

Antitone \mathbf{L} -bonds*

Jan Konecny

Data Analysis and Modeling Lab, Department of Computer Science
Faculty of Science, Palacky University, Olomouc, Czech Republic
`jan.konecny@upol.cz`

Abstract. L -bonds represent relationships between fuzzy formal contexts. We study these intercontextual structures w.r.t. antitone Galois connections in fuzzy setting. Furthermore, we define direct \triangleleft -product and \triangleright -product of two formal fuzzy contexts and show conditions under which a fuzzy bond can be obtained as an intent of the product. This extends our previous work on isotone fuzzy bonds.

1 Introduction

Formal Concept Analysis (FCA) [10] is an exploratory method of analysis of relational data. The method identifies some interesting clusters (formal concepts) in a collection of objects and their attributes (formal context) and organizes them into a structure called concept lattice. Formal Concept Analysis in fuzzy setting [3] allows us to work with graded data.

In the present paper, we deal with intercontextual relationships in FCA in fuzzy setting. Particularly, our approach originated in relation to [16] on the notion of Chu correspondences between formal contexts, which led to obtaining information about the structure of \mathbf{L} -bonds. In [15] we studied properties of \mathbf{L} -bonds w.r.t. isotone concept-forming operators.

The present paper concerns with \mathbf{L} -bonds with antitone character; We describe their properties and explain how these \mathbf{L} -bonds relate to the structures studied in [16]. In addition, we also focus on the direct products of two formal fuzzy contexts and show conditions under which a bond can be obtained as an intent of the product.

The paper is structured as follows: in Section 2 we recollect some notions used in this paper; in Section 3 we define the \mathbf{L} -bonds and direct products, and describe their properties. Our conclusions and related further research are summarized in Section 4.

2 Preliminaries

In this section, we recall some basic notions used in the paper.

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2.1 Residuated Lattices and Fuzzy Sets

We use complete residuated lattices as basic structures of truth-degrees. A complete residuated lattice [3,12,21] is a structure $\mathbf{L} = \langle L, \wedge, \vee, \otimes, \rightarrow, 0, 1 \rangle$ such that

- (i) $\langle L, \wedge, \vee, 0, 1 \rangle$ is a complete lattice, i.e. a partially ordered set in which arbitrary infima and suprema exist;
- (ii) $\langle L, \otimes, 1 \rangle$ is a commutative monoid, i.e. \otimes is a binary operation which is commutative, associative, and $a \otimes 1 = a$ for each $a \in L$;
- (iii) \otimes and \rightarrow satisfy adjointness, i.e. $a \otimes b \leq c$ iff $a \leq b \rightarrow c$.

0 and 1 denote the least and greatest elements. The partial order of \mathbf{L} is denoted by \leq . Throughout this paper, \mathbf{L} denotes an arbitrary complete residuated lattice.

Elements a of L are called truth degrees. Operations \otimes (multiplication) and \rightarrow (residuum) play the role of a (truth functions of) “fuzzy conjunction” and “fuzzy implication”. Furthermore, we define the complement of $a \in L$ as

$$\neg a = a \rightarrow 0. \quad (1)$$

L-sets and L-relations An \mathbf{L} -set (or fuzzy set) A in a universe set X is a mapping assigning to each $x \in X$ some truth degree $A(x) \in L$. The set of all \mathbf{L} -sets in a universe X is denoted L^X , or \mathbf{L}^X if the structure of \mathbf{L} is to be emphasized.

The operations with \mathbf{L} -sets are defined componentwise. For instance, the intersection of \mathbf{L} -sets $A, B \in L^X$ is an \mathbf{L} -set $A \cap B$ in X such that $(A \cap B)(x) = A(x) \wedge B(x)$ for each $x \in X$, etc. An \mathbf{L} -set $A \in L^X$ is also denoted $\{A(x)/x \mid x \in X\}$. If for all $y \in X$ distinct from x_1, x_2, \dots, x_n we have $A(y) = 0$, we also write

$$\{A(x_1)/x_1, A(x_2)/x_2, \dots, A(x_n)/x_n\}.$$

An \mathbf{L} -set $A \in L^X$ is called crisp if $A(x) \in \{0, 1\}$ for each $x \in X$. Crisp \mathbf{L} -sets can be identified with ordinary sets. For a crisp A , we also write $x \in A$ for $A(x) = 1$ and $x \notin A$ for $A(x) = 0$. An \mathbf{L} -set $A \in L^X$ is called empty (denoted by \emptyset) if $A(x) = 0$ for each $x \in X$. For $a \in L$ and $A \in L^X$, the \mathbf{L} -sets $a \otimes A, a \rightarrow A, A \rightarrow a$, and $\neg A$ in X are defined by

$$(a \otimes A)(x) = a \otimes A(x), \quad (2)$$

$$(a \rightarrow A)(x) = a \rightarrow A(x), \quad (3)$$

$$(A \rightarrow a)(x) = A(x) \rightarrow a, \quad (4)$$

$$\neg A(x) = A(x) \rightarrow 0. \quad (5)$$

An a -complement is an \mathbf{L} -set A which satisfies $(A \rightarrow a) \rightarrow a = A$.

Binary \mathbf{L} -relations (binary fuzzy relations) between X and Y can be thought of as \mathbf{L} -sets in the universe $X \times Y$. That is, a binary \mathbf{L} -relation $I \in L^{X \times Y}$ between a set X and a set Y is a mapping assigning to each $x \in X$ and each $y \in Y$ a truth degree $I(x, y) \in L$ (a degree to which x and y are related by I).

$V \subseteq L^X$ is called an \mathbf{L} -closure system if

- V is closed under left \rightarrow -multiplication (or \rightarrow -shift), i.e. for every $a \in L$ and $C \in V$ we have $a \rightarrow C \in V$,
- V is closed under intersection, i.e. for $C_j \in V$ ($j \in J$) we have $\bigcap_{j \in J} C_j \in V$.

$V \subseteq L^X$ is called an \mathbf{L} -interior system if

- V is closed under left \otimes -multiplication, i.e. for every $a \in L$ and $C \in V$ we have $a \otimes C \in V$,
- V is closed under union, i.e. for $C_j \in V$ ($j \in J$) we have $\bigcup_{j \in J} C_j \in V$.

Relational products We use three relational product operators, \circ , \triangleleft , and \triangleright , and consider the corresponding products $R = S \circ T$, $R = S \triangleleft T$, and $R = S \triangleright T$ (for $R \in L^{X \times Z}$, $S \in L^{X \times Y}$, $T \in L^{Y \times Z}$). In the compositions, $R(x, z)$ is interpreted as the degree to which the object x has the attribute z ; $S(x, y)$ as the degree to which the factor y applies to the object x ; $T(y, z)$ as the degree to which the attribute z is a manifestation (one of possibly several manifestations) of the factor y . The composition operators are defined by

$$(S \circ T)(x, z) = \bigvee_{y \in Y} S(x, y) \otimes T(y, z), \quad (6)$$

$$(S \triangleleft T)(x, z) = \bigwedge_{y \in Y} S(x, y) \rightarrow T(y, z), \quad (7)$$

$$(S \triangleright T)(x, z) = \bigwedge_{y \in Y} T(y, z) \rightarrow S(x, y). \quad (8)$$

Note that these operators were extensively studied by Bandler and Kohout, see e.g. [13]. They have natural verbal descriptions. For instance, $(S \circ T)(x, z)$ is the truth degree of the proposition “there is factor y such that y applies to object x and attribute z is a manifestation of y ”; $(S \triangleleft T)(x, z)$ is the truth degree of “for every factor y , if y applies to object x then attribute z is a manifestation of y ”. Note also that for $L = \{0, 1\}$, $S \circ T$ coincides with the well-known composition of binary relations.

We will need following lemma.

Lemma 1 ([3]). *For $R \in L^{W \times X}$, $S \in L^{X \times Y}$, $T \in L^{Y \times Z}$ we have*

$$R \triangleleft (S \triangleleft T) = (R \circ S) \triangleleft T \quad \text{and} \quad R \triangleright (S \circ T) = (R \triangleright S) \triangleright T.$$

2.2 Formal Concept Analysis in the Fuzzy Setting

An \mathbf{L} -context is a triplet $\langle X, Y, I \rangle$ where X and Y are (ordinary) sets and $I \in L^{X \times Y}$ is an \mathbf{L} -relation between X and Y . Elements of X are called objects, elements of Y are called attributes, I is called an incidence relation. $I(x, y) = a$ is read: “The object x has the attribute y to degree a .” An \mathbf{L} -context is usually depicted as a table whose rows correspond to objects and whose columns correspond to attributes; entries of the table contain the degrees $I(x, y)$.

Concept-forming operators induced by an \mathbf{L} -context $\langle X, Y, I \rangle$ are the following operators: First, the pair $\langle \uparrow, \downarrow \rangle$ of operators $\uparrow : L^X \rightarrow L^Y$ and $\downarrow : L^Y \rightarrow L^X$ is defined by

$$A^\uparrow(y) = \bigwedge_{x \in X} A(x) \rightarrow I(x, y), \quad B^\downarrow(x) = \bigwedge_{y \in Y} B(y) \rightarrow I(x, y). \quad (9)$$

Second, the pair $\langle \wedge, \vee \rangle$ of operators $\wedge : L^X \rightarrow L^Y$ and $\vee : L^Y \rightarrow L^X$ is defined by

$$A^\wedge(y) = \bigvee_{x \in X} A(x) \otimes I(x, y), \quad B^\vee(x) = \bigwedge_{y \in Y} I(x, y) \rightarrow B(y), \quad (10)$$

Third, the pair $\langle \wedge, \vee \rangle$ of operators $\wedge : L^X \rightarrow L^Y$ and $\vee : L^Y \rightarrow L^X$ is defined by

$$A^\wedge(y) = \bigwedge_{x \in X} I(x, y) \rightarrow A(x), \quad B^\vee(x) = \bigvee_{y \in Y} B(y) \otimes I(x, y), \quad (11)$$

for $A \in L^X$, $B \in L^Y$. When we need to emphasize that a pair of concept-forming operators is induced by a particular \mathbf{L} -relation we write it as a subscript, for instance we write \uparrow_I instead of just \uparrow .

Furthermore, denote the corresponding sets of fixed points by $\mathcal{B}^{\uparrow\downarrow}(X, Y, I)$, $\mathcal{B}^{\wedge\vee}(X, Y, I)$, and $\mathcal{B}^{\wedge\vee}(X, Y, I)$, i.e.

$$\begin{aligned} \mathcal{B}^{\uparrow\downarrow}(X, Y, I) &= \{\langle A, B \rangle \in L^X \times L^Y \mid A^\uparrow = B, B^\downarrow = A\}, \\ \mathcal{B}^{\wedge\vee}(X, Y, I) &= \{\langle A, B \rangle \in L^X \times L^Y \mid A^\wedge = B, B^\vee = A\}, \\ \mathcal{B}^{\wedge\vee}(X, Y, I) &= \{\langle A, B \rangle \in L^X \times L^Y \mid A^\wedge = B, B^\vee = A\}. \end{aligned}$$

The sets of fixpoints are complete lattices [1,11,20], called \mathbf{L} -concept lattices associated to I , and their elements are called formal concepts.

For a concept lattice $\mathcal{B}^{\Delta\nabla}(X, Y, I)$, where $\mathcal{B}^{\Delta\nabla}$ is either of $\mathcal{B}^{\uparrow\downarrow}$, $\mathcal{B}^{\wedge\vee}$, or $\mathcal{B}^{\wedge\vee}$, denote the corresponding sets of extents and intents by $\text{Ext}^{\Delta\nabla}(X, Y, I)$ and $\text{Int}^{\Delta\nabla}(X, Y, I)$. That is,

$$\begin{aligned} \text{Ext}^{\Delta\nabla}(X, Y, I) &= \{A \in L^X \mid \langle A, B \rangle \in \mathcal{B}^{\Delta\nabla}(X, Y, I) \text{ for some } B\}, \\ \text{Int}^{\Delta\nabla}(X, Y, I) &= \{B \in L^Y \mid \langle A, B \rangle \in \mathcal{B}^{\Delta\nabla}(X, Y, I) \text{ for some } A\}. \end{aligned}$$

The operators induced by an \mathbf{L} -context and their sets of fixpoints have been extensively studied, see e.g. [1,2,4,11,20].

3 L-bonds

This section introduces antitone \mathbf{L} -bonds, namely a-bonds and c-bonds, and describes their properties.

Definition 1. (a) An a-bond from \mathbf{L} -context $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle$ to \mathbf{L} -context $\mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ is an \mathbf{L} -relation $\beta \in L^{X_1 \times Y_2}$ s.t.

$$\text{Ext}^{\uparrow\downarrow}(X_1, Y_2, \beta) \subseteq \text{Ext}^{\wedge\vee}(X_1, Y_1, I_1) \text{ and } \text{Int}^{\uparrow\downarrow}(X_1, Y_2, \beta) \subseteq \text{Int}^{\uparrow\downarrow}(X_2, Y_2, I_2).$$

(b) A c-bond from \mathbf{L} -context $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle$ to \mathbf{L} -context $\mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ is an \mathbf{L} -relation $\beta \in L^{X_1 \times Y_2}$ s.t.

$$\text{Ext}^{\uparrow\downarrow}(X_1, Y_2, \beta) \subseteq \text{Ext}^{\uparrow\downarrow}(X_1, Y_1, I_1) \text{ and } \text{Int}^{\uparrow\downarrow}(X_1, Y_2, \beta) \subseteq \text{Int}^{\wedge\vee}(X_2, Y_2, I_2).$$

Remark 1. 1) The terms—a-bond and c-bond—were chosen to match with notions of a-morphism and c-morphism [7,14,9]. We show in Theorem 2 that the a-bonds and c-bonds are in one-to-one correspondence of a-morphisms and c-morphisms, respectively, on sets of intents of associated concept lattices.

2) Note that all considered sets of extents and intents in Definition 1 are \mathbf{L} -closure systems. From this point of view, the condition of subsethood is natural.

Theorem 1. (a) $\beta \in L^{X_1 \times Y_2}$ is an a-bond between $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle$ and $\mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ iff there exist \mathbf{L} -relations $S_i \in L^{Y_1 \times Y_2}$ and $S_e \in L^{X_1 \times X_2}$, such that

$$\beta = I_1 \triangleleft S_i = S_e \triangleleft I_2. \quad (12)$$

(b) $\beta \in L^{X_1 \times Y_2}$ is a c-bond between $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle$ and $\mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ iff there exist \mathbf{L} -relations $S_i \in L^{Y_1 \times Y_2}$ and $S_e \in L^{X_1 \times X_2}$, such that

$$\beta = I_1 \triangleright S_i = S_e \triangleright I_2. \quad (13)$$

Proof. Follows from results in [9]. \square

3.1 Morphisms

This section explains correspondence of \mathbf{L} -bonds with morphisms of \mathbf{L} -interior/ \mathbf{L} -closure spaces. First, we recall notions of c-morphisms and a-morphisms. These morphisms were previously studied in [7,9,14].

Definition 2. (a) A mapping $h : V \rightarrow W$ from an \mathbf{L} -interior system $V \subseteq L^X$ into an \mathbf{L} -closure system $W \subseteq L^Y$ is called an a-morphism if

- $h(a \otimes C) = a \rightarrow h(C)$ for each $a \in L$ and $C \in V$;
- $h(\bigvee_{k \in K} C_k) = \bigwedge_{k \in K} h(C_k)$ for every collection of $C_k \in V$.

An a-morphism $h : V \rightarrow W$ is called an extendable a-morphism if h can be extended to an a-morphism of L^X to L^Y , i.e. if there exists an a-morphism $h' : L^X \rightarrow L^Y$ such that for every $C \in V$ we have $h'(C) = h(C)$.

(b) A mapping $h : V \rightarrow W$ from an \mathbf{L} -closure system $V \subseteq L^X$ into an \mathbf{L} -closure system $W \subseteq L^Y$ is called a c-morphism if it is a \rightarrow - and \bigwedge -morphism and it preserves a-complements, i.e. if

- $h(a \rightarrow C) = a \rightarrow h(C)$ for each $a \in L$ and $C \in V$;
- $h(\bigwedge_{k \in K} C_k) = \bigwedge_{k \in K} h(C_k)$ for every collection of $C_k \in V$ ($k \in K$);
- if C is an a-complement then $h(C)$ is an a-complement.

A c -morphism $h : V \rightarrow W$ is called an extendable c -morphism if h can be extended to a c -morphism of L^X to L^Y , i.e. if there exists a c -morphism $h' : L^X \rightarrow L^Y$ such that for every $C \in V$ we have $h'(C) = h(C)$.

In this paper we consider only extendable $\{a,c\}$ -morphisms.

Theorem 2. (a) The a -bonds between $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle$ and $\mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ are in one-to-one correspondence with

- a -morphisms from $\text{Int}^{\cup}(X_1, Y_1, I_1)$ to $\text{Int}^{\downarrow}(X_2, Y_2, I_2)$;
- c -morphisms from $\text{Ext}^{\downarrow}(X_2, Y_2, I_2)$ to $\text{Ext}^{\cup}(X_1, Y_1, I_1)$.

(b) The c -bonds between $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle$ and $\mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ are in one-to-one correspondence with

- c -morphisms from $\text{Int}^{\downarrow}(X_1, Y_1, I_1)$ to $\text{Int}^{\vee}(X_2, Y_2, I_2)$;
- a -morphisms from $\text{Ext}^{\vee}(X_2, Y_2, I_2)$ to $\text{Ext}^{\downarrow}(X_1, Y_1, I_1)$.

Proof. Follows from Theorem 1 and results in [9,14].

Theorem 3. (a) The system of all a -bonds is an \mathbf{L} -closure system.

(b) The system of all c -bonds is an \mathbf{L} -closure system.

Proof. (a) Consider a collection of a -bonds β_i . By Theorem 1 the β_i s are in the form $\beta_i = I_1 \triangleleft S_i = S_e \triangleleft I_2$. We have

$$\begin{aligned} \bigcap_{j \in J} \beta_j &= \bigcap_{j \in J} (I_1 \triangleleft S_{i_j}) = I_1 \triangleleft \left(\bigcap_{j \in J} S_{i_j} \right) \\ &= \bigcap_{j \in J} (S_{e_j} \triangleleft I_2) = \left(\bigcup_{j \in J} S_{e_j} \right) \triangleleft I_2; \\ a \rightarrow \beta &= a \rightarrow (I_1 \triangleleft S_i) = I_1 \triangleleft (a \rightarrow S_i) \\ &= a \rightarrow (S_e \triangleleft I_2) = (a \otimes S_e) \triangleleft I_2. \end{aligned}$$

Thus, $\bigcap_{j \in J} \beta_j$ and $a \rightarrow \beta$ are a -bonds. Proof of (b) is similar. \square

3.2 Direct Products

In this part, we focus on direct products of \mathbf{L} -contexts related to a -bonds and c -bonds.

Definition 3. Let $\mathbb{K}_1 = \langle X_1, Y_1, I_1 \rangle, \mathbb{K}_2 = \langle X_2, Y_2, I_2 \rangle$ be \mathbf{L} -contexts.

(a) A direct \triangleleft -product of \mathbb{K}_1 and \mathbb{K}_2 is defined as the \mathbf{L} -context $\mathbb{K}_1 \boxtimes \mathbb{K}_2 = \langle X_2 \times Y_1, X_1 \times Y_2, \Delta \rangle$ with $\Delta(\langle x_2, y_1 \rangle, \langle x_1, y_2 \rangle) = I_1(x_1, y_1) \rightarrow I_2(x_2, y_2)$ for all $x_1 \in X_1, x_2 \in X_2, y_1 \in Y_1, y_2 \in Y_2$.

(b) A direct \triangleright -product of \mathbb{K}_1 and \mathbb{K}_2 is defined as the \mathbf{L} -context $\mathbb{K}_1 \boxplus \mathbb{K}_2 = \langle X_2 \times Y_1, X_1 \times Y_2, \Delta \rangle$ with $\Delta(\langle x_2, y_1 \rangle, \langle x_1, y_2 \rangle) = I_2(x_2, y_2) \rightarrow I_1(x_1, y_1)$ for all $x_1 \in X_1, x_2 \in X_2, y_1 \in Y_1, y_2 \in Y_2$.

The following theorem shows that $\mathbb{K}_1 \boxplus \mathbb{K}_2$ (resp. $\mathbb{K}_1 \boxminus \mathbb{K}_2$) induces a-bonds (resp. c-bonds) as its intents.

Theorem 4. (a) The intents of $\mathbb{K}_1 \boxplus \mathbb{K}_2$ w.r.t $\langle \uparrow, \downarrow \rangle$ are a-bonds from \mathbb{K}_1 to \mathbb{K}_2 , i.e for each $\phi \in L^{X_2 \times Y_1}$, ϕ^\uparrow is an a-bond from \mathbb{K}_1 to \mathbb{K}_2 .
(b) The intents of $\mathbb{K}_1 \boxminus \mathbb{K}_2$ w.r.t $\langle \uparrow, \downarrow \rangle$ are c-bonds from \mathbb{K}_1 to \mathbb{K}_2 , i.e for each $\phi \in L^{X_2 \times Y_1}$, ϕ^\uparrow is a c-bond from \mathbb{K}_1 to \mathbb{K}_2 .

Proof. (a) For $\phi \in L^{X_2 \times Y_1}$ we have

$$\begin{aligned}
\phi^\uparrow(x_1, y_2) &= \bigwedge_{\langle x_2, y_1 \rangle \in X_2 \times Y_1} \phi(x_2, y_1) \rightarrow \Delta(\langle x_2, y_1 \rangle, \langle x_1, y_2 \rangle) \\
&= \bigwedge_{x_2 \in X_2} \bigwedge_{y_1 \in Y_1} \phi(x_2, y_1) \rightarrow (I_1(x_1, y_1) \rightarrow I_2(x_2, y_2)) \\
&= \bigwedge_{x_2 \in X_2} \bigwedge_{y_1 \in Y_1} (I_1(x_1, y_1) \rightarrow (\phi(x_2, y_1) \rightarrow I_2(x_2, y_2))) \\
&= \bigwedge_{y_1 \in Y_1} (I_1(x_1, y_1) \rightarrow \bigwedge_{x_2 \in X_2} (\phi(x_2, y_1) \rightarrow I_2(x_2, y_2))) \\
&= \bigwedge_{y_1 \in Y_1} (I_1(x_1, y_1) \rightarrow \bigwedge_{x_2 \in X_2} (\phi^\top(y_1, x_2) \rightarrow I_2(x_2, y_2))) \\
&= \bigwedge_{y_1 \in Y_1} I_1(x_1, y_1) \rightarrow (\phi^\top \triangleleft I_2)(y_1, y_2) \\
&= (I_1 \triangleleft (\phi^\top \triangleleft I_2))(x_1, y_2) \\
&= ((I_1 \circ \phi^\top) \triangleleft I_2)(x_1, y_2).
\end{aligned}$$

Thus ϕ^\uparrow is an a-bond by Theorem 1. Proof of (b) is similar. \square

Not all a-bonds are intents of the direct product as the following examples shows.

Example 1. Consider \mathbf{L} -context $\mathbb{K} = \langle \{x\}, \{y\}, \{^{0.5}/\langle x, y \rangle\} \rangle$ with \mathbf{L} being the three-element Łukasiewicz chain. Obviously, $\{^{0.5}/\langle x, y \rangle\}$ is an a-bond from \mathbb{K} to \mathbb{K} . We have $\mathbb{K} \boxplus \mathbb{K} = \langle \{\langle x, y \rangle\}, \{\langle x, y \rangle\}, \{\langle x, y \rangle, \langle x, y \rangle\} \rangle$. The only intent of $\mathbb{K} \boxplus \mathbb{K}$ is $\{\langle x, y \rangle\}$; thus the a-bond $\{^{0.5}/\langle x, y \rangle\}$ is not among its intents.

Example 2. Consider following \mathbf{L} -context with \mathbf{L} being three-element Łukasiewicz chain.

$$\mathbb{K}_1 = \overline{\begin{matrix} 0 & 0 & 0 & \frac{1}{2} \\ 1 & 0 & \frac{1}{2} & \frac{1}{2} \end{matrix}} \quad \mathbb{K}_2 = \overline{\begin{matrix} 0 & 1 & 1 \\ 1 & 1 & 1 \\ \frac{1}{2} & \frac{1}{2} & 1 \end{matrix}}$$

There are 11 a-bonds from \mathbb{K}_1 to \mathbb{K}_2 , but $\mathbb{K}_1 \boxplus \mathbb{K}_2$ has only 9 concepts; see Figure 1.

Since the definition of direct \triangleleft -product and direct \triangleright -product differ only in the direction of residuum, we can make the following corollary.

Theorem 5. *Let the double negation law hold true in \mathbf{L} . the strong antitone \mathbf{L} -bonds from $\langle X_1, Y_1, I_1 \rangle$ to $\langle X_2, Y_2, I_2 \rangle$ are exactly a -bonds from $\langle X_1, Y_1, \neg I_1 \rangle$ to $\langle X_2, Y_2, I_2 \rangle$; and c -bonds from $\langle X_1, Y_1, I_1 \rangle$ to $\langle X_2, Y_2, \neg I_2 \rangle$.*

Note that the incidence relation Δ in direct product $\mathbb{K}_1 \boxtimes \mathbb{K}_2$ then becomes

$$\Delta(\langle x_2, y_1 \rangle, \langle x_1, y_2 \rangle) = \neg I_1(x_1, y_1) \rightarrow I_2(x_2, y_2);$$

that is in agreement with results in [16]. Similarly, the incidence relation Δ in direct product $\mathbb{K}_1 \boxtimes \mathbb{K}_2$ becomes

$$\Delta(\langle x_2, y_1 \rangle, \langle x_1, y_2 \rangle) = I_1(x_1, y_1) \rightarrow \neg I_2(x_2, y_2).$$

Using an alternative notion of complement. The mutual reducibility of concept-forming operators (9)–(11) does not hold generally. In [8], we proposed a new notion of complement of \mathbf{L} -relation to overcome that. Using this notion we showed that each for each $I \in L^{X \times Y}$, one can define $\neg I \in L^{X \times (Y \times L)}$ as

$$\neg I(x, \langle y, a \rangle) = I(x, y) \rightarrow a,$$

and obtain

$$\text{Ext}^{\uparrow\downarrow}(X, Y \times L, \neg I) = \text{Ext}^{\text{ou}}(X, Y, I)$$

and, similarly,

$$\text{Int}^{\uparrow\downarrow}(X, Y \times L, \neg I) = \text{Int}^{\wedge\vee}(X \times L, Y, (\neg I^{\text{T}})^{\text{T}})$$

Unfortunately, the opposite direction holds true only for those \mathbf{L} -contexts $\langle X, Y, I \rangle$ whose set $\text{Ext}^{\uparrow\downarrow}(X, Y, I)$ (resp. $\text{Int}^{\uparrow\downarrow}(X, Y, I)$) is a c -closure system [7]; i.e. an \mathbf{L} -closure system generated by a system of all a -complements of some $\mathcal{T} \subseteq L^X$.

Theorem 6. *If $\text{Ext}^{\uparrow\downarrow}(X_1, Y_1, I_1)$ is a c -closure system, the strong antitone \mathbf{L} -bonds from $\langle X_1, Y_1, I_1 \rangle$ to $\langle X_2, Y_2, I_2 \rangle$ are exactly a -bonds from $\langle X_1, Y_1 \times L, \neg I_1 \rangle$ to $\langle X_2, Y_2, I_2 \rangle$. If $\text{Int}^{\uparrow\downarrow}(X_2, Y_2, I_2)$ is a c -closure system, the antitone \mathbf{L} -bonds from $\langle X_1, Y_1, I_1 \rangle$ to $\langle X_2, Y_2, I_2 \rangle$ are exactly c -bonds from $\langle X_1, Y_1, I_1 \rangle$ to $\langle X_2 \times L, Y_2, (\neg I_2^{\text{T}})^{\text{T}} \rangle$.*

We omit further details due to the lack of space.

4 Conclusions and Further Research

We studied bonds between fuzzy contexts related to mutually different types of concept-forming operators and their relationship to antitone fuzzy bonds.

Our future research includes:

- Covering the \mathbf{L} -bonds described above and isotone \mathbf{L} -bonds in [15] by a general framework. The isotone and antitone concept-forming operators are one type of operators in [6,5]; also in [18].
- Generalizing the described theory to bond \mathbf{L} -contexts which each use different residuated lattice as the structure of truth-degrees. Results described in [17] seem to be promising for this goal.

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