From Algebra to Logic: There and Back Again The Story of a Hierarchy*-*

(Invited Paper)

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Abstract. This is a survey about a collection of results about a (double) hierarchy of classes of regular languages, which occurs in a natural fashion in a number of contexts. One of these occurrences is given by an alternated sequence of deterministic and co-deterministic closure operations, starting with the piecewise testable languages. Since these closure operations preserve varieties of languages, this defines a hierarchy of varieties, and through Eilenberg's variety theorem, a hierarchy of pseudovarieties (classes of finite monoids that are defined by pesudo-identities). The point of this excursion through algebra is that it provides reasonably simple decision algorithms for the membership problem in the corresponding varieties of languages. Another interesting point is that the hierarchy of pseudo-varieties bears a formal resemblance with another hierarchy, the hierarchy of varieties of idempotent monoids, which was much studied in the 1970s and 1980s and is by now well understood. This resemblance provides keys to a combinatorial characterization of the different levels of our hierarchies, which turn out to be closely related with the so-called rankers, a specification mechanism which was introduced to investigate the two-variable fragment of the first-order theory of the linear order. And indeed the union of the varieties of languages which we consider coincides with the languages that can be defined in that fragment. Moreover, the quantifier alternation hierarchy within that logical fragment is exactly captured by our hierarchy of languages, thus establishing the decidability of the alternation hierarchy.

There are other combinatorial and algebraic approaches of the same logical hierarchy, and one recently introduced by Krebs and Straubing also establishes decidability. Yet the algebraic operations involved are seemingly very different, an intriguing problem. . .

Formal language theory historically arose from the definition of models of computation (automata, grammars, etc) and relied for its first step on combinatorial reasoning, especially combinatorics on words. Very quickly however, algebra and

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logic were identified as powerful tools for the classification of rational languages, e.g. with the definition of the syntactic monoid of a language and Büchi's theorem on monadic second-order logic. It did not take much time after that to observe that, conversely, formal language theory is itself a tool for algebra and logic.

The results which we will present are an illustration of this back-and-forth movement between languages, algebra and logic. They deal with a hierarchy of classes of rational languages which arises in different contexts and turned out to solve a problem in logic, namely the decidability of t[he](#page-2-0) [qu](#page-2-1)antifier alternation hierarchy within the two-variable fragment of first-order logic $\text{FO}^2[\lt]$.

The full picture uses a collection of results in logic, combinatorics on words and algebra which were obtained independently of the quantifier alternation hierarchy by various authors over several decades.

Let \mathcal{R}_0 be the c[la](#page-2-2)[ss o](#page-2-3)f piecewise testable languages, which is natural from a combinatorial and automata-theoretic point of view, and corresponds to the first level of the quantifier alternation hierarchy within $\text{FO}^2[<]$ (and within $\text{FO}[\lt]$) as well). This class i[s r](#page-2-4)ather simp[le](#page-2-5) [and](#page-2-6) [re](#page-3-0)asonably well understood, see [1, 10]. We first consider the hierarchies of classes of languages obtained from \mathcal{R}_0 by alternatingly closing it under deterministic and co-deterministic closure: we let $\mathcal{L}_0 = \mathcal{R}_0, \mathcal{R}_{k+1}$ (resp. \mathcal{L}_{k+1}) be the deterministic (resp. co-deterministic) closure of \mathcal{L}_k (resp. \mathcal{R}_k).

Results from the 1970s and 1980s [9, [12\]](#page-2-7) show that the classes \mathcal{R}_k and \mathcal{L}_k are varieties (whether a language *L* belongs to one of these classes depends only on its syntactic monoi[d\)](#page-2-8) and describe the corresponding varieties of finite monoids \mathbf{R}_k and \mathbf{L}_k . Results from the 1960s [5] (see also [6, 11, 17]) shows that their membership problems are decidable and they form [an](#page-2-9) infinite hierarchy.

A first view of the structure of the lattice formed by these varieties can be obtained by using purely algebraic results from the 1970s on a seemingly different hierarchy, that of varieties of idempotent monoids [2]. The theory of the latter varieties is particularly well underst[ood,](#page-3-1) and one can exhibit for each of them structurally elegant identities [and](#page-2-10) solutions of the word problem (of the corresponding relatively free object) [3].

To completely elucidate the structure of the lattice generated by the \mathcal{R}_k and \mathcal{L}_k , Kufleitner and Weil introduced the [not](#page-2-11)ion of condensed rankers [8]. These are a rather natural extension of the algorithm to solve the word problem in th[e re](#page-2-12)latively free idempo[ten](#page-3-2)t monoids and have natural connections with deterministic and codeterministic products. But they are also – and foremost – a variant of the rankers introduced by Weiss and Immerman [18] (following the turtle programs of Schwentick, Thérien and Vollmer $[13]$) to characterize the levels of the quantifier alternation hierarchy of $\text{FO}^2[\langle \cdot]$. As a result one can show that the *k*-th level of this hierarchy coincides with the intersection $\mathcal{R}_{k+1} \cap \mathcal{L}_{k+1}$, thus proving the decidability of each level of the hierarchy [7].

The story does not end there: using algebraic methods similar to those described in his book [15], Straubing showed [16] that the *k*-th level of the quantifier alternation hierarchy of $FO^2\leq$ is the variety of languages whose syntactic monoid is in the *k*-th term of the sequence given by $V_1 = J$ and $V_{n+1} = V_n \square J$. Here J is the class of J -trivial monoids, which characterizes piecewise testable languages by Simon's theorem [14] and \square denotes the two-sided block product, the bilateral version of the more classical wreath product. Then Straubing and Krebs showed that every one of these classes of finite monoids is decidable [4], thus providing an alternate proof of the decidability of the quantifier alternation hierarchy, but also giving an alternative characterization of the classes V_k : finite monoid M is in V_k if and only if it sits in both \mathbf{R}_{k+1} and \mathbf{L}_{k+1} .

The coincidence of these two very differently defined hierarchies raises an intriguing question: what connects the block product with the alternate operation of deterministic and co-deterministic closure?. . .

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