the barriers to building new and innovative applications. The possibility of combining the expressiveness of Choco, the extensibility of JADE and our powerful Dynamic Distributed Constraint Reasoning Add-On bring a strong added value in the development of innovative applications based on Constraints Programming paradigm. The JChoc platform presented in this paper has been designed to support extensions: security, cryptography. In our experiments, We have implemented ABT protocol and solved the Meeting Scheduling problem (MSP) in a real distributed environment. In a dynamic environment, we have solved dynamic problems with DynABT, to show the benefits of our platform that supports the dynamic aspect. We found that, by using this platform, we can adopt easily any proposed protocol for solving distributed constraint problem even if environment changes dynamically. Future investigations are focusing on enhancing the platform, by adding other layers that will allow users to implement their multi-robots applications easily. The platform will allow the communication between robots in a dynamically changing environment.

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