

# Using Agents for Dynamic Components Redeployment and Replication in Distributed Systems

Nadim Obeid and Samih Al-Areqi

**Abstract.** Availability is one of the important criteria that affect the usefulness and efficiency of a distributed system. It mainly depends on how the components are deployed on the available hosts. In this paper, we present a generic agent-based monitor approach that supports the dynamic component redeployment and replication mechanisms which were presented in Avala and E-Avala. Avala and E-Avala were proposed to improve availability in large and distributed component-based systems via redeployment and replication. By reifying the interaction between the system and components, agents can detect when it is necessary to change the configuration and whether redeployment or replication is more appropriate.

**Keywords:** Distributed Systems, Agents, Availability, Redeployment.

## 1 Introduction

Distributed Systems (DS) have to face the problem of disconnected operations. In addition to the fact that the initial deployment architecture may not be very suitable, it is difficult to predict, at design time, the applications which the DS has to deal with. Therefore, finding and maintaining a desirable (e.g. availability) deployment architecture that satisfies a given set of constraints is a challenging problem. This is due to the facts that (1) there are many parameters which influence the selection of an appropriate deployment architecture (2) the space of possible architectures is large and (3) there may be constant need to change locations of components to meet changing requirements. This leads to some problems such as availability, dependency management [15]], and dynamic configuration. Hence,

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mechanisms such as components replication and redeployment may be necessary in order to improve availability and reliability [1, 2, 10, 12, 5].

In this paper, we present a generic agent-based monitor approach that supports the dynamic component redeployment and replication mechanisms which were presented in Avala [12] and E-Avala [3]. E-Avala improves on Avala by (1) considering positive and negative dependencies among components and (2) implementing replication taking into consideration negative dependencies. By reifying the interaction between the system and application components, agents can detect when it is necessary to change the configuration and whether redeployment or replication is more appropriate.

In section 2 we discuss Avala and E-Avala. In section 3 we present the agent-based redeployment approach. In Section 4 we discuss previous approaches to replication and redeployment.

## 2 Avala and E-Avala

In this section we give a brief presentation of Avala [6] and E-Avala [2]. Let  $h_1, h_2, \dots, h_k$  ( $1 \leq k$ ) stands for hosts,  $MEM(h_i)$  be the memory of  $h_i$ .  $C_1, \dots, C_n$  ( $1 \leq n$ ) stands for components,  $MEM(C_i)$  be the memory of  $C_i$  and  $FREQ(C_i, C_j)$  be the frequency between components  $C_i$  and  $C_j$ . The Avala algorithm [6] starts by ranking all hardware nodes and software components as follows:

$$IHR_i = \sum_{j=1}^k REL(h_i, h_j) + MEM(h_i) \quad (1)$$

The ranking of software components is performed as follows:

$$ICR_i = d * \sum_{j=1}^n FREQ(C_i, C_j) + \frac{E}{MEM(C_i)} \quad (2)$$

Where  $d$  denotes contributions of host memory and  $E$  contributions of event size of interactions between  $C_i$  and  $C_j$ .

The next software component to be assigned to  $h$ , is the one with the smallest memory requirement and which would contribute maximally to the availability function if placed on  $h$ . The Component Rank (CR) is calculated as follows:

$$CR(C_i, h) = D_1(C_i, h, n) + D_2(C_i, h) \quad (3)$$

where  $D_1(C_i, h, n) = d * \sum_{j=1}^n FREQ(C_i, MC_j) * REL(h, f(MC_j))$

and  $D_2(C_i, h, n) = \frac{E}{MEM(C_i)}$

and  $MC_j$  is a shorthand for mapped  $C_j$ ,  $f(MC_j)$  is a function that determines the hosts of mapped components,  $REL(h, f(MC_j))$  is a function that determines the reliability between selected host  $h$ , and hosts of mapped components.

Host Rank (HR) is calculated as follows:

$$HR(h_i) = \sum_{j=1}^m REL(h_i, MH_j) + MEM(h_i) \quad (4)$$

where  $m$  is number of hosts that are already selected.

E-Avala [2] employs the notion  $\text{Depend}(C_i, C_j)$ , not present in Avala, as follows:

$$\text{Depend}(C_i, C_j) = \begin{cases} 1 & \text{if } C_i \text{ depends on } C_j \\ -1 & \text{if } C_j \text{ depends on } C_i \end{cases} \quad (5)$$

Furthermore, E-Avala takes into consideration whether or not is a need for data consistency check regarding a  $C_i$  as shown below in (6):

$$\text{Consis}(C_i) = \begin{cases} 1 & \text{if } C_i \text{ does requires data consistency} \\ 0 & \text{Otherwise} \end{cases} \quad (6)$$

Let  $h$  be the selected host,  $l$  is the level of dependency for system configuration, determined by the designer, and  $nm$  be number of mapped components (i.e., already been assigned to selected hosts), E-Avala uses the same equations of Avala to calculate the intial ranking and distribution. It improves on Avala by employing two additional functions: RCR (resp. Consis-RCR) that compute Replicate Component Rank without (resp. with) consideration for data consistency.

$$\text{RCR}(C_i, h, n, nm) = D_3(C_i, h, n) + D_1((C_i, h, nm)) \quad (7)$$

where 
$$D_3(C_i, h, n) = \sum_{p=1}^n \text{Depend}(C_p, C_i) + \frac{l + 2E}{MEM(C_i)}$$

$$\text{Consis-RCR}(C_i, h, n, nm) = D_3(C_i, h, n) * (1 - \text{Consis}(C_i)) + D_1(C_i, h, nm) \quad (8)$$

E-Avala makes a comparison between the selected components for redeployment determined by CR (cf. (3)) and those to be replicated determined by RCR (cf. (5)). The selected component will be the one with the highest value of CR and RCR and that satisfies the constraints of memory, *Loc*, and *Colloc* with respect to the current host  $h$  and components which are already assigned. This process is be repeated until  $h$  is saturated. The performance of Avala and E-Avala is discussed in [2].

### 3 Agent-Based Redeployment

Agents are specialized autonomous problem solving entities that are suitable for problem solving in DS [6, 8, 9, 10, 11, 12]. The use of agents enables us (1) to keep track of the communication cost, (2) to mange dynamic reconfiguration while the system is operational and (3) to choose the better mechanism (e.g., re-deployment or replication) to maintain availability at minimal cost.

Let  $H_R$  (resp.  $H_T$ ) stands for the host of the requesting component,  $C_R$ , (resp. target component  $C_T$ ). We employ two kinds of Agents: (1) Comp-Agent (CPA), which has the required information about its host's components and has the ability to monitor any frequent interactions between a component on its host and components on other hosts and (2) Comm-Agent (CMA), which manages the communications with the other Host's CMAs. Let  $\text{Cor}(C_R, C_T)$  stand for the cost of request between  $C_R$  and  $C_T$ . When  $\text{Cor}(C_R, C_T)$  becomes high (e.g., above a certain threshold), CPA of  $H_T$  will negotiate with the CPA of  $H_R$  (through CMAs of  $H_T$  and  $H_R$ ) in order to agree on one of the following options : (1) redeploying  $C_R$  in

$H_T$ , (2) redeploying  $C_T$  in  $H_R$ , (3) replicating  $C_R$  in  $H_T$ , (4) replicating  $C_T$  in  $H_R$  or (5) no change. Assuming  $H_R \neq H_T$ ,  $Cor(C_R, C_T)$  can be defined as follows:

$$Cor(C_R, C_T) = \text{freq}(C_R, C_T) * \text{eventsize}(C_R, C_T) / \text{reliable}(H_R, H_T) \quad (9)$$

where  $\text{freq}(C_R, C_T)$  represents the frequency of interaction between  $C_R$  and  $C_T$ ,  $\text{eventsize}(C_R, C_T)$  denotes the size of interactions between  $C_R$  and  $C_T$ ,  $\text{reliable}(H_R, H_T)$  is the reliability between  $H_R$  and  $H_T$ . Fig. 1 and Fig. 2 show the agents and negotiations algorithms.

```

Agent_algorithm (hosts, comps)
{
  A= Component Agent of component CT;
  B= Component Agent of component CR;

  If ( Cor(CR,CT)=high))
  If (Agent Negotiation (A) ==false)
  {
    Agent Negotiation (A) =true
    Result=Negotiation (A, B)
    If (Result==result1)
      Redeployment (CT)
    If (Result==result2)
      Replication (CR)
    If (Result==result3)
      Redeployment (CT)
    If (Result==result4)
      Replication (CT)
    If (Result==result5)
      None
  }
}

```

**Fig. 1** Agent Algorithm

```

Result=Negotiation ( Agent A, Agent B)
{
  Exchange information (A,B)
  If (Redeployment (CT) ==true)
    Result 1=availability (Redeployment (CT))
  If (Replication (CR) =true)
    Result 2=availability (Replication (CR))
  If (Redeployment (CT)=true)
    Result 3=availability Redeployment (CT)
  If (Replication (CT) =true)
    Result 4=availability (Replication (CT))
  Else
    Result 5=current availability
  Result=max (result1, result2, result3, result4)
  Return result;
}

```

**Fig. 2** Agent Negotiation

We have made some improvement on the **DeSi** simulator [6] in order to simulate interactions between any two components on different hosts. We generate a deployment architecture that consists of 10 components, 3 hosts with their software agents and with  $\text{availability}=.8122$  distributed as follows:

$\text{Host0} = \{0,4,7\}$ ,  $\text{Host1} = \{2,6,3,8\}$  and  $\text{Host3} = \{0,5,1,9\}$

the input value are as in Table 1.

**Table 1** Input Values

Input Parameter	Value	Input Parameter	Value
Number of Component	10	Min host reliability	0
No. of hosts	3	Max host reliability	1
Min comp memory (in KB)	2	Min comp event size (in KB)	.01
Max comp memory (in KB)	8	Max comp event size (in KB)	10
Min host memory (in KB)	15	Min host bandwidth (in KB/S)	30
Max host memory (in KB)	30	Max host bandwidth (in KB/S)	100
Min comp frequency (in events/s)	0	Level of dependency	3
Max comp frequency (in events/s)	10		

Let  $C_i^R$  where  $0 \leq i \leq 9$  and  $C_j^T$  where  $0 \leq j \leq 9$  be two operating components and let  $Rep(i)$  (resp.  $Red(i)$ ) denote replicating  $C_i^R$  (resp. redeploying) on host of  $C_j^T$ . To test the viability of the algorithms, we execute several scenarios.

The values, as generated by the simulator (of an E-deployment architecture), which effect the agents' negotiation results, are shown below: frequency values between components and their memory size in Fig. 3, dependency values in Fig. 4 and reliability values between hosts and their memory size in Table 2.

We now consider two Scenarios. In the first,  $C_4$  makes requests frequently to  $C_6$  (cf. Table 3). The result (cf. Fig. 5) is to replicate  $C_4$  as there are many components dependent on it, and it provides better availability. We could not replicate  $C_6$  because there is a need for data consistency and it depends on two components in its host. In the second,  $C_3$  makes frequents requests to  $C_9$  (Table 4). The result (cf. Fig. 6) is that either mechanism is possible. Redeploying  $C_3$  will improve availability because it has more interaction and both positive and negation dependency relations with components in the host of  $C_9$ .

0	1	2	3	4	5	6	7	8	9
0.0	5.682	0.0	5.479	5.660	0.0	4.776	0.0	0.182	0.327
5.682	0.0	1.179	6.086	0.0	2.971	0.007	3.969	4.365	0.0
0.0	1.179	0.0	9.802	3.060	0.0	5.494	2.090	0.177	0.299
5.479	6.086	9.802	0.0	6.510	0.619	9.443	0.0	8.226	0.578
5.660	0.0	3.060	6.510	0.0	5.306	2.059	2.932	5.883	4.236
0.0	2.971	0.0	0.619	5.306	0.0	1.465	0.0	1.015	0.0
4.776	0.007	5.494	9.443	2.059	1.465	0.0	1.015	2.352	0.0
0.0	3.969	2.090	0.0	2.932	0.0	0.0	0.0	0.0	0.0
0.182	4.365	0.177	8.226	5.883	4.613	2.352	0.0	0.0	6.085
0.327	0.0	0.299	0.578	4.236	0.0	0.0	0.0	6.085	0.0
4.28	6.12	3.97	4.80	7.66	4.24	3.87	4.72	6.69	6.79

Fig. 3 Components Frequency

0	1	2	3	4	5	6	7	8	9
0.0	5.662	0.0	5.479	5.660	0.0	4.776	0.0	0.0	0.0
-5.66	0.0	0.0	6.086	0.0	0.0	0.0	0.0	4.365	0.0
0.0	0.0	0.0	9.802	0.0	0.0	5.494	0.0	0.0	0.0
-5.66	-5.90	-9.80	0.0	6.510	0.0	9.443	0.0	8.226	0.578
0.0	0.0	0.0	-6.51	0.0	5.306	0.0	0.0	5.883	4.236
0.0	0.0	0.0	0.0	-5.30	0.0	0.0	0.0	4.613	0.0
-4.77	0.0	-5.49	-9.44	0.0	0.0	0.0	0.0	0.0	0.0
0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0
0.0	-4.36	0.0	-8.23	-5.89	-4.61	0.0	0.0	0.0	6.085
0.0	0.0	0.0	-0.57	-4.23	0.0	0.0	0.0	0.0	0.0

Fig. 4 Component Dependency

Table 2 Host Reliability/Memory

Host No.	0	1	2
0	1	.49	.38
1	.49	1	.94
3	.38	.94	1
Host MEM	21	27	22

Table 3 Component properties

Comp. properties	Comp (4)	Comp (6)
Comp. memory size	7,6kb	3,8
Free host size	4,5 kb	12.5 kb
Positive dependency	→0,3	→0,2,3
Data Consistency	0	1

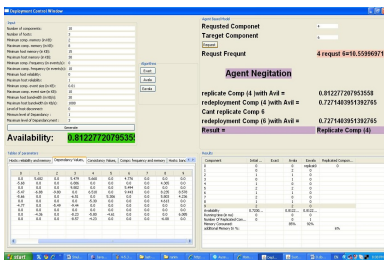


Fig. 5 Senario 1 Results

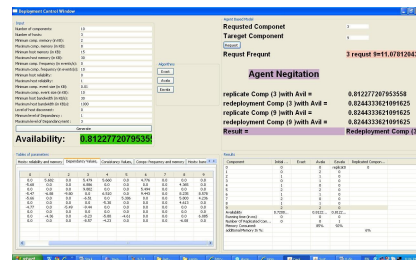


Fig. 6 Senario 2 Results

**Table 4** Scenario 1 Component properties

<b>Component Properties</b>	<b>Component (3)</b>	<b>Component (9)</b>
Component memory size	4.8 kb	6.8
Free Host size	4.5kb	12.5kb
Positive dependency	→0,1,2	→4,5,8
Negative dependency	4,6,8,9→	
Data consistency	0	0

## 4 Previous Work and Concluding Remark

Several approaches that support the replication of components in DSs have been proposed. However, only a few address redeployment. In [4], Dock is proposed. It employs mobile agents to perform deployment tasks among hosts. It differs from our approach in that it is more concerned with the practical issues of implementing deployment rather than extracting parameters and evaluating deployment architectures. In [5], a constraint-based deployment approach is presented. It addresses the deployment of hierarchical components on heterogeneous dynamic networks. In [3], MARP employs mobile agents to coordinate the updates made to replications maintained at different servers to ensure consistency.

In this paper, we present a generic agent-based monitor approach that supports the dynamic component redeployment and replication mechanisms which were presented in Avala [6] and improved in E-Avala [2]. Some of the issues that need to be addressed include: (1) dealing with functional consistency among components, (2) expanding the solution to include additional parameters such as components structure representation.

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