

THE KUHN-TUCKER THEOREM IN CONCAVE PROGRAMMING

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1. Introduction

In order to solve problems of constrained extrema, it is customary in the calculus to use the method of the Lagrangian multiplier. Let us, for example, consider a problem: maximize $f(x_1, \dots, x_n)$ subject to the restrictions $g_k(x_1, \dots, x_n) = 0$ ($k = 1, \dots, m$). First, formulate the so-called Lagrangian form

$$\varphi(x, y) = f(x_1, \dots, x_n) + \sum_{k=1}^m y_k g_k(x_1, \dots, x_n)$$

where unknown y_1, \dots, y_m are called the Lagrangian multipliers. Then solutions x_1, \dots, x_n are found among extreme points of $\varphi(x, y)$, with unrestricted x and y , which in turn are characterized as the solutions of

$$\varphi_{x_i}(x, y) = \frac{\partial f}{\partial x_i} + \sum_k y_k \frac{\partial g_k}{\partial x_i} = 0 \quad (i = 1, \dots, n),$$

$$\varphi_{y_k}(x, y) = g_k(x_1, \dots, x_n) = 0 \quad (k = 1, \dots, m).$$

This method, although not necessarily true without certain qualifications, has been found to be useful in many particular problems of constrained extrema.

The method is with a suitable modification applied to solve the programming problems also where we are concerned with maximizing a function $f(x_1, \dots, x_n)$ subject to the restrictions $x_i \geq 0$ ($i = 1, \dots, n$) and $g_k(x_1, \dots, x_n) \geq 0$ ($k = 1, \dots, m$). Kuhn and Tucker [2] first proved that under some qualifications, concave programming is reduced to finding a saddle-point of the Lagrangian form $\varphi(x, y)$. This Kuhn-Tucker Theorem was further elaborated by Arrow and Hurwicz [1] so that non-concave programming may be handled. In the present chapter we shall, under different qualifications, prove the Kuhn-Tucker Theorem for concave programming.

2. Maximum Problem and Saddle-Point Problem

Let $g(x) = \langle g_1(x), \dots, g_m(x) \rangle$ (see fn. 1, p. 2) be an m -dimensional vector-valued function and $f(x)$ be a real-valued function, both defined for non-negative vectors $x = \langle x_1, \dots, x_n \rangle$.

Consider the following

MAXIMUM PROBLEM. Find a vector \bar{x} that maximizes

$$(1) \quad f(x)$$

subject to the restrictions

$$(2) \quad x \geq 0, \quad g(x) \geq 0.$$

A vector x will be called *feasible* if it satisfies (2), and a feasible vector \bar{x} maximizing $f(x)$ subject to (2) will be called an *optimum* vector of, or a *solution* to, the problem.

Associated with the Maximum Problem, the *Lagrangian form* is defined by

$$(3) \quad \varphi(x, y) = f(x) + y \cdot g(x),$$

where¹

$$x = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \geq 0 \text{ and } y = \begin{pmatrix} y_1 \\ \vdots \\ y_m \end{pmatrix} \geq 0.$$

A pair of vectors (\bar{x}, \bar{y}) is called a *saddle-point* of $\varphi(x, y)$ in $x \geq 0, y \geq 0$, if

$$(4) \quad \bar{x} \geq 0, \quad \bar{y} \geq 0,$$

$$(5) \quad \varphi(x, \bar{y}) \leq \varphi(\bar{x}, \bar{y}) \leq \varphi(\bar{x}, y) \text{ for all } x \geq 0 \text{ and } y \geq 0,$$

which may be written as follows:

$$(6) \quad \varphi(\bar{x}, \bar{y}) = \min_{y \geq 0} \max_{x \geq 0} \varphi(x, y) = \max_{x \geq 0} \min_{y \geq 0} \varphi(x, y).$$

SADDLE-POINT PROBLEM. Find a saddle-point (\bar{x}, \bar{y}) of $\varphi(x, y) = f(x) + y \cdot g(x)$.

3. Saddle-Point Implies the Optimality

We are interested in the reduction of a maximum problem to the saddle-point problem of the associated Lagrangian form. First, a proposition will be noted which is true without any qualification on f and g , whenever there exists a saddle-point.

THEOREM 1. If (\bar{x}, \bar{y}) is a saddle-point of $\varphi(x, y)$ in $x \geq 0, y \geq 0$, then \bar{x} is an optimum vector of the maximum problem.

PROOF. Substituting (3) into (5), we have

$$(7) \quad f(x) + \bar{y} \cdot g(x) \leq f(\bar{x}) + \bar{y} \cdot g(\bar{x}) \leq f(\bar{x}) + y \cdot g(\bar{x})$$

for all $x \geq 0, y \geq 0$.

¹ For two vectors $x = \langle x_1, \dots, x_n \rangle$ and $u = \langle u_1, \dots, u_n \rangle$, we shall, as usual, define

$$x \geq u \text{ if } x_i \geq u_i, \quad (i = 1, \dots, n),$$

$$x \geq u \text{ if } x \geq u \text{ and } x \neq u,$$

$$x > u \text{ if } x_i > u_i, \quad (i = 1, \dots, n),$$

and $x \cdot u$ stands for the inner product

$$x \cdot u = \sum_{i=1}^n x_i u_i.$$

Since the right-hand inequality holds for any $y \geq 0$, it follows that $g(\bar{x})$ cannot have a negative component, and $\bar{y} \cdot g(\bar{x})$ must be zero:

$$g(\bar{x}) \geq 0, \quad \bar{y} \cdot g(\bar{x}) = 0.$$

Thus the left-hand inequality of (7) may be written as

$$(8) \quad f(x) + \bar{y} \cdot g(x) \leq f(\bar{x}) \quad \text{for all } x \geq 0.$$

Since, for any feasible vector x we have $\bar{y} \cdot g(x) \geq 0$, it follows that $f(x) \leq f(x) + \bar{y} \cdot g(x) \leq f(\bar{x})$, which shows that \bar{x} is optimum, q.e.d.

4. The Kuhn-Tucker Theorem

Now a question naturally arises whether, given an optimum vector \bar{x} , it is possible to find a vector \bar{y} for which (\bar{x}, \bar{y}) is a saddle-point of $\varphi(x, y)$. This, of course, is not true in general, e.g., for convex programming (i.e., where the maximand is a convex function). The following simple example shows that it does not hold even for concave programming:

$$f(x) = x, \quad g(x) = -x^2.$$

Regarding concave programming, however, the reduction is shown to be possible provided f and g satisfy certain regularity conditions, e.g., the Kuhn-Tucker Constraint Qualification.² We shall give sufficient conditions which make the reduction possible.

THEOREM 2. *Suppose that $f(x)$ and $g(x)$ are concave functions on $x \geq 0$, and $g(x)$ satisfies the following condition (due to M. Slater [3]):³*

(9) *There exists an $x^0 \geq 0$ such that $g(x^0) > 0$.*

Then a vector \bar{x} is optimum if, and only if, there is a vector $\bar{y} \geq 0$ such that (\bar{x}, \bar{y}) is a saddle-point of $\varphi(x, y)$.

PROOF. Let \bar{x} be optimum. We shall, in the $(m+1)$ -dimensional vector space, define A and B by

$$A = \left\{ \begin{pmatrix} z_0 \\ z \end{pmatrix}; \begin{pmatrix} z_0 \\ z \end{pmatrix} \leq \begin{pmatrix} f(x) \\ g(x) \end{pmatrix} \text{ for some } x \geq 0 \right\},$$

$$B = \left\{ \begin{pmatrix} z_0 \\ z \end{pmatrix}; \begin{pmatrix} z_0 \\ z \end{pmatrix} > \begin{pmatrix} f(\bar{x}) \\ 0 \end{pmatrix} \right\}.$$

Since $f(x)$ and $g(x)$ are concave, the set A is convex. Since \bar{x} is opti-

² See Kuhn and Tucker [2], p. 483.

³ A seemingly weaker condition: (9') For any $u > 0$, there exists a vector $x \geq 0$ such that $u \cdot g(x) > 0$ is due to S. Karlin. The condition, however, is equivalent to the Slater's condition (9). For the proof, see Chapter 5, pp. 109-10 of the present volume.

mum, A and B have no vector in common. Therefore, by the lemma on the separation of convex sets, there is a non-zero vector $\langle v_0, v \rangle \neq 0$ such that

$$(10) \quad v_0 z_0 + v \cdot z \leq v_0 u_0 + v \cdot u \quad \text{for all } \begin{pmatrix} z_0 \\ z \end{pmatrix} \in A, \begin{pmatrix} u_0 \\ u \end{pmatrix} \in B .$$

By the definition of B , (10) implies $\langle v_0, v \rangle \geq 0$. Since $\langle f(\bar{x}), 0 \rangle$ is on the boundary of B , we also have, by the definition of A ,

$$(11) \quad v_0 f(x) + v \cdot g(x) \leq v_0 f(\bar{x}) \quad \text{for all } x \geq 0 .$$

We have $v_0 > 0$. Otherwise, we have $v \geq 0$ and $v \cdot g(x) \leq 0$ for all $x \geq 0$, which contradicts (9).

Let $\bar{y} = v/v_0$. Then

$$(12) \quad \bar{y} \geq 0 ,$$

$$(13) \quad f(x) + \bar{y} \cdot g(x) \leq f(\bar{x}) \quad \text{for all } x \geq 0 .$$

Putting $x = \bar{x}$ in (13), we have $\bar{y} \cdot g(\bar{x}) \leq 0$. On the other hand, we have

$$(14) \quad g(\bar{x}) \geq 0 .$$

Hence

$$(15) \quad \bar{y} \cdot g(\bar{x}) = 0 .$$

Relations (13), (14), and (15) show that (\bar{x}, \bar{y}) is a saddle-point of $\varphi(x, y)$ in $x \geq 0, y \geq 0$, q.e.d.

5. A Modification of the Kuhn-Tucker Theorem

Slater's condition (9), however, excludes the case where part of the second half of restriction (2) is

$$h(x) \geq 0 \quad \text{and} \quad -h(x) \geq 0 ,$$

for linear $h(x)$. In order to make the reduction possible for such cases, we have to modify the Kuhn-Tucker Theorem.

Let sub-sets I and II of $\{1, \dots, m\}$ be defined by

$$(16) \quad I = \{k; g_k(x) = 0 \quad \text{for all feasible } x\} ,$$

and

$$(17) \quad II = \{1, \dots, m\} - I .$$

We shall assume that

$$(18) \quad g_k(x) \text{ is linear in } x, \text{ for } k \in I .$$

(19) For any i , there is a feasible vector x^i such that

$$x_i^i > 0 .$$

Then we have as a modification of Theorem 2 the following :

THEOREM 3. *Suppose that $f(x)$, $g(x)$ are concave, and $g(x)$ satisfies (18) and (19). Then a vector \bar{x} is optimum if, and only if, there is a vector \bar{y} such that (\bar{x}, \bar{y}) is a saddle-point of $\varphi(x, y)$ in $x \geq 0$ and $y \geq 0$ (II).⁴*

PROOF. It is obvious that, if (\bar{x}, \bar{y}) is a saddle-point of $\varphi(x, y)$ in $x \geq 0$ and $y \geq 0$ (II), then \bar{x} is optimum.

In order to prove the converse we may assume that

$$(20) \quad \frac{dg_k}{dx}, \quad k \in I, \text{ are linear independent.}$$

Let \bar{x} be optimum. We consider two sets A and B defined by

$$A = \left\{ \begin{pmatrix} z_0 \\ z \\ u \end{pmatrix}; \quad z_0 \leq f(x), \quad z_I = g_I(x), \quad z_{II} \leq g_{II}(x), \quad u \leq x, \quad \text{for some } x \right\},$$

$$B = \left\{ \begin{pmatrix} z_0 \\ z \\ 0 \end{pmatrix}; \quad z_0 > f(\bar{x}), \quad z = 0 \text{ (I)}, \quad z > 0 \text{ (II)} \right\}.$$

Then A and B are convex, and have no point in common. Therefore, there is a vector $\langle v_0, v, w \rangle \neq 0$ such that

$$(21) \quad v_0 \geq 0, \quad v \geq 0 \text{ (II)}, \quad w \geq 0,$$

and

$$(22) \quad v_0 f(x) + v \cdot g(x) + w \cdot x \leq v_0 f(\bar{x}) \quad \text{for all } x.$$

It now suffices to prove that $v_0 > 0$. If we had assumed that $v_0 = 0$, then

$$(23) \quad v \cdot g(x) + w \cdot x \leq 0 \quad \text{for all } x.$$

For any $k \in II$, there is a feasible vector x^k such that $g_k(x^k) > 0$. Hence

$$(24) \quad v_k = 0 \quad \text{for } k \in II.$$

By (19) and (23),

$$(25) \quad w = 0.$$

Using (24) and (25), the inequality (23) may be written as follows :

$$(26) \quad v_I \cdot g_I(x) \leq 0 \quad \text{for all } x, \quad \text{where } v_I \neq 0.$$

Since we have assumed that $g_I(x)$ is linear, (26) implies

⁴ The notation $y \geq 0$ (II) means $y_k \geq 0$ for all $k \in II$. The point (\bar{x}, \bar{y}) is said to be a saddle-point of $\varphi(x, y)$ in $x \geq 0$ and $y \geq 0$ (II) if $\bar{x} \geq 0$, $\bar{y} \geq 0$ (II) and (5) holds for any $x \geq 0$ and $y \geq 0$ (II).

$$(27) \quad v_i \cdot \frac{\partial g_i}{\partial x} = 0, \quad \text{where } v_i \neq 0,$$

which contradicts (20), q.e.d.

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